

DM → + note for mid-term

Propositional Logic

A proposition is a declarative sentence.

* shakib all Hasan is a cricketer. → True

□ Question is not proposition

□ Order is not proposition

□ $n+1 \geq y$ → Predicate

Truth value $\begin{cases} \text{False} \\ \text{True} \end{cases}$

□ I am lying → Paradox

কখন মিথ্যা বলার মতি কিন্তু যা বলে তা সত্য বা মিথ্যা
কাজে $T \rightarrow F \rightarrow T \rightarrow F$ cause Paradox

□ conjunction → AND → \wedge ($P \wedge Q$)

Today is Friday (and) It is raining

P

$$2^n = 2^2 = 4$$

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

□ But সত্যলভ
conjunction

□ , - , ... কমা
সত্যলভ conjune

* conjunction to get value of not 20 20
that's why output is True
Like AND operation means AND

Disjunction \rightarrow OR $\rightarrow \vee (p \vee q)$

* At-least one proposition is True

Today is Friday (or) It is raining

p

q

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

* Let (p) shakib all hasan is a cricketers

(or) (q) shakib khon is a cricketers

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

☐ Negation $\rightarrow \neg p$

(p) Shakib Al Hasan is a footballer \rightarrow False

$\neg(p)$ Shakib Al Hasan is not a footballer \rightarrow True

p	$\neg p$
1	0
0	1

\rightarrow

p	$\neg p$
T	F
F	T

$$1+3 \neq 4$$

☐ The rules of writing a truth table

* Start to from input of right

* A T and a F will go down along \rightarrow (1)

The column like this.

* 2nd 2 T and 2 F will go down along \rightarrow (2)

The column like this

* 3rd 4 T and 4 F will go down along \rightarrow (3)

The column like this

p	q	r
T	T	T
T	T	F
T	F	T
T	F	F
F	T	T
F	T	F
F	F	T
F	F	F

Truth table

* input output $\rightarrow ((p \wedge a) \vee r)$

p	q	r	$p \wedge a$	$((p \wedge a) \vee r)$
T	T	T	T	T
T	T	F	T	T
T	F	T	F	T
T	F	F	F	F
F	T	T	F	T
F	T	F	F	F
F	F	T	F	T
F	F	F	F	F

(v) * $(p \vee \neg a) \wedge r$

Exchange \rightarrow or \rightarrow and

p	q	r	$(\neg a)$	$(p \vee \neg a)$	$(p \vee \neg a) \wedge r$
T	T	T	F	T	T
T	T	F	F	T	F
T	F	T	T	T	T
T	F	F	T	T	F
F	T	T	F	F	F
F	T	F	F	F	F
F	F	T	T	T	T
F	F	F	T	T	F

Exclusive OR $\rightarrow (\oplus)$

* * Can't do the same work, two or more person on things, same time.

* two wicket keeper can't play in the same match

* two goal keeper can't playing the same match

* Quinton de Kock or Heinrich Kullsen will play in today's match

or but not both
 $p \oplus q$

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

* Mushfiquur Rahim or Liton Das will play in today's match

p	s	$p \oplus s$
T	T	F
T	F	T
F	T	T
F	F	F

* Or... not both

"students who have taken calculus or computers science, but not both, can take the class"

P	q	$P \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

* coffee or tea comes with dinner.

* Experience with c++ or Java is required

* Lunch includes soup or salad

* you can pay using us. dollars or euros

2nd class

conditional state ment

* "If P then Q"

P: today is holiday

Q: The storey is closed

If today is holiday then the storey is closed

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

*** statement $P \rightarrow Q = F$
T F शरु गारल $P \rightarrow Q = F$

* If it rains, I will stay home.

we can also say "It is raining implies that I will stay home."

$P \rightarrow Q$

* If P then Q

* P implies Q

* Q if P

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

* If $\frac{1+1=2}{T}$ then $\frac{2+2=5}{F}$ F

If $\frac{1+1=3}{F}$ then $\frac{2+2=4}{T}$ T

If $\frac{1+1=3}{F}$ then dog can fly T

If $\frac{1+1=2}{T}$ then dog can fly F

* "If p, q"

"If it is below freezing, it is also snowing"

$P \rightarrow Q$

☐ (যদি) তাহলে (সufficient)
তাহলে (necessary)

☐ If p then q

$\Rightarrow q$ if p

* q unless $\neg p$

maria will ^{find} get a job good job unless she

does not learn discrete mathematics p \rightarrow q

\Rightarrow q unless $\neg p$ $\neg p$

Items of Implies

* If then

* when

* "q whenever p"

* "q is necessary for p"

* "q follow from p"

* "q when p"

* "q unless $\neg p$ "

** \rightarrow is sufficient

" only if " is necessary

sufficient \Rightarrow necessary
symbol of Implies

☐ sufficient condition

If you will pass only if you study.

