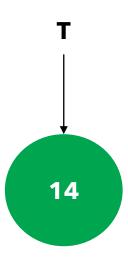
#### Alberi binari

Liceo G.B. Brocchi - Bassano del Grappa (VI) Liceo Scientifico - opzione scienze applicate Giovanni Mazzocchin

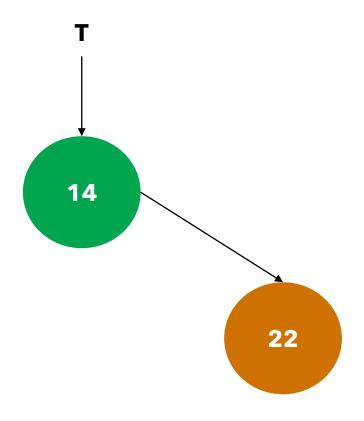
```
T = nil
bst_insert(T, 14)
```

**caso base**: si crea un nuovo nodo e lo si fa puntare da T

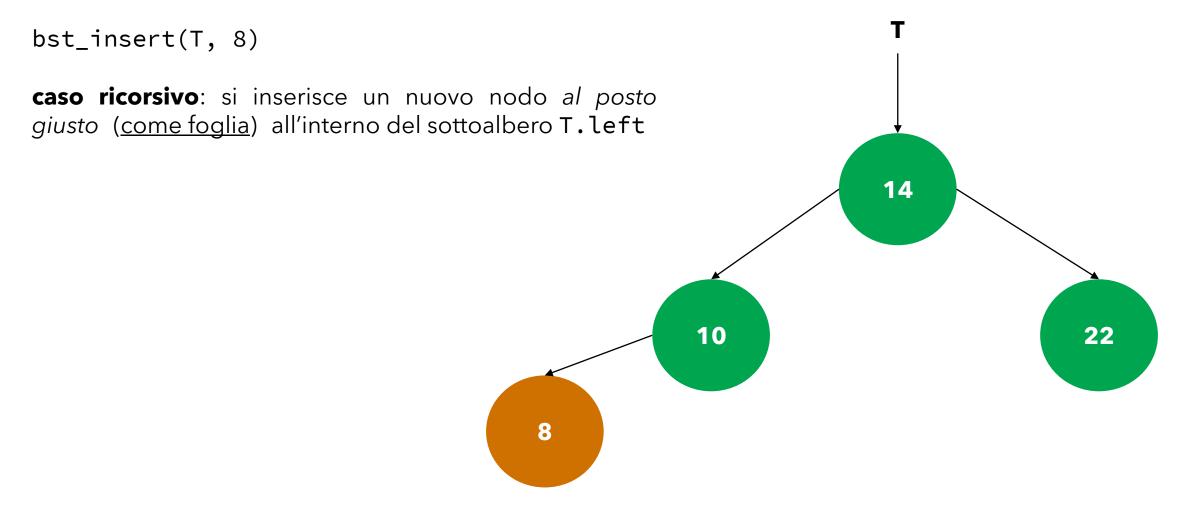


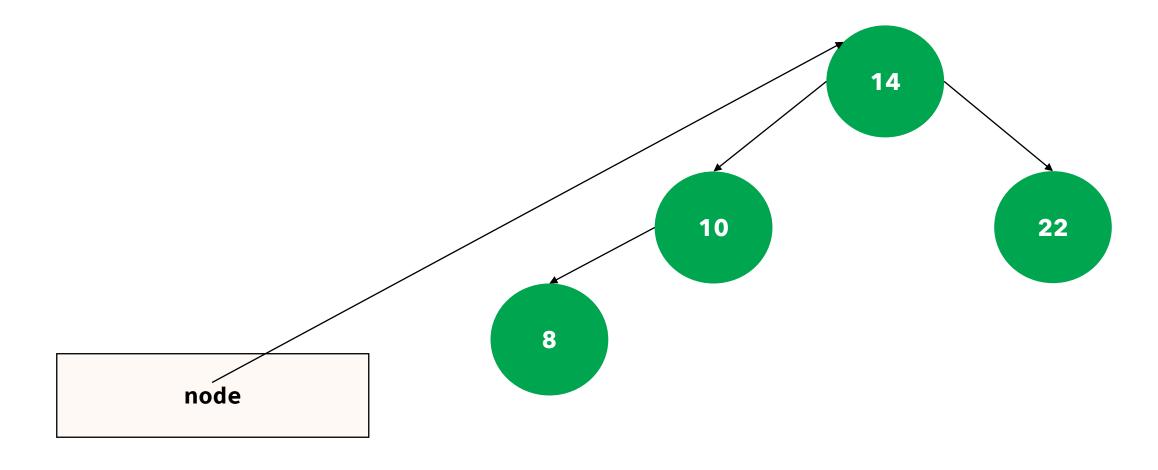
bst\_insert(T, 22)

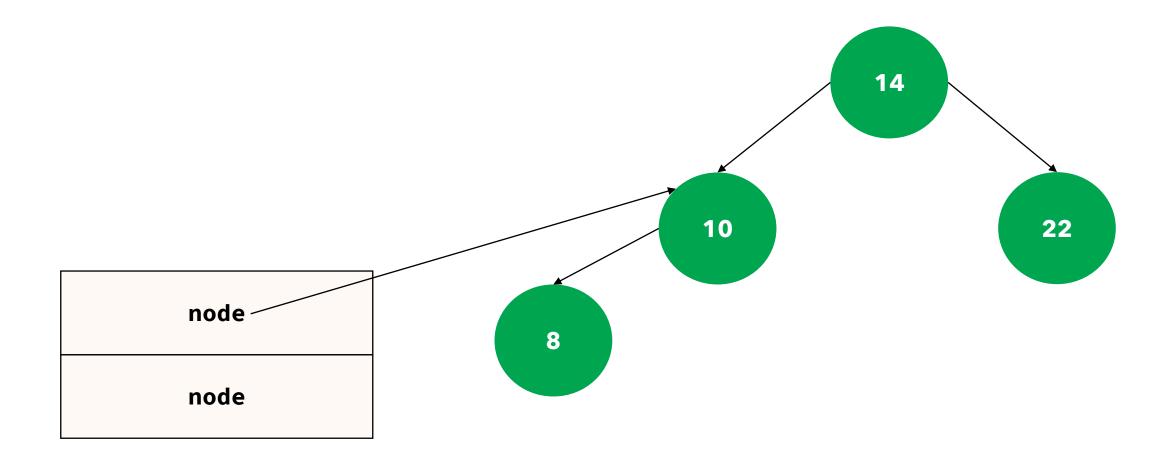
caso ricorsivo: si inserisce un nuovo nodo al posto
giusto (come foglia) all'interno del sottoalbero
T.right

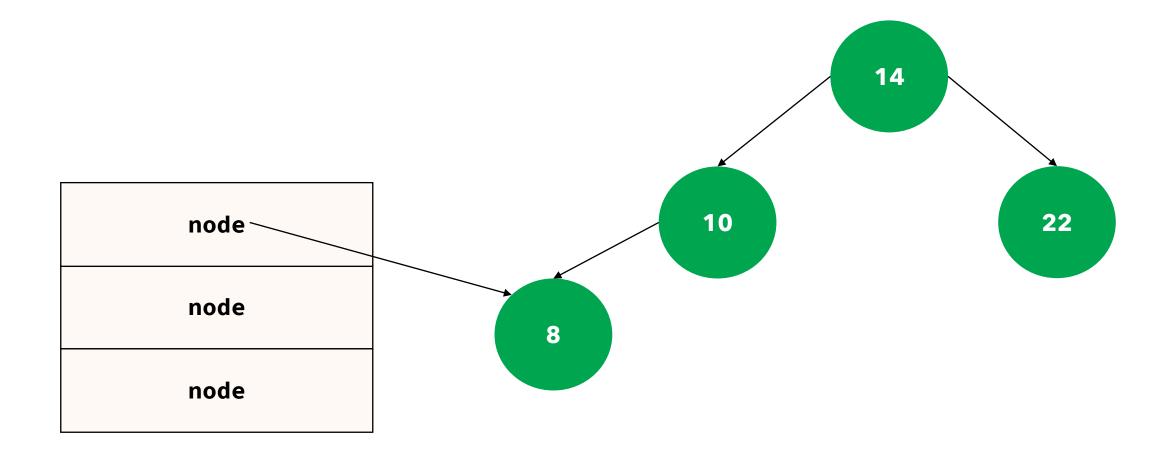


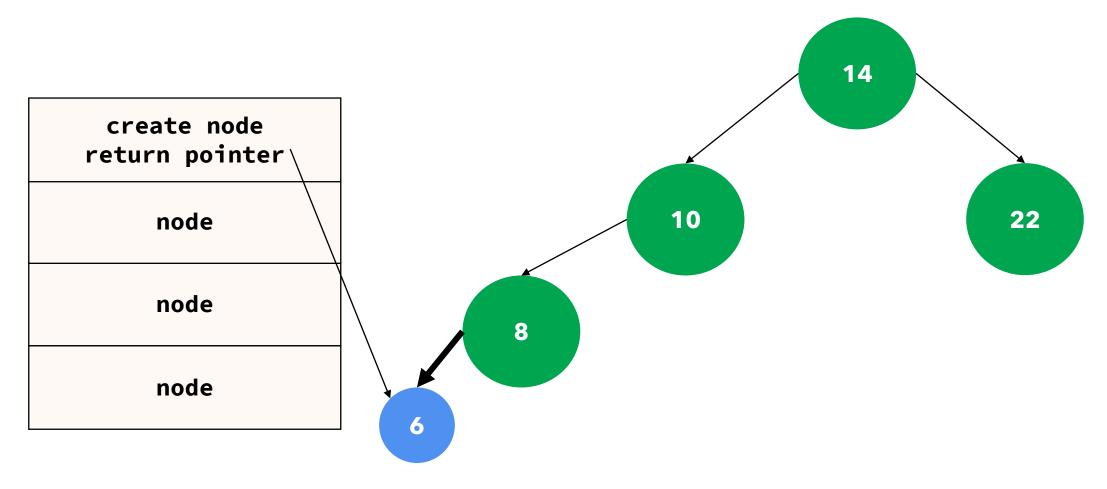
bst\_insert(T, 10) caso ricorsivo: si inserisce un nuovo nodo al posto giusto (come foglia) all'interno del sottoalbero T.left 14 10 22

