

Today in Cryptography (5830)

Review of modes of operation & active attacks

Message authentication

CBC-MAC

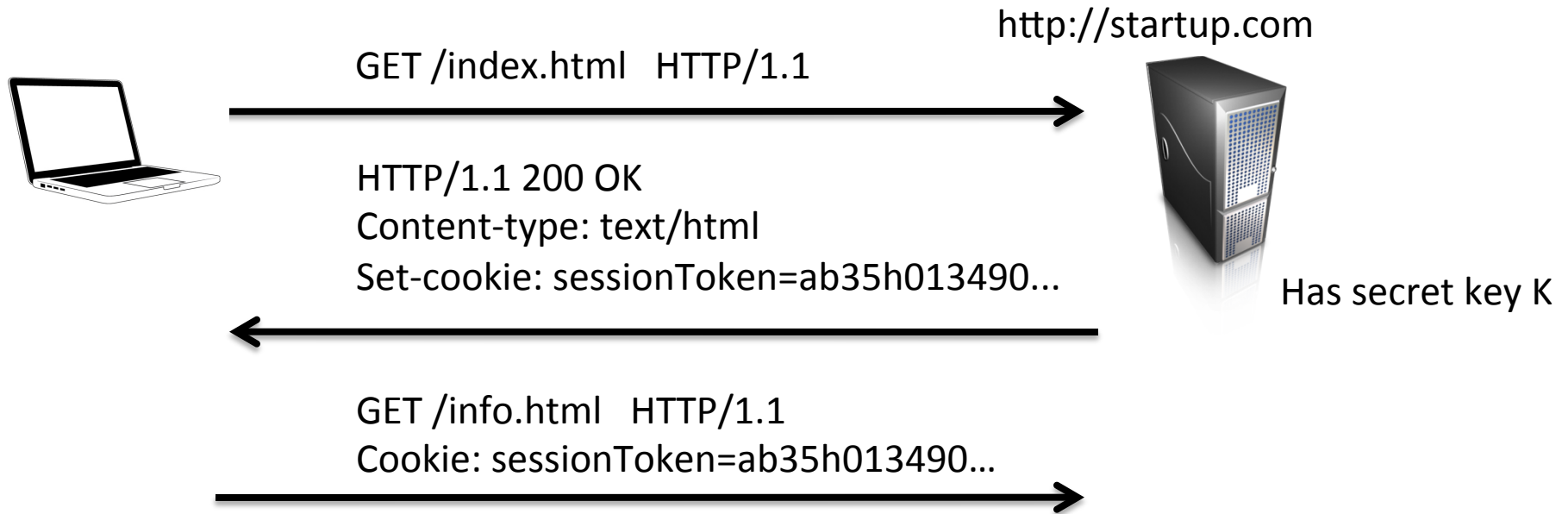
Attacks against bad CBC-MAC implementations

Variable-length secure CBC-MAC

Review

- **Goal:** secure (length-extending) encryption
- What we have so far:
 - Block cipher modes of operation (CBC, CTR)
 - Insecurity against active attacks
 - Bit flip “mauling” attacks against CTR
 - Padding oracle attacks against CBC
- We need another tool:
 authenticity mechanisms

