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ESE 344

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ESE 344 HW 1

QUESTION 1:

1a.

Homers algorithm in an O notation is a runtime of $O(n)$, because it iterates from $i=n$ down to , and is a $n+1$ iteration structure.

1b.

One implementation in pseudo code for a naive polynomial evaluation algorithm could be:

```
naive-eval-funct(A, n, x)
  p = 0
  for i = 0 to n
    term = 1
    for j = 1 to i
      term = term * x
    p = p + A[i] * term
  return p
```

For each i in the loop, the total computer time sums to

$$\sum_{i=0}^n O(i) = O(n^2)$$

For our comparison we can see that homers is the MORE efficient solution due to its runtime.

1c.

Let's look at the loop-invariant proof and prove the termination

I. At initialization we have the state of

$$p = \sum_{k=0}^{n-(i+1)} A[k + i + 1] x^k$$

Where $i=n$, there are NO terms ($n-(n+1)=-1$)

$p=0$

Invariant will hold in this case.

II. During maintenance we have

$P = A[i] + x * (\text{previous } p)$

Matches required form for following iteration

Invariant still holds.

III. During the end or termination of this algorithm

Loop ends when $i = -1$

Then we substitute for -1 and get

$P(x) = \text{Summation}(n-(-1+1) \text{ } k=0 \text{ } A[k+(-1+1)] x^k$

$= \text{summation } n-k=0 \text{ } A[k] x^k$, and this is the polynomial we actually wanted

Horners does correctly evaluate what we want.

QUESTION 2:

Given the array [2; 3; 8; 6; 1]

2a.

The inversions for the above array would be for any pair of indices where $i < j$ and $A[i] > A[j]$

(1,5) because $2 > 1$

(2,5) because $3 > 1$

(3,4) because $8 > 6$

(3,5) because $8 > 1$

(4,5) because $6 > 1$

2b.

The set of numbers amongst all sets Z, that has the highest amount of inversions is in fact the “reverse sorted array” in which as the indices increase you grow closer to 1, starting with N and ending in 1. Like such:

$\langle n, n-1, \dots, 2, 1 \rangle$

Total inversions for this case will be :

$$(n-1) + n-2 + (\dots) + 1 = \frac{n(n-1)}{2}$$

2c.

There are 3 cases for this algorithm given 3 different extremes:

If the array is already sorted : it will run in $O(n)$ time

If the array is the reverse order array : it will run in $O(n^2)$ time.

Under normal conditions the array will run in a random amount of time and the performance will be affected determined by $n+1$,

2d.

The following code pseudo snippet could be an option for determining that performance

start:

```
Function: int merge_then_count (A, p, q, r);
```

```
Count_num_inversions(A, p, r):  
    if p >= r:  
        return 0  
    q = floor((p + r) / 2)  
    leftInv = Count_num_inversions(A, p, q)  
    rightInv = Count_num_inversions(A, q + 1, r)  
    splitInv = Merge_then_count(A, p, q, r)  
    return leftInv + rightInv + splitInv
```

```
Merge_then_count(A, p, q, r):  
    Let L = A[p ... q]  
    Let R = A[q+1 ... r]  
    i = 0, j = 0, inv = 0  
    for k from p to r:  
        if i >= length(L):  
            A[k] = R[j]  
            j = j + 1  
        else if j >= length(R):  
            A[k] = L[i]  
            i = i + 1  
        else if L[i] <= R[j]:  
            A[k] = L[i]  
            i = i + 1  
        else:  
            A[k] = R[j]  
            inv = inv + (length(L) - i)  
            j = j + 1  
    return inv
```

The runtime for this deterministic algorithm will in fact be the same amount of runtime as a normal merge sort, since it is a modified $\Theta(n \log n)$ algorithm.