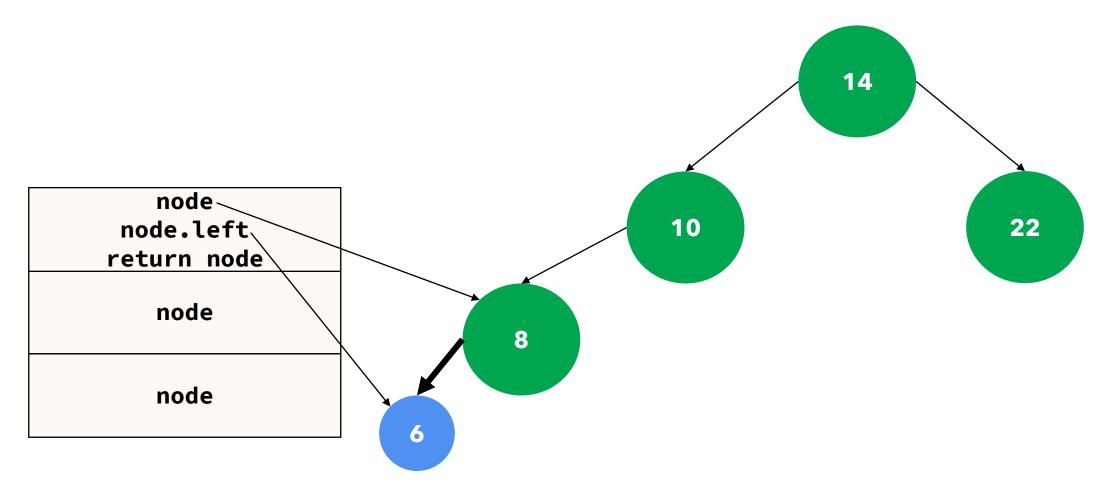


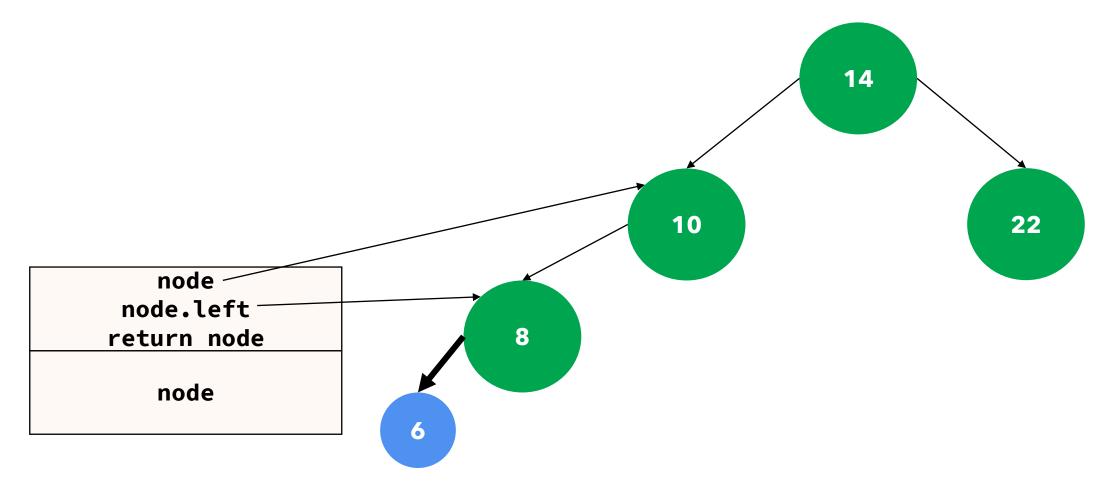


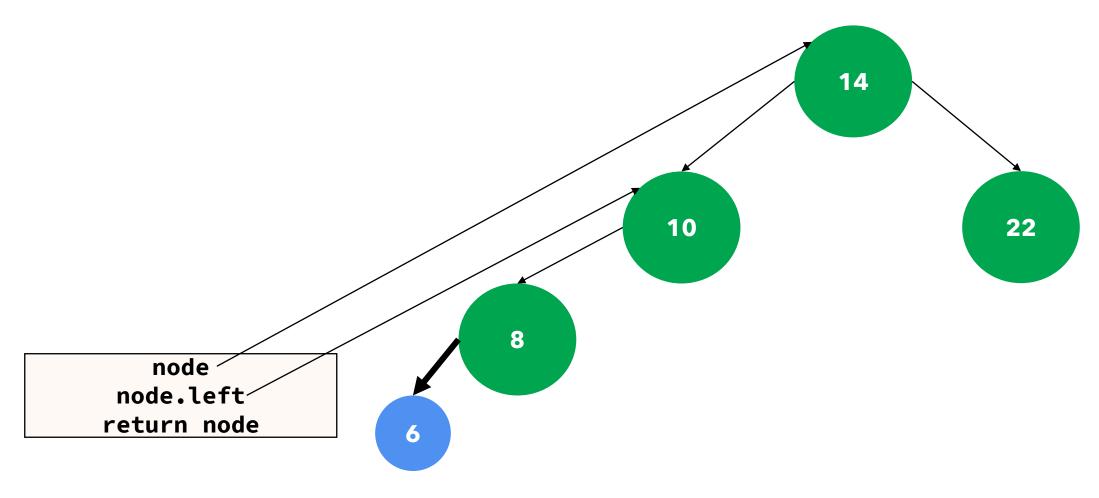




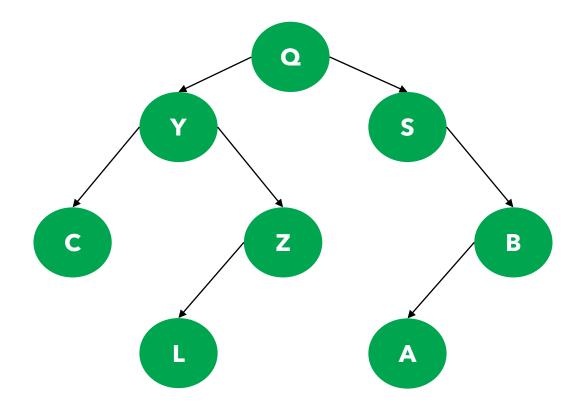


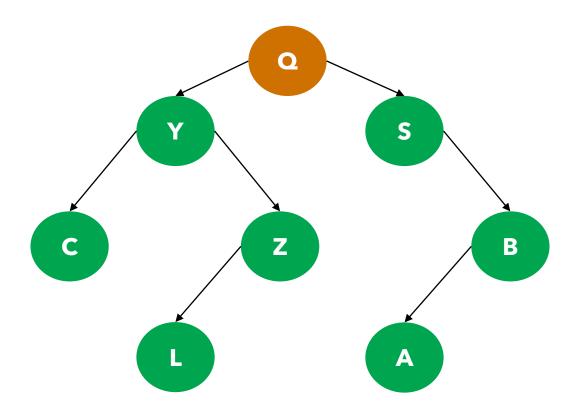


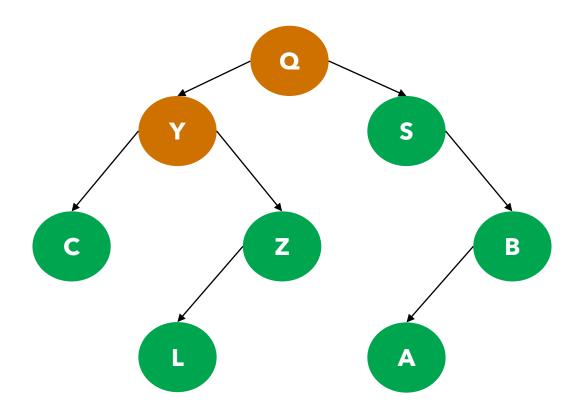


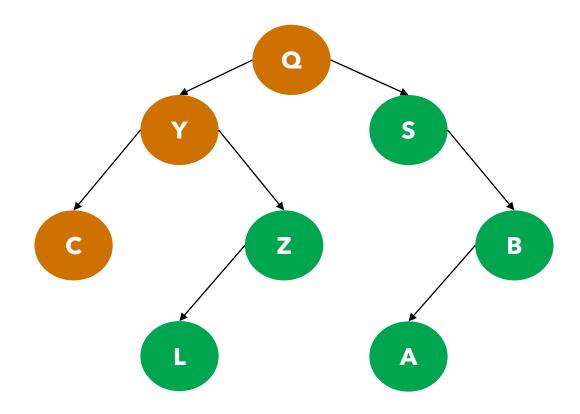


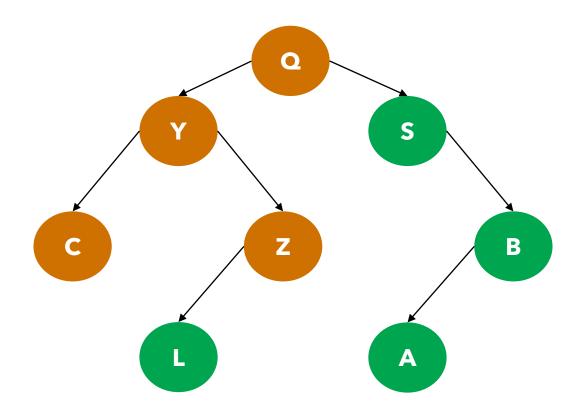
- Nella <u>ricerca in profondità</u> (depth-first), per ogni nodo vengono visitati ricorsivamente in profondità prima il sottoalbero sinistro e poi il sottoalbero destro
- Questo significa che prima di tornare indietro ad un nodo T, uno dei due sottoalberi di T deve essere visitato completamente
- Per tornare indietro fino al punto giusto e non ripetere gli stessi percorsi serve uno stack. Ovviamente, la ricorsione ci aiuta molto, perché si basa sul call stack del programma
- Esistono diverse versioni della ricerca in profondità:
  - visita **preorder**
  - visita **inorder**
  - visita <u>postorder</u>

