

# Title

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## Abstract

## 1 Introduction

## 2 Well-formed rewriting steps

We fix a commutative field  $\mathbb{K}$  as well as a well-founded ordered set  $(X, <)$ . We denote by  $\mathbb{K}X$  the vector space spanned by  $X$ : an element  $v \in \mathbb{K}X$  is a finite formal linear combination of elements of  $X$  with coefficients in  $\mathbb{K}$ . In particular, for every  $v \in \mathbb{K}X$ , there exists a unique finite set  $\text{supp}(v) \subset X$ , called the *support* of  $v$ , such that

$$v = \sum_{x \in \text{supp}(v)} \lambda_x x \text{ and } x \in \text{supp}(v) \Rightarrow \lambda_x \neq 0. \quad (1)$$

We denote by  $\text{supp}(v)^c = X \setminus \text{supp}(v)$ . The sum of  $u = \sum \lambda_x x$  and  $v = \sum \mu_x x$  equals  $\sum (\lambda_x + \mu_x) x$  and the product of  $\lambda \in \mathbb{K}$  by  $v$  equals  $\sum (\lambda \lambda_x) x$ . We extend the order  $<$  into the multiset order, still written  $<$ , on  $\mathbb{K}X$ : we have  $u < v$  if for every  $x \in \text{supp}(u) \cap \text{supp}(v)^c$ , there exists  $y \in \text{supp}(v) \cap \text{supp}(u)^c$  such that  $y > x$ .

We fix a set  $R \subseteq X \times \mathbb{K}X$  which represents rewrite rules of the form  $x \xrightarrow{R} r$ . The set  $R$  induces the rewriting relation on  $\mathbb{K}X$ , still written  $\xrightarrow{R}$ , defined as follows:

$$\sum \lambda_x x + v \xrightarrow{R} \sum \lambda_x r_x + v, \quad (2)$$

whenever  $\lambda_x \neq 0$ ,  $x \xrightarrow{R} r_x \in R$  and  $x \notin \text{supp}(v)$ .

**Definition 2.1.** A *local strategy* for  $R$  is the choice, for every  $x \in X$  not minimal for  $<$ , of a rewriting rule  $h_x = x \xrightarrow{R} r_x$  such that  $r_x < x$ .

Suppose chosen such a local strategy  $h$ . Any vector  $v$  can be decomposed in a unique way as  $\sum \lambda_x x + v'$ , where  $y \in \text{supp}(v')$  implies that  $y$  is minimal for  $<$ , and  $x \in \text{supp}(v) \cap \text{supp}(v')^c$  is not. We define a rewriting relation  $\xrightarrow{h}$  as follows:

$$\sum \lambda_x x + v' \xrightarrow{h} \sum \lambda_x r_x + v', \quad (3)$$

where for every  $x$ ,  $h_x = x \xrightarrow{R} r_x$ .

**Definition 2.2.** A vector  $v$  is said to be a  $h$ -normal form if it is a normal form for  $\xrightarrow{h}$ .

**Example 2.3.** Let  $X = \{x, y, z, t\}$ ,  $x \xrightarrow{R} y$ ,  $y \xrightarrow{R} z + t$ ,  $z \xrightarrow{R} y - t$ . Note that this is not terminating since we have the infinite loop  $y \xrightarrow{R} z + t \xrightarrow{R} (y - t) + t = y$ . We choose the order  $x > y > z$ , and the following distinguished rewrite rules:  $h_x = x \xrightarrow{h} y$ ,  $h_y = y \xrightarrow{h} z + t$ . Then the  $R$ -normal forms are the  $\lambda_t t$ , while the  $h$ -normal form are all the  $\lambda_t t + \lambda_z z$ .

**Lemma 2.4.** Let  $v$  be a vector in  $\mathbb{K}X$ . Either  $v$  is minimal for  $<$ , or there exists  $v' < v$  such that  $v \xrightarrow{h} v'$ . In particular,  $h$ -normal forms are precisely the minimal elements of  $\mathbb{K}X$  for  $<$ .

For each  $v \in \mathbb{K}X$ , there exists at most one  $v'$  such that  $v \xrightarrow{h} v'$ , and  $\xrightarrow{h}$  is compatible with the termination order  $<$ . As a consequence, any  $v \in \mathbb{K}X$  is sent by multiple applications of  $\xrightarrow{h}$  to a unique  $h$ -normal form that we denote by  $H(v)$ . This defines a map  $H : \mathbb{K}X \rightarrow \mathbb{K}X$ .

