# 第五课 二叉树与图

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## 内容概述

#### 1.5道经典二叉树与图的相关题目

预备知识:二叉树基础知识

例1:路径之和2(medium) (二叉树深搜)

例2:最近的公共祖先(medium)(二叉树性质)

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预备知识:图的基础知识

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2.详细讲解题目解题方法、代码实现

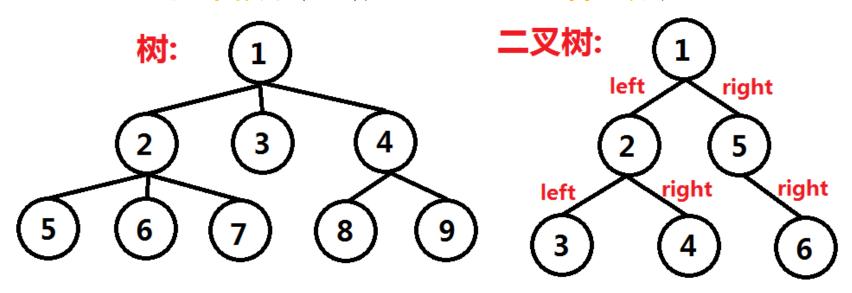
## 预备知识:二叉树定义

#### 树是n(n>=0)个节点的有限集,且这些节点满足如下关系:

- (1)有且仅有一个节点没有父结点,该节点称为树的根。
- (2)除根外,其余的每个节点都有且仅有一个父结点。
- (3)树中的每一个节点都构成一个以它为根的树。

#### **二叉树**在满足树的条件时,满足如下条件:

每个节点最多有两个孩子(子树),这两个子树有左右之分,次序不可颠倒。



```
预备知识:二叉树构造
#include <stdio.h>
                   //二叉树数据结构
struct TreeNode {
                    //数据域val
    int val;
                                                  right
                                         left
   TreeNode *left;
   TreeNode *right; //left, right左右子树指针
   TreeNode(int x) : val(x), left(NULL), right(NULL) {}
};
void preorder print(TreeNode *node, int layer) {
   if (!node) {
                  //正在遍历的节点 //当前节点的层数
       return:
   for (int i = 0; i < layer; i++) {</pre>
       printf("----"); //根据层数,打印相应数量的'-'
   printf("[%d]\n", node->val);
   preorder_print(node->left, layer + 1); //遍历左子树,层数+1
   preorder_print(node->right, layer + 1); //遍历右子树, 层数+1
int main(){
                                                           [1]
   TreeNode a (1);
   TreeNode b(2);
   TreeNode C(5);
   TreeNode d(3);
                                                                          [3]
   TreeNode e(4);
                                                                          [4]
   TreeNode f(6);
   a.left = &b;
   a.right = &c;
   b.left = &d;
   b.right = &e;
   c.right = &f;
   preorder print(&a, 0);
   return 0;
```

## 预备知识:二叉树的深度遍历

```
void traversal(TreeNode *node) {
   if (!node) {
      return;
   }
```

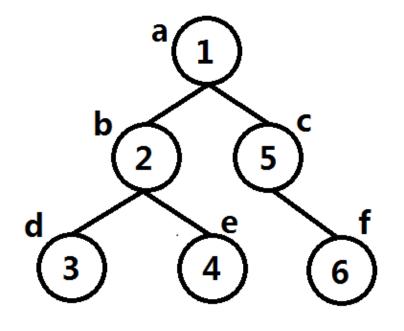
#### 此时访问node称为前序遍历

traversal (node->left);

#### 此时访问node称为中序遍历

traversal(node->right);

#### 此时访问node称为后序遍历



前序遍历: a(1), b(2), d(3), e(4), c(5), f(6)

中序遍历: d(3), b(2), e(4), a(1), c(5), f(6)

后序遍历: d(3), e(4), b(2), f(6), c(5), a(1)

## 预备知识:二叉树的遍历课堂练习

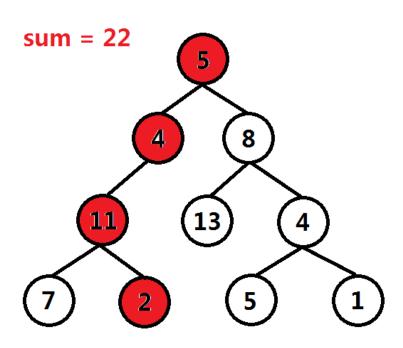
```
[1]
void traversal printl(TreeNode *node,int layer) {
    if (!node) {
                                                                              --[2]
        return;
                                                                                        [3]
    }
    traversal printl(node->left, layer + 1);
                                                                                        Г41
    for (int i = 0; i < layer; i++) {</pre>
                                                                             ---[5]
        printf("----");
                                                                                       -[6]
    }
    printf("[%d]\n", node->val);
    traversal print1(node->right, layer + 1);
void traversal print2(TreeNode *node,int layer){
                                                                                        [3]
    if (!node) {
        return;
                                                                                        [4]
    traversal print2(node->left, layer + 1);
                                                                         [1]
    traversal print2(node->right, layer + 1);
    for (int \bar{i} = 0; i < layer; i++) {
                                                                            ---[5]
        printf("----");
                                                                                      -[6]
    printf("[%d]\n", node->val);
void traversal print3(TreeNode *node,int layer) {
                                                                                        [3]
    if (!node) {
        return;
                                                                                        Г41
    }
                                                                            ----[2]
    for (int i = 0; i < layer; i++) {</pre>
        printf("----");
                                                                                     -- [6]
    printf("[%d]\n", node->val);
                                                                                 [5]
    traversal print3(node->left, layer + 1);
    traversal print3(node->right, layer + 1);
```

## 预备知识:二叉树的遍历讲解

```
void traversal print1(TreeNode *node,int layer) {
                                                                          [1]
    if (!node) {
                                                                                -[2]
        return;
                                                                                      -[3]
    traversal print1(node->left, layer + 1);
                                                                                       [4]
    for (int i = 0; i < layer; i++) {
                                                                             ---[5]
        printf("----");
                                                                                       -[6]
    printf("[%d]\n", node->val);
    traversal print1(node->right, layer + 1);
void traversal print2(TreeNode *node, int layer) {
                                                                                       131
    if (!node) {
        return;
                                                                                       -[4]
    traversal print2(node->left, layer + 1);
    traversal print2(node->right, layer + 1);
                                                                         [1]
    for (int i = 0; i < layer; i++) {</pre>
                                                                                - [5]
        printf("----");
                                                                                       [6]
    printf("[%d]\n", node->val);
void traversal print3(TreeNode *node,int layer) {
    if (!node) {
        return;
                                                                                       Г41
    for (int i = 0; i < layer; i++) {
        printf("----");
                                                                                      -[6]
    printf("[%d]\n", node->val);
                                                                                -[5]
    traversal print3(node->left, layer + 1);
    traversal print3(node->right, layer + 1);
```

#### 例1:路径之和2

给定一个二叉树与整数sum,找出所有从根节点到叶结点的路径,这些路径上的节点值累加和为sum。

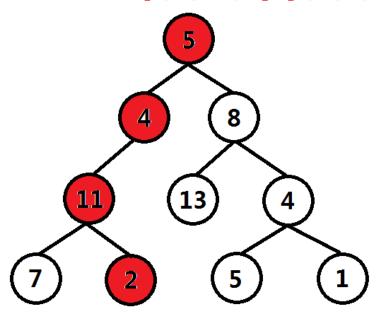


#### 选自 LeetCode 113. Path Sum II

https://leetcode.com/problems/path-sum-ii/description/

难度:Medium

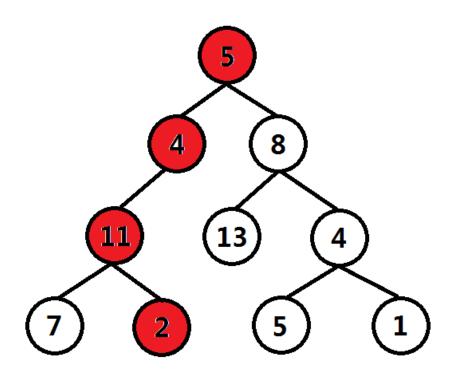
## 例1:思考



#### 思考:

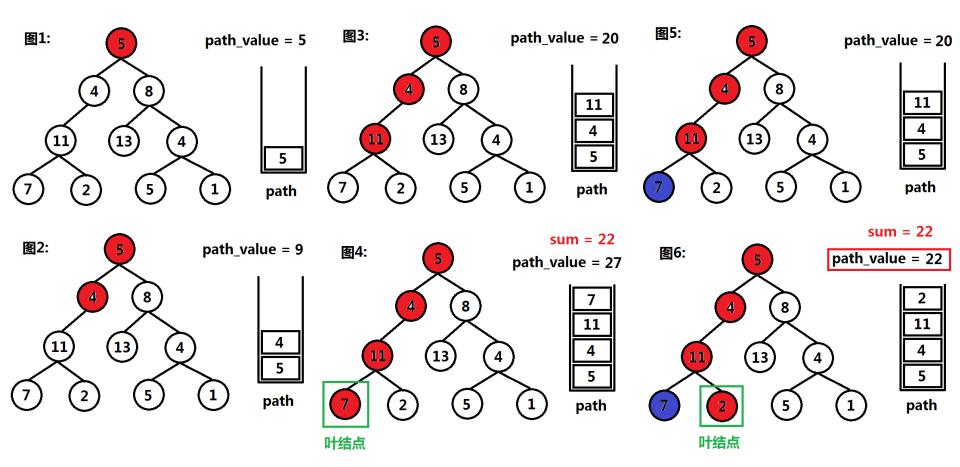
- 1.使用何种数据结构存储遍历路径上的节点?
- 2.在树的前序遍历时做什么?后序遍历时做什么?
- 3.如何判断一个节点为叶结点?当遍历到叶结点时应该做什么?