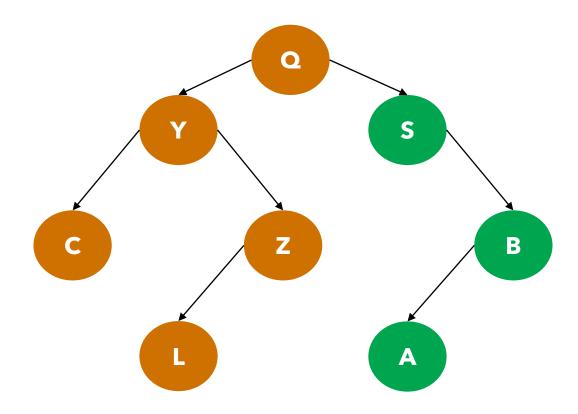


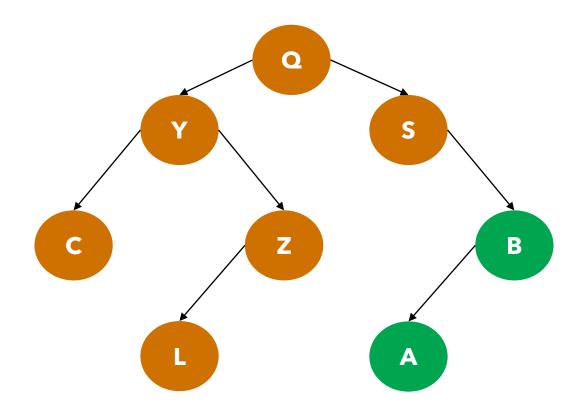


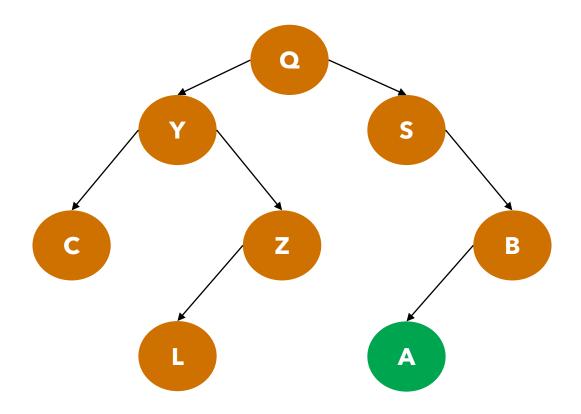


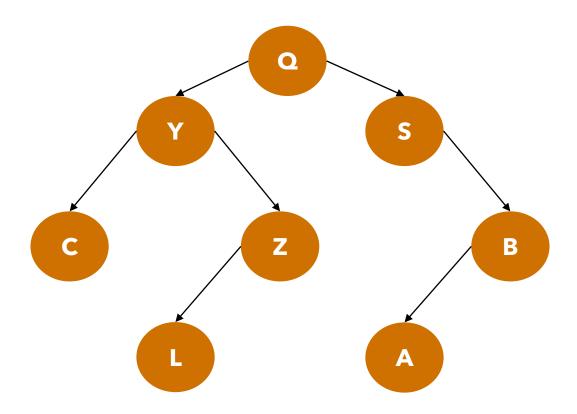








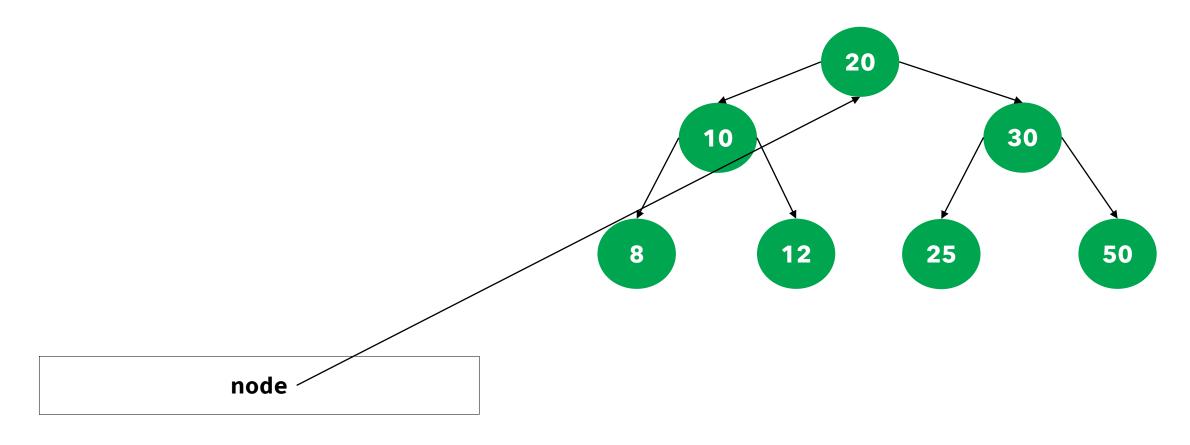


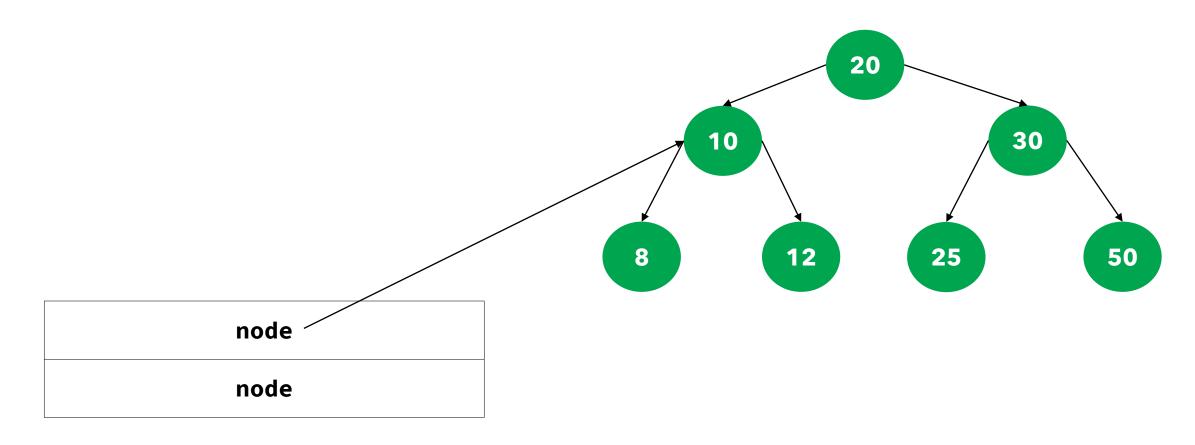


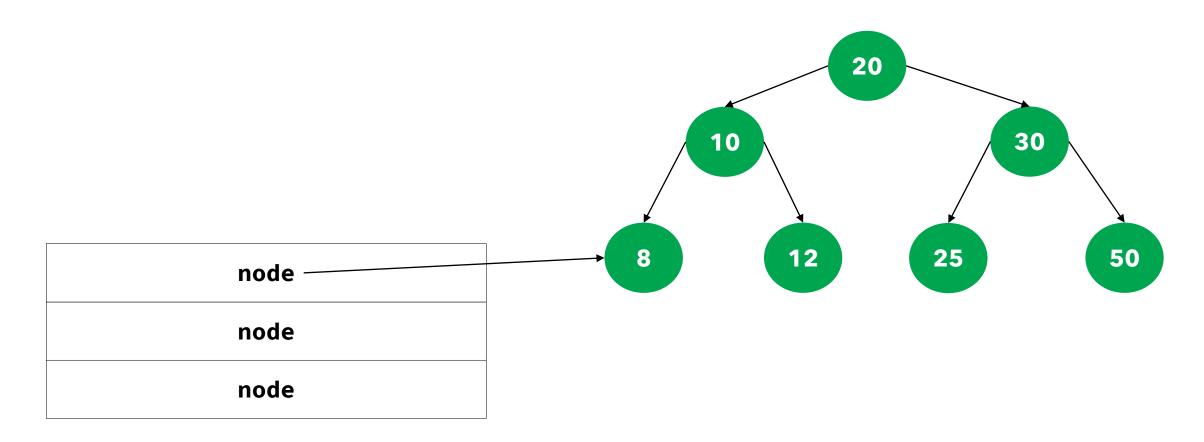
```
inorder_tree_walk(T):
```

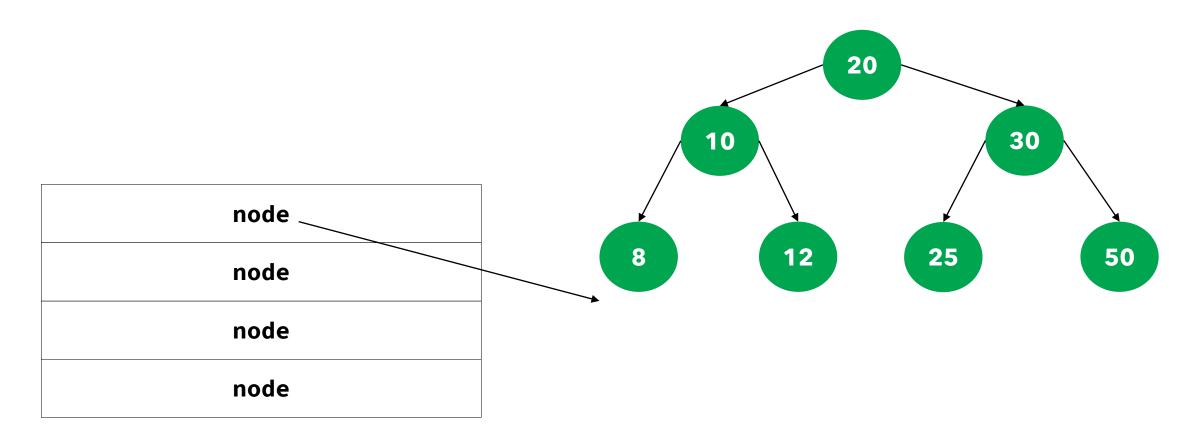
- if T is empty, do nothing
- otherwise:
  - call inorder\_tree\_walk on T.left
  - open T (e.g.: print T.key)
  - call inorder\_tree\_walk on T.right

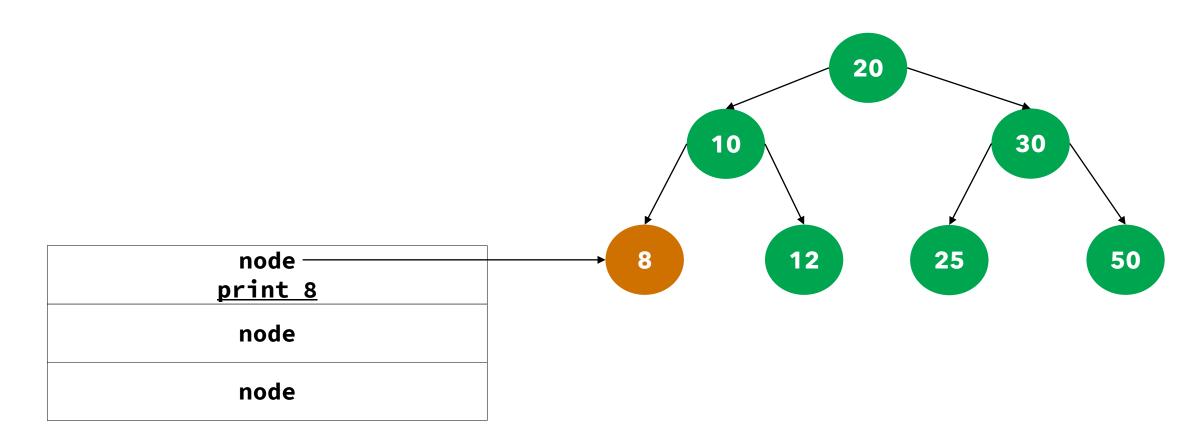
• NB: su un BST stampa le chiavi in ordine ascendente