**Proposition 2.5.** The map  $H$  is a linear projector.

*Proof.* The  $h$ -normal forms are closed under sums, so that  $H(H(v)) = H(v)$  for every  $v$ , that is  $H$  is a projector. Moreover, if  $u \xrightarrow{h} u'$  and  $v \xrightarrow{h} v'$ , then we have  $u + v \xrightarrow{h} u' + v'$ . Iterating  $\xrightarrow{h}$ , we get  $H(u + v) = H(H(u) + H(v)) = H(u) + H(v)$ .  $\square$

### 3 A confluence criterion

**Definition 3.1.** We say that  $R$  is  $h$ -confluent if for every rewrite rule  $x \xrightarrow{R} v \in R$ , we have  $x - v \in \ker(H)$ .

**Example 3.2.** Let us take the same example as in 2.3. Then  $H(x) = H(y) = z + t$ , with  $H(z) = z = H(y - t)$ , and so  $R$  is  $h$ -confluent. Replacing the rule  $z \xrightarrow{R} y - t$  by  $z \xrightarrow{R} y$ , we get  $H(z) = z$  and  $H(y) = z + t$ , so  $R$  is not  $h$ -confluent anymore.

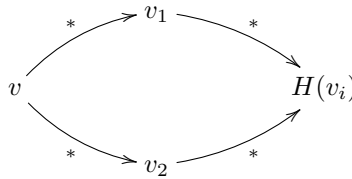
**Proposition 3.3.** If  $R$  is  $h$ -confluent, then  $u \xleftarrow{R}^* v$  if and only if  $u - v \in \ker(H)$ .

*Proof.* The relation  $\xleftarrow{R}^*$  is the closure of  $\xrightarrow{R}$  under transitivity, symmetry and sum. Since the relation  $u - v \in \ker(H)$  is closed under these operations, we get one implication.

Reciprocally, if  $u - v \in \ker(H)$  then by definition of  $H$  we have  $u \xleftarrow{h}^* v$ , and in particular  $u \xleftarrow{R}^* v$ .  $\square$

**Proposition 3.4.** If  $R$  is  $h$ -confluent, then  $\xrightarrow{R}$  is confluent.

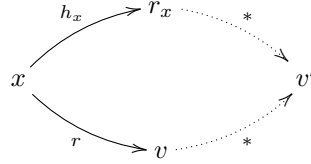
*Proof.* Let  $v, v_1, v_2 \in \mathbb{K}X$  be such that  $v \xrightarrow{R}^* v_i$ , for  $i = 1, 2$ . From Proposition 3.3,  $v_1 - v_2$  belongs to  $\ker(H)$ , that is  $H(v_1) = H(v_2)$ . Hence, we get



□

In Theorem 3.6, we introduce a confluence criterion when  $R$  satisfies 3.1. For that, we assume that  $R$  is equipped with a well-founded order  $\prec$  satisfying the following decreasingness property:

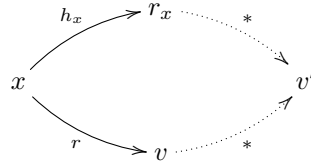
**Definition 3.5.** We say that  $R$  is *locally  $h$ -confluent* if for every non minimal  $x \in X$  and  $r = x \xrightarrow[R]{\quad} v$ , then letting  $h_x = x \xrightarrow[h]{\quad} r_x$ , we have the confluence diagram:



where each rewriting step occurring in the dotted arrows are strictly smaller than  $r$  for  $\prec$ .

**Theorem 3.6.** *If  $R$  is locally  $h$ -confluent, then  $R$  is  $h$ -confluent. In particular,  $\xrightarrow[R]{\quad}$  is confluent.*

*Proof.* We reason by induction on  $r$ . Looking at the square corresponding to  $r$ :



We have  $H(x) = H(r_x)$  by definition of  $H$ , and  $H(r_x) = H(v') = H(v)$  by induction hypothesis, which concludes the proof. □

**Remark 3.7.** Let us show how Newman's Lemma fits as a particular case of our set up. Suppose that  $R$  is locally confluent and terminating. We define the relation  $x > y$  on  $X$  as the transitive closure of the relation  $x \xrightarrow[R]{\quad} u$  and  $y$  appears in  $u$ : this is a well-founded relation. By definition if  $x \in X$  is not minimal for  $>$  then  $x$  is not an  $R$ -normal form. Let us fix an arbitrary rewriting step  $h_x = x \xrightarrow[h]{\quad} r_x$ . By definition of  $>$ ,  $r_x < x$  so that  $h$  is a local strategy. Ordering the rewrite rules by their sources makes  $R$  locally  $h$ -confluent, and Theorem 3.6 finally shows that  $R$  is confluent.