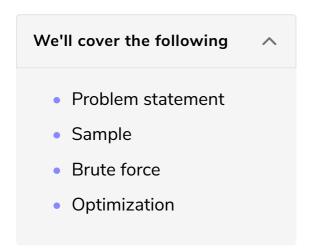
Solved Problem - Subarray Sum

In this lesson, we'll see an effective way to compute sum of an subarray.



Problem statement

Given an array, A, of N integers, answer Q queries of the type (l,r) - sum of the elements in the subarray A[l...r]. Print the sum for each query.

Input format

The first line contains two space-separated integers N and Q $(1 \le N, Q \le 10^6)$.

The second line contains N integers representing the array A[] $(1 \le A[i] \le 10^6)$.

The next Q lines each contain a pair of integers l and r.

Output format

Print Q integers and answer to the queries.

Sample

Input

```
5 3
1 2 4 8 16
1 5
2 3
3 5
```

Brute force

The brute force method would be to loop over the subarray in the query and sum all the elements.

I am skipping the code for this solution because it is trivial.

The time for each query would be O(N), the total complexity of the solution would be O(Q*N). This means it is not good enough for the given constraints.

Optimization

First, let's discuss what the prefix sum array is.

The prefix sum array sum[] of an array A[] is defined as

$$sum[i] = \sum_{k=1}^{i} A[k]$$

or, ith element of sum[] is sum of first i elements of A[]

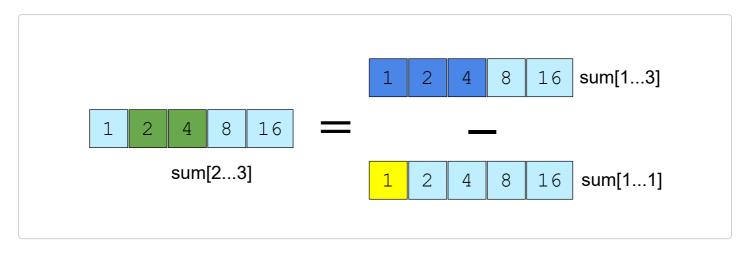
We can use the prefix sum array to find the sum of a subarray in O(1) time.

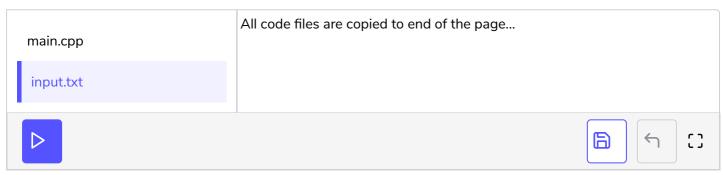
$$egin{aligned} sum[i...j] &= A[i] + A[i+1] + ... + A[j] \ &= (A[1] + A[2] + ... + A[j]) - (A[1] + A[2] + ... + A[i-1]) \ &= sum[j] - sum[i-1] \end{aligned}$$

From preprocessing to computing the sum[] array takes O(N) time. Each query is just O(1).

So the total time complexity is O(N+Q).

See the below illustration for a better understanding.





In the next chapter, we'll start with Sieve of Eratosthenes, an efficient algorithm to generate primes.

Code Files Content !!!


```
#include
#include
#include
using namespace std;

int main() {
   ifstream cin("input.txt");

   int N, Q;
   cin >> N >> Q;
   vector A(N);
   for (int i = 0: i < N: i++) cin>>A[i]:
```

```
vector sum(N);
 sum[0] = A[0];
 for(int i = 1; i < N; i++)</pre>
   sum[i] = sum[i - 1] + A[i];
 for(int i = 0; i < Q; i++) {
   int 1, r;
   cin >> 1 >> r;
   l--; r--; // Convert to 0-based indexing
   int ans = sum[r];
   if (1 > 0)
     ans -= sum[l - 1];
   cout << ans << " ";
 }
 return 0;
}
| input.txt [1]
5 3
1 2 4 8 16
1 5
2 3
3 5
*******************************
```