

- State whether the following statements are TRUE or FALSE. Justify your answer. (4 marks)
 - A good algorithm for scheduling of hard real-time tasks tries to complete each task in the shortest possible time.
 - Soft real-time tasks do not have any associated time bounds.
 - With EDF a higher utilization is possible than with RMS.
 - The LL-bound ($LL(n)=n(2^{1/n}-1)$) and HB-bound ($HB(n): \prod(U_i+1) \leq 2$) are examples of sufficient conditions for fixed-priority preemptive scheduling (FPPS). Give at least three assumptions on which these bounds are based.
- Draw the full scheduling tree for the following set of non-preemptive tasks and mark the branches that are pruned by the Bratley's algorithm. (3 marks)

	J1	J2	J3	J4
a_i	0	4	2	6
C_i	6	2	4	2
D_i	18	4	7	4

- Consider the following taskset. (5 marks)

Name	Priority	Computation Time	Resources
τ_1	1	C_1	Ra, Rb
τ_2	2	C_2	none
τ_3	3	C_3	Ra
τ_4	4	C_4	Rb

The priorities are fixed with lower numbers representing higher priorities. The tasks need resources that they reserve through a regular locking mechanism. For example, task τ_1 needs resources Ra and Rb for its operations and acquires them in that order.

- Discuss two scenarios for priority inversion.
 - Assume that we use the priority inheritance protocol to resolve this. What is the maximum blocking time for τ_1 expressed in the given computation times? And what is this value if the priority ceiling protocol is used?
- How can the following taskset be optimally scheduled (using fixed priority scheduling)? Is this task set schedulable? (6 marks)

	Time Period	Computation Time	Relative Deadline
τ_1	8	4	8
τ_2	10	2	7
τ_3	30	5	30

Now we include the Maximum Blocking time also for each task as shown below.

	Time Period	Computation Time	Blocking Time	Relative Deadline
τ_1	8	4	2	8
τ_2	10	2	2	7
τ_3	30	5	0	30

Is the task set schedulable? If not, what change in deadline parameter will make it schedulable.

- Consider two jobs $J1 (r1, C1, d1)$ and $J2 (r2, C2, d2)$ where $J1$ precedes $J2$ (i.e., $J1 \rightarrow J2$), but $J2$ arrives before $J1$ and has an earlier deadline. How can the two tasks be scheduled so that the timing and precedence constraints are met. The tasks are preemptable. Assume that $C1+C2 < d2$. (3 marks)
- Formulate processor demand function for 'n' preemptive periodic tasks with deadlines less than periods. (4 marks)