

ArrayLists & LinkedLists Analysis

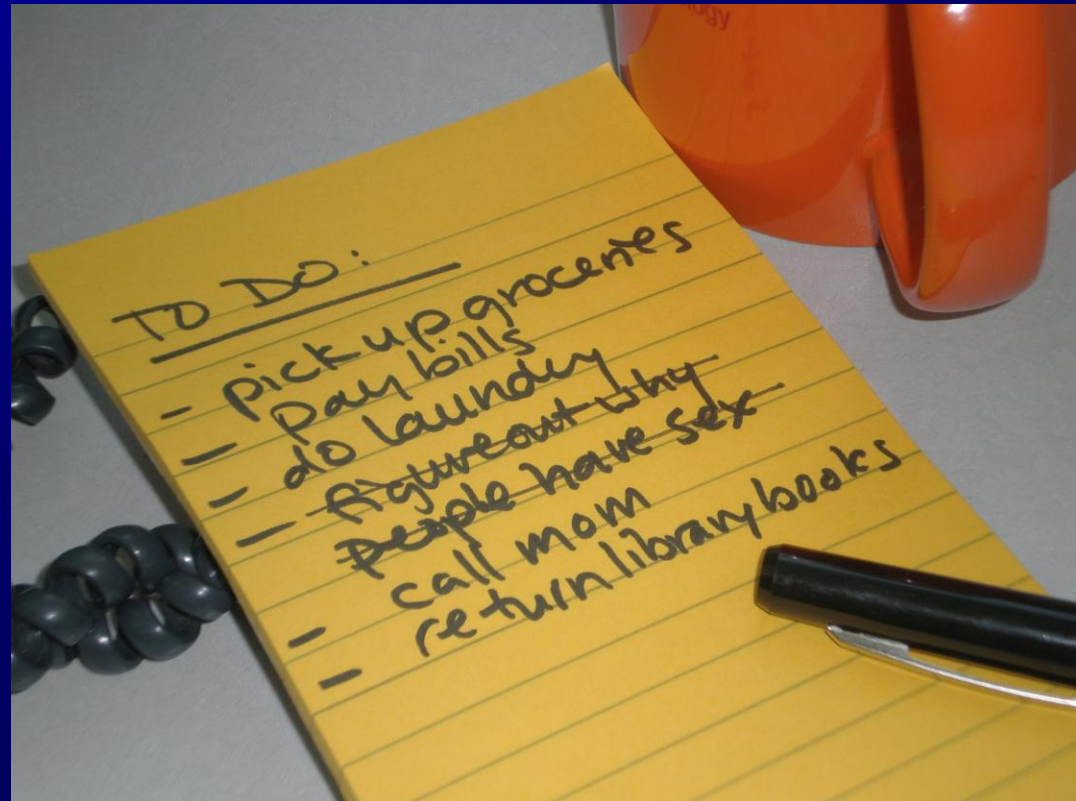
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- ArrayList Class (again)
- Linked Lists (again)
- Algorithm analysis (continuation)
- Amortized runtime

Lists as Data Structures

- *Lists* are ordered collections of data
 - They can contain duplicate values
 - They do have indexes 0 .. length-1



Java's List<E> interface

- List<E> represents a sequence of element that can be accessed by index. It may have duplicate values.
- Java has two implementations of List<E> interface:
 - ArrayList<E>
 - LinkedList<E>
- The only difference, besides the implementation, is the run times of their methods.

Methods in the ArrayList class

- `add(object)`, `add(index, object)`
 - adds element (an object) to the end of the list
 - adds element (an object) at specified index, relocating the rest
- `set(index, object)`
 - changes object at specified element
- `get()`
 - returns an element (an object)
- `remove(index)`, `remove(obj)`
 - deletes an element given a value or an index
- `size()`
- `contains(obj)`
 - searches for obj

ArrayList running times

- `size()`
 - $O(1)$. Why?

- `add(E object)`
 - $O(1)$. Why?

- `remove(int index)`
 - Expected is $O(n)$. Why?
 - Best is $O(1)$. Why?

