

# I2I-2 Functional Verification Example Exercises

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## 1 Treecount

The following OCaml functions are given:

```
1 type node = Leaf of int | Inner of node * node
2
3 let rec countleaves t = match t with
4   | Leaf _ -> 1
5   | Inner (l, r) -> countleaves l + countleaves r
6
7 let rec countinner t = match t with
8   | Leaf _ -> 0
9   | Inner (l, r) -> 1 + countinner l + countinner r
```

Prove that the equality

```
1 count_leaves t = count_inner t + 1
```

holds.

### Solution

- Case 1:  $t = \text{Leaf}$

$$\begin{aligned} \text{countleaves } t & \stackrel{\text{countleaves}}{=} \text{match } t \text{ with Leaf } \rightarrow 1 \mid \text{Inner } (l, r) \rightarrow \\ & \quad \text{countleaves } l + \text{countleaves } r \\ & \stackrel{\text{def } t, \text{match}}{=} 1 \\ & \stackrel{\text{arith}}{=} 0 + 1 \\ & \stackrel{\text{match, countinner}}{=} \text{countinner } t + 1 \end{aligned}$$

- Case 2:  $t = \text{Inner } (a, b)$

$$\begin{aligned} \text{countleaves } t & \stackrel{\text{countleaves}}{=} \text{match } t \text{ with Leaf } \rightarrow 1 \mid \text{Inner } (l, r) \rightarrow \\ & \quad \text{countleaves } l + \text{countleaves } r \\ & \stackrel{\text{def } t, \text{match}}{=} \text{countleaves } a + \text{countleaves } b \\ & \stackrel{\text{I.H.}}{=} (\text{countinner } a + 1) + (\text{countinner } b + 1) \\ & \stackrel{\text{arith}}{=} (1 + \text{countinner } a + \text{countinner } b) + 1 \\ & \stackrel{\text{match, countinner, def}}{=} \text{countinner } t + 1 \end{aligned}$$

## 2 Fun With Fold

The following OCaml functions are given:

```
1 let rec fl f a l = match l with [] -> a
2   | x::xs -> fl f (f a x) xs
3 let rec fr f l a = match l with [] -> a
4   | x::xs -> f x (fr f xs a)
5 let rec rev_map f l a = match l with [] -> a
6   | x::xs -> rev_map f xs (f x :: a)
```

Prove that the equality

```
1 fl (+) 0 (rev_map (fun x -> x * 2) l []) = fr (fun x a -> a + 2 * x) l 0
```

holds.

### Solution

To prove the given equality, we first generalize and try to prove that

$$\text{fl } (+) \text{ b } (\text{rev\_map } (\text{fun } x \rightarrow x * 2) \text{ l } a) = \text{fr } (\text{fun } x \text{ a} \rightarrow a + 2 * x) (\text{l @ a}) \text{ b}$$

We prove the equality by induction on the length of  $l$

- Base case:  $l = []$

$$\text{fl } (+) \text{ 0 } (\text{rev\_map } (\text{fun } x \rightarrow x * 2) \text{ l } []) \stackrel{\text{rev\_map, def l, match}}{=} \text{fl } (+) \text{ b } a$$

### 3 Even list

The following OCaml functions are given:

```

1 let rec el a lst = match lst with
2   | h::_::t -> el (h::a) t
3   | [h] -> el (h::a) []
4   | [] -> a
5
6 let rec de i a lst = match lst with
7   | h::t -> de (i+1) (if i mod 2 = 0 then h::a else a) t
8   | [] -> a

```

Prove that the equality

```

1 de 0 [] 1 = el [] 1

```

holds.

### Solution

To prove that

$$\text{de } 0 \ [] \ 1 = \text{el } [] \ 1$$

we generalize and try to prove

$$\text{de } (2*i) \ a \ l = \text{el } a \ l$$

Induction on the list  $l$

- Base case 1:  $l = []$

$$\begin{array}{lcl} \text{de } (2*i) \ a \ l & \stackrel{\text{de, def } l, \text{ match}}{=} & a \\ & \stackrel{\text{match, def } l, \text{ el}}{=} & \text{el } a \ l \end{array}$$

- Base case 2:  $l = [h] = h::[]$

$$\begin{array}{lcl} \text{de } (2*i) \ a \ l & \stackrel{\text{de, def } l}{=} & \text{match } h::[] \text{ with } | h::t \rightarrow \text{de } (2*i+1) \ (\text{if } 2*i \bmod 2 = 0 \text{ then } h::a \text{ else } a) \ t \ | \ [] \rightarrow a \\ & \stackrel{\text{match}}{=} & \text{de } (2*i+1) \ (\text{if } 2*i \bmod 2 = 0 \text{ then } h::a \text{ else } a) \ [] \\ & \stackrel{\text{if}}{=} & \text{de } (2*i+1) \ (h::a) \ [] \\ & \stackrel{\text{de, match}}{=} & h::a \\ & \stackrel{\text{match, el}}{=} & \text{el } (h::a) \ [] \\ & \stackrel{\text{match, el}}{=} & \text{el } a \ [h] \\ & \stackrel{\text{def } l}{=} & \text{el } a \ l \end{array}$$

- Induction step:  $l = h1::h2::t$

```

de (2*i) a l      de,def 1      match h1::h2::t with | h::t -> de (2*i+1) (if 2*i
                    ==          mod 2 = 0 then h::a else a) t | [] -> a
                    match
                    ==      de (2*i+1) (if 2*i mod 2 = 0 then h1::a else a)
                    ==      h2::t
if,de,match      de (2*i+2) (if 2*i+1 mod 2 = 0 then h2::h1::a
                    ==      else h1::a) t
                    if
                    ==      de (2*i+2) (h1::a) t
                    I.H.
                    ==      el (h1::a) t
match,el      el a (h1::h2::t)
                    ==
                    def 1
                    ==      el a l)

```

## 4 Bigotree

The following OCaml functions are given:

```
1 type node = Empty | Inner of node * int * node
2
3 let rec insertintree v t = match t with
4   | Empty -> Inner (Empty, v, Empty)
5   | Inner (l, u, r) -> if v > u then
6     Inner (l, u, insertintree v r)
7   else
8     Inner (insertintree v l, u, r)
9
10 let rec totree a lst = match lst with
11   | [] -> a
12   | h::t -> insertintree h (totree a t)
13
14 let rec tolist t = match t with
15   | Empty -> []
16   | Inner (l, v, r) -> tolist l @ [v] @ tolist r
17
18 let rec insert n lst = match lst with
19   | [] -> [n]
20   | h::t -> if n > h then
21     h::(insert n t)
22   else
23     n::h::t
24
25 let rec sort lst = match lst with
26   | [] -> []
27   | h::t -> insert h (sort t)
```

Prove that the equality

```
1 tolist (totree Empty 1) = sort 1
```

holds.

### Solution

There is no solution.

To prove that

$$\text{tolist } (\text{totree Empty } 1) = \text{sort } 1$$

we need to first prove that **Lemma 1**:

$$\text{tolist } t = \text{sort } (\text{tolist } t)$$

holds. We do this by induction on  $t$

- Case 1:  $t = \text{Empty}$

$$\text{tolist } t$$