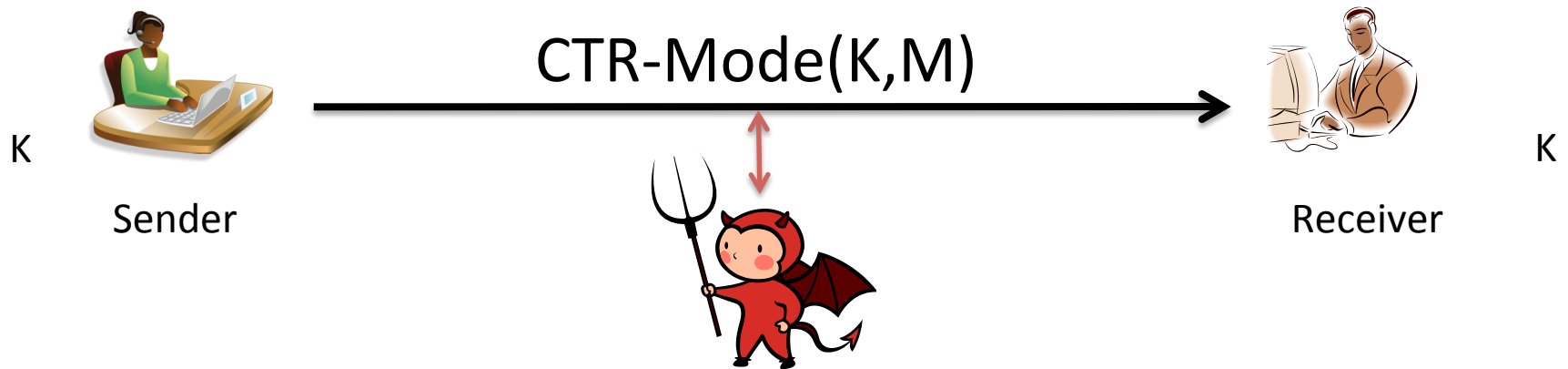
Malleability example: Encrypted cookies



abc35h013490... = CTR-Mode(K, "admin=0")

Malicious client can simply flip a few bits to change admin=1

More generally:

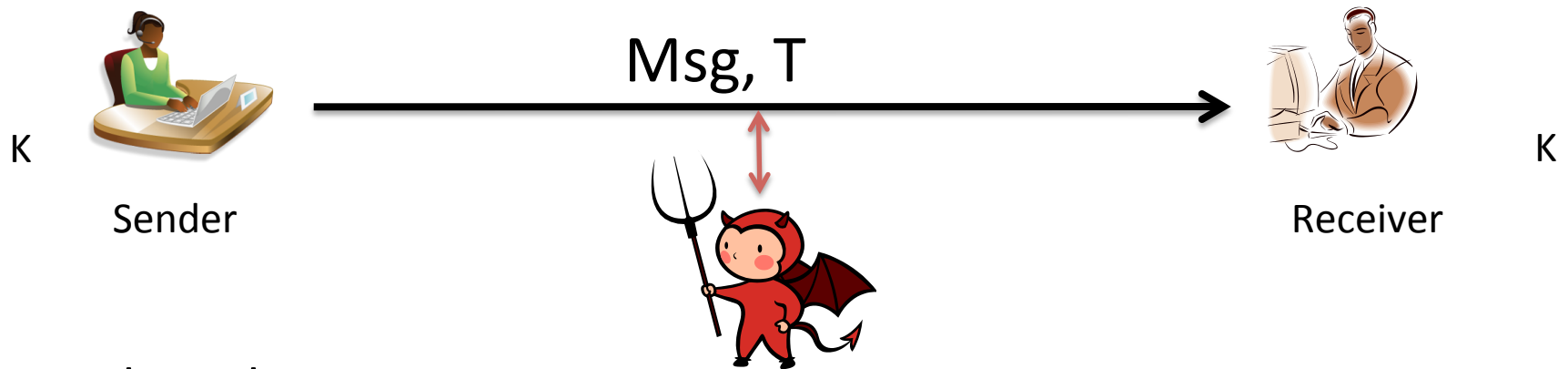


Attacker has read/write access to communications channel

The strategy:

Arrange so that that all bits received can be validated as having come from sender (the person with key **K**)

The tool: Message authentication



Two algorithms:

- (1) $\text{Tag}(K, \text{Msg})$ outputs a tag T
- (2) $\text{Verify}(K, \text{Msg}, T)$ outputs 0/1 (invalid / valid)

