

# dOpt Compulsory Assignment 3 Re-handin

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## 1 Proof

We want to prove that one of two languages are in  $\mathbf{P}$  iff the other one is as well. These languages are decided by the same function, but the input to the deciding function is represented differently (both are still strings  $x \in \{0, 1\}^*$ ) and the two representations are *good* and *polynomially equivalent*.

Starting off, we can assume that one of the languages, lets say  $L_1$ , is in  $\mathbf{P}$ :

$L_1 \in \mathbf{P}$

Under this assumption, the language is decided by a Turing machine  $T_1$  in polynomial time. Now to continue our proof we need to show that  $L_2$  is also in  $\mathbf{P}$ . From our assumption, we know that  $\pi_1$  is a good representation and we want to map it to  $\pi_2$ . However, we cannot assume that  $x \in \pi_2(S)$  just because they are polynomially equivalent.

Since  $\pi_2$  is also supposedly a good representation we can create a Turing machine  $T_2$  that decides the language it represents in polynomial time. Being

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certain of a good representation AND knowing it is polynomially equivalent with  $pi_1$  from the description, we can perform the mapping

$$\pi_1(x) = r_1(\pi_2(x)).$$

Since every input to the mapping was a good representation (and thus decidable by a Turing machine) and it can be mapped back to  $\pi_1$  used by  $L_1$ , which was definitely in  $\mathbf{P}$  from our assumption, which translates back to  $L_2 \in \mathbf{P}$  iff it is decided by a TM  $T_2$  and  $L_1$  is decided by a TM  $T_1$ .  $\square$

## 2 Notes

So the idea is that if  $L_1 \in \mathbf{P}$  we can utilise a mapping between the polynomially equivalent representations. We know that  $\pi_1$  is a good representation from our assumption (and our constructed Turing machine). For the mapping to make sense,  $\pi_2$  must also be a good representation, which is why we construct the second Turing machine that hopefully decides the language  $L_2$ . This can be used to confirm that  $L_2 \in \mathbf{P}$  if and only if our assumptions are true (i.e. that  $\pi_2$  is a good representation and that  $L_1 \in \mathbf{P}$ ).

This was how I interpreted the proof made using Turing machines.