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# 1 Infos

You can get many infos there. <https://github.com/bashardudin/GraphsAndFlows>

## 2 Recap

### shortest path

- single-pair (src dst) shortest path problem
- single src any destination : shortest path problem
- Any pair shortest path pb

### algorithms

- Bellman-Ford [general no negative valued edges]
- Dijkstra [positive weights]
- Bellman [no cycles]
- Floyd warshall

### 2.1 Safest path

Aim : safest path along which to go.

To use shortest paths algorithms for it is about applying shortest paths to graphs having  $-\log(\text{above weights})$ .

### 2.2 Finding paths

To find critical paths going back from ending point you look for predecessors whose earliest starting date is earliest starting point 'E' - time task takes. And you go through recursively.

### margins

It is the difference between earliest starting date of task and the latest (not delaying project).

### free margin

Is the difference between earliest starting date & latest one not modifying the earliest starting date of any successors.

### 2.3 A bit of formalism

Let  $c : R_+$  be the capacity in a given graph  $G$  (having  $s$  and  $t$ )

Then a flow on  $G$  is a function  $f : A \rightarrow R_+$

Satisfying : (1)  $\forall a \in A, f(a) \leq c(a)$

(2)  $\forall i \in V \setminus \{s, t\} :$

$$\sum_{a \rightarrow i} f(a) = \sum_{i \rightarrow a} f(a)$$

$$a \rightarrow i \quad i \rightarrow a$$

$$a \in A \quad a \in A$$

## 2.4 Exercice 6 (exercise sheet)

1)