- How does Java implements the ArrayList Class?

ArrayList running times

- `get(int index)` and `set(int index, E object)`
 - Expected?
 - Worst?
 - Best?

- `add(int index, E object)`
 - Expected?
 - Worst?
 - Best?

Methods in the LinkedList class

- `add(object)`, `add(index, object)`
 - adds element to the end of the list
- `clear()`
 - removes all the elements in the list
- `get(index)`
 - returns an element
- `remove(index)`, `remove(obj)`,
 - deletes an element given a value or an index
- `size()`
- `set(index, obj)`
 - changes the value of a element given an index
- `contains(obj)`
 - searches for obj

Other Methods in a LinkedList class

- `addFirst()`, `addLast()`
 - adds an element
- `getFirst()`, `getLast()`
 - returns an element
- `removeFirst()`, `removeLast()`,
 - deletes an element

LinkedList running times

- `size()`
 - $O(1)$. Why?

- `add(E object)`
 - $O(1)$. Why?

- `remove(int index)`
 - Expected is $O(n)$. Why?
 - Best is $O(1)$. Why?

- How does Java implements the LinkedList Class?

LinkedList running times

- `get(int index)` and `set(int index, E object)`
 - Expected?
 - Worst?
 - Best?

- `add(int index, E object)`
 - Expected?
 - Worst?
 - Best?

- Why do the *ArrayList* class and the *LinkedList* class in Java have the same methods?
 - Is this a coincidence?

Expanding the size of an array

- When the array is full, we create a new array that is larger and copy all the elements to the new array.
 - What's the cost of doing so?

```
private void expand() {  
    int biggerSize = -----;  
    E[] largerArray = (E[]) new Object[biggerSize];  
    for (int i = 0; i < size; i++) {  
        largerArray[i] = values[i];  
    }  
    values = largerArray;  
}
```

Expanding the size of an array

- What size should we make *biggerSize* ?

n	size	copies
1	1	0
2	2	1
3	4	1 + 2 (1 copy when you expand from 1 to 2, and 2 copies when you expand from 2-4)
5	8	1 + 2 + 4
...		
n	n	$1 + 2 + 4 + \dots + n = n*(n+1)/2 = \mathbf{O(n^2)}$

Expanding the size of an array

- Suppose INITIAL_CAPACITY is 10. How many copies would you make to add n elements?

n	size	copies
10	10	0
20	20	10
30	30	10 + 20
40	40	10 + 20 + 30
...		
$10*k$	$10*k$	$10 + 20 + 30 + \dots + 10*(k-1) = 10*(1+2+3+\dots+n/10) =$ $10*(n/10)*((n/10)+1)/2 = \mathbf{O(n^2)}$

It doesn't matter how big INITIAL_CAPACITY is. The cost of n add() operations is $O(n^2)$.

On average, the cost of a single add() operation is $O(n)$.

Better idea ...

- Double the size of the array for every expand

n	size	copies	
1	1	0	
2	2	1	
4	4	$1 + 2 = 3$	$= O(n)$
8	8	$1 + 2 + 4 + 8 = 15 = 2*n-1$	$= O(n)$
16	16	$1 + 2 + 4 + 8 + 16 = 31 = 2*n-1$	$= O(n)$
32	32	$1 + 2 + 4 + 8 + 16 + 32 = 63 = 2*n-1$	$= O(n)$

Amortized runtimes

- Amortized runtime is the expected runtime per operation of a sequence of n operations.
- It is not the same as the average runtime.

Example: $\text{add}(E)$ is $O(n)$ because we next to expand the array quite often. Now, let's add and delete using $\text{remove}(\text{index})$. Then we don't need to expand as often. Thus, the amortize time of a single call to $\text{add}(E)$ is $O(1)$!

Readings

- Java API for ArrayList class
- Java API for LinkedList class
- Java API for List interface

- Watch the video *Linked List Efficiency Analysis*
(Linked List Unit folder)

- Exam I next Thursday Oct. 8
- No homework due next week
- No quiz next week
- Exam review session
 - This weekend
 - TBA on Piazza