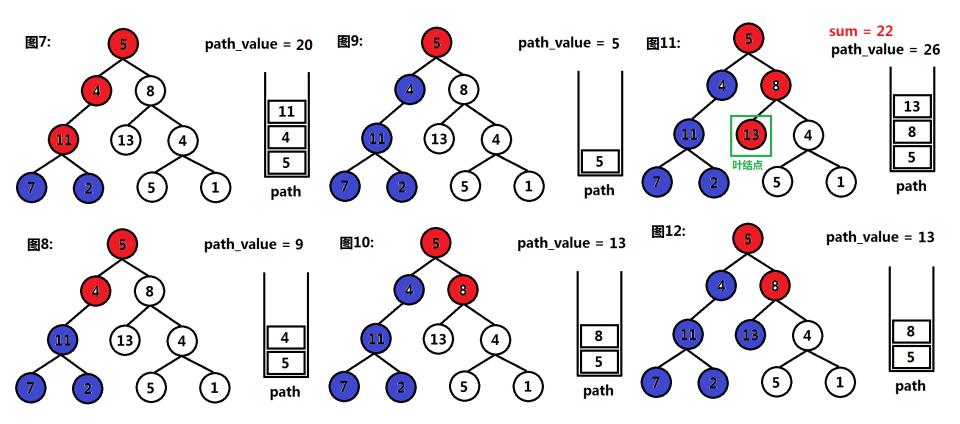
- 1.从根节点深度遍历二叉树,先序遍历时,将该节点值存储至**path栈**中(vector实现),使用 path\_value<mark>累加</mark>节点值。
- 2.当遍历至**叶结点**时,检查**path\_value值**是否为sum,若为sum,则将path **push** 进入result结果中。
- 3.在后续遍历时,将该节点值从path栈中弹出,path\_value减去节点值。

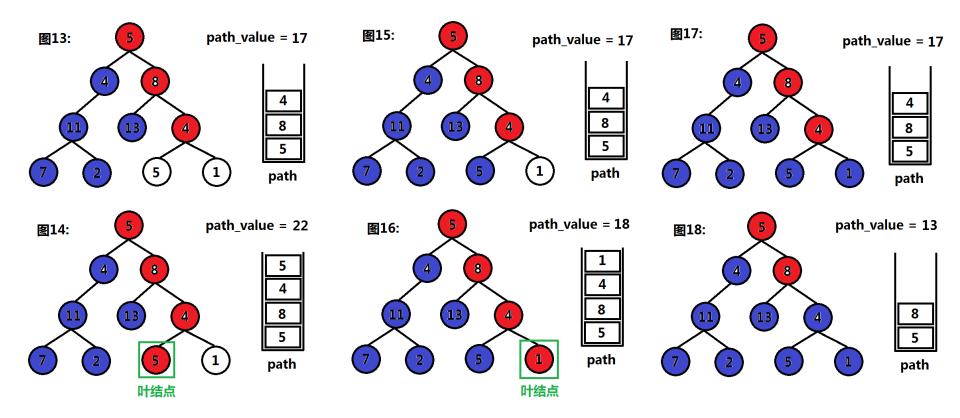


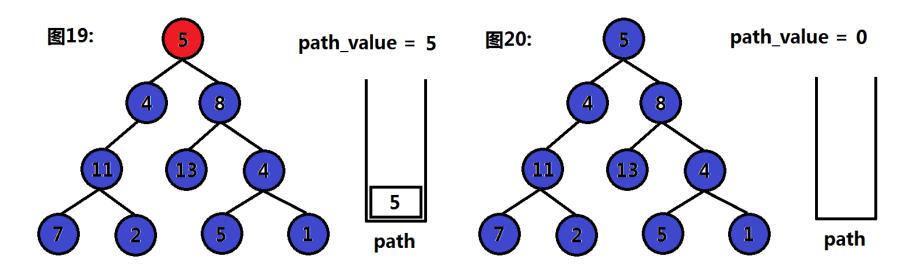
path\_value = 22 path\_value ?? sum











#### 所有路径:

[5, 4, 11, 7] path\_value = 27

[5, 4, 11, 2] path\_value = 22

[5, 8, 13] path\_value = 26

[5, 8, 4, 5] path\_value = 22

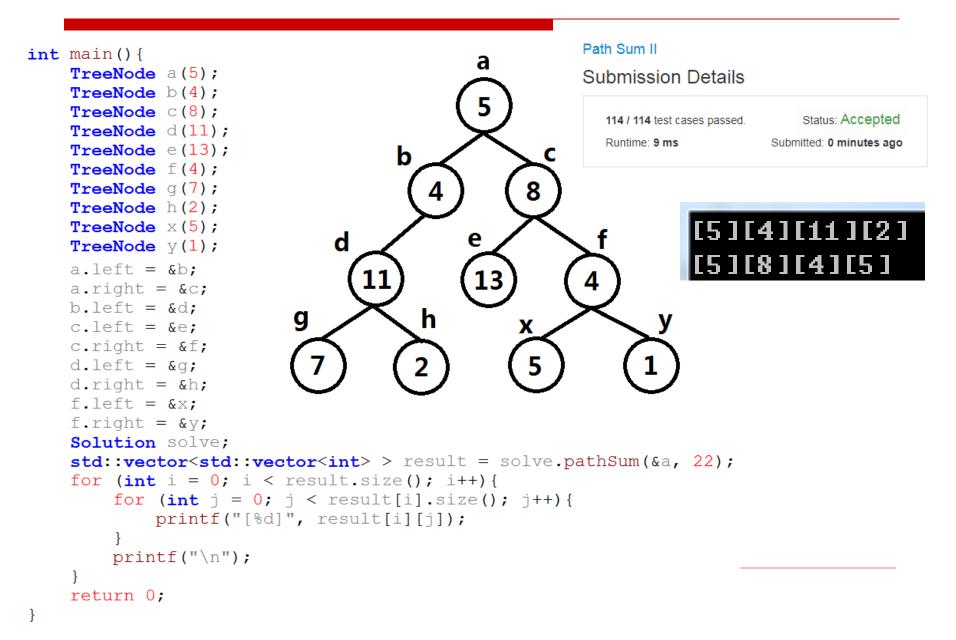
[5, 4, 11, 1] path\_value = 21

```
class Solution {
public:
   std::vector<std::vector<int> > pathSum(TreeNode* root, int sum) {
       std::vector<std::vector<int> > result; //存储满足条件路径的数组
       std::vector<int> path; //路径栈与路径值
       int path value = 0;
       preorder (root, path_value, sum, path, result); 例1:课堂练习
       return result;
private:
    void preorder(TreeNode *node, int &path value, int sum,
               std::vector<int> &path,
               std::vector<std::vector<int> > &result) {
       if (!node) {
           return;
                    //遍历一个节点即更新一次路径值
                                               3分钟填写代码,
       path value += node->val;
                                               有问题随时提出!
       if
           result.push back(path);
                                 //满足 ?? 条件时,将path添加至结果数组
       preorder(node->left, path value, sum, path, result);
       preorder(node->right, path value, sum, path, result);
       path.pop_back(); //遍历完成后,将该节点送路径栈中弹出
};
```

## 例1:实现

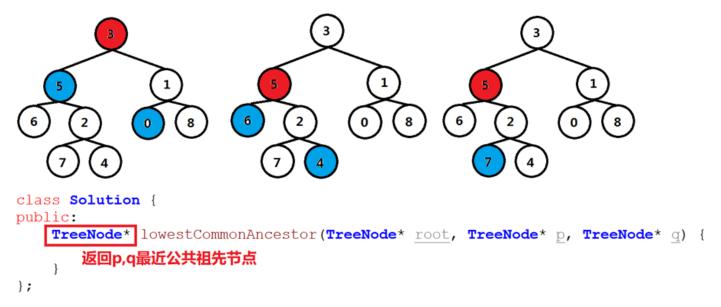
```
sum = 22
                                                                      path_value = 22
void preorder (TreeNode *node, int &path value, int sum,
                                                              11
           std::vector<int> &path,
           std::vector<std::vector<int> > &result) {
   if (!node) {
       return;
                 //遍历一个节点即更新一次路径值
   path value += node->val;
     path.push_back(node->val);
                                                            path
          !node->left && !node->right && path value == sum ) {
   if
       result.push_back(path); //满足??条件时,将path添加至结果数组
   preorder(node->left, path value, sum, path, result);
   preorder(node->right, path value, sum, path, result);
                                                                                   5
     path_value -= node->val;
    path.pop_back(); //遍历完成后,将该节点送路径栈中弹出
                                                                      叶结点
```

## 例1:测试与leetcode提交结果



### 例2:最近的公共祖先

已知二叉树,求二叉树中给定的两个节点的最近公共祖先。 最近公共祖先: 两节点v与w的最近公共祖先u,满足在树上最低(离根最远),且v,w两个节点都是u的子孙。



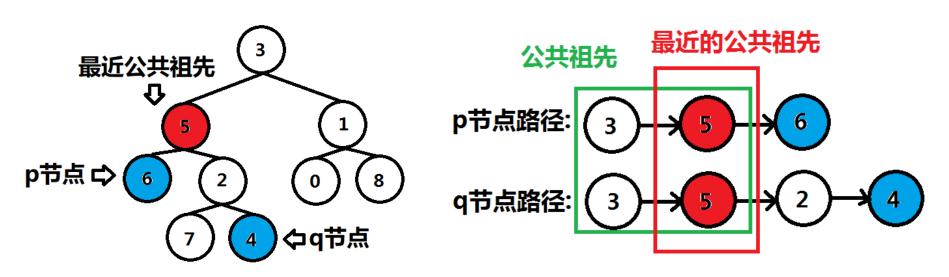
#### 选自 LeetCode 236. Lowest Common Ancestor of a Binary Tree

https://leetcode.com/problems/lowest-common-ancestor-of-a-binary-tree/description/

难度:Medium

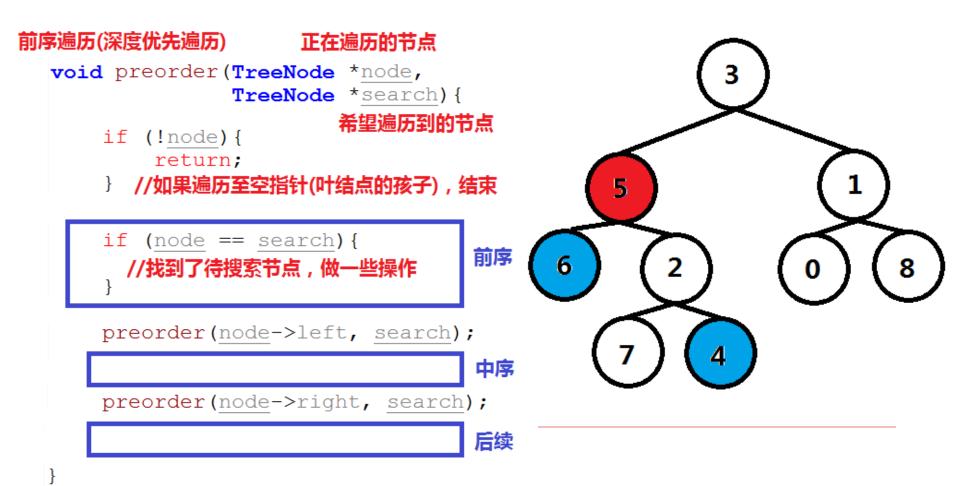
## 例2:思考与分析

- 1.两个节点的公共祖先一定在从根节点,至这两个节点的路径上。
- 2.由于求公共祖先中的最近公共祖先,那么即同时出现在这两条路径上的**离根节点最远**的节点(或离两个最近)。
- 3.最终算法即:求p节点路径, q节点路径, 两路径上最后一个相同的节点。



