

## Cryptography and Network Security

### FAT – F1 Key

1)  $RC_4 - P = [4 \ 5 \ 6 \ 7] \quad K = [2 \ 3 \ 4 \ 5]$

a) Initialization

State Vector,  $S = [0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7]$

Temporary Vector,  $T = [2 \ 3 \ 4 \ 5 \ 2 \ 3 \ 4 \ 5]$

b) Initial Permutation

$i=0, j=0, S = [2 \ 1 \ 0 \ 3 \ 4 \ 5 \ 6 \ 7]$

$i=1, j=2, S = [2 \ 6 \ 0 \ 3 \ 4 \ 5 \ 1 \ 7]$

$i=2, j=6, S = [2 \ 6 \ 0 \ 3 \ 4 \ 5 \ 1 \ 7]$

$i=3, j=2, S = [2 \ 6 \ 3 \ 0 \ 4 \ 5 \ 1 \ 7]$

$i=4, j=2, S = [4 \ 6 \ 3 \ 0 \ 4 \ 5 \ 1 \ 7]$

$i=5, j=0, S = [5 \ 6 \ 3 \ 0 \ 2 \ 5 \ 1 \ 7]$

$i=6, j=0, S = [5 \ 6 \ 3 \ 0 \ 2 \ 4 \ 1 \ 7]$

$i=7, j=5, S = [5 \ 7 \ 3 \ 0 \ 2 \ 1 \ 4 \ 7]$

c) Stream Generation

Iteration 1:  $i=0, j=0, S = [5 \ 6 \ 3 \ 0 \ 2 \ 1 \ 4 \ 7], k=1, C=5$

Iteration 2:  $i=1, j=7, S = [5 \ 6 \ 3 \ 0 \ 2 \ 1 \ 4 \ 7], k=4, C=1$

Iteration 3:  $i=2, j=2, S = [5 \ 6 \ 3 \ 0 \ 2 \ 1 \ 4 \ 7], k=3, C=5$

Iteration 4:  $i=3, j=8, S = [5 \ 6 \ 0 \ 3 \ 2 \ 1 \ 4 \ 7], k=2, C=5$

2a) Extended Euclidean Algorithm

$$17y \equiv 1 \pmod{43}$$

i	Dividend	Divisor	Quotient	Remainder	$x_i$	$y_i$
-1				43	1	0
0				17	0	1
1	43	17	2	9	$x_1 = x_0 - q_1 x_{-1} = 1 - 2 \times 0 = 1$	$y_1 = y_0 - q_1 y_{-1} = 0 - 2 \times 1 = -2$
2	17	9	1	8	$x_2 = x_0 - q_2 x_1 = 0 - 1 \times 1 = -1$	$y_2 = y_0 - q_2 y_1 = 1 - 1 \times (-2) = 3$
3	9	8	1	1	$x_3 = x_1 - q_3 x_2 = 1 - 1 \times (-1) = 2$	$y_3 = y_1 - q_3 y_2 = -2 - 1 \times 3 = -5$
4	8	1	8	0		

$$x = 2, y = -5 = 38$$

$$17 \times 38 \equiv 1 \pmod{43}$$

$$\therefore \boxed{y = 38}$$

2b)

	$a^1$	$a^2$	$a^3$	$a^4$	$a^5$	$a^6$	$a^7$	$a^8$	$a^9$	$a^{10}$	$a^{11}$	$a^{12}$	$a^{13}$	$a^{14}$	$a^{15}$	$a^{16}$
$a=3$	3	9	10	13	5	15	11	16	14	8	7	4	12	2	6	1
$\log_3(a)$	16	14	1	12	5	15	11	10	2	3	7	13	4	9	6	8

a) Digraph cipher or Playfair cipher.

Key: SECURITY

Message: HELLOWORLD

Playfair Matrix

S	E	C	U	R
H	T	Y	A	B
D	F	G	I	K
L	M	N	O	P
Q	V	W	X	Z

Encryption:

HE LX LO WO RL DX  
FU OQ MP XN SP HQ

4.  $p=17, q=11, e=7, M=88$

$$1) n = p \times q = 17 \times 11 = 187$$

$$2) \phi(n) = (p-1)(q-1) = 16 \times 10 = 160$$

$$3) \text{Compute } d' = e^{-1} \bmod 160$$

$$d = 23$$

$$4) \text{Signature generation } S = M^d \bmod n \\ = 88^{23} \bmod 187 \\ S = 11$$

$$5) \text{Verification } M = S^e \bmod n \\ = 11^7 \bmod 187 \\ = 88$$

Here verified

6) ECC -  $E_{11}(1,6)$ ,  $G = (2,7)$ ,  $n_B = 2$ ,  $P_m = (10,9)$ ,  $k = 2$

a) Public key  $P_U = n_B G = 2(2,7) = (2,7) + (2,7)$

$$P_U = (5,2)$$

b) Encryption  $C = [C_1^{C_2}, (M + kP_U)]$

$$C_1 = (5,2)$$

$$C_2 = \text{Infinity}$$