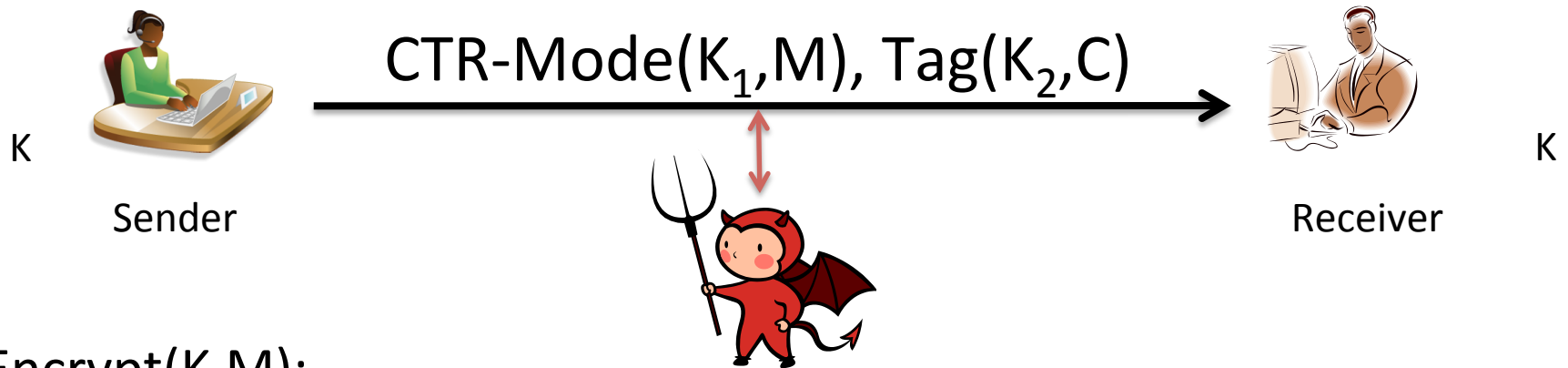
Correctness: $\text{Verify}(K, \text{Msg}, \text{Tag}(K, \text{Msg})) = 1$ always

Security: No computationally efficient attacker can forge tags for a new message even when attacker gets

$(\text{Msg}_1, T_1), (\text{Msg}_2, T_2), \dots, (\text{Msg}_q, T_q)$

for messages of his choosing and reasonably large q .

Composing encryption and authentication



Encrypt(K, M):

Use secret keys K_1 and K_2 . These can be derived from K if needed

$$K_1 = \text{AES}(K, 0^n) \quad K_2 = \text{AES}(K, 1^n)$$

$$C = \text{CTR-Mode}(K_1, M)$$

$$T = \text{Tag}(K_2, C)$$

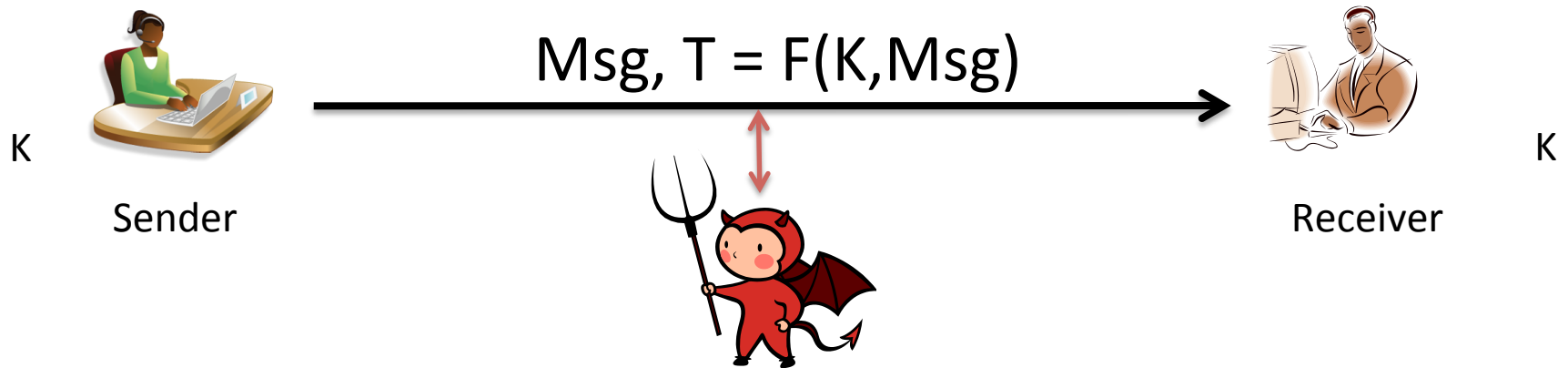
Output $C || T$

Decrypt($K, C || T$)