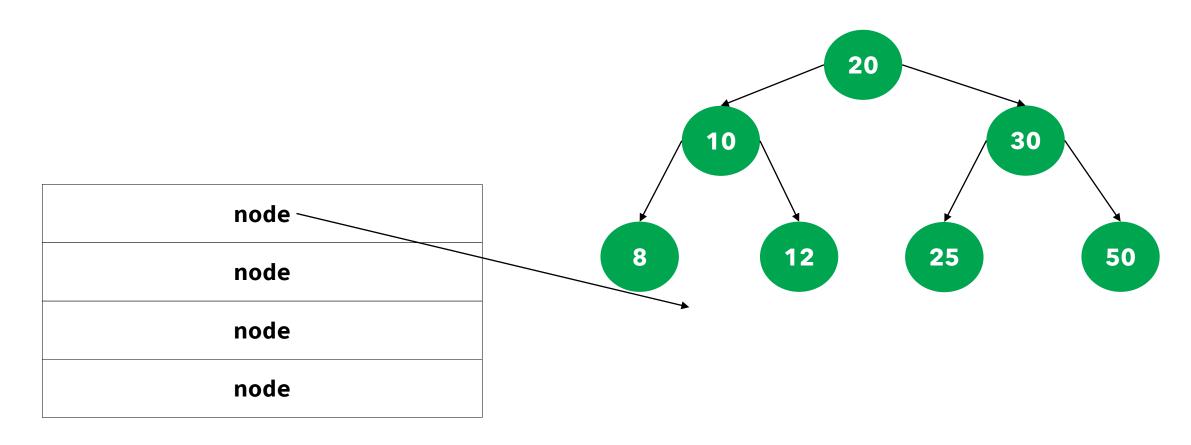


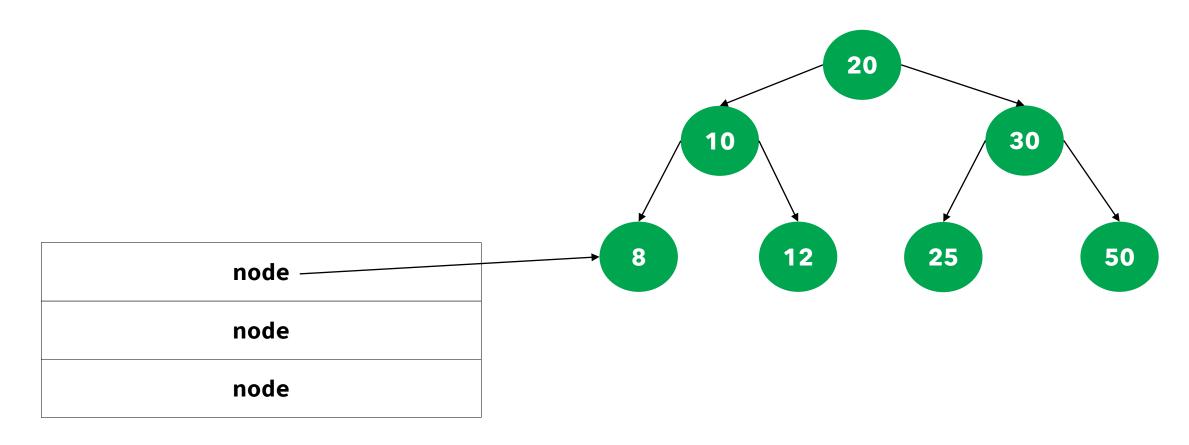


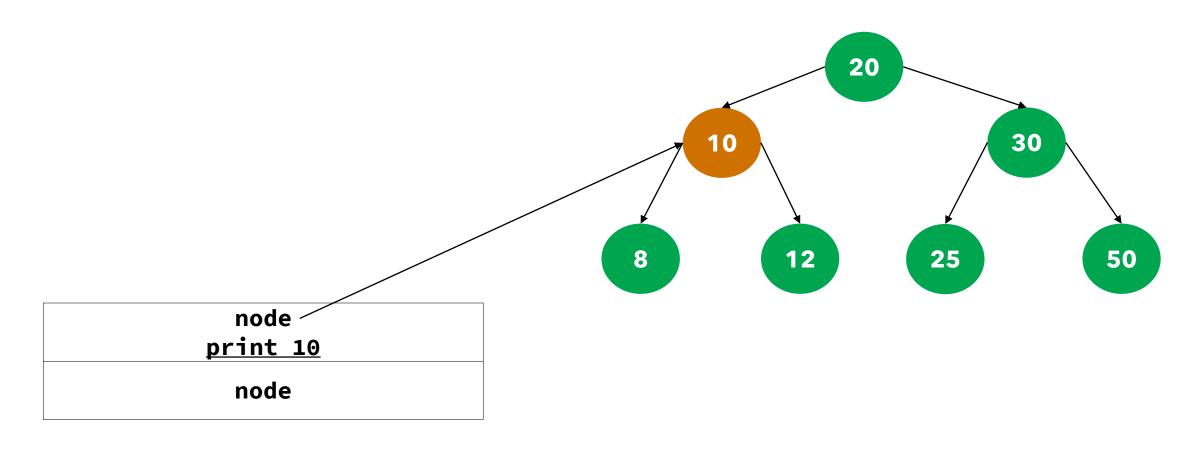


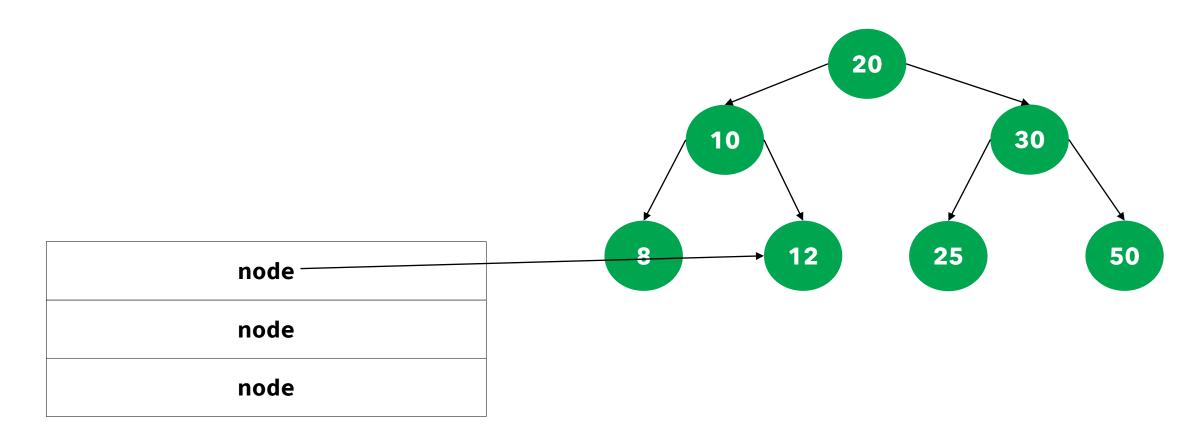


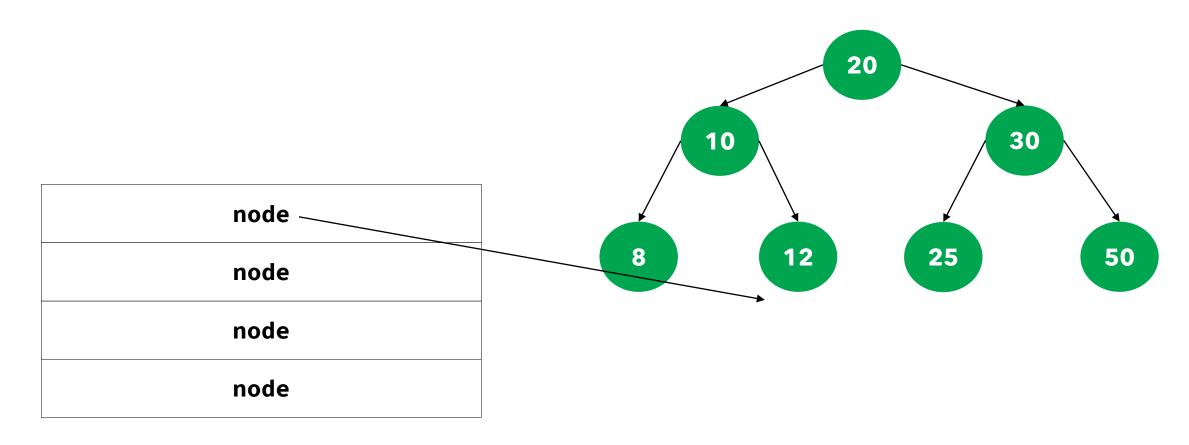


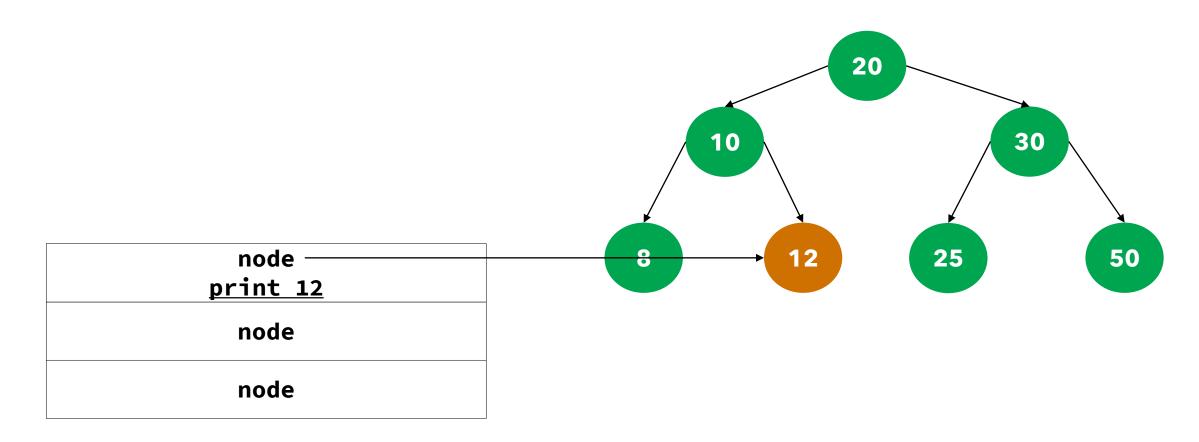


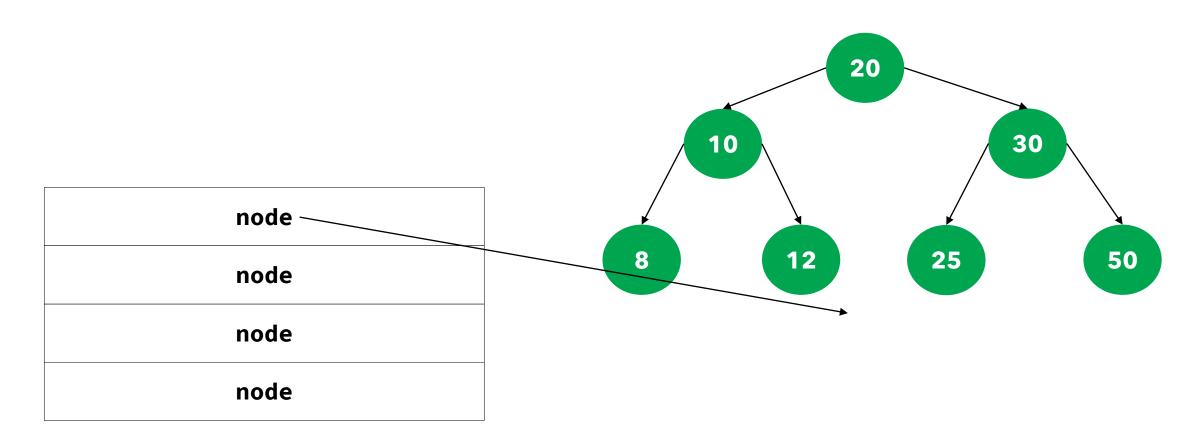


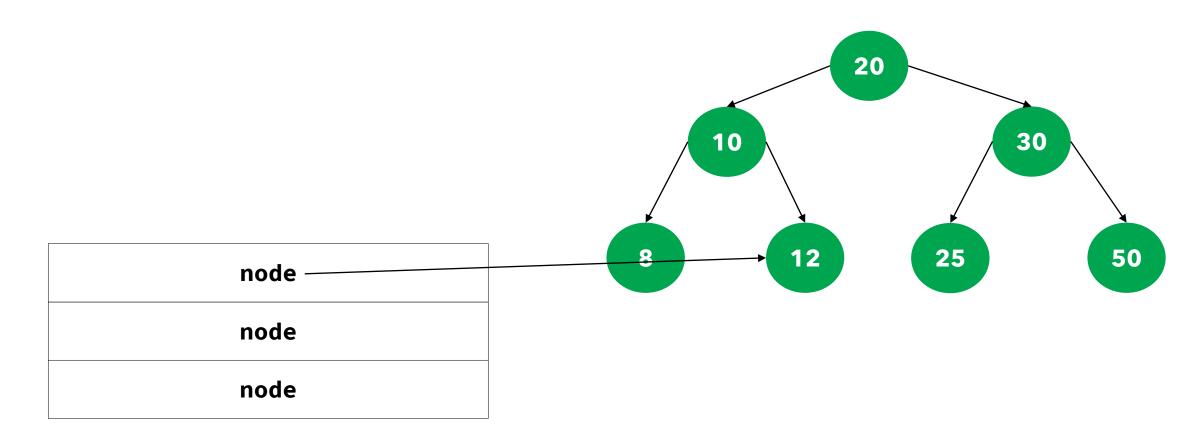


