# Chinese Remainder Theorem and Powers of an Element

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 ,  $i\neq j\to gcd(n_i,n_j)=1$  (pairwise relatively prime) define mapping

$$cr: \mathbb{Z}_n \to \mathbb{Z}_{n_1} \times \dots \mathbb{Z}_{n_k}$$

$$cr(a) = (a \bmod n_1, \dots, a \bmod n_k)$$

We (and only we) call cr(a) the *chinese remainder representation* of a

Lemma 25. Mapping cr is bijective.

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*Proof.* • by construction of inverse function. Given

$$(a_1,\ldots,a_k)\in\mathbb{Z}_{n_1}\times\ldots\mathbb{Z}_{n_k}$$

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$$m_i = n/n_i$$
 for  $i = 1, \dots, k$ 

$$m_i = n_1 \dots n_{i-1} n_{i+1} \dots n_k$$

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•  $gcd(n_i, m_1) = 1$  and lemma 24  $\rightarrow$ 

 $m_i x \equiv 1 \mod n_i$  has unique solution  $m_i^{-1}$ 

$$c_i = m_i(m_i^{-1} \bmod n_i)$$

$$a \equiv (a_1c_1. + \ldots + a_kc_k) \bmod n$$

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• claim:  $a \equiv a_i \mod n_i$  for all i.

$$i \neq j \rightarrow c_j = m_j (m_j^{-1} \bmod n_j) \equiv 0 \bmod n_i$$

$$c_i \equiv 1 \mod n_i$$



Observe

$$cr(c_i) = (0, ..., 0, 1, 0 ... 0)$$
 with 1 at position i

$$a \equiv a_i c_i \bmod n_i$$

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and

$$a, x \in \mathbb{Z}$$

then

$$x \equiv a \mod n_i \text{ for all } i \in [1:k] \quad \leftrightarrow \quad x \equiv a \mod n$$

# example:

$$a \equiv 2 \mod 5$$

$$a \equiv 3 \mod 13$$

$$a \mod 65 = ?$$

$$a_1 = 2$$
,  $n_1 = m_2 = 5$   
 $a_2 = 3$ ,  $n_2 = m_1 = 13$   
 $13^{-1} \equiv 2 \mod 5$ ,  $5^{-1} \equiv 8 \mod 13$ 

$$c_1 = 13(2 \mod 5) = 26$$
  
 $c_2 = 5(8 \mod 13) = 40$ 

$$a \equiv 2 \cdot 26 + 3 \cdot 40 \mod 65$$
$$\equiv 52 + 120 \mod 65$$
$$\equiv 42 \mod 65$$

Consider group

$$(Z_n^*, \cdot_n)$$
 ,  $Z_n^* = \{ a \in \mathbb{Z}_n : gcd(a, n) = 1 \}$ 

$$a^{(i)} = a^i \mod n$$
 ,  $\langle a \rangle = \{ a^i \mod n : i \in \mathbb{N} \}$  ,  $ord(a) = |\langle a \rangle|$ 

i = 0?

$$e = 1$$
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$$\langle 3 \rangle = \{1, 2, 3, 4, 5, 6\}$$
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recall  $|\mathbb{Z}_n^*| = \varphi(n)$  and lemma 17:

If  $(S, \circ)$  is a finite group with identity e, then  $a^{(|S|)} = e$  for all  $a \in S$ .

#### **Euler's theorem**

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#### Fermat's theorem

**Lemma 30.** *If p is prime, then* 

$$a^{p-1} \equiv 1 \mod p \quad \text{for all } a \in \mathbb{Z}_p^*$$

Proof.

$$p \text{ prime } \rightarrow \varphi(p) = p-1$$

# primitive roots, generators

If  $g \in \mathbb{Z}_n^*$  and  $ord(g) = |Z_n^*|$ , then g is called a *primitive root* or *generator* of  $Z_n^*$ .

e.g. 3 is generator of  $\mathbb{Z}_7^*$  and 2 is not.

 $\mathbb{Z}_n^{*}$  is *cyclic* iff it has a generator.

**Lemma 31.** The only values n > 1, for which  $\mathbb{Z}_n^*$  is cyclic are

 $2, 4, p^e, 2p^e$  for p prime and  $e \in \mathbb{N}$ 

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• assume  $x \equiv y \mod \varphi(n)$ :

$$x = y + k\varphi(n)$$
 with  $k \in \mathbb{Z}$ 

$$g^{x} \equiv g^{y+k\varphi(n)} \mod n$$

$$\equiv g^{y} \cdot (g^{\varphi(n)})^{k} \mod n$$

$$\equiv g^{y} \cdot 1^{k} \mod n \quad \text{(Euler's theorem)}$$

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• assume  $g^x \equiv g^y \mod n$ 

$$|\langle g \rangle| = |\mathbb{Z}_n^*| = \varphi(n)$$

recall lemma 16: with  $|\langle a \rangle| = t$  sequence  $a^{(0)}, a^{(1)}, \dots$  is periodic with period t, i.e. for all i, j:

$$i \equiv j \mod t \leftrightarrow a^{(i)} = a^{(j)}$$

Hence with  $t = \varphi(n)$ :

$$x \equiv y \bmod \varphi(n)$$

# square roots of 1 modulo n:

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$$y - 1 = \forall p \mid (x-1) \land p \mid (x+1)$$

$$(5 - \forall c) p = 2$$

$$5 - \forall c = 2 < 1$$

$$6 \land 1 \land 2 \land 3 \land 4$$

•  $p^{*} \not | (x-1)$ 

$$gcd(p^e, (x-1)) = 1$$
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x is a nontrivial square root of 1 modulo n iff  $x^2 \equiv 1 \mod n$  and  $x \notin \{-1, 1\}$ . e.g. 6 is nontrivial square root of 1 modulo 35.

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nontrivial square root mod n proves that n is not a prime

**goal:** compute  $a^c \mod n$ . Let

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$$i-1 \rightarrow i$$

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$$= a^{2\langle b[k-1:k-1-(i-1)] \rangle + b[k-1-i]} \mod n$$

$$= C(i-1)^2 \cdot a^{b[k-1-i]} \mod n$$

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**complexity:**  $O(\log b)$  arithmetic operations mod n.