

## Numerical Linear Algebra

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October 19, 2022



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# Well posed problem, ill conditioned problem, condition Number

- Recap of Previous Lecture
- Perturbations in right hand side and coefficients
- Error sources
- Number systems
- ► Floating point
- ► Q & A

## Recap of Previous Lecture

- ► III conditioned linear system and matrix properties
- Condition number of a matrix
- ► Properties of *Cond(A)*
- Perturbations in right hand side
- Perturbations in coefficients

#### Theorem 5.1

(Right perturbation theorem)
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Then the following holds true:

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$$\frac{1}{cond(A)} \frac{\|\delta b\|}{\|b\|} \leq \frac{\|\delta x\|}{\|x\|} \leq cond(A) \frac{\|\delta b\|}{\|b\|}$$

### Theorem 5.2

(Left perturbation theorem) Suppose

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► Formula (gen.perturbation theorem)

$$\frac{\|\delta x\|}{\|x\|} \leq \frac{\operatorname{cond}(A)}{1 - \operatorname{cond}(A)\frac{\|\delta A\|}{\|A\|}} \left(\frac{\|\delta A\|}{\|A\|} + \frac{\|\delta b\|}{\|b\|}\right)$$

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max. amplification factor for relative errorrs

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Analysis similar to right perturbation theorem on previous slide

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► Hilbert matrix

$$a_{ij} = \frac{1}{i+j-1}$$
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Vandermonde matrix

$$a_{ij} = v_i^{n-j}, v \in \mathbb{R}, \quad cond_2(A_{5 \times 5, v_i = i}) = 2.617 \cdot 10^4$$

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- 1.  $(I M)^{-1}$  exists
- 2.  $\|(I-M)^{-1}\| \le \frac{1}{1-\|M\|}$
- 3.  $(I-M)^{-1} = \sum_{k=0}^{\infty} M^k$

## Proof.

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$$\|(I-M)x\| = \|x-Mx\|$$

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► 
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- $\sum_{k=0}^{j} M^k \sum_{k=0}^{j} M^{k+1} = I M^{j+1}$
- $\blacktriangleright \parallel -M \parallel <1 \Rightarrow \parallel M^{j+1} \parallel \rightarrow_{j \rightarrow \infty} 0$

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$$\blacktriangleright \ \|-M\|<1\Rightarrow \|M^{j+1}\|\to_{j\to\infty} 0$$

$$\downarrow \downarrow$$

$$\|S_j(I-M)-I\| = \|M^{j+1}\| \to_{j\to\infty} 0$$

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$$\Downarrow$$

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THM [?].3 
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Then the following holds true:

$$\frac{\|\delta x\|}{\|x\|} \leq \frac{cond(A)}{1-cond(A)\frac{\|\delta A\|}{\|A\|}} (\frac{\|\delta A\|}{\|A\|} + \frac{\|\delta b\|}{\|b\|})$$

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$$||A^{-1}\delta b|| \le ||A^{-1}|| ||\delta b|| \frac{||Ax||}{||Ax||} \le K(A) \frac{||\delta b||}{||b||} ||x||$$

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### Sources of errors, 1

#### Main sources of errors in numerical computations

- Rounding errors (arithmetic errors)
  - consequence of finite precision arithmetic
  - unavoidable
- Uncertainty in data, may arise in several ways:
  - errors in measuring physical quantities
  - errors from earlier computations
  - from wrong mathematical models of reality
  - **...**
- Truncation errors (discretization errors, approximation errors)

# Sources of errors, 2

- ► Rounding errors (arithmetic errors)
  - consequence of finite precision arithmetic
  - unavoidable

#### Example 5.11

```
In [14]: x=0.001
In [15]: print((1-math.cos(x))/(x*x))
0.49999995832550326
In [16]: x=0.000001
In [17]: print((1-math.cos(x))/(x*x))
0.5000444502911705
In [18]: x=0.00000000001
In [19]: print((1-math.cos(x))/(x*x))
0.0
```

Figure: Arithmetic error, true value is 0.5

▶ Binary, base = 2

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- ► Octal, base = 8

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Example 5.12

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### Example 5.12

#### Decimal system

▶ Base: 10

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### Example 5.12

#### Decimal system

▶ Base: 10

► Symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9

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### Example 5.12

- ▶ Base: 10
- ► Symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
- ► Integer:  $2021 = 2 * 10^3 + 0 * 10^2 + 2 * 10^1 + 1 * 10^0$

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- n-digit number:

$$d_{n-1}d_{n-2}...d_1d_0 = d_{n-1}*10^{n-1} + d_{n-2}*10^{n-2} + ... + d_1*10^1 + d_0*10^0$$

- ▶ Binary, base = 2
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- ► n-digit integral and m-digit fractional part:
  - $d_{n-1}d_{n-2}...d_1d_0.d_{-1}d_{-2}...d_{-m} = d_{n-1} * 10^{n-1} + ... + d_1 * 10^1 + d_0 * 10^0 + d_{-1} * 10^{-1} + d_{-2} * 10^{-2} + ... + d_{-m} * 10^{-m}$

Example 5.13

Binary system

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Binary system ▶ Base: 2

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$$(d_{n-1}d_{n-2}...d_1d_0)_2 = (d_{n-1}*2^{n-1} + d_{n-2}*2^{n-2} + ... + d_1*2^1 + d_0*2^0)_{10}$$

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- ightharpoonup 0.2<sub>10</sub> = 0.00110011[0011]...<sub>2</sub>