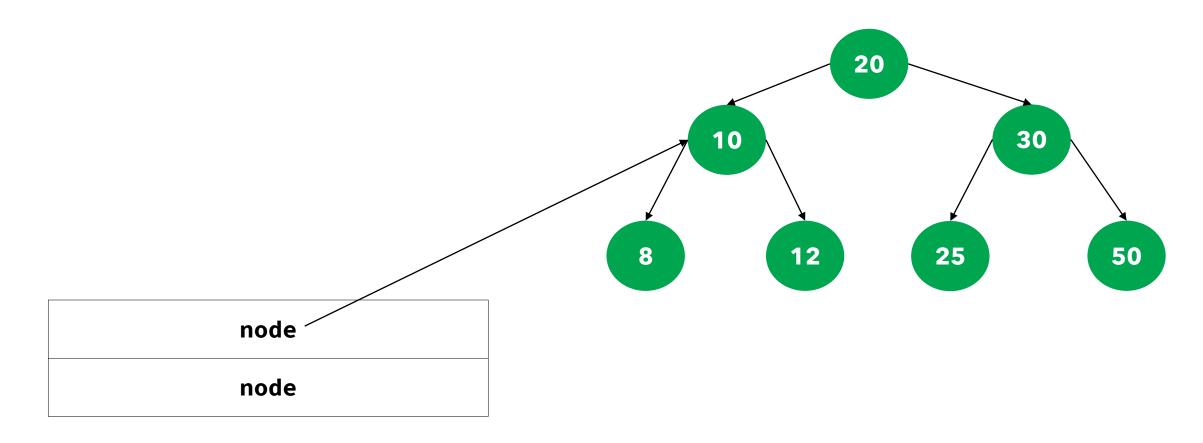


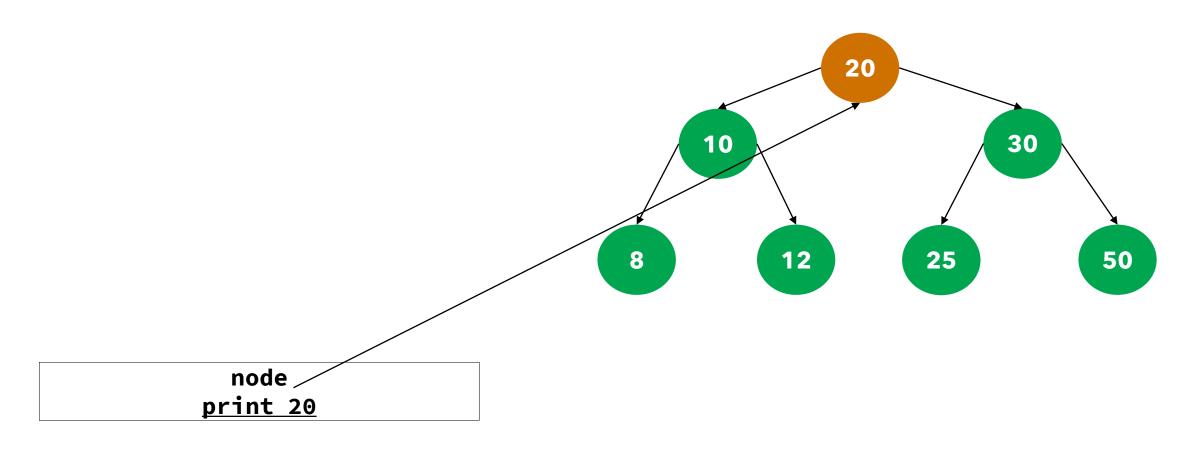




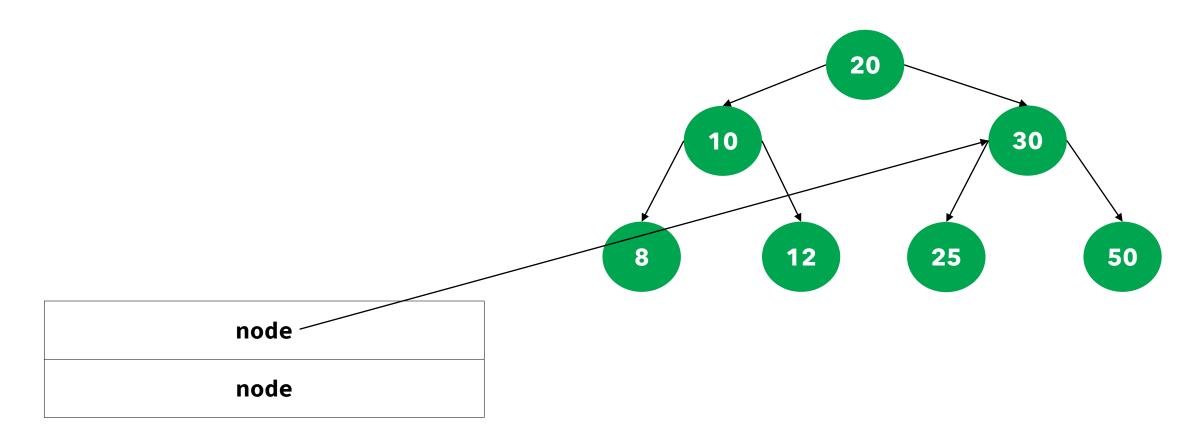


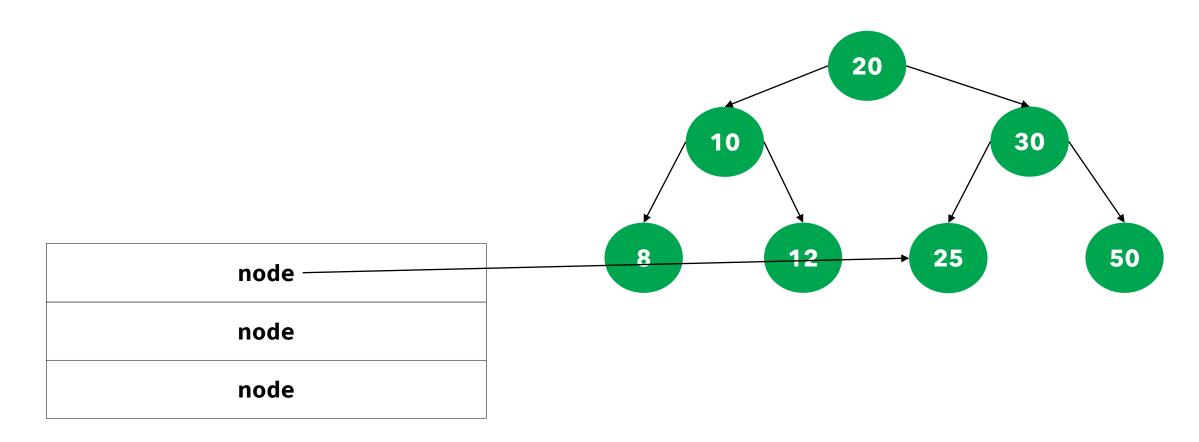


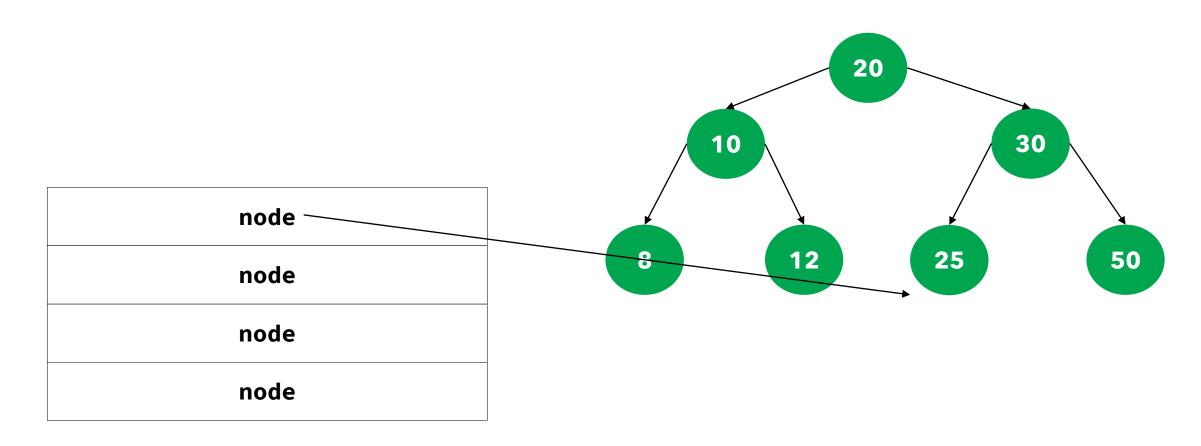


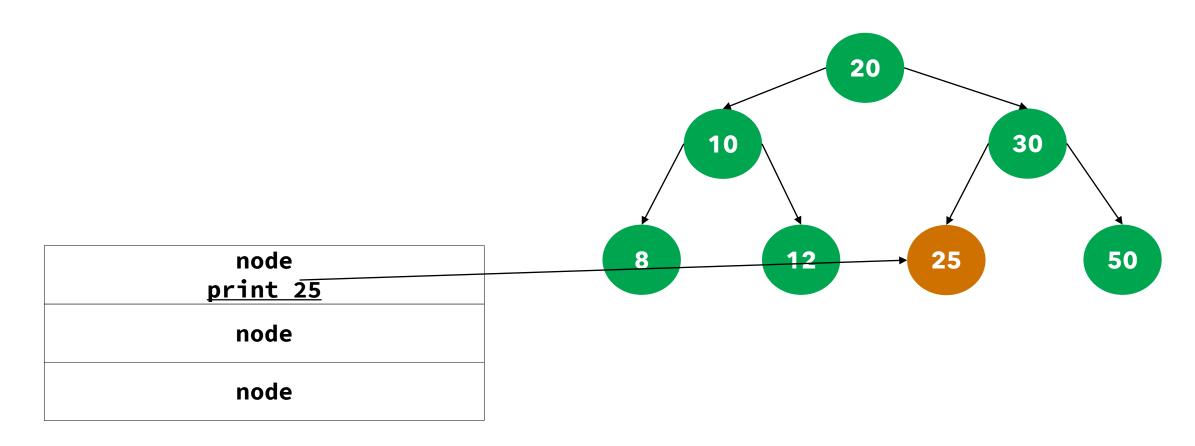


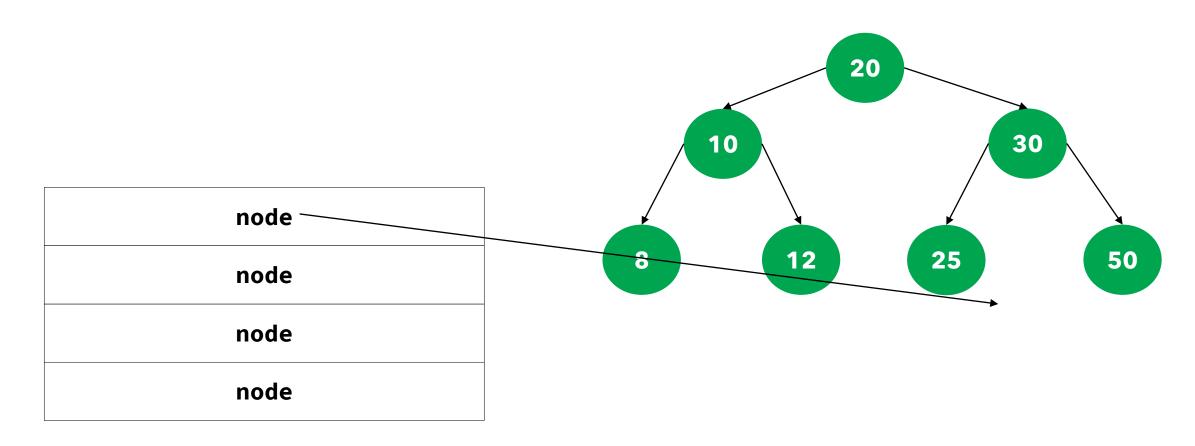
17/11/2023

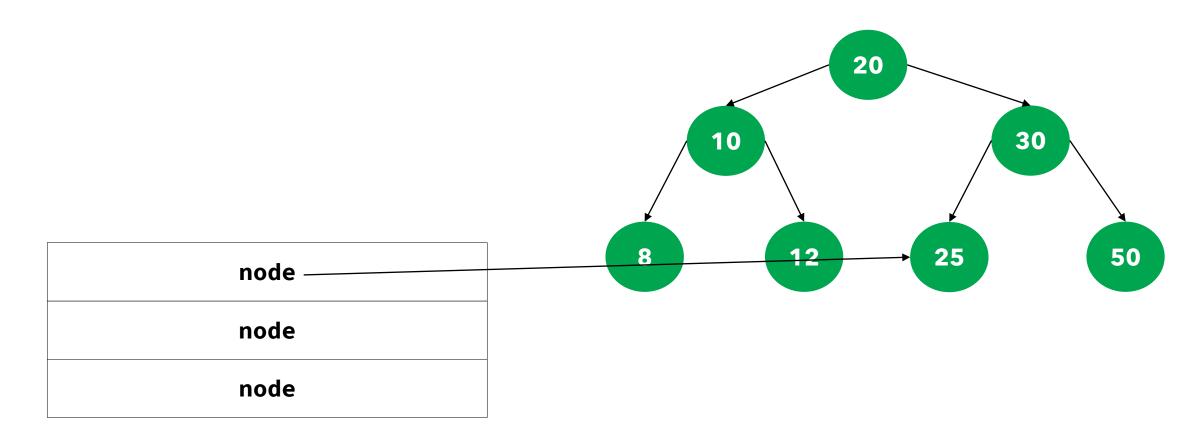