## 例2:求根节点至某节点路径(深度搜索)

- 1.从根节点遍历(搜索)至该节点,找到该节点后就结束搜索。
- 2.将遍历过程中遇到的节点按照顺序存储起来,这些节点即路径节点。



## 例2:求根节点至某节点路径(栈存储路径)

图2

前序(后序)

中序

图3

图4

前序(后序)

图1

#### 6 记录节点p路径的过程: 深度遍历节点(3) 深度遍历节点(2) 深度遍历节点(1) 深度遍历节点(2) (2)前序 前序 前序(后序) 中序 (4)(3)图6 图7 图8 图5 (5) 2 5 5 5 深度遍历节点(4) 深度遍历节点(6) 深度遍历节点(5) 深度遍历节点(4)

前序

#### 例2:求根节点至某节点路径(实现,课堂练习)

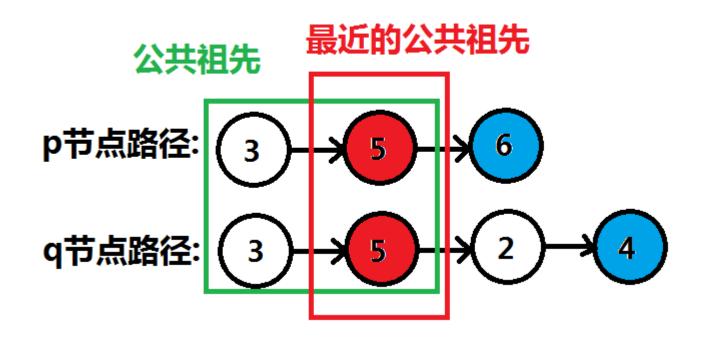
```
void preorder (TreeNode* node, //正在遍历的节点
           TreeNode *search, //待搜索节点
           std::vector<TreeNode*> &path, //遍历时的节点路径栈
           std::vector<TreeNode*> &result, //最终搜索到节点search的路径结果
           int &finish) { //记录是否找到节点search的变量 , 未找到时是0 , 找到为1
                  」){ //当node为空或已找到search节点直接返回 , 结束搜索
  if (!node | | 1
      return:
  path.push_back(node); //先序遍历时,将节点压入path栈
  if (node == search) {
      finish = 1; //当找到search节点后,标记finish变量
  preorder(node->left, search, path, result, finish); //深度遍历node左孩子
  preorder(node->right, search, path, result, finish); //深度遍历node右孩子
                     3分钟时间填写代码,
                      有问题随时提出!
```

## 例2:求根节点至某节点路径(实现)

```
void preorder (TreeNode* node, //正在遍历的节点
            TreeNode *search, //待搜索节点
            std::vector<TreeNode*> &path, //遍历时的节点路径栈
            std::vector<TreeNode*> &result, //最终搜索到节点search的路径结果
            int &finish) { //记录是否找到节点search的变量 , 未找到时是0 , 找到为1
   if (!node || finish ) {
                      //当node为空或已找到search节点直接返回,结束搜索
      return:
   path.push back(node); //先序遍历时,将节点压入path栈
   if (node == search) {
                     //当找到search节点后,标记finish变量
      finish = 1;
      result=path
                     //将当前的path存储到result中
  preorder(node->left, <u>search</u>, <u>path</u>, <u>result</u>, finish); //深度遍历node左孩子
  preorder(node->right, search, path, result, finish); //深度遍历node右孩子
    path.pop_back()
                     //结束遍历node时,将node节点弹出path栈
```

## 例2:求两路径上最后一个相同的节点

- 1.求出较短路径的长度n。
- 2.同时遍历p节点的路径与q节点的路径,遍历n个节点,最后一个发现的相同节点,即最近公共祖先。



# 例2:整体代码(实现,课堂练习)

```
class Solution {
public:
   TreeNode* lowestCommonAncestor(TreeNode* root, TreeNode* p, TreeNode* q) {
       std::vector<TreeNode*> path; //声明遍历用的临时栈
       std::vector<TreeNode*> node _p_path; //存储p节点路径
       std::vector<TreeNode*> node q path; //存储q节点路径
       int finish = 0; //记录是否完成搜索的变量finish
       path.clear();
                     //清空path、finish , 计算q节点路径
       finish = 0:
       preorder(root, q, path, node q path, finish);
                                                          中时间填写代码,
       int path len = 0; //较短路径的长度
                                                   有问题随时提出!
       if
          path len = node p path.size();
       else{
          path len = node q path.size();
       TreeNode *result = 0; 同时遍历根到p,q两个节点的的路径上的节点
       for (int i = 0; i < path len; i++) {</pre>
                            3
           if
              result = node p path[i];
       return result;
};
```

# 例2:整体代码(实现)

```
class Solution {
public:
    TreeNode* lowestCommonAncestor(TreeNode* root, TreeNode* p, TreeNode* q) {
       std::vector<TreeNode*> path; //声明遍历用的临时栈
       std::vector<TreeNode*> node p_path; //存储p节点路径
       std::vector<TreeNode*> node q path; //存储q节点路径
       int finish = 0; //记录是否完成搜索的变量finish
          preorder(root, p, path, node_p_path, finish);
        path.clear();
                      //清空path、finish , 计算q节点路径
        finish = 0:
       preorder(root, q, path, node q path, finish);
       int path len = 0; //较短路径的长度
            node p path.size() < node q path.size()</pre>
       if
           path len = node p path.size();
       else{
           path len = node q path.size();
       TreeNode *result = 0; 同时遍历根到p,q两个节点的的路径上的节点
       for (int i = 0; i < path len; i++) {</pre>
              node_p_path[i] == node_q_path[i]
           i f
               result = node p path[i];
       return result:
};
```

## 例2:测试与leetcode提交结果

```
int main(){
    TreeNode a (3);
                     构造样例中的树
    TreeNode b (5);
    TreeNode C(1);
    TreeNode d(6);
    TreeNode e(2);
    TreeNode f(0);
    TreeNode x(8);
    TreeNode y(7);
    TreeNode Z(4);
    a.left = &b;
    a.right = &c;
    b.left = &d;
    b.right = &e;
    c.left = &f;
    c.right = &x;
    e.left = &v;
    e.right = &z;
```

```
lowestCommonAncestor = 3
lowestCommonAncestor = 5
lowestCommonAncestor = 5
请按任意键继续. . .
```

Lowest Common Ancestor of a Binary Tree

#### Submission Details

```
31 / 31 test cases passed. Status: Accepted
Runtime: 19 ms Submitted: 1 minute ago
```

```
Solution solve;
TreeNode *result = solve.lowestCommonAncestor(&a, &b, &f);
printf("lowestCommonAncestor = %d\n", result->val);
result = solve.lowestCommonAncestor(&a, &d, &z);
printf("lowestCommonAncestor = %d\n", result->val);
result = solve.lowestCommonAncestor(&a, &b, &y);
printf("lowestCommonAncestor = %d\n", result->val);
return 0;
```

### 例3:二叉树转链表

给定一个二叉树,将该二叉树就地(in-place)转换为单链表。单链表

中节点顺序为二叉树前序遍历顺序。 //树节点的数据结构 struct **TreeNode** { //单链表仍使用该数据结构 , int val; 即left=NULL,right = next TreeNode \*left; TreeNode \*right; TreeNode(int x) : val(x), left(NULL), right(NULL) {} }; class Solution { public: void flatten(TreeNode \*root) } ;

#### 选自 LeetCode 114. Flatten Binary Tree to Linked List

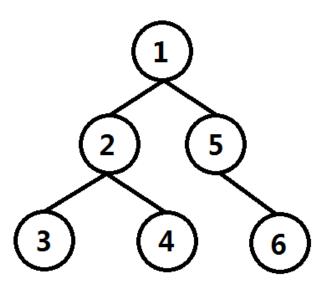
https://leetcode.com/problems/flatten-binary-tree-to-linked-list/description/

难度:Medium

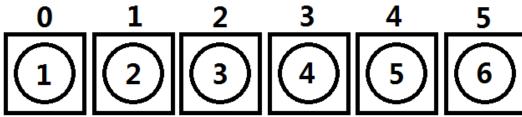
### 例3:思考

前序遍历二叉树,将节点指针push进入vector,顺序遍历vector中的节点,链接相邻两节点,形成链单链表。(投机取巧) 该方法虽然可通过题目,但不满足就地(in-place)转换的条件。

若就地(in-place)转换应该如何做?



vector:



按顺序遍历vector,将前面的节点左指针置空 右指针与后面的节点相连。

### 例3:方法1课堂练习

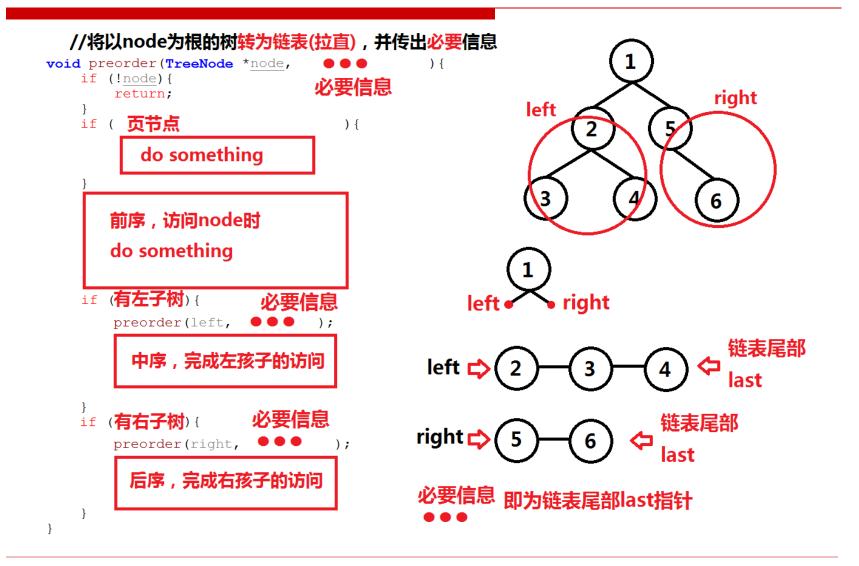
```
#include <vector>
                                            3分钟填写代码,
class Solution {
public:
                                             有问题随时提出!
   void flatten(TreeNode *root) {
        std::vector<TreeNode *> node vec;
       preorder(root, node vec);
       for (int i = 1; i < node vec.size(); i++){</pre>
private:
   void preorder(TreeNode *node, std::vector<TreeNode *> &node vec) {
        if (!node) {
           return:
        preorder(node->left, node vec);
        preorder(node->right, node vec);
};
```

