

Solution Review: Find All Permutations of a String

In this lesson, we will review the solution to the challenge from the previous lesson.

We'll cover the following ^

- Solution
- Explanation
- Time complexity

Solution

```
def permutations(str):
    if str == "": # base case
        return [""]
    permutes = []
    for char in str:
        subpermutes = permutations(str.replace(char, "", 1))    # recursive step
        for each in subpermutes:
            permutes.append(char+each)
    return permutes

def main():
    print (permutations("abc"))

main()
```

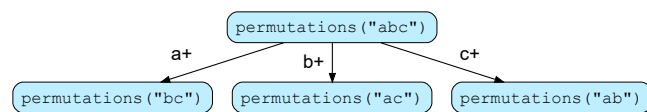


Explanation

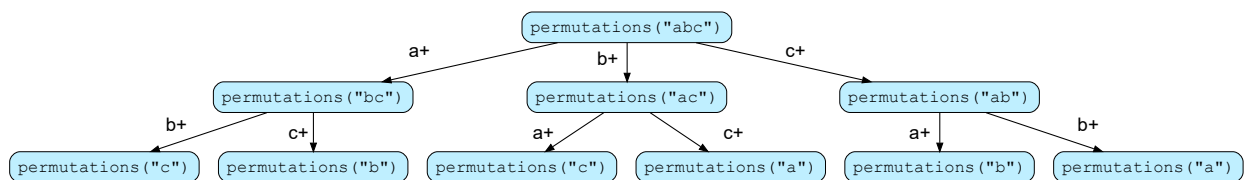
Just as we learned in previous lessons, the key to acing recursion is focusing on one step at a time. Try to think of this problem in this way: you already have every possible arrangement of every possible subsequence of the string `str`. All you need to do now is prepend all the characters to their corresponding list of substrings. Let's see a visualization of this before we jump into the code.

permutations("abc")

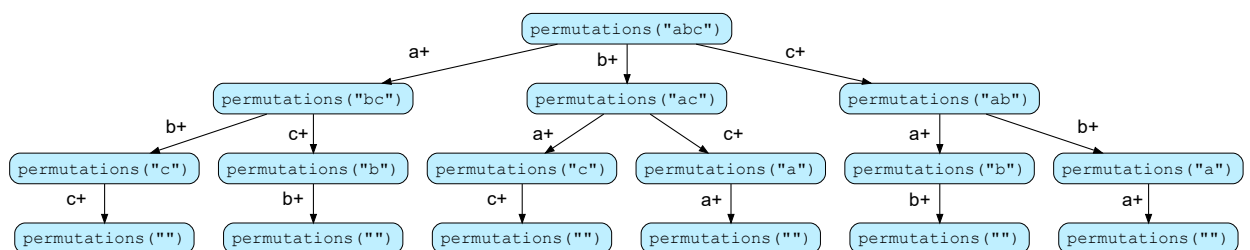
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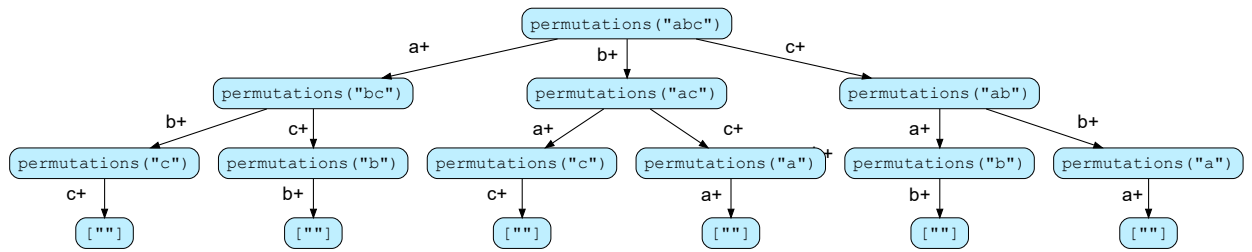
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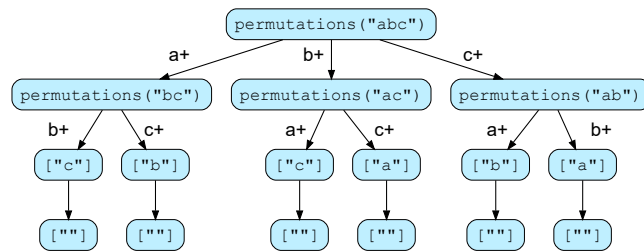
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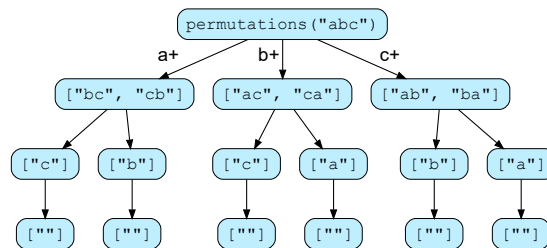
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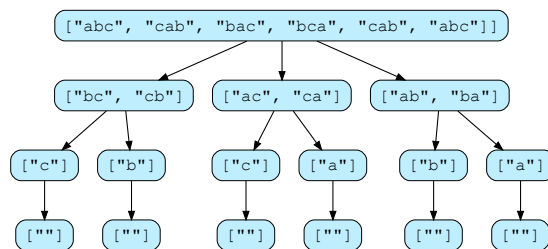
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So, what is happening here? At each step, we make as many recursive calls to `permutations` as there are characters in our `str`. For each character, `char`, we find permutations of `str` excluding that character, and simply prepend the character to the results. This is the crux of this algorithm, highlighted in code at *lines 5-8*. We should also be careful about our termination condition. Our algorithm has only one base case of when the string is empty, and it returns a list containing an empty string (*lines 2-3*). You could have had the terminating condition as a string of length 1, but then our algorithm would not have worked for empty strings. This highlights the importance of choosing base cases for a recursive algorithm.

Time complexity

Look at the visualization again and notice how at the first step we had three branches, then two branches per branch, and then finally one. Which is $3 \times 2 \times 1 = 3!$. If we were to generalize this to n , our time complexity would be $O(n!)$. This makes sense because for an n character long string, there are $n!$ total permutations.

Hopefully, this was a good exercise for you. We have another one waiting in the next lesson.