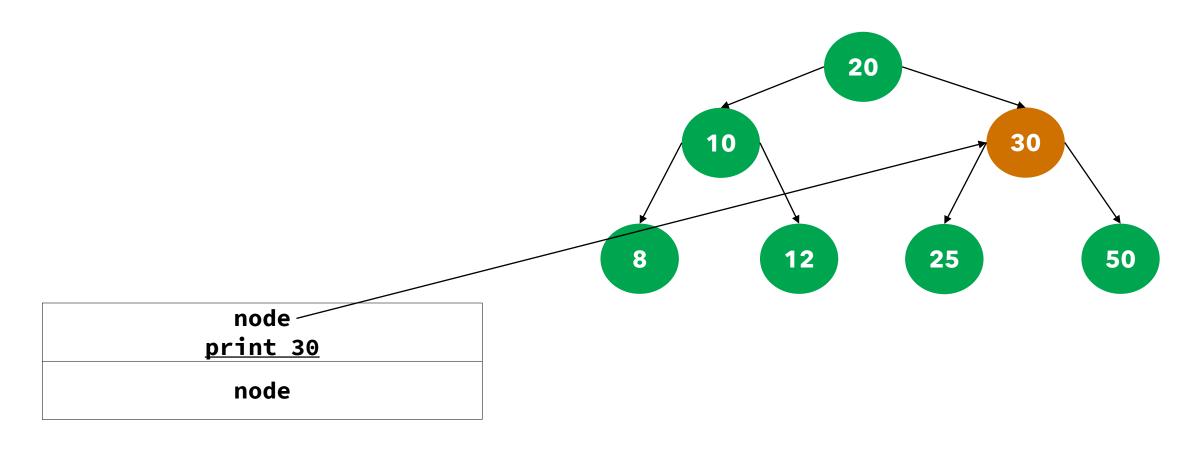


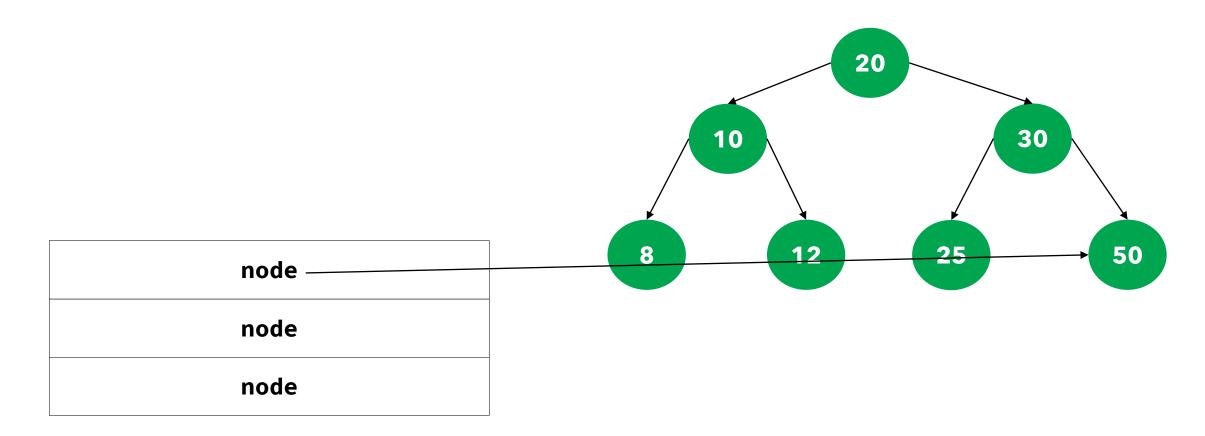


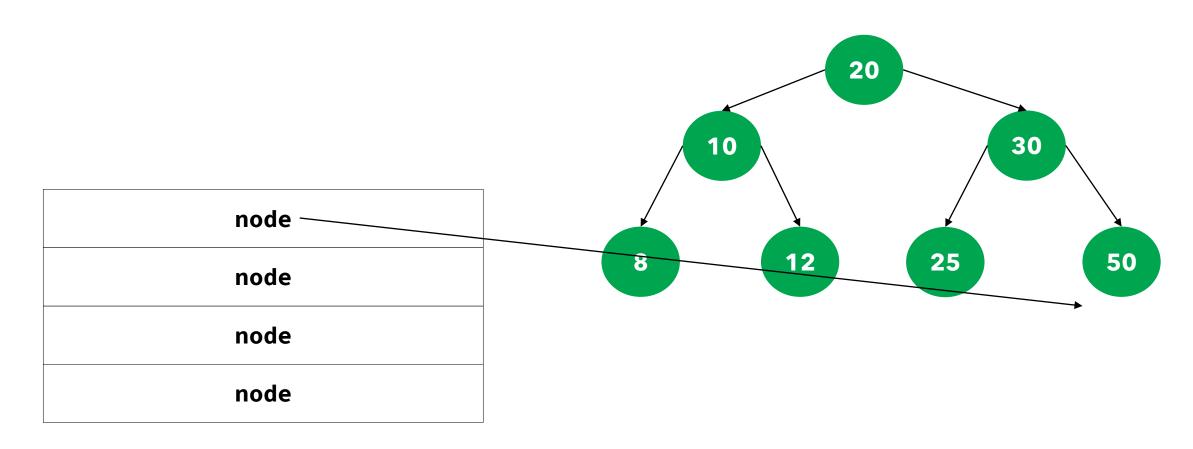


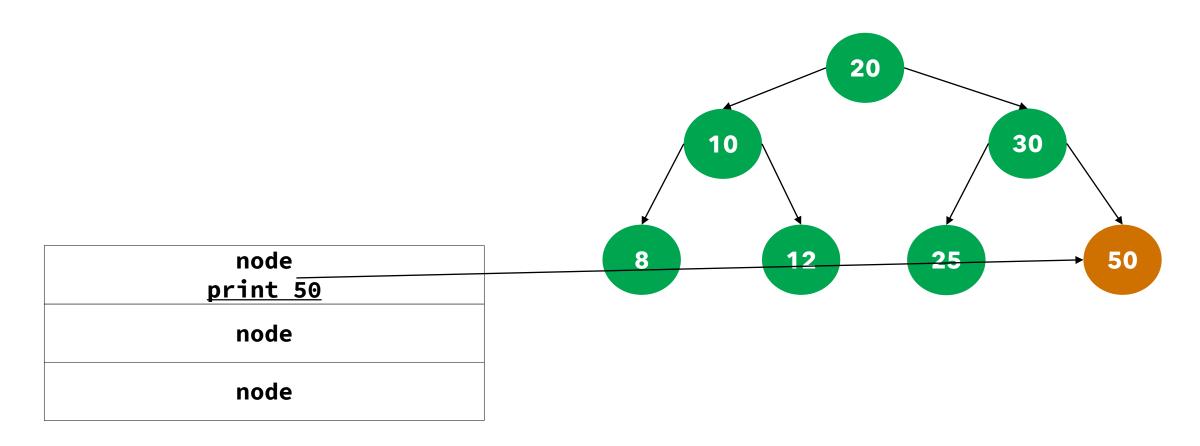


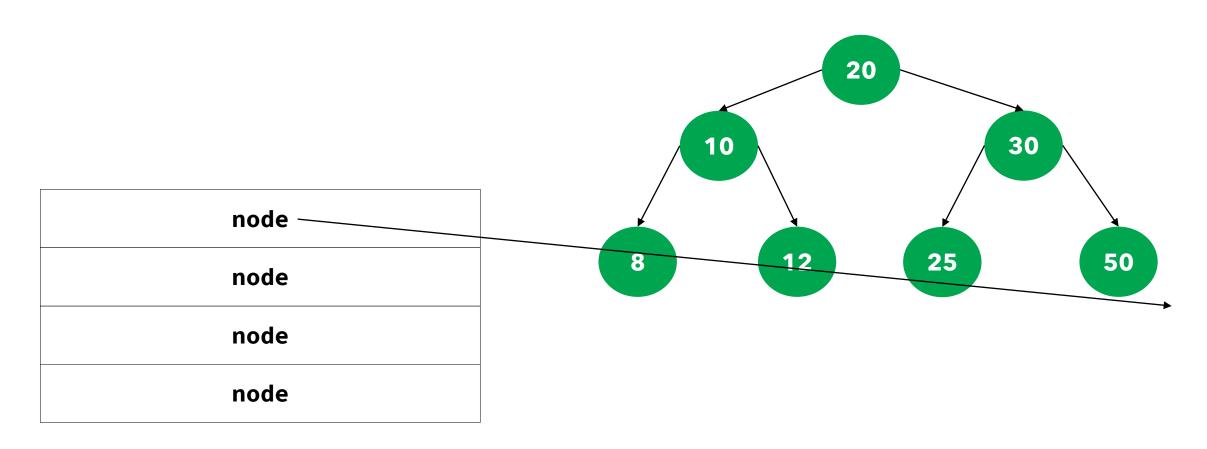


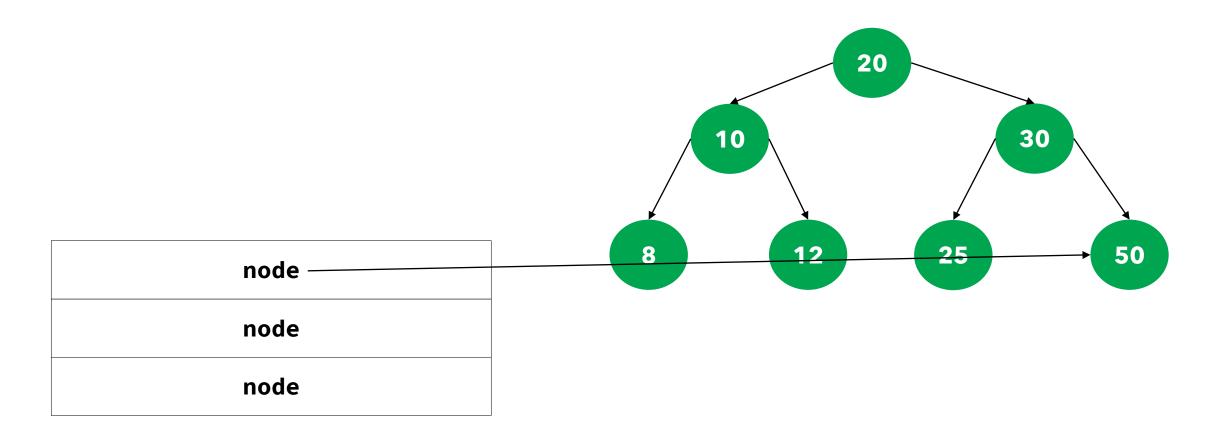


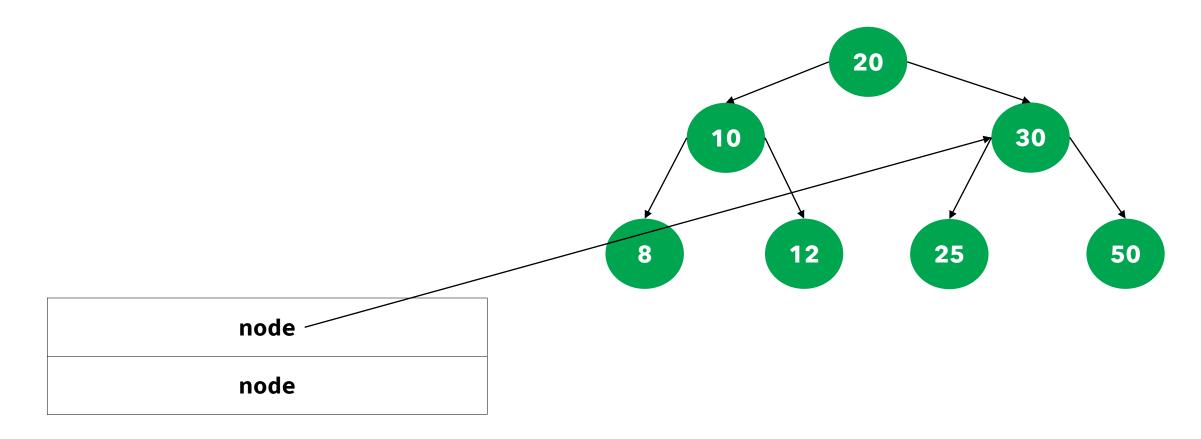


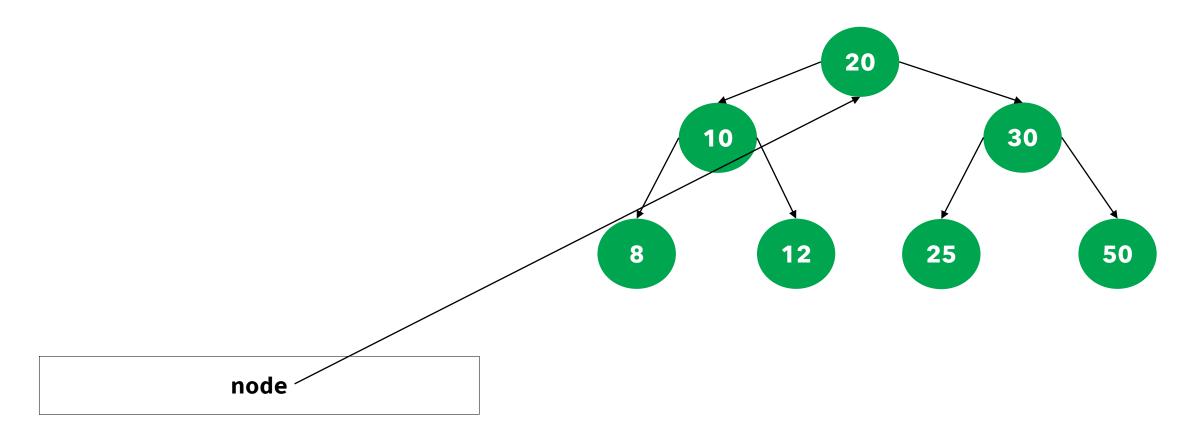