## 例3:方法1实现

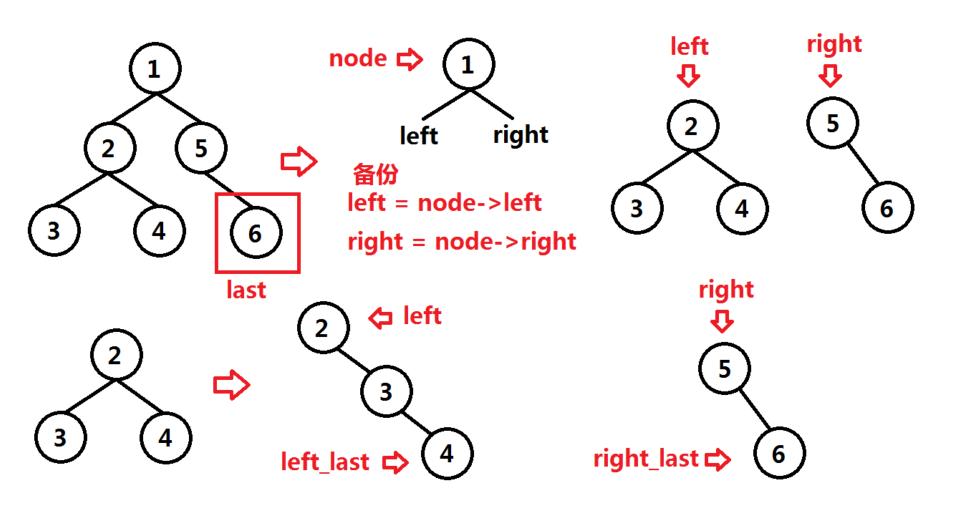
```
left
#include <vector>
                                              right
class Solution {
public:
                                                  NULL
   void flatten(TreeNode *root) {
        std::vector<TreeNode *> node vec;
        preorder(root, node vec);
        for (int i = 1; i < node vec.size(); i++){</pre>
              node_vec[i-1]->left = NULL;
            node_vec[i-1]->right = node_vec[i];
private:
    void preorder(TreeNode *node, std::vector<TreeNode *> &node vec) {
        if (!node) {
            return:
              node_vec.push_back(node);
        preorder(node->left, node vec);
        preorder(node->right, node vec);
};
```

node\_vec[i-1] node\_vec[i]

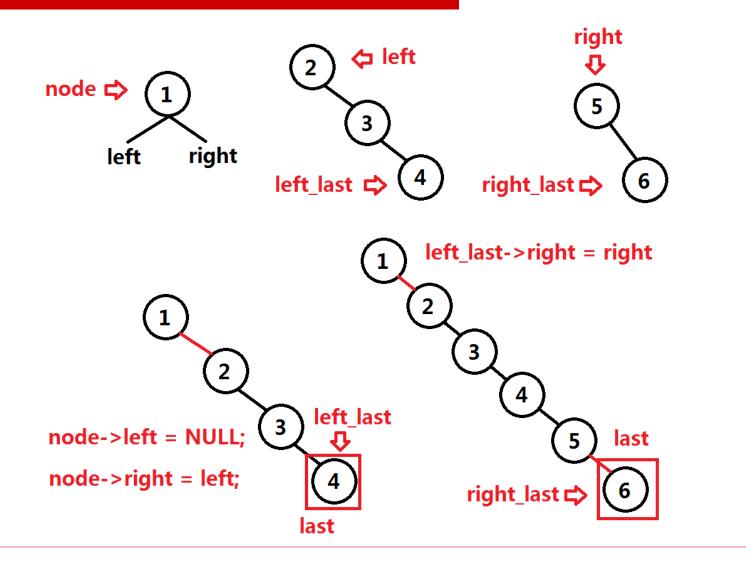
# 例3:算法思路(方法2整体)



# 例3:算法思路(拆解并解决子问题)



## 例3:算法思路(解决当前问题)



```
例3:方法2课堂练习
class Solution {
public:
   void flatten(TreeNode *root) {
      TreeNode *last = NULL;
                          //当前子树的先序遍历的
      preorder(root, last);
                          最后一个节点, 传引用会传出
              // 当前的节点
private:
   void preorder(TreeNode *node, TreeNode *&last) {
      if (!node) {
         return;
      if
         return;
                              //备份左右指针
      TreeNode *left = node->left;
      TreeNode *right = node->right;
      TreeNode *left last = NULL;
                             //左右子树最后一个节点
      TreeNode *right last = NULL;
      if (left) {
                              //若有左子树,递归将左子树转换单链表
         preorder(left, left last);
                              //左指针赋空
         node->left = NULL;
                                                    5分钟填写代码,
                  3
                                                    有问题随时提出!
         //若有右子树,递归将右子树转换单链表
      if (right) {
         preorder(right, right_last); //若node找到左子树最后一个节点
         if (left last) {
                                    (有左子树)
                  5
```

```
class Solution {
                                                          例3:方法2实现
public:
    void flatten(TreeNode *root) {
        TreeNode *last = NULL;
                               //当前子树的先序遍历的
       preorder(root, last);
                               最后一个节点, 传引用会传出
                 // 当前的节点
private:
    void preorder(TreeNode *node, TreeNode *&last) {
        if (!node) {
           return;
             !node->left && !node->right
        if
             last = node;
           return:
       TreeNode *left = node->left;
                                     //备份左右指针
       TreeNode *right = node->right;
       TreeNode *left last = NULL;
                                   //左右子树最后一个节点
       TreeNode *right last = NULL;
       if (left) {
                                    //若有左子树,递归将左子树转换单链表
           preorder(left, left last);
                                    //左指针赋空
           node->left = NULL;
            node->right = left;
           last = left last;
                            //将该节点的last保存为左子树的last
       }
                                 //若有右子树,递归将右子树转换单链表
       if (right) {
           preorder (right, right_last); //若node找到左子树最后一个节点
           if (left last) {
                                           (有左子树)
                 left last->right = right;
               last = right_last;
};
```

#### 例3:测试与leetcode提交结果

printf("\n");

return 0:

```
int main(){
                                                   [1][2][3][4][5][6]
    TreeNode a (1);
    TreeNode b(2);
    TreeNode C(5);
    TreeNode d(3);
    TreeNode e(4);
    TreeNode f(6);
    a.left = &b;
    a.right = &c;
    b.left = &d;
    b.right = &e;
    c.right = &f;
    Solution solve;
    solve.flatten(&a);
    TreeNode *head = &a;
                                          Flatten Binary Tree to Linked List
    while (head) {
                                          Submission Details
         if (head->left) {
              printf("ERROR\n");
                                                                         Status: Accepted
                                             225 / 225 test cases passed.
         printf("[%d]", head->val);
                                             Runtime: 6 ms
                                                                 Submitted: 3 hours, 4 minutes ago
         head = head->right;
```

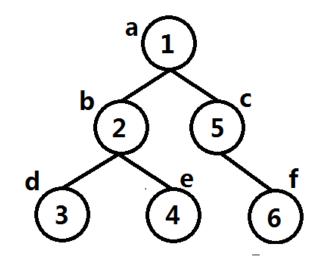
### 课间休息10分钟

# 有问题提出!

#### 预备知识:二叉树层次遍历

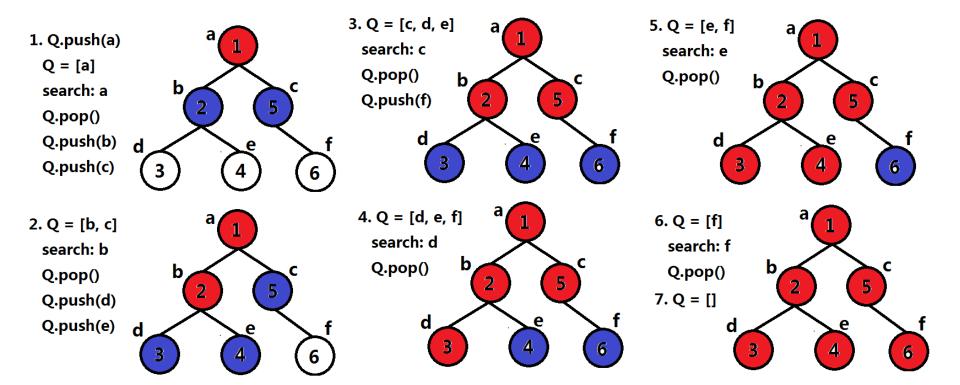
二叉树<mark>层次遍历</mark>,又称为<mark>宽度优先搜索</mark>,按树的层次依次访问树的结点。层次遍历使用<mark>队列</mark>对遍历节点进行**存储**,先进入队列的结点,<mark>优先</mark>遍历拓展其左孩子与右孩子。

设置队列 Q 将根节点 push 进Q while(Q不空){ 取出队列头部节点node 对node访问 将node的左、右孩子push进队列 }



层次遍历: a(1), b(2), c(5), d(3), e(4), f(6)

# 预备知识:二叉树层次遍历



# 预备知识:二叉树层次遍历,课堂练习

```
#include <stdio.h>
                                                                           int main(){
#include <vector>
                                                                               TreeNode a(1);
#include <queue>
                                                                               TreeNode b(2);
                                                                               TreeNode c(5);
struct TreeNode {
                   //二叉树数据结构
                                                                               TreeNode d(3);
    int val;
                                                                               TreeNode e(4);
   TreeNode *left;
                                                                               TreeNode f(6);
   TreeNode *right;
                                                                               a.left = &b;
   TreeNode(int x) : val(x), left(NULL), right(NULL) {}
                                                                               a.right = &c;
};
                                                                               b.left = &d;
void BFS print (TreeNode* root) { //宽度优先搜索二叉树
                                                                               b.right = &e;
                                                                               c.right = &f;
    std::queue<TreeNode *> Q;
                                                                               BFS print(&a);
                                                                               return 0;
    while (
        TreeNode *node = Q.front();
        printf("[%d]\n", node->val);
        if
            Q.push (node->left);
        if (node->right) {
                                         层次遍历:
                                         a(1), b(2), c(5), d(3), e(4), f(6)
```

