

# dOpt Compulsory Assignment 3

Christian Zhuang-Qing Nielsen\*

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## 1 Proof

Since  $\pi_1$  and  $\pi_2$  are good representations, they must necessary run in polynomial time and since they are equivalent it is possible to reduce one of them from the other and vice versa.  $f$  is a simple function that run in constant time (basically converts '1' to 'yes' and '0' to 'no'). Since:

- $L_1$  is based upon the functionality of  $\pi_1$ ,
- $f$  does the same in both languages,
- $\pi_1 \leq \pi_2$  (per definition of being polynomially equivalent)

$L_1$  must also be reducible from  $L_2$  and vice versa (i.e.  $L_1 \leq L_2$ ). We know that  $L_2 \in P$  as it was restricted by  $\pi_2$  which is bounded by polynomial time and thus in  $P$  as well. We can then finally use the property of downward closure of  $P$  on our languages:

$$L_1 \leq L_2 \wedge L_2 \in P \implies L_1 \in P$$

to make the implication that if  $L_2$  is in  $P$  then  $L_1$  must also necessarily be in  $P$  and vice versa. Similarly if  $L_1$  or  $L_2$  is not in  $P$  (and  $f$  is trivial), then the inner  $\pi$ -function is *not* a good representation and both languages must not be

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\*201504624, christian@czn.dk

in  $P$ , which follows from the closure. This gives us that  $L_1 \in P \Leftrightarrow L_2 \in P$ , if  $\pi_1$  and  $\pi_2$  are good representations.  $\square$