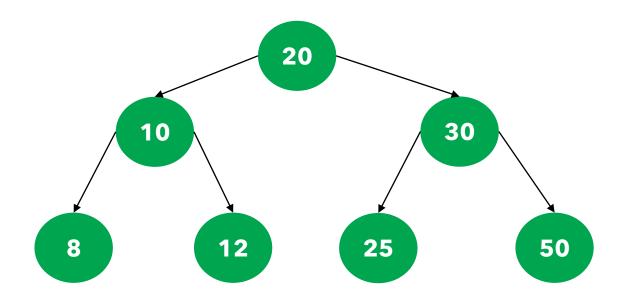






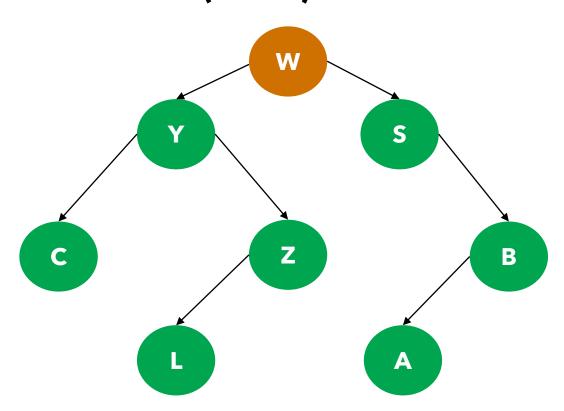






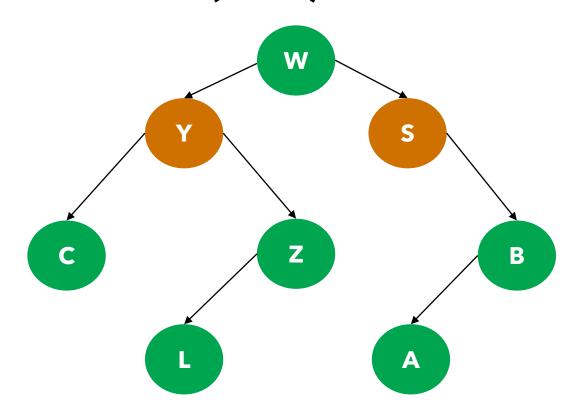
```
postorder_tree_walk(T):
    if T is empty, do nothing
    otherwise:
        call postorder_tree_walk on T.left
        call postorder_tree_walk on T.right
        open T (e.g.: print T.key)
```