# 预备知识:二叉树层次遍历,实现

```
#include <stdio.h>
                                                                  int main(){
#include <vector>
                                                                      TreeNode a (1);
#include <queue>
                                                                      TreeNode b(2);
                                                                      TreeNode C(5);
struct TreeNode {
                   //二叉树数据结构
                                                                      TreeNode d(3);
    int val;
                                                                      TreeNode e(4);
   TreeNode *left;
                                                                      TreeNode f(6);
   TreeNode *right;
                                                                      a.left = &b;
   TreeNode(int x) : val(x), left(NULL), right(NULL) {}
                                                                      a.right = &c;
} ;
                                                                      b.left = &d;
                                                                      b.right = &e;
void BFS_print(TreeNode* root) { //宽度优先搜索二叉树
                                                                      c.right = &f;
    std::queue<TreeNode *> Q;
                                                                      BFS print(&a);
      Q.push(root);
                                                                      return 0;
    while (!Q.empty()) {
        TreeNode *node = Q.front();
         Q.pop();
        printf("[%d]\n", node->val);
             node->left
            O.push (node->left);
        if (node->right) {
                                         层次遍历:
            Q.push(node->right);
                                         a(1), b(2), c(5), d(3), e(4), f(6)
```

[1] [2] [5] [3] [4] [6]

#### 例4:侧面观察二叉树

给定一个二叉树,假设从该二叉树的右侧观察它,将观察到的节点按照从上到下的顺序输出。

```
struct TreeNode {
   int val;
   TreeNode *left;
   TreeNode *right;
   TreeNode (int x):
      val(x), left(NULL), right(NULL) {}
};
class Solution {
public:
   std::vector<int> rightSideView(TreeNode* root) {
   }
};
   [1, 3, 4, 6]
```

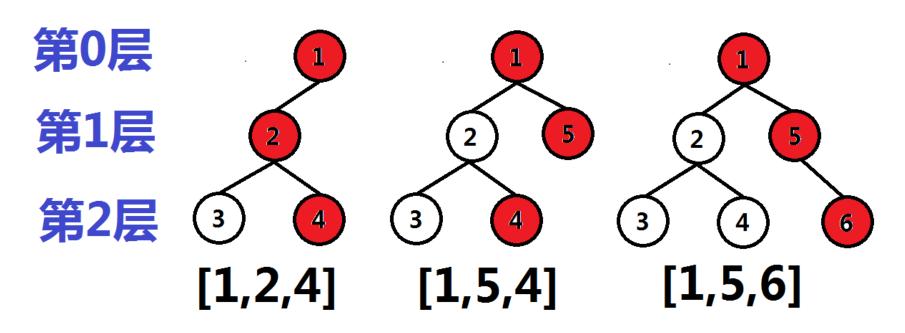
#### 选自 LeetCode 199. Binary Tree Right Side View

https://leetcode.com/problems/binary-tree-right-side-view/description/

难度:Medium

#### 例4:思考与分析

从二叉树的右侧观察它,将观察到的节点按照从上到下的顺序输出,就是求层次遍历二叉树,每个层中的最后一个节点。

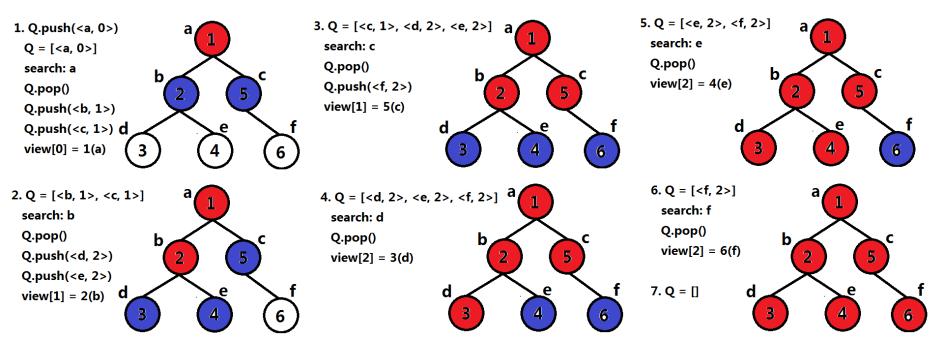


思考:层次遍历时,如何记录每一层中出现的最后一个节点?

#### 例4:算法思路

层次遍历时,将节点与层数绑定为pair, 压入队列时, 将节点与层数同时压入队列, 并记录每一层中出现的最后一个节点。 在层次遍历中, 每一层中的最后一个节点最后遍历到. 随时更

新对每层的最后一个节点即可。



```
class Solution {
                                                  例4:课堂练习
public:
   std::vector<int> rightSideView(TreeNode* root) {
       std::vector<int> view; //按层遍历的最后一个节点
       std::queue<std::pair<TreeNode *, int> > Q;
                       //宽度优先搜索队列<节点,层数>
                                                         钟填写代码,
       if (root) {
          Q.push(std::make pair(root, 0));
                                                    有问题随时提出!
              //根节点非空时,将<root, 0> push进入队列
       while(!Q.empty()){
          TreeNode *node = Q.front().first; //搜索节点
          int depth = Q.front().second; //待搜索节点的层数
          Q.pop();
           if
              view.push back(node->val);
          else{
           if
              Q.push(std::make pair(node->left
           if (node->right) {
              Q.push(std::make pair(node->right,
       return view;
};
```

### 例4:实现

```
1. Q.push(<a, 0>)
class Solution {
public:
                                                                           Q = [ <a, 0 > ]
    std::vector<int> rightSideView(TreeNode* root) {
                                                                            search: a
        std::vector<int> view; //按层遍历的最后一个节点
                                                                            Q.pop()
        std::queue<std::pair<TreeNode *, int> > Q;
                            //宽度优先搜索队列<节点,层数>
                                                                           Q.push(<b, 1>)
        if (root) {
                                                                            Q.push(<c, 1>) d
             Q.push(std::make pair(root, 0));
                                                                            view[0] = 1(a)
                 //根节点非空时,将<root, 0> push进入队列
        while(!Q.empty()){
             TreeNode *node = Q.front().first; //搜索节点
                                                                         2. Q = \{ \langle b, 1 \rangle, \langle c, 1 \rangle \}
             int depth = Q.front().second; //待搜索节点的层数
                                                                           search: b
             Q.pop();
                                                                           Q.pop()
             if ( view.size() == depth
                                                                           Q.push(<d, 2>)
                 view.push back(node->val);
                                                                           Q.push(<e, 2>)
             else{
                                                                           view[1] = 2(b) d
                   view[depth] = node->val;
                   node->left
             if
                  Q.push(std::make pair(node->left, depth + 1 ));
                                                                          3. Q = [\langle c, 1 \rangle, \langle d, 2 \rangle, \langle e, 2 \rangle] a
                                                                           search: c
             if (node->right) {
                                                         depth + 1 ));
                  Q.push(std::make pair(node->right,
                                                                           Q.pop()
                                                                           Q.push(<f, 2>)
         return view;
                                                                           view[1] = 5(c)
};
```

#### 例4:测试与leetcode提交结果

```
int main(){
                                                            \Gamma 1 \ 1
    TreeNode a (1);
    TreeNode b(2);
    TreeNode c(5);
    TreeNode d(3);
    TreeNode e(4);
    TreeNode f(6);
    a.left = &b;
    a.right = &c;
    b.left = &d;
    b.right = &e;
    c.right = &f;
    Solution solve;
    std::vector<int> result = solve.rightSideView(&a);
    for (int i = 0; i < result.size(); i++) {
        printf("[%d]\n", result[i]);
    return 0;
                       Binary Tree Right Side View
                       Submission Details
```

互联网新技术在线教育领航者

210 / 210 test cases passed.

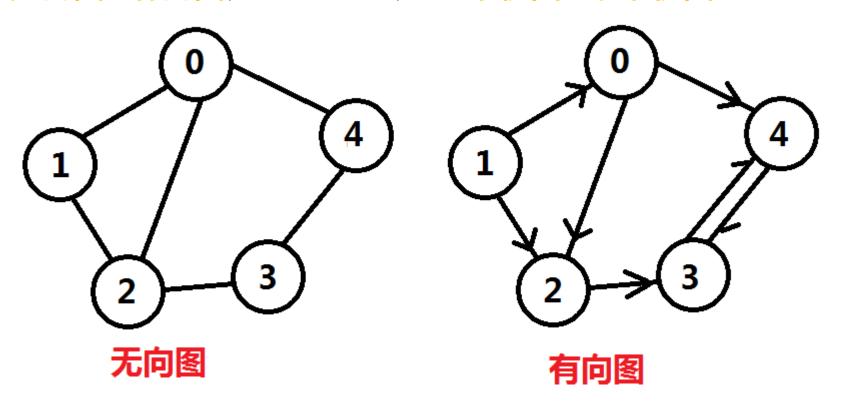
Runtime: 3 ms

Status: Accepted

Submitted: 0 minutes ago

### 预备知识:图的定义

图(Graph)是由顶点的有穷非空集合和顶点之间边的集合组成,通常表示为: G(V, E), 其中, G表示一个图, V是图G中顶点的集合, E是图G中边的集合。 图分无向图与有向图, 根据图的边长, 又分带权图与不带权图。