P

Q

☐ শুধুমাত্র Q এর জন্যে only if আছে তাহলে
P এর জন্যে তা হবে নিশ্চয়

$$P \rightarrow Q$$

যায যগাটা sufficient, কিন্তু অপ্রাপ্য ক্রিয়া বলা
এর দ্বারা

☐ Necessary condition

you will pass only if you study

P

Q

$$P \rightarrow Q$$

P	Q	$P \rightarrow Q$
T	F	F
T	T	T
F	F	T
F	T	T

☐ Bioconditional statement ($\Rightarrow = \leftrightarrow$)

* (if and only if) হাফেল্ড Bio conditional

* you can take the flight if and only if you

buy a ticket \rightarrow necessary and sufficient

9

$$\Rightarrow p \leftrightarrow q$$

* p is necessary and sufficient for a

* 'p itt q"

- * "p iff q"
- * if p then q, and conversely

□ $\text{value} = \text{True}$ 23

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

$$1+1=2 \text{ only if } 2+2=4 = \text{F}$$

$1+2=4$ only if $2+2=4 = F$

$$1+2=3 \text{ only if } 2+2=5 = F$$
$$1 + 1 = 3 \text{ only if } 2 + 2 = 5 = T$$

☐ converse, contrapositive, Inverse

* converse $q \rightarrow p$

If it rains, I will stay home $p \rightarrow q$
 p q

If I will stay home, it is raining $q \rightarrow p$
 q p

* contrapositive. $\neg p \rightarrow \neg q$

If I do not stay home, it is not raining
 $\neg q$ $\neg p$

$\Rightarrow p \rightarrow q \rightarrow \neg p \rightarrow \neg q$

*** ☐. original and contrapositive are meaning same

☐ Inverse $\Rightarrow \neg p \rightarrow \neg q$

If it does not rain, I will not stay home
 $\neg p$ $\neg q$

$\Rightarrow \neg p \rightarrow \neg q$

p	q	$\neg p$	$\neg q$
T	T	F	F
T	F	F	T
F	T	T	F
F	F	T	T

Problem

Exercise

$p \equiv$ You get an A on the final exam.

$q \equiv$ You do every exercise.

$r \equiv$ You get an A in this class.

① $p \wedge \neg q$

② $p \wedge q \wedge r$

③ $r \rightarrow p$

□ proposition to English sentence.

Let p and q be the propositions

"swimming at the new Jersey is called allowed"
and "sharks have been spotted near the shore."

a/ $p \rightarrow q$ = If swimming at the new Jersey is allowed,
then sharks have been spotted near the shore.

b/ $p \leftrightarrow q$ = swimming at the new Jersey is allowed
if and only if sharks have been spotted
near the shore.

3rd clS
04-2-23

\rightarrow

Logical Equivalence.

$$P \rightarrow Q \equiv \neg P \vee Q$$

P	Q	$\neg P$	$P \rightarrow Q$	$\neg P \vee Q$
T	T	F	T	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T

$$(P \wedge Q) \rightarrow R \equiv (P \rightarrow R) \wedge (Q \rightarrow R)$$

P	Q	R	$P \wedge Q$	$(P \wedge Q) \rightarrow R$	$P \rightarrow R$	$Q \rightarrow R$	$(P \rightarrow R) \wedge (Q \rightarrow R)$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	F
T	F	T	F	T	T	T	T
T	F	F	F	T	F	T	F
F	T	T	F	T	T	T	T
F	T	F	F	T	T	F	F
F	F	T	F	T	T	T	T
F	F	F	F	T	T	T	T

not same.



Law

□ Identity

$$p \wedge T \equiv p$$

$$p \vee F \equiv p$$

□ Domination

$$p \vee T \equiv T$$

$$p \wedge F \equiv F$$

(v) □ Idempotent

$$p \wedge p \equiv p$$

$$p \vee p \equiv p$$

□ Double Negation

$$\neg(\neg p) \equiv p$$

□ Commutative

$$p \vee q \equiv q \vee p$$

$$p \wedge q \equiv q \wedge p$$

□ Associative

$$(p \vee q) \vee r \equiv p \vee (q \vee r)$$

$$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$$

□ Distributive

$$(p \wedge q) \vee (q \wedge s) \equiv (p \vee r) \wedge (p \vee s) \wedge (q \vee r) \wedge (q \vee s)$$

