



Searching(1)

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outline

- ▣ Searching
- ▣ Searching on unsorted arrays
- ▣ Searching on sorted arrays
- ▣ Binary search

Searching

- ❑ **Search can be viewed abstractly as a process to determine if an element with a particular value is a member of a particular set.**
- ❑ **The more common view of searching is an attempt to find the record within a collection of records that has a particular key value, or those records in a collection whose key values meet some criterion such as falling within a range of values.**

Searching: formal definition

- Suppose that we have a collection L of n records of the form

$(k_1, I_1), (k_2, I_2), \dots, (k_n, I_n)$

where I_j is information associated with key k_j from record j for $1 \leq j \leq n$. Given a particular key value K , the search problem is to locate a record (k_j, I_j) in L such that $k_j = K$ (if one exists). Searching is a systematic method for locating the record (or records) with key value $k_j = K$.

Searching

- A **successful search** is one in which a record with key $k_j = K$ is found.
- An **unsuccessful search** is one in which no record with $k_j = K$ is found (and no such record exists).
- An **exact-match query** is a search for the record whose key value matches a specified key value.
- A **range query** is a search for all records whose key value falls within a specified range of key values.

Searching Algorithms

- ❑ **We can categorize search algorithms into three general approaches:**
 - 1. Sequential and list methods.
 - 2. Direct access by key value (hashing).
 - 3. Tree indexing methods.

Searching on Unsorted Arrays

□ Sequential search algorithm

- the simplest form of search
- Sequential search on an unsorted list requires $\Theta(n)$ time in the worst case.

□ Basic algorithm:

Get the search criterion (**key**)

Get the first record from the file

While (**(record != key) and (still more records)**)

 Get the next record

End_while

□ When do we know that there wasn't a record in the file that matched the key?



Sequential Search on unsorted arrays: implementation

```
Int SeqSearch(int A[],int n,int K)
{  int i=n;
   A[0]=K;
   while(A[i]!=K) i--;
   return i;
}
```


Searching on sorted Arrays

□ Basic algorithm:

Get the search criterion (key)

Get the first record from the file

While ((record < key) and (still more records))

 Get the next record

End_while

If (record = key)

 Then success

 Else there is no match in the file

End_else

□ When do we know that there wasn't a record in the file that matched the key?

Unsorted vs. Sorted

- ❑ **Observation:** the search is faster on an sorted list only when the item being searched for is not in the list.
- ❑ **Also, keep in mind that the list has to first be placed in order for the ordered search.**
- ❑ **Conclusion:** the efficiency of these algorithms is roughly the same.
- ❑ **So, if we need a faster search, we need a completely different algorithm.**
- ❑ **How else could we search an ordered file?**

Binary search

- ❑ **If we have an ordered list and we know how many things are in the list (i.e., number of records in a file), we can use a different strategy.**
- ❑ **The binary search gets its name because the algorithm continually divides the list into two parts.**

How a Binary Search Works



Always look at the center value. Each time you get to discard half of the remaining list.

Is this fast ?