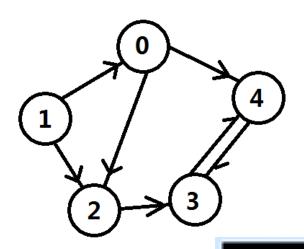


#### 预备知识:图的构造与表示(临接矩阵)

```
Graph:
#include <stdio.h>
                                                       0101
int main(){
   const int MAX N = 5; //一共5个顶点
   int Graph [MAX N] [MAX N] = {0}; //使用邻接矩阵表示
   Graph[0][2] = 1;
   Graph[0][4] = 1;
                    //将图连通 , 且不带权
   Graph[1][0] = 1;
   Graph[1][2] = 1; //一般用临接矩阵表示稠密图
   Graph[2][3] = 1;
   Graph[3][4] = 1;
   Graph[4][3] = 1;
   printf("Graph:\n");
   for (int i = 0; i < MAX N; i++) {
       for (int j = 0; j < MAX N; j++) {
           printf("%d ", Graph[i][j]);
       printf("\n");
   return 0;
```

```
#include <stdio.h>
#include <vector>
                  //图的邻接表数据结构
struct GraphNode{
                  //图的顶点的值
                                  //相邻节点指针数组
    int label;
    std::vector<GraphNode *> neighbors;
    GraphNode(int x) : label(x) {};
} ;
int main(){
    const int MAX N = 5;
    GraphNode *Graph[MAX N]; //5个顶点
    for (int i = 0; i < MAX N; i++) {
       Graph[i] = new GraphNode(i);
                 //添加边
    Graph[0]->neighbors.push back(Graph[2]);
    Graph[0]->neighbors.push back(Graph[4]);
   Graph[1]->neighbors.push back(Graph[0]);
   Graph[1]->neighbors.push back(Graph[2]);
   Graph[2]->neighbors.push back(Graph[3]);
   Graph[3]->neighbors.push back(Graph[4]);
    Graph[4]->neighbors.push back(Graph[3]);
   printf("Graph:\n");
    for (int i = 0; i < MAX N; i++) {</pre>
        printf("Label(%d) : ", i);
        for (int j = 0; j < Graph[i]->neighbors.size(); j++){
            printf("%d ", Graph[i]->neighbors[j]->label);
        printf("\n");
   for (int i = 0; i < MAX N; i++) {</pre>
       delete Graph[i];
   return 0;
```

# 预备知识:图的构造 与表示(临接表)

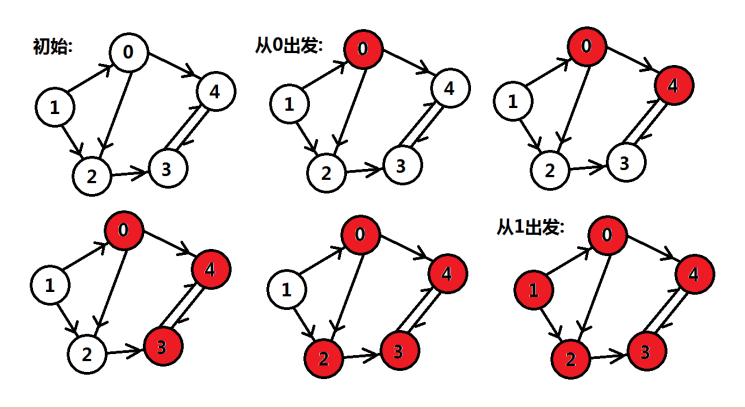


#### 有向图

```
Graph:
Label(0) : 2 4
Label(1) : 0 2
Label(2) : 3
Label(3)
Labe 1(4)
```

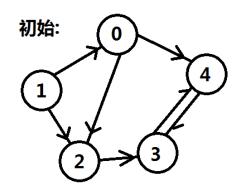
### 预备知识:图的深度优先遍历

从图中**某个顶点v**出发,首先访问该顶点,然后**依次**从它的各个未被访问的**邻接点**出发<mark>深度优先搜索</mark>遍历图,直至图中所有和v有**路径相通且未被访问**的顶点都被访问到。 若此时尚有其他顶点**未被访问**到,则另选一个未被访问的顶点作起始点,重复上述过程,直至图中<mark>所有顶点</mark>都被访问到为止。

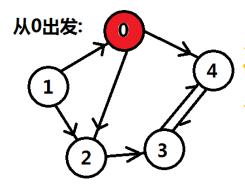


```
#include <stdio.h>
#include <vector>
                   //图的邻接表数据结构
                   //图的顶点的值
struct GraphNode{
                                    //相邻节点指针数组
    int label;
    std::vector<GraphNode *> neighbors;
    GraphNode(int x) : label(x) {};
void DFS graph(GraphNode *node, int visit[]){
    printf("%d ", node->label);
    for (int i = 0; i < node->neighbors.size(); i++) {
        if (
            DFS graph(node->neighbors[i], visit);
int main() {
                            //创建图的顶点
    const int MAX N = 5;
    GraphNode *Graph[MAX N];
    for (int i = 0; i < MAX N; i++) {</pre>
        Graph[i] = new GraphNode(i);
                //添加图的边,注意添加边的顺序
   Graph[0]->neighbors.push back(Graph[4]);
    Graph[0] -> neighbors.push back(Graph[2]);
   Graph[1] -> neighbors.push back(Graph[0]);
    Graph[1] ->neighbors.push back(Graph[2]);
   Graph[2] -> neighbors.push back(Graph[3]);
   Graph[3] -> neighbors.push back(Graph[4]);
   Graph[4]->neighbors.push back(Graph[3]);
   int visit[MAX N] = {0}; //标记已访问的顶点
   for (int i = \overline{0}; i < MAX N; i++) {
        if
           printf("From label(%d) : ", Graph[i]->label);
            DFS graph(Graph[i], visit);
           printf("\n");
   for (int i = 0; i < MAX N; i++){</pre>
       delete Graph[i];
```

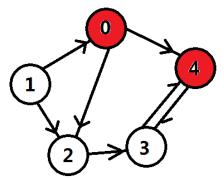
return 0;



预备知识: 图的深度 优先遍历 ,课堂练习



**3分钟**填写代码, <mark>有问题随时提</mark> 出!



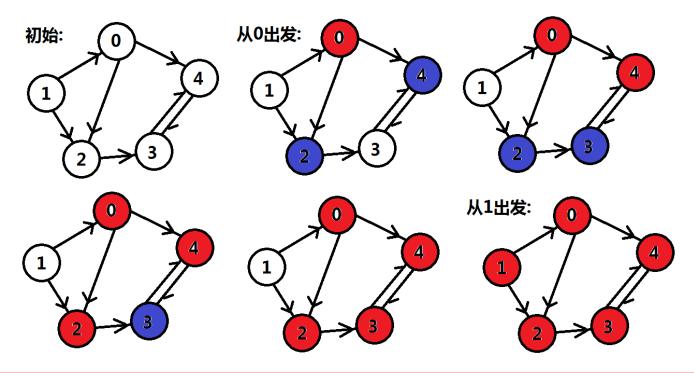
From label(0) : 0 4 3 2 From label(1) : 1

```
#include <stdio.h>
#include <vector>
                  //图的邻接表数据结构
                                                          初始:
                  //图的顶点的值
struct GraphNode{
                                  //相邻节点指针数组
    int label;
    std::vector<GraphNode *> neighbors;
   GraphNode(int x) : label(x) {};
                                                            1
void DFS graph(GraphNode *node, int visit[]){
      visit[node->label] = 1;
                                      //标记已访问的顶点
   printf("%d ", node->label); //访问相邻的且没有被访问的顶点
    for (int i = 0; i < node->neighbors.size(); i++) {
            visit[node->neighbors[i]->label] == 0
                                                          从0出发:
           DFS graph(node->neighbors[i], visit);
int main(){
                          //创建图的顶点
   const int MAX N = 5;
   GraphNode *Graph[MAX N];
   for (int i = 0; i < MAX N; i++) {</pre>
       Graph[i] = new GraphNode(i);
               //添加图的边,注意添加边的顺序
   Graph[0]->neighbors.push back(Graph[4]);
   Graph[0]->neighbors.push back(Graph[2]);
   Graph[1] ->neighbors.push back(Graph[0]);
   Graph[1]->neighbors.push back(Graph[2]);
   Graph[2]->neighbors.push back(Graph[3]);
   Graph[3]->neighbors.push back(Graph[4]);
   Graph[4] ->neighbors.push back(Graph[3]);
                                                            1
   int visit[MAX N] = {0}; //标记已访问的顶点
   for (int i = \overline{0}; i < MAX N; i++) {
       if (
                             ) { //顶点没有被标记才会访问
              visit[i] == 0
           printf("From label(%d) : ", Graph[i]->label);
           DFS graph(Graph[i], visit);
           printf("\n");
                                         From label(0)
                                                                        5
   for (int i = 0; i < MAX N; i++) {</pre>
       delete Graph[i];
                                                  label(1)
   return 0;
```

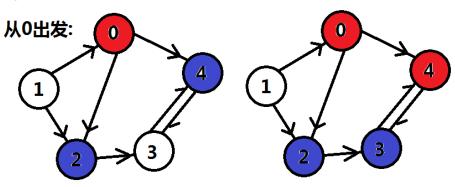
预备知识: 图的深度 优先遍历, 实现

### 预备知识:图的宽度优先遍历

从图中某个顶点v出发,在访问了v之后依次访问v的各个未曾访问过的邻接点,然后分别从这些邻接点出发依次访问它们的邻接点,并使得"先被访问的顶点的邻接点先于后被访问的顶点的邻接点被访问",直至图中所有已被访问的顶点的邻接点都被访问到。如果此时图中尚有顶点未被访问,则需要另选一个未曾被访问过的顶点作为新的起始点,重复上述过程,直至图中所有顶点都被访问到为止。

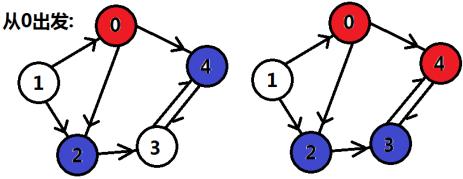


#### 预备知识:图的宽度优先遍历,课堂练习



```
int main(){
                            //创建图的顶点
    const int MAX N = 5;
    GraphNode *Graph[MAX N];
    for (int i = 0; i < MAX N; i++) {</pre>
        Graph[i] = new GraphNode(i);
                //添加图的边,注意添加边的顺序
    Graph[0]->neighbors.push back(Graph[4]);
    Graph[0]->neighbors.push back(Graph[2]);
    Graph[1]->neighbors.push back(Graph[0]);
    Graph[1]->neighbors.push back(Graph[2]);
    Graph[2]->neighbors.push back(Graph[3]);
    Graph[3]->neighbors.push back(Graph[4]);
   Graph[4]->neighbors.push back(Graph[3]);
   int visit[MAX N] = {0}; //标记已访问的顶点
   for (int i = 0; i < MAX N; i++) {</pre>
       if (visit[i] == 0){
           printf("From label(%d) : ", Graph[i]->label);
            BFS graph (Graph [i], visit);
            printf("\n");
   for (int i = 0; i < MAX N; i++) {</pre>
       delete Graph[i];
   return 0;
```