□ De Morgan's Law

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

☐ Negation

$$p \vee \neg p \equiv T$$

$$p \wedge \neg p = F$$

☐ Absorption

$$p \vee (p \wedge q) \equiv p$$

$$p \wedge (p \vee q) \equiv p$$

☐ p only if q

if p then q

$$p \rightarrow q$$

☐ p if q

if q then p

$$q \rightarrow p$$

Golden Law

☐ conditional

$$p \rightarrow q \equiv \neg p \vee q$$

☐ contrapositive

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

☐ Biconditional

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

*** Q. proved that $(p \wedge q) \rightarrow (p \vee q)$ is a tautology

p	q	$p \wedge q$	$p \vee q$	$(p \wedge q) \rightarrow (p \vee q)$
T	T	T	T	T
T	F	F	T	T
F	T	F	T	T
F	F	F	F	T

is a tautology

* ** tautology means is every element is true

$$\begin{aligned}
 \neg(p \rightarrow q) &\equiv p \wedge \neg q \\
 &\equiv \text{L.H.S} = \neg(p \rightarrow q) \\
 &\equiv \neg(\neg p \vee q) \\
 &\equiv \neg(\neg p) \wedge \neg q \\
 &\equiv p \wedge \neg q \\
 &\equiv \text{R.H.S}
 \end{aligned}$$

prove that

$(p \rightarrow q) \wedge (q \rightarrow p) \rightarrow (p \rightarrow r)$ is a tautology

$$\equiv \neg((p \rightarrow q) \wedge (q \rightarrow p)) \vee (p \rightarrow r)$$

$$\equiv \neg(p \rightarrow q) \vee \neg(q \rightarrow p) \vee (p \rightarrow r)$$

$$\equiv \neg(\neg p \vee q) \vee \neg(\neg q \vee p) \vee (\neg p \vee r)$$

$$\equiv (\neg(\neg p) \wedge \neg q) \vee (\neg(\neg q) \wedge \neg p) \vee \neg p \vee r$$

$$\equiv (p \wedge \neg q) \vee (q \wedge \neg p) \vee \neg p \vee r$$

$$\equiv (p \wedge \neg q) \vee \neg p \vee (q \wedge \neg p) \vee r$$

$$\equiv ((p \vee \neg p) \wedge (\neg q \vee \neg p)) \vee ((q \vee r) \wedge (\neg p \vee r))$$

$$\equiv (\top \wedge (\neg q \vee \neg p)) \vee ((q \vee r) \wedge \top)$$

$$\equiv (\neg q \vee \neg p) \vee (q \vee r)$$

$$\equiv \neg q \vee \neg p \vee q \vee r$$

$$\equiv \neg q \vee q \vee \neg p \vee r$$

$$\equiv \top \vee p \vee r$$

$$\equiv \top$$

\therefore it's a tautology

$$*** (p \wedge q) \rightarrow r \equiv (p \rightarrow r) \wedge (q \rightarrow r)$$

$$\equiv \neg(p \wedge q) \vee r$$

$$\equiv (\neg p \vee \neg q) \vee r$$

$$\equiv \neg p \vee \neg q \vee r \vee r$$

Logically not possible
possibly by the Truth
Table.

now

$$(p \rightarrow r) \wedge (q \rightarrow r)$$

$$\equiv (\neg p \vee r) \wedge (\neg q \vee r)$$

$$\equiv ((\neg p \wedge \neg q) \vee (\neg p \wedge r) \vee (r \wedge \neg q) \vee (r \wedge r))$$

$$\equiv (\neg(p \vee q) \vee (\neg p \wedge r) \vee (r \wedge \neg q) \vee r)$$

$$\equiv \neg(p \vee q) \vee$$

$$\equiv \neg p \vee r \vee \neg q \vee r$$

$$\equiv (p \rightarrow r) \vee (q \rightarrow r)$$

$$\equiv$$

$$\neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q$$

L.H.S

$$\equiv \neg(p \leftrightarrow q)$$

$$\equiv \neg((p \rightarrow q) \wedge (q \rightarrow p))$$

$$\equiv \neg(p \rightarrow q) \vee \neg(q \rightarrow p)$$

$$\equiv \neg(\neg p \vee q) \vee \neg(\neg q \vee p)$$

$$\equiv (\neg(\neg p) \wedge \neg q) \vee (\neg(\neg q) \wedge \neg p)$$

$$\equiv (p \wedge \neg q) \vee (q \wedge \neg p)$$

$$\equiv (p \vee q) \wedge (p \vee \neg p) \wedge (\neg q \vee q) \wedge (\neg q \vee \neg p)$$

$$\equiv (p \vee q) \wedge T \wedge T \wedge (\neg q \vee \neg p)$$

$$\equiv (p \vee q) \wedge (\neg q \vee \neg p)$$

now

$$p \leftrightarrow \neg q$$

$$\equiv (p \rightarrow \neg q) \wedge (\neg q \rightarrow p)$$

$$\equiv (\neg p \vee \neg q) \wedge (\neg(\neg q) \vee p)$$

$$\equiv (\neg q \vee \neg p) \wedge (q \vee p)$$

$$\equiv (q \vee p) \wedge (\neg q \vee \neg p)$$