If $\text{Tag}(K_2, C, T) \neq 1$ then Return error

Return $\text{CTR-Mode}(K_1, C)$

Message authentication using pseudorandom functions (PRFs)

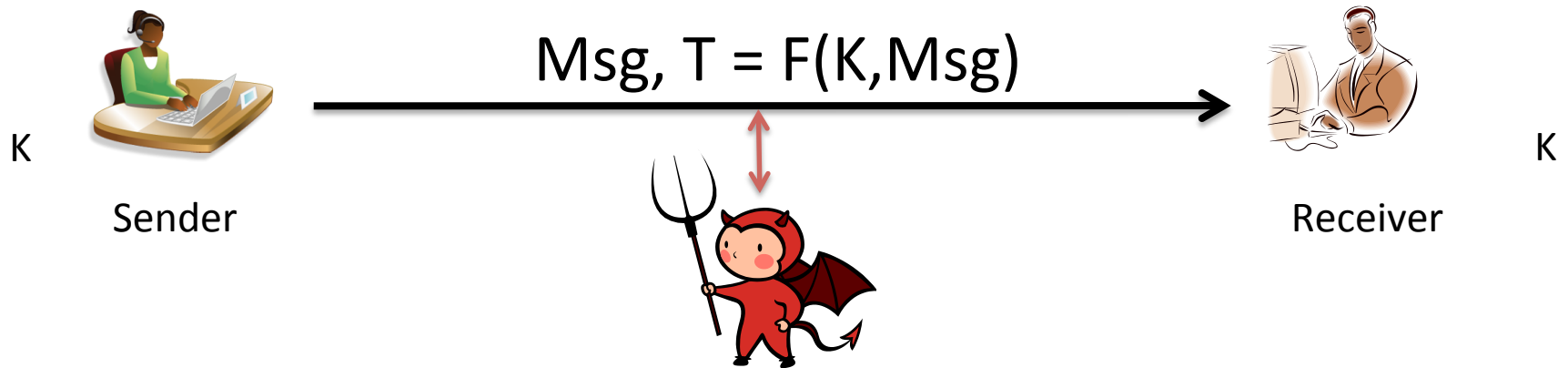


$$Tag(K, M) = F(K, Msg)$$

If F behaves like random function (to those w/o K , this will be secure.

What was example of a good PRF?

Message authentication using pseudorandom functions (PRFs)



$$\text{Tag}(K, M) = F(K, \text{Msg})$$

If F behaves like random function (to those w/o K , this will be secure.

Deterministic message authentication scheme is often called *message authentication code (MAC)* and tag called MAC

What was example of a good PRF?

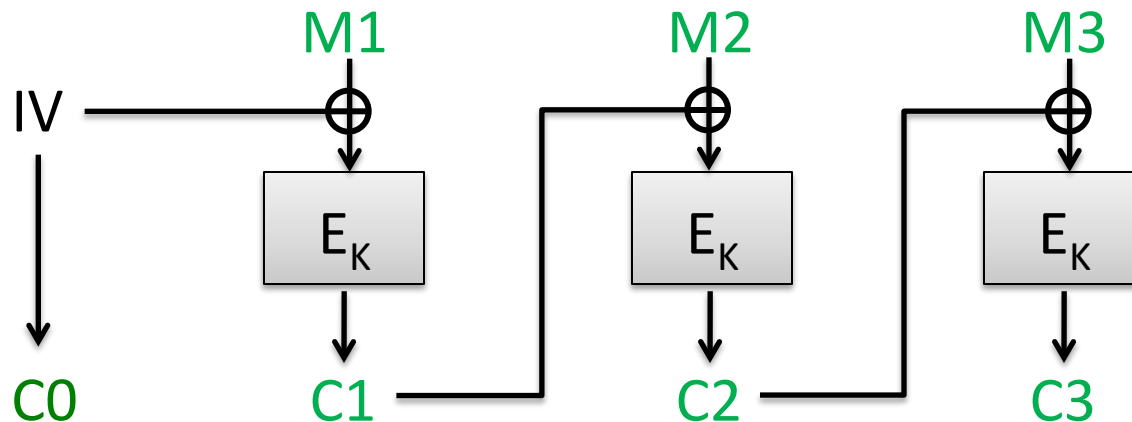
Recall CBC mode

Ciphertext block chaining (CBC)