#### 预备知识:图的宽度优先遍历,实现



```
int main(){
                           //创建图的顶点
   const int MAX N = 5;
   GraphNode *Graph[MAX N];
    for (int i = 0; i < MAX N; i++) {
       Graph[i] = new GraphNode(i);
                //添加图的边,注意添加边的顺序
   Graph[0]->neighbors.push back(Graph[4]);
   Graph[0] ->neighbors.push back(Graph[2]);
   Graph[1]->neighbors.push back(Graph[0]);
   Graph[1]->neighbors.push back(Graph[2]);
   Graph[2]->neighbors.push back(Graph[3]);
   Graph[3]->neighbors.push back(Graph[4]);
   Graph[4]->neighbors.push back(Graph[3]);
   int visit[MAX_N] = {0}; //标记已访问的顶点
   for (int i = 0; i < MAX N; i++) {
       if (visit[i] == 0){
           printf("From label(%d) : ", Graph[i]->label);
           BFS graph(Graph[i], visit);
           printf("\n");
   for (int i = 0; i < MAX N; i++) {</pre>
       delete Graph[i];
   return 0;
```

#### 例5:课程安排

已知有n个课程,标记从0至n-1,课程之间是有依赖关系的,例如希望完成A课程,可能需要先完成B课程。已知n个课程的依赖关系,求是否可以将n个课程全部完成。

#### 选自 LeetCode 207. Course Schedule

https://leetcode.com/problems/course-schedule/description/

难度:Medium

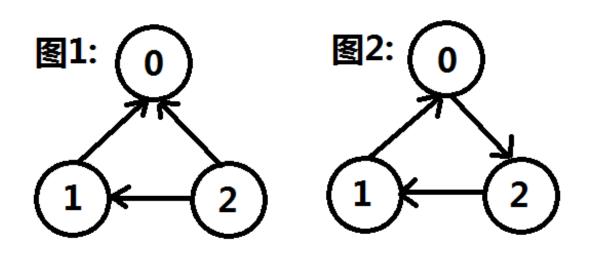
#### 例5:分析

n个课程,它们之间有m个依赖关系,可以看成<mark>顶点</mark>个数为n,<mark>边</mark>个数为m的有向图。

图1:n = 3, m = [[0, 1], [0, 2], [1, 2]]; 可以完成。

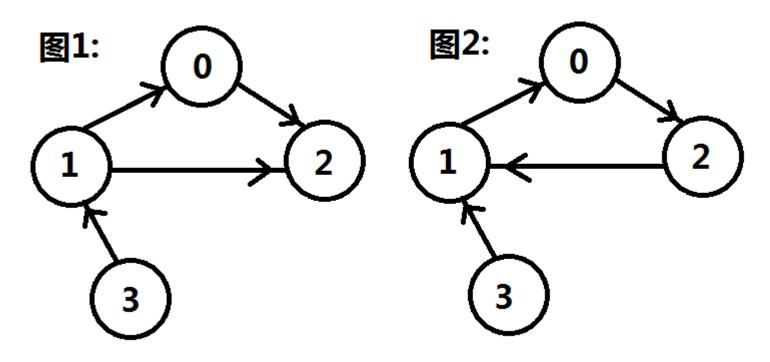
图2:n = 3, m = [[0, 1], [1, 2], [2, 0]]; 不可以完成。

故,若**有向图无环**,则可以完成全部课程,否则不能。问题转换成,构建图,并判断图是否有环。

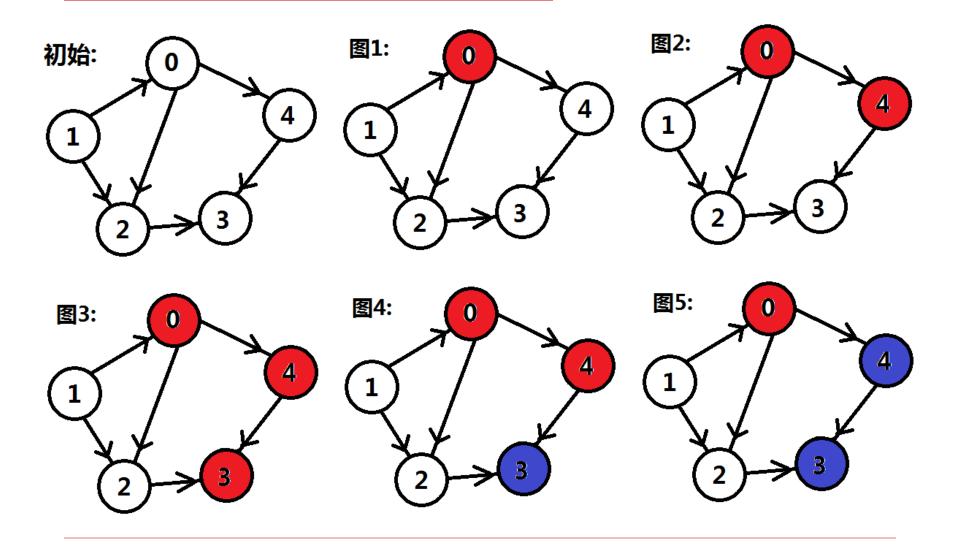


# 例5:方法1,深度优先搜索

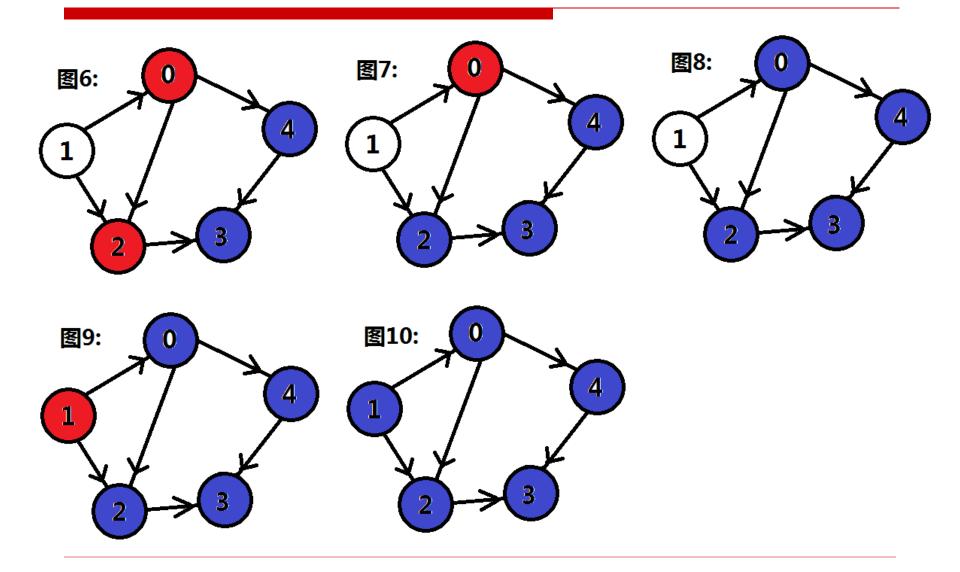
在深度优先搜索时,如果正在搜索某一顶点(还未退出该顶点的递归深度搜索),又回到了该顶点,即证明图有环。如下图:



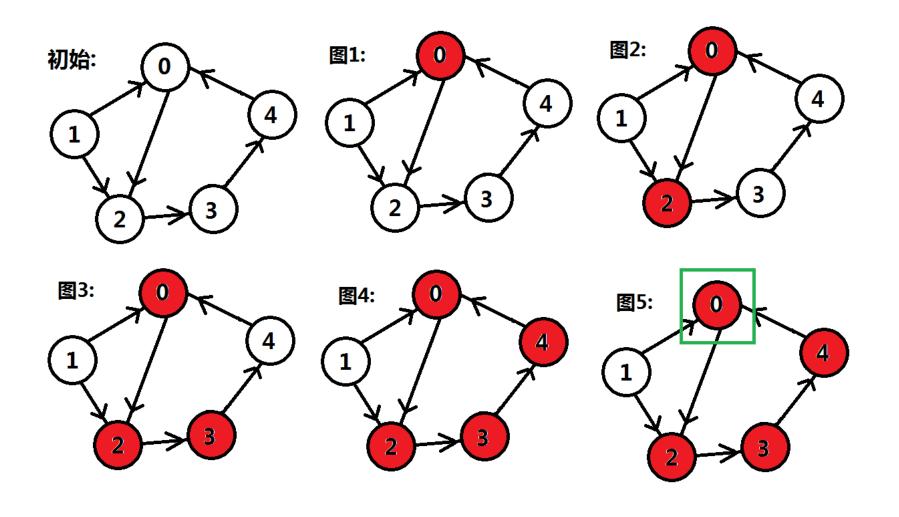
# 例5:算法思路(无环)



# 例5:算法思路(无环)



# 例5:算法思路(有环)

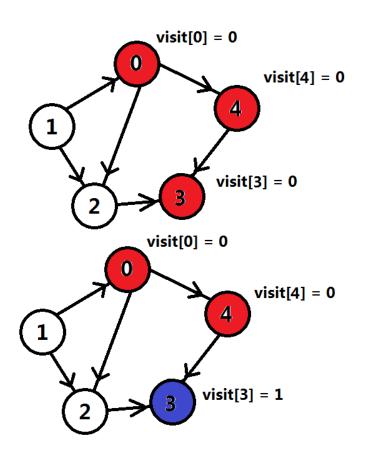


# 例5:方法1,调用代码

```
class Solution {
public:
                                 pair<课程1,课程2>课程1依赖课程2
   bool canFinish(int numCourses,
       std::vector<std::pair<int, int> >& prerequisites) {
       std::vector<GraphNode*> graph; //邻接表
       std::vector<int> visit; //节点访问状态,-1没有访问过,0代表正在访问,1代表已完成访问
       for (int i = 0; i < numCourses; i++) {
           graph.push_back(new GraphNode(i)); //创建图的节点,并赋访问状态为空
           visit.push back(-1);
                                    //创建图,连接图的顶点
       for (int i = 0; i < prerequisites.size(); i++){</pre>
           GraphNode *begin = graph[prerequisites[i].second];
           GraphNode *end = graph[prerequisites[i].first];
           begin->neighbors.push back(end);
                                              //课程2指向课程1
       for (int i = 0; i < graph.size(); i++) {
           if (visit[i] == -1 && !DFS graph(graph[i], visit)) {
               return false:
                              //如果节点没访问过,进行DFS,如果DFS遇到环,
       for (int i = 0; i < numCourses; i++) {</pre>
           delete graph[i];
       return true; //返回可以完成
};
```