Binary search: Implementation

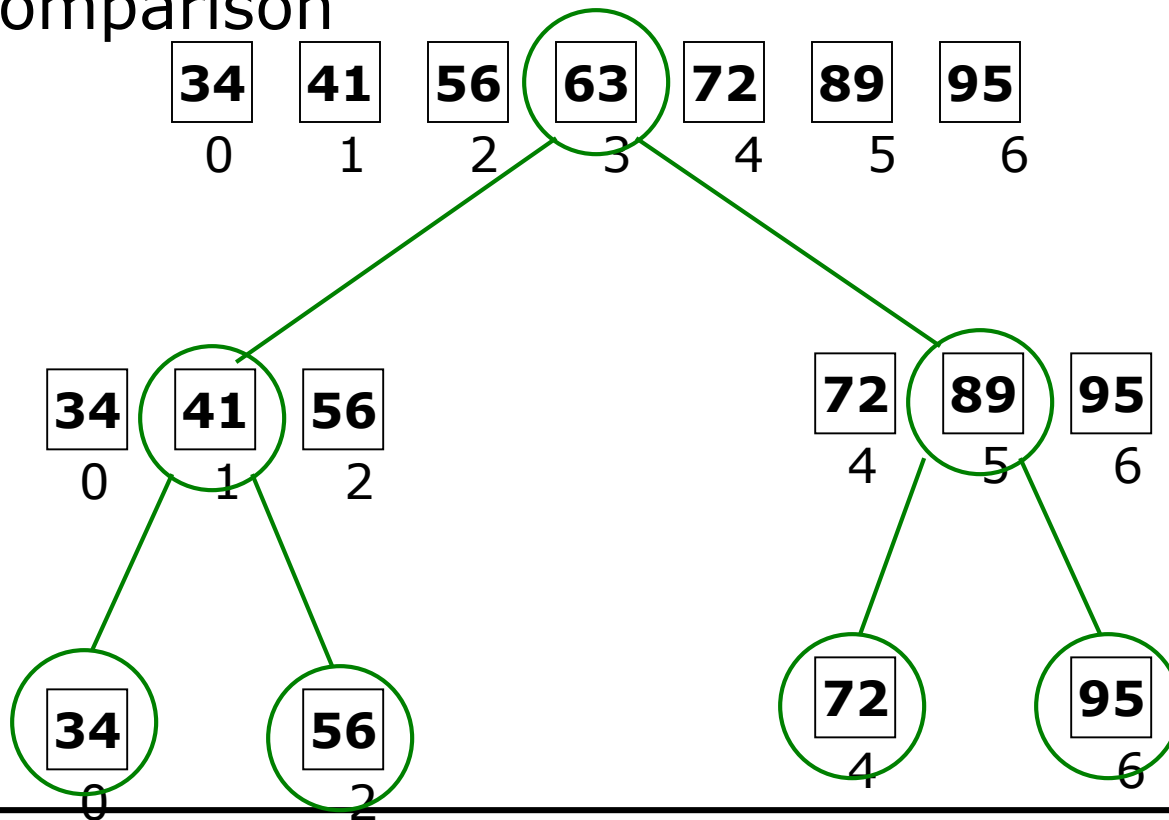
```
// Return the position of an element in sorted array "A" of
// size "n" with value "K". If "K" is not in "A", return
// the value "n".
int binary(int A[], int n, int K) {
    int l = -1;
    int r = n;           // l and r are beyond array bounds
    while (l+1 != r) {    // Stop when l and r meet
        int i = (l+r)/2;  // Check middle of remaining subarray
        if (K < A[i]) r = i;    // In left half
        if (K == A[i]) return i; // Found it
        if (K > A[i]) l = i;    // In right half
    }
    return n; // Search value not in A
}
```

Binary search: non-recursive Implementation

```
int BinSearch1(int A[ ], int n, int k)
{
    low=1; high=n;
    while (low<=high)
    {
        mid=(low+high)/2;
        if (k<A[mid]) high=mid-1;
        else if (k>A[mid]) low=mid+1;
        else return mid;
    }
    return 0;
}
```

Binary Search Tree

Binary Search algorithm of an array of *sorted* items reduces the search space by one half after each comparison



How Fast is a Binary Search?

- ❑ **Worst case: 11 items in the list took 4 tries**
- ❑ **How about the worst case for a list with 32 items ?**
 - 1st try - list has 16 items
 - 2nd try - list has 8 items
 - 3rd try - list has 4 items
 - 4th try - list has 2 items
 - 5th try - list has 1 item

How Fast is a Binary Search? (con't)

List has 250 items

1st try - 125 items

2nd try - 63 items

3rd try - 32 items

4th try - 16 items

5th try - 8 items

6th try - 4 items

7th try - 2 items

8th try - 1 item

List has 512 items

1st try - 256 items

2nd try - 128 items

3rd try - 64 items

4th try - 32 items

5th try - 16 items

6th try - 8 items

7th try - 4 items

8th try - 2 items

9th try - 1 item



What's the Pattern?

- ❑ List of 11 took 4 tries
 - ❑ List of 32 took 5 tries
 - ❑ List of 250 took 8 tries
 - ❑ List of 512 took 9 tries
-
- ❑ $32 = 2^5$ and $512 = 2^9$
 - ❑ $8 < 11 < 16$ $2^3 < 11 < 2^4$
 - ❑ $128 < 250 < 256$ $2^7 < 250 < 2^8$



A Very Fast Algorithm!

- How long (worst case) will it take to find an item in a list 30,000 items long?

$$2^{10} = 1024$$

$$2^{11} = 2048$$

$$2^{12} = 4096$$

$$2^{13} = 8192$$

$$2^{14} = 16384$$

$$2^{15} = 32768$$

- So, it will take only 15 tries!
- Time complexity: $\Theta(\log n)$



References

- **Data Structures and Algorithm Analysis Edition 3.2 (C++ Version)**
 - P.57 Example 3.1
 - P.74-75

End