Pad message M to M_1, M_2, M_3, \dots where each block M_i is n bits

Choose random n -bit string IV

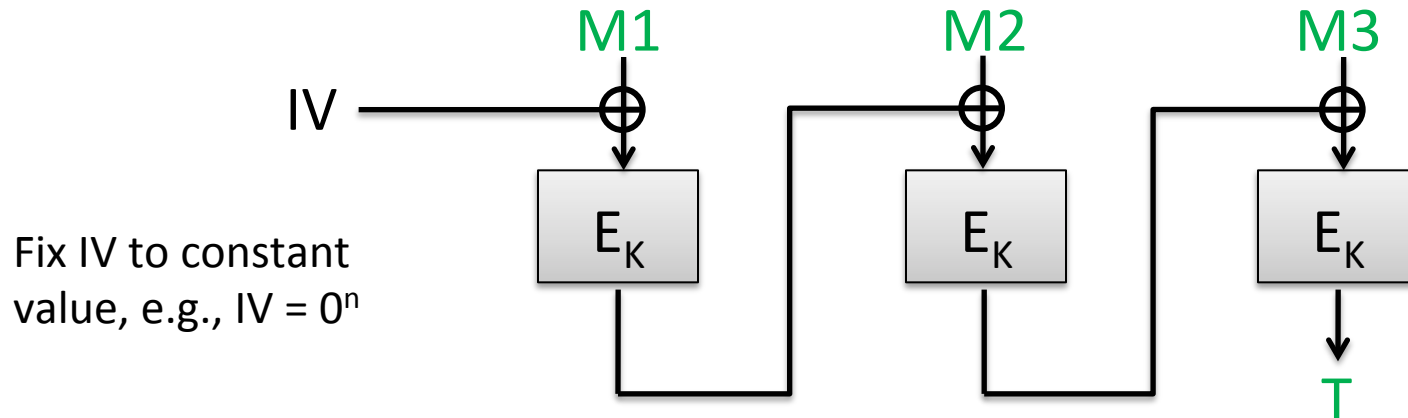
Then:



Can we convert this into variable-message-length PRF?

CBC-MAC

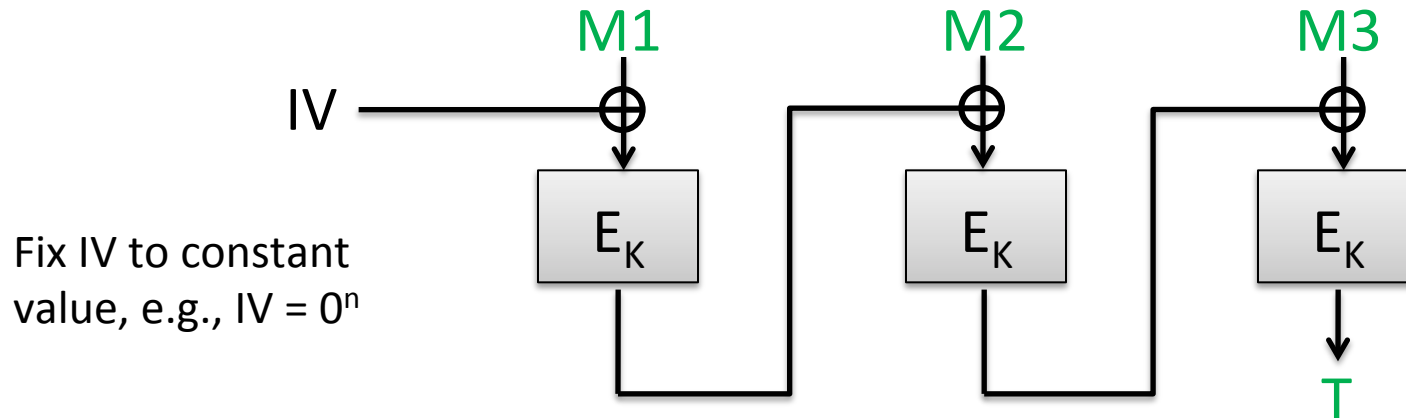
Message authentication code (MAC)



