

CMPSC122 Lab Section 004L

Lab 12

Lab Proj 12 – Algorithm Analysis

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Date: April 29th, 2025

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1 Objective (or Abstract)

Objective of this lab report:

1. Analyze each algorithm's time complexity theoretically using Big-O notation.
2. Compare their actual performance through practical testing on varying input sizes.
3. Determine the most efficient algorithm for real-world applications

2 Introduction

In this case study, I will provide a report on how different algorithms calculate the maximum sum for a contiguous subset in a larger set. There will be three algorithms analyzed in this report: Brute force, divide and conquer using the recursion method, and linear scan using dynamic programming.

I will first provide an analysis of each algorithm, as well as a hypothesis of how well each algorithm will run

3 Procedure

Equipment used for the experiment:

- Laptop Asus ROG Zephyrus G16 (2024) GU605
 - o OS: Windows 11 Home - 23H2
 - o CPU: Intel(R) Core(TM) Ultra 9 185H
 - 16 Cores, 22 Threads
 - 5.1 GHz – 115W
 - o GPU: NVIDIA GeForce RTX 4070 Laptop GPU
 - VRAM: 8GB GDDR6X
 - o RAM: 16 GB
 - DDR5 – 7467 MT/s

During the testing period, the program (in Visual Studio Code) was reported by Task Manager to use roughly $10 \pm 2\%$ of the CPU and 520 ± 10 MB of RAM

4 Discussion

1. Hypothesis

Brute force (Blue)

This algorithm (Blue) uses a nested loop to check all possible continuous subsets of numbers.

Code analysis:

- The outer loop runs from $i = 0$ to $i = \text{size} - 1$.

- The inner loop runs from $j = 0$ to $j = \text{size} - i - 1$, computing the sum of the subsequence starting at i and of length $j + 1$.
- The maximum sum encountered is stored in `currentMax`

Big-O analysis:

- The outer loop runs $O(n)$ times.
 - The inner loop runs $O(n)$ times in the worst case (when $i = 0$, it runs n times).
- ⇒ In conclusion, time complexity is $O(n^2)$

Divide and conquer (Green)

This algorithm (Green) uses a recursive divide-and-conquer approach.

Code analysis:

1. Base case
 - If `start > end`, return 0.
 - If `start == end`, return `max(0, array[start])`
2. Divide
 - Split the array into left and right halves around `mid = (start + end) / 2`.
 - Recursively compute the max sublist sum for the left and right halves.
3. Recursive
 - Compute the max sum that crosses the midpoint by:
 - Extending leftwards from `mid` to `start` to find the best left suffix.
 - Extending rightwards from `mid + 1` to `end` to find the best right prefix.
 - The max of the left half, right half, and crossing sum is returned.

Big-O analysis:

- The recurrence relation is: $T(n) = 2T(n/2) + O(n)$. (The $O(n)$ term comes from the two linear scans to compute the crossing sum)
- ⇒ Time complexity is $O(n \log n)$

Linear scan (Red)

The Red function computes the maximum subset sum in one pass.

Code analysis:

- If `maxSumEndingAt[i-1]` is negative, start a new sublist at i (since adding a negative reduces the sum).
- Otherwise, extend the previous sublist by including `array[i]`.
- Update `currentMax` if the current sublist sum is greater.

Big-O analysis:

- The loop runs exactly $n-1$ times from $i = 1$ to $i = \text{size} - 1$.
- ⇒ Total time complexity is $O(n)$

2. Expected performance

ALGORITHM	APPROACH	TIME COMPLEXITY	# OPERATIONS TO FINISH N = 64000
BLUE	Brute force	$O(n^2)$	$64000^2 \approx 4 \text{ Billion}$
GREEN	Recursion	$O(n \log n)$	$64000 * 16 \approx 1 \text{ Million}$
RED	Linear	$O(n)$	64000

3. Actual running time

The number of loops is set at 20, which will test $n = 500, 1000, 2000, 4000, \dots, 262144000$.

ALGORITHM	TIME COMPLEXITY	TOTAL TIME TAKEN (SECONDS)
BLUE	$O(n^2)$	1009.494 *
GREEN	$O(n \log n)$	34.211
RED	$O(n)$	1.248

* The Blue algorithm is only tested for 12 loops ($n = 500, 1000, 2000, \dots, 1024000$) due to time constraints, but a comprehensive conclusion can still be surmised.

5 Conclusion

- The Blue algorithm ($O(n^2)$) is substantially slower compared to the other two algorithms and is extremely slow for large inputs. This matches with the time complexity where for each time n doubles, the time it takes to calculate quadruple.
 - ⇒ Not suitable for real-world use beyond very small datasets
- The Green algorithm ($O(n \log n)$) is significantly faster than the Blue algorithm but still slightly slower than the Red algorithm, because the $\log n$ function grows incredibly slowly, slower than n (in the Blue algorithm), but it is still more than 1 (in the Red algorithm). This is a viable but still generally not a recommended choice for calculating.
 - ⇒ Usable but not recommended.
- The Red algorithm ($O(n)$) is the best-performing algorithm in practice, matching theoretical expectations for a linear function; matching the expectation for a linear time complexity.
 - ⇒ Ideal for real-world use, even for very large datasets.