QUESTION 3:

3a: Coding Exercise

SEE CPP file attached:

Visual Studio COPY PASTE

```
#include <iostream>
#include <vector>
#include <unordered_map>
#include <algorithm>
using namespace std;

//relative sort function
vector<int> relativeSortArray(const vector<int>& arr1, const vector<int>& arr2) {
    unordered_map<int, int> freq;
    //frequency count loop to use for later
    for (int num : arr1) {
        freq[num]++;
    }

    vector<int> result;
    //using counts from before insert into a new array having been counted.
    for (int num : arr2) {
        if (freq.count(num)) {
            //insertion
            result.insert(result.end(), freq[num], num);
            //move onto next one and remove from addition queue
            freq.erase(num);
        }
    }

    //array 2 needs a secondary vector for holding
    vector<int> extras;
    for (auto pair : freq) {
        // Insert the value "pair.second" times.
        extras.insert(extras.end(), pair.second, pair.first);
    }

    // Sort the extras in ascending order with built in sort function
    sort(extras.begin(), extras.end());
}
```

```

        // Append extras
        result.insert(result.end(), extras.begin(), extras.end());
        return result;
    }

int main() {
    cout<< " EXPECTED BEHAVIOR : Example 1: \n \n Input: arr1 =
[2,3,1,3,2,4,6,7,9,2,19], arr2 = [2,1,4,3,9,6] \n Output: [2,2,2,1,4,3,3,9,6,7,19] \n
Example 2: \n Input: arr1 = [28,6,22,8,44,17], arr2 = [22,28,8,6] \n Output:
[22,28,8,6,17,44]";

    // Example 1 run
    vector<int> arr1_1 = {2, 3, 1, 3, 2, 4, 6, 7, 9, 2, 19};
    vector<int> arr2_1 = {2, 1, 4, 3, 9, 6};
    vector<int> sorted1 = relativeSortArray(arr1_1, arr2_1);

    cout << "\n\nNOW FOR SORTING PROOFS: \n";


    cout << "Sorted arr1 for Example 1: ";
    for (int num : sorted1)
        cout << num << " ";
    cout << "\n";

    // Example 2 run
    vector<int> arr1_2 = {28, 6, 22, 8, 44, 17};
    vector<int> arr2_2 = {22, 28, 8, 6};
    vector<int> sorted2 = relativeSortArray(arr1_2, arr2_2);

    cout << "Sorted arr1 for Example 2: ";
    for (int num : sorted2)
        cout << num << " ";
    cout << "\n";

    return 0;
}

```



The screenshot shows a C++ IDE with a dark theme. The top bar displays navigation icons, a search bar containing 'ESE344', and window management icons. Below the top bar, the file explorer shows two files: 'Hw1_demo.cpp' and 'sort.h'. The main editor window displays the code in 'Hw1_demo.cpp' with line numbers 1 through 8. The code includes standard headers and defines a 'relativeSortArray' function. The terminal at the bottom shows the command to compile and run the program, followed by the expected behavior and the actual output for two examples.

```
#include <iostream>
#include <vector>
#include <unordered_map>
#include <algorithm>
using namespace std;

//relative sort function
vector<int> relativeSortArray(const vector<int>& arr1, const vector<int>& arr2) {
```

PROBLEMS OUTPUT DEBUG CONSOLE **TERMINAL** PORTS

(base) chrisnielsen@Chriss-MacBook-Pro ESE344 % g++ -std=c++17 /Users/chrisnielsen/Documents/ESE344/Hw1_demo.cpp -o Hw1_demo && ./Hw1_demo

EXPECTED BEHAVIOR : Example 1:

Input: arr1 = [2,3,1,3,2,4,6,7,9,2,19], arr2 = [2,1,4,3,9,6]

Output: [2,2,2,1,4,3,3,9,6,7,19]

Example 2:

Input: arr1 = [28,6,22,8,44,17], arr2 = [22,28,8,6]

Output: [22,28,8,6,17,44]

NOW FOR SORTING PROOF:

Sorted arr1 for Example 1: 2 2 2 1 4 3 3 9 6 7 19

Sorted arr1 for Example 2: 22 28 8 6 17 44

(base) chrisnielsen@Chriss-MacBook-Pro ESE344 % g++ -std=c++17 /Users/chrisnielsen/Documents/ESE344/Hw1_demo.cpp -o Hw1_demo && ./Hw1_demo