Turns out this is (provably) a good PRF
if only K used only on same-length messages

CBC-MAC

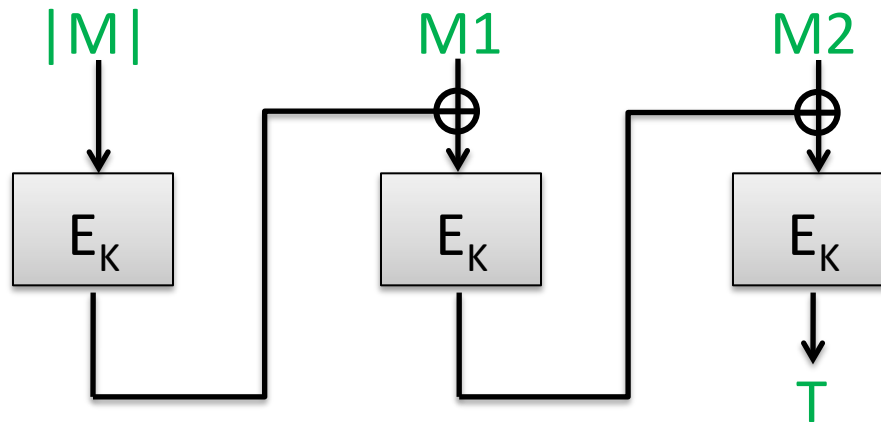
Message authentication code (MAC)



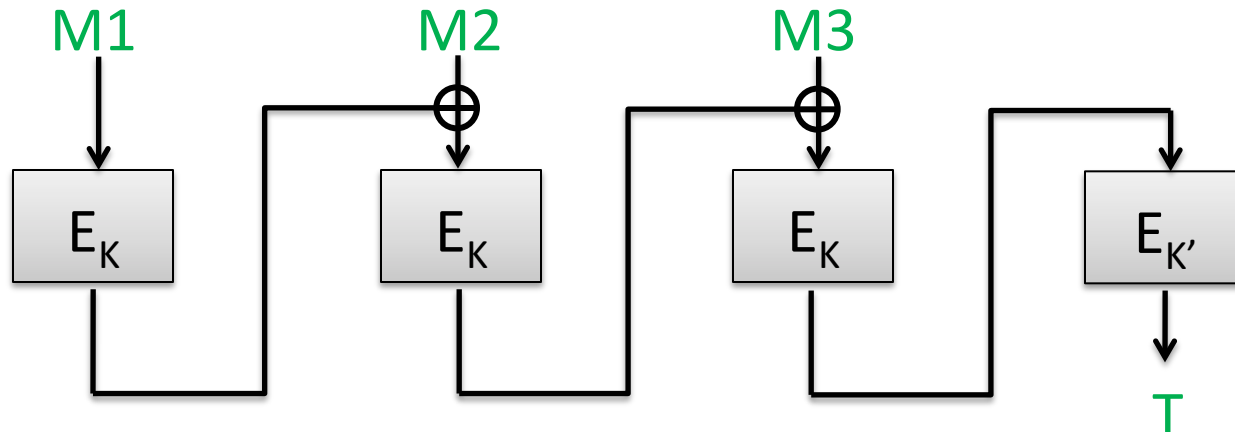
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Variable-message-length CBC-MAC

- Prepend message length



- Encrypted CBC-MAC



Discussion exercise

- We used hazmat interface for CBC mode in HW1 to implement Feistel round function.
- Does this realize a secure implementation of CBC-MAC?