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- ► Symbols: 0,1
- ► Integer:  $1010_2 = 1 * 2^3 + 0 * 2^2 + 2^1 + 1 * 2^0 = 10_{10}$
- ► Real:  $10.10_2 = 1 * 2^1 + 0 * 2^0 + 1 * 2^{-1} + 0 * 2^{-2} = 2.5_{10}$
- n-digit number:

$$(d_{n-1}d_{n-2}...d_1d_0)_2 = (d_{n-1}*2^{n-1} + d_{n-2}*2^{n-2} + ... + d_1*2^1 + d_0*2^0)_{10}$$

n-digit integral and m-digit fractional part:

$$(d_{n-1}d_{n-2}...d_1d_0.d_{-1}d_{-2}...d_{-m})_2 = (d_{n-1}*2^{n-1}+...+d_1*2^1+d_0*2^0+d_{-1}*2^{-1}+d_{-2}*2^{-2}+...+d_{-m}*2^{-m})_{10}$$

- ▶ Limitation: can only exactly represent numbers  $x/2^k$
- ightharpoonup 0.2<sub>10</sub> = 0.00110011[0011]...<sub>2</sub>

Octal system

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  - ► Base: 8

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- Can you define number system with arbitrary base r?

### Number systems, 3

- Octal system
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  - ► Symbols: 0, 1, 2, 3, 4, 5, 6, 7
  - ▶ Representing numbers and conversion to other number system?
- Hexadecimal system
  - ▶ Base: 16
  - Symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F
  - Representing numbers and conversion to other number system?
- Can you define number system with arbitrary base r?
- ▶ Is conversion to and from decimal system possible?

► Most computers use binary number system

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Definition 5.14

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#### Definition 5.14

Floating point number  $\tilde{x} = (-1)^s (\sum_{i=1}^m d_{-i}\beta^{-i})\beta^e$ 

 $\triangleright$   $\beta$  - base

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### Example 5.16

 $0.1 \cdot 10^{-2}$  normalized,  $0.01 \cdot 10^{-1}$  - unnormalized

#### Definition 5.17

- ightharpoonup eta base
- ► s -sign,  $d_{-i} \in \{0, ..., \beta 1\}, i = 1, ..., m$
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#### Definition 5.18

IEEE floating point standard

► Single precision: 
$$\begin{pmatrix} \text{sign mantissa exponent} \\ 1 & 23 & 8 \end{pmatrix}$$

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#### Definition 5.18

IEEE floating point standard

- ► Single precision:  $\begin{pmatrix} \text{sign mantissa exponent} \\ 1 & 23 & 8 \end{pmatrix}$
- Double precision:  $\begin{pmatrix} sign & mantissa & exponent \\ 1 & 52 & 11 \end{pmatrix}$

- ▶ Floating point number  $\tilde{x} = (-1)^s (\sum_{i=1}^m d_{-i}\beta^{-i})\beta^e$
- ► IEEE floating point single precision:  $\begin{pmatrix} sign & mantissa & exponent \\ 1 & 23 & 8 \end{pmatrix}$

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#### Definition 5.19

**Overflow**: caused by a floating point number whose exponent is larger then permissible range  $e_{max}$ 

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### Example 5.21

Underflow:

$$\beta = 10, m = 2, e_{min} = -3, e_{max} = 3$$

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### Example 5.21

Underflow:

- $\beta = 10, m = 2, e_{min} = -3, e_{max} = 3$
- ►  $a = 0.3 \cdot 10^{-3}, b = 0.2 \cdot 10^{-3}, a \cdot b = 0.3 \cdot 10^{-3} \cdot 0.2 \cdot 10^{-3} = 0.06 \cdot 10^{-6} = 0.6 \cdot 10^{-7}$

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#### Underflow:

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- ▶ Exponent out of range: -7 < -3

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$$ightharpoonup$$
  $a = 0.4 \cdot 10^2, b = 0.3 \cdot 10^2, a \cdot b = 0.4 \cdot 10^2 \cdot 0.3 \cdot 10^2 = 0.12 \cdot 10^4 = 0.12 \cdot 10^4$ 

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- ► Exponent out of range: 4 > 3

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- ▶ Underflow is a problem: for most system result is 0
- ► Example: Ariane 5 explosion due to overflow

# Floating point system, 6 Rounding and Chopping

### **Rounding and Chopping**

▶ Not all real numbers are machine representable

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in m-digit arithmetics  $d_{-m-1}$  and all other further digits are thrown away

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## Example 5.24

 $\pi = 3.141596$ 

▶ Two-digit arithmetic  $fI(\pi) = 0.31 \cdot 10^{-2}$ 

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## Example 5.24

 $\pi = 3.141596$ 

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### Machine precision, significant numbers

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#### Definition 5.26

Suppose x is real number and  $\tilde{x}$  is its approximation.  $\tilde{x}$  approximates x to s significant digits if s is **largest nonnegative integer** for which relative error satisfies the inequality:

$$\frac{\left|x-\tilde{x}\right|}{\left|x\right|}<5\cdot10^{-s}$$

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$$ightharpoonup x = 1.31, \ddot{x} = 1.3, |x - \ddot{x}| = 0.01, \frac{|x - \ddot{x}|}{|x|} = 0.007635$$

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- $ightharpoonup 7.635 \cdot 10^{-3} < 5 \cdot 10^{-2}$ , agree up to two significant digits

Q & A