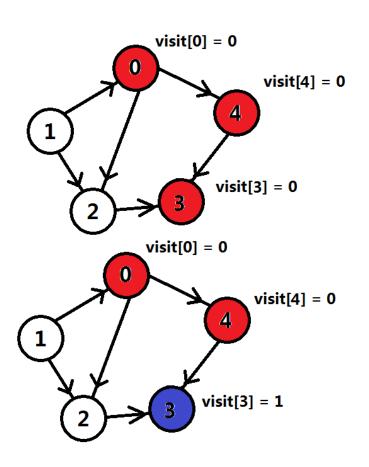
# 例5:方法1,课堂练习

```
#include <vector>
struct GraphNode{
    int label;
    std::vector<GraphNode *> neighbors;
    GraphNode(int x) : label(x) {};
};
     //节点访问状态,-1没有访问过,0代表正在访问,1代表已完成访问
bool DFS graph(GraphNode *node, std::vector<int> &visit) {
    for (int i = 0; i < node->neighbors.size(); i++) {
        if
            if (DFS graph(node->neighbors[i], visit) == 0) {
                           3
        else if (visit[node->neighbors[i]->label] == 0){
    return true;
```



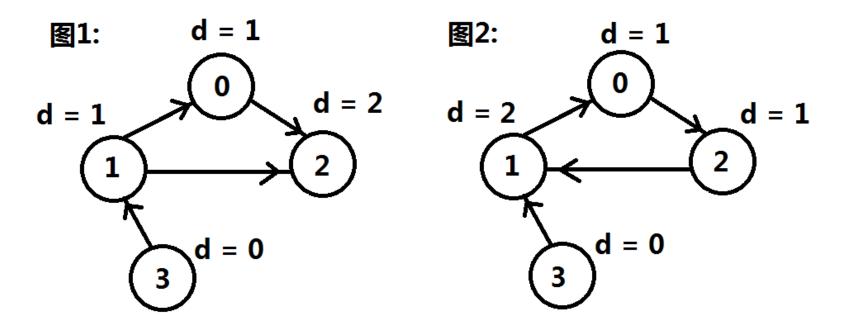
# 例5:方法1,实现

```
#include <vector>
struct GraphNode{
    int label;
    std::vector<GraphNode *> neighbors;
    GraphNode(int x) : label(x) {};
};
     //节点访问状态,-1没有访问过,0代表正在访问,1代表已完成访问
bool DFS graph(GraphNode *node, std::vector<int> &visit) {
         visit[node->label] = 0;
    for (int i = 0; i < node->neighbors.size(); i++){
              visit[node->neighbors[i]->label] == -1
        if (
            if (DFS graph(node->neighbors[i], visit) == 0) {
                    return false:
        else if (visit[node->neighbors[i]->label] == 0) {
                    return false;
     visit[node->label] = 1;
    return true;
```

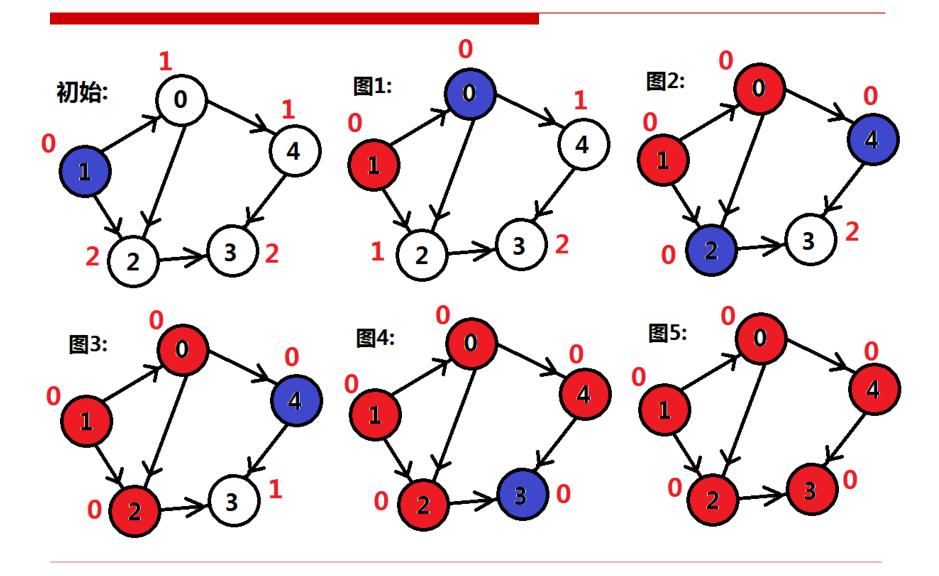


#### 例5:方法2, 拓扑排序(宽度优先搜索)

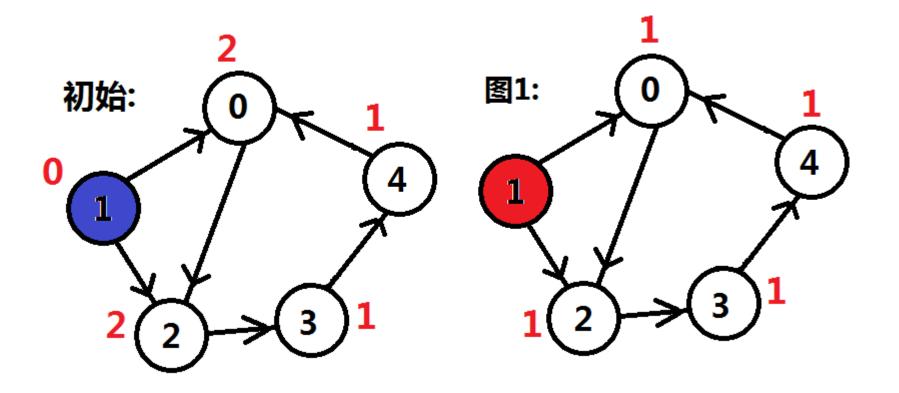
在**宽度优先搜索**时,只将入度为0的点添加至队列。当完成一个顶点的搜索(从队列取出),它指向的所有顶点入度都减1,若此时某顶点入度为0则添加至队列,若完成宽度搜索后,所有的点入度都为0,则图无环,否则有环。



# 例5:算法思路(无环)

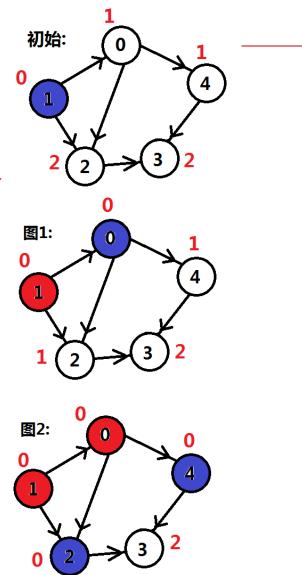


# 例5:算法思路(有环)



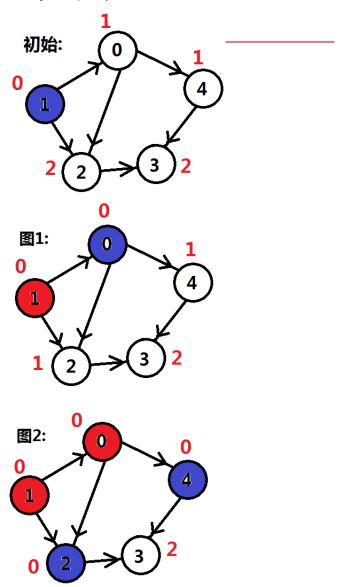
#### 例5:方法2,课堂练习

```
class Solution {
public:
    bool canFinish(int numCourses,
        std::vector<std::pair<int, int> >& prerequisites) {
         std::vector<GraphNode*> graph;
         std::vector<int> degree; //入度数组
        for (int i = 0; i < numCourses; i++) {</pre>
            degree.push back(0);
            graph.push back(new GraphNode(i));
        for (int i = 0; i < prerequisites.size(); i++){</pre>
            GraphNode *begin = graph[prerequisites[i].second];
            GraphNode *end = graph[prerequisites[i].first];
            begin->neighbors.push back(end);
                                                //入度++,即pair<课程1,课程2>
                                                课程1的入度++
        std::queue<GraphNode *> Q;
        for (int i = 0; i < numCourses; i++) {</pre>
            if
                Q.push(graph[i]);
        while(!Q.empty()){
            GraphNode *node = Q.front();
            Q.pop();
            for (int i = 0; i < node->neighbors.size(); i++){
                if (
                     Q.push (node->neighbors[i]);
        for (int i = 0; i < graph.size(); i++){</pre>
            delete graph[i];
        for (int i = 0; i < degree.size(); i++) {</pre>
            if (degree[i]){
        return true;
};
```



# 例5:方法2,实现

```
class Solution {
public:
    bool canFinish(int numCourses,
         std::vector<std::pair<int, int> >& prerequisites) {
         std::vector<GraphNode*> graph;
         std::vector<int> degree; //入度数组
         for (int i = 0; i < numCourses; i++) {
            degree.push back(0);
            graph.push back(new GraphNode(i));
         for (int i = 0; i < prerequisites.size(); i++){</pre>
            GraphNode *begin = graph[prerequisites[i].second];
             GraphNode *end = graph[prerequisites[i].first];
             begin->neighbors.push back(end);
                                                //入度++,即pair<课程1,课程2>
              degree[prerequisites[i].first]++;
                                                 课程1的入度++
        std::queue<GraphNode *> Q;
        for (int i = 0; i < numCourses; i++) {</pre>
            if
                    degree[i] == 0
                Q.push(graph[i]);
        while(!Q.empty()){
            GraphNode *node = Q.front();
            Q.pop();
            for (int i = 0; i < node->neighbors.size(); i++) {
                   degree[node->neighbors[i]->label]--;
                      degree[node->neighbors[i]->label] == 0
                     Q.push(node->neighbors[i]);
        for (int i = 0; i < graph.size(); i++){</pre>
            delete graph[i];
        for (int i = 0; i < degree.size(); i++){</pre>
            if (degree[i]){
                   return false:
        return true;
} ;
```



### 例5:测试与leetcode提交结果

```
int main() {
    std::vector<std::pair<int, int> > prerequisites;
    prerequisites.push_back(std::make_pair(1, 0));
    prerequisites.push_back(std::make_pair(2, 0));
    prerequisites.push_back(std::make_pair(3, 1));
    prerequisites.push_back(std::make_pair(3, 2));
    Solution solve;
    printf("%d\n", solve.canFinish(4, prerequisites));
    return 0;
}

Course Schedule
```

Submission Details

37 / 37 test cases passed. Status: Accepted
Runtime: 13 ms Submitted: 39 minutes ago

#### 结束

# 非常感谢大家!

林沐