



RISC-V Capacity and Bandwidth QoS Register Interface

RISC-V CMQRI Task Group

Version 0.1, 11/2022: This document is in development. Assume everything can change. See
<http://riscv.org/spec-state> for details.

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Preamble



This document is in the [Development state](#)

Assume everything can change. This draft specification will change before being accepted as standard, so implementations made to this draft specification will likely not conform to the future standard.

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Chapter 1. Introduction

Quality of Service (QoS) is the minimal end-to-end performance that is guaranteed in advance by a service level agreement (SLA) to an workload. The performance may be measured in the form of metrics such as instructions per cycle (IPC), latency of servicing work, etc.

Various factors such as the available cache capacity, memory bandwidth, interconnect bandwidth, CPU cycles, system memory, etc. affect the performance in a computing system that runs multiple workloads concurrently. Further when there is arbitration for shared resources, the prioritization of the workloads requests against other competing requests may also affect the performance of the workload. Such interference due to resource sharing may lead to unpredictable workload performance [1].

When multiple workloads are running concurrently on modern processors with large core counts, multiple cache hierarchies, and multiple memory controllers, the performance of an workload becomes less deterministic or even non-deterministic as the performance depends on the behavior of all the other workloads in the machine that contend for the shared resources leading to interference. In many deployment scenarios such as public cloud servers the workload owner may not be in control of the type and placement of other workloads in the platform.



A workload may be a single application, a group of application, a virtual machine, a group of virtual machines, or some combination of above.

System software can control some of these resources available to the workload such as the number of hardware threads made available for execution, the amount of system memory allocated to the workload, the number of CPU cycles provided for execution, etc.

System software needs additional tools to control interference to a workload and thereby reduce the variability in performance experienced by one workload due to other workload's cache capacity usage, memory bandwidth usage, interconnect bandwidth usage, etc. through a resource allocation extension. The resource allocation extension enables system software to reserve capacity and/or bandwidth to meet the performance goals of the workload. Such controls enable improving the utilization of the system by colocating workloads while minimizing the interference caused by one workload to another [2].

Effective use of the resource allocation extension requires hardware to provide a resource monitoring extension by which the resource requirements of a workload needed to meet a certain performance goal can be characterized. A typical use model involves profiling the resource usage of the workload using the resource monitoring extensions and to establish resource allocations for the workload using the resource allocation extensions.

The allocation may be in the form of capacity or bandwidth depending on the type of resource. For caches and directories, the resource allocation is in the form of storage capacity. For interconnects and memory controllers the resource allocation is in the form of bandwidth.

Workloads typically generate a different types of accesses to shared resources. For example, some of the requests may be for accessing instructions and others may be to access data operated on by the workload. Certain data accesses may not have temporal locality whereas others may have a

high probability of reuse. In some cases it is desirable to provide differentiated treatment to each type of access by providing unique resource allocation to each access type.

RISC-V Capacity and Bandwidth Controller QoS Register Interface (CBQRI) extension specifies:

1. QoS identifiers to identify workloads that originate requests to the shared resources. These QoS identifiers include an identifier for resource allocation configuration and an identifier for the monitoring counters used to monitor resource usage. These identifiers accompany each request made by the workload to the shared resource. [Chapter 2](#) specifies the mechanism to associate the identifiers with workloads.
2. Access-type identifiers to accompany request to access a shared resource to allow differentiated treatment of each access-type (e.g., code vs. data, etc.). The access-types are defined in [Chapter 2](#)
3. Register interface for capacity allocation in capacity controllers such as shared caches and directories. The capacity allocation register interface is specified in [Chapter 3](#).
4. Register interface for capacity usage monitoring. The capacity usage monitoring register interface is specified in [Chapter 3](#).
5. Register interface for bandwidth allocation in bandwidth controllers such as interconnect and memory controllers. The bandwidth allocation register interface is specified in [Chapter 4](#).
6. Register interface for bandwidth usage monitoring. The bandwidth usage monitoring register interface is specified in [Chapter 4](#).

The capacity and bandwidth controller register interfaces for resource allocation and usage monitoring are defined as memory-mapped registers. Each controller that supports the CBQRI extension provides a set of registers that are located at a range of physical address space that is a multiple of 4-KiB and the lowest address of the range is aligned to 4-KiB. The memory-mapped registers may be accessed using naturally aligned 4-byte or 8-byte memory accesses. The controller behavior for register accesses where the address is not aligned to the size of the access, or if the access spans multiple registers, or if the size of the access is not 4-bytes or 8-bytes, is **UNSPECIFIED**. The atomicity of access to an 8-byte register is **UNSPECIFIED**. The implementation may observe the 8-byte access as two 4-byte accesses. A 4-byte access to a register must be single-copy atomic.