- La <u>ricerca in ampiezza</u> è una ricerca per livelli
- Prima si visitano tutti i nodi al livello 0, poi tutti quelli al livello 1, poi tutti quelli al livello 2 etc...
- Visiteremo ciascun livello da sinistra a destra
- Per tenere traccia dei livelli nell'ordine corretto serve una queue



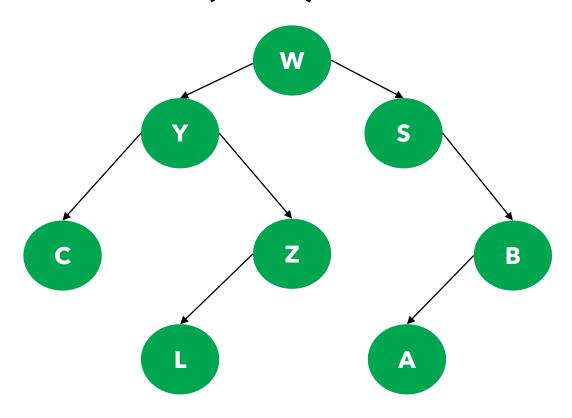
```
Q: []
enqueue(W);
```

**Q:** [W]



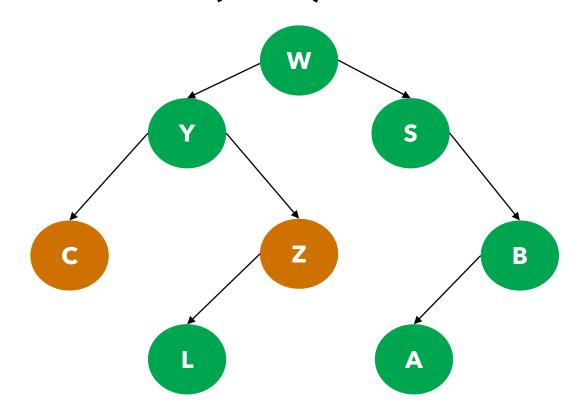
```
enqueue(W.left);
enqueue(W.right);
```

Q: [W, Y, S]



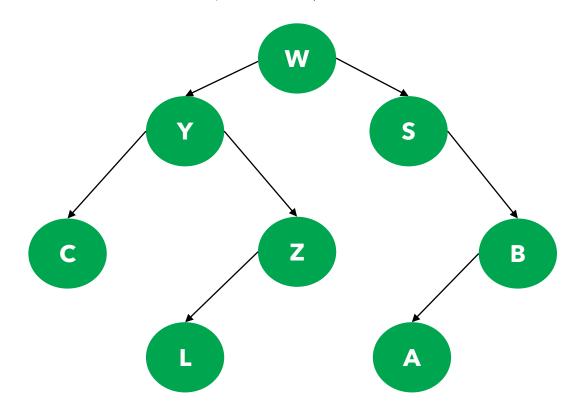
dequeue();

Q: [Y, S]



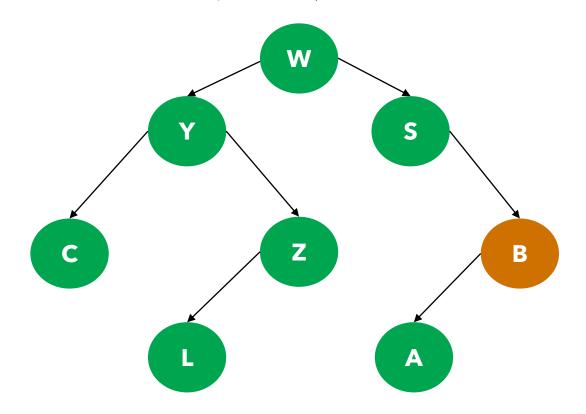
```
enqueue(Y.left);
enqueue(Y.right);
```

Q: [Y, S, C, Z]



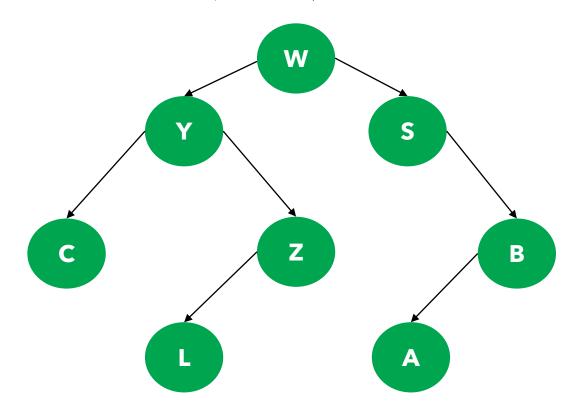
dequeue();

Q: [S, C, Z]



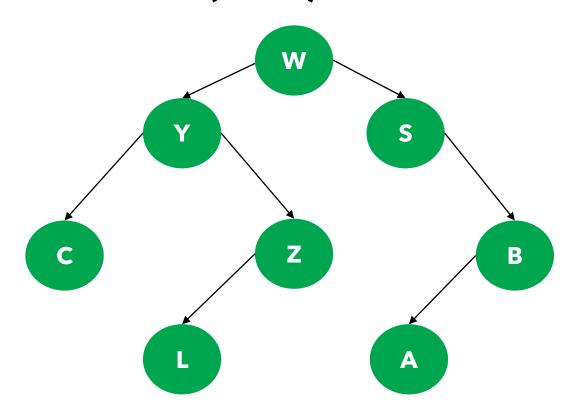
enqueue(S.right);

Q: [S, C, Z, B]



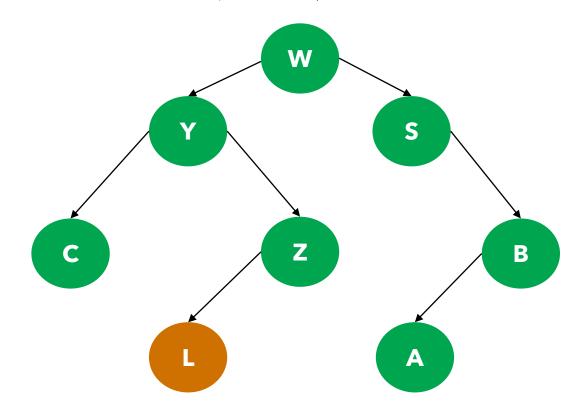
dequeue();

Q: [C, Z, B]



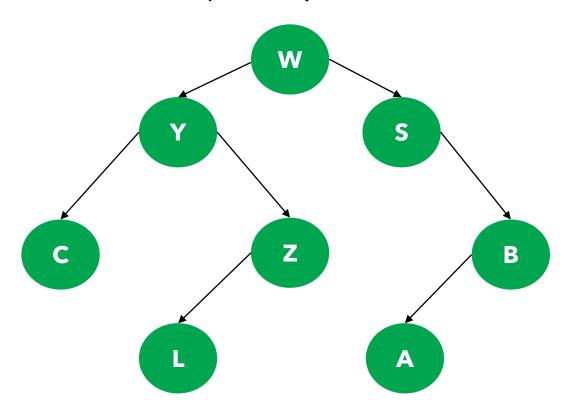
dequeue();

Q: [Z, B]



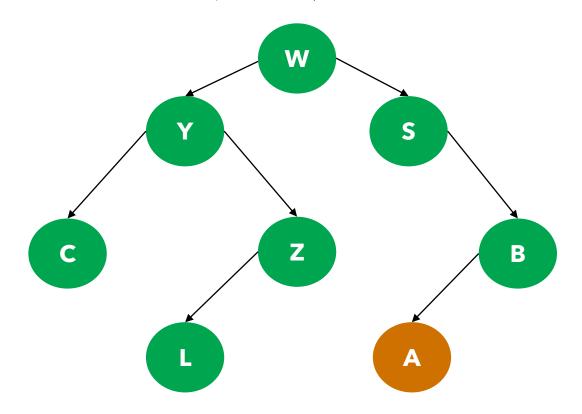
enqueue(Z.left);

Q: [Z, B, L]



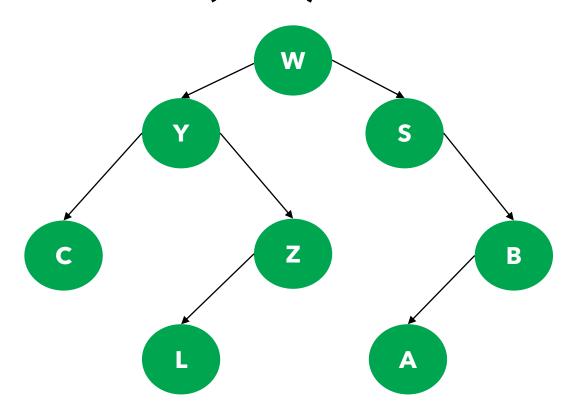
dequeue();

Q: [B, L]



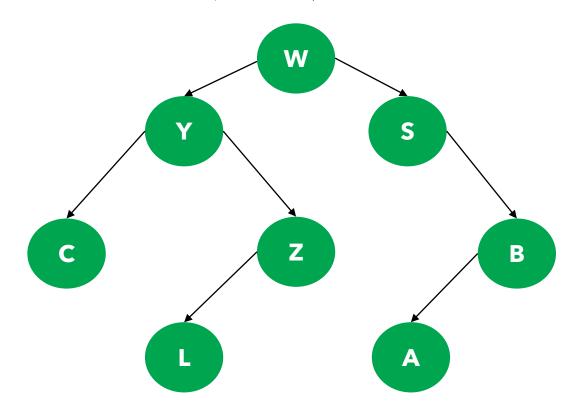
enqueue(B.left);

Q: [B, L, A]



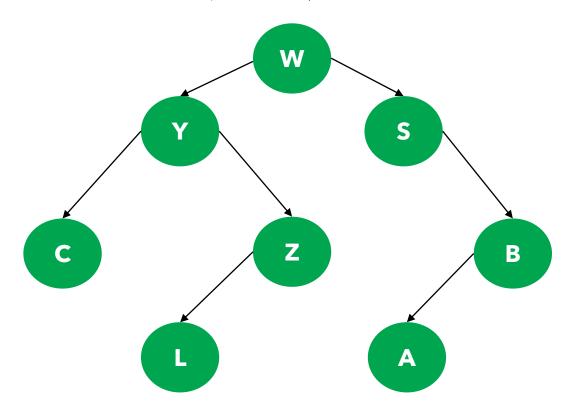
dequeue();

Q: [L, A]



dequeue();

**Q:** [A]



dequeue();

**Q:** []