If an implementation may observe a 8-byte register access as two 4-byte accesses then such implementations must preserve the semantics of the 8-byte access and must cause any effects of the access must only occur after both accesses have been observed.

The controller registers have little-endian byte order (even for systems where all harts are big-endian-only).



Big-endian-configured harts that make use of the register interface are expected to implement the **REV8** byte-reversal instruction defined by the Zbb extension. If **REV8** is not implemented, then endianness conversion may be implemented using a sequence of instructions.

A controller may support a subset of capabilities defined by CBQRI. When a capability is not supported the registers and/or fields used to configure and/or control such capabilities are

hardwired to 0. Each controller supports a capabilities register to enumerate the supported capabilities.

Chapter 2. QoS Identifiers

Monitoring or allocation of resources requires a way to identify the originator of the request to access the resource.

The CBQRI extension provides a mechanism by which a workload can be associated with a resource control ID (**RCID**) and a monitoring counter ID (**MCID**) that accompany each request made by the workload to shared resources.

To provide differentiated services to workloads, the CBQRI extension defines a mechanism to configure resource usage limits, in the form of capacity or bandwidth, and optionally a priority for an **RCID** in the resource controllers that control accesses to such shared resources.

To monitor the resource utilization by a workload the CBQRI extension defines a mechanism to configure counters identified by the **MCID** to count events in the resource controllers that control accesses to such shared resources.

2.1. Associating **RCID** and **MCID** with requests

2.1.1. RISC-V hart initiated requests

The **sqoscfg** CSR is a 32-bit S/HS-mode read/write register to configure a resource control ID (**RCID**) and a monitoring counter ID (**MCID**). The **RCID** and **MCID** accompany each request made by the hart to shared resources such as interconnects, caches, memory, etc.

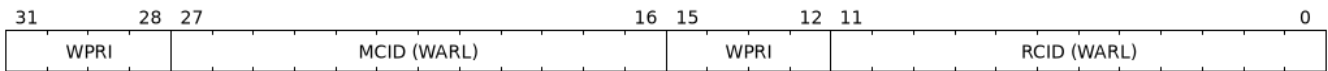


Figure 1. **sqoscfg** register for RV32 and RV64



The type of request made to the shared resource controller depends on the type of shared resource. In case of resources such as caches or memory these may be a memory access request. In case of resources such as CPU dispatch slots and retirement bandwidth the request may be to allocate such resources for execution.

The **RCID** in the request is used by the resource controllers to determine the resource allocations (e.g., cache occupancy limits, memory bandwidth limits, etc.) to enforce. The **MCID** in the request is used by the resource controllers to identify the ID of a counter to monitor resource usage (e.g., cache occupancy, memory bandwidth, etc.).

Access to **sqoscfg** when **V=1** causes a virtual-instruction exception.



At reset it is suggested that the **RCID** field of **sqoscfg** be set to 0 as typically the resource controllers in the SoC default to a reset behavior of associating all capacity or bandwidth to the **RCID** value of 0.

The value of **MCID** at reset, unlike the **RCID**, does not affect functional behavior. Implementations may choose a convenient legal value for the **MCID** reset value.

The **RCID** and **MCID** configured in the **sqoscfg** CSR apply to all privilege modes of software execution on that hart.

2.1.2. Device initiated requests

Devices may be configured with an **RCID** and **MCID** for requests originated from the device if the device implementation supports such capability. The method to configure the QoS identifiers into devices is **UNSPECIFIED**.

Where the device does not natively support being configured with an **RCID** and **MCID**, the implementation may provide a shim at the device interface that may be configured with the **RCID** and **MCID** that are associated with requests originating from the device. The method to configure such QoS identifiers into a shim is **UNSPECIFIED**.

If the system supports an IOMMU, then the IOMMU may be configured with the **RCID** and **MCID** to associate requests from the device with the QoS identifiers. The RISC-V IOMMU specification may support a future standard extension to support this configuration.

2.2. Access-type (**AT**)

In some usages, in addition to providing differentiated service among workloads, the ability to differentiate between resource usage for accesses made by the same workload may be required. For example, the capacity allocated in a shared cache for code storage may be differentiated from the capacity allocated for data storage and thereby avoid code from being evicted from such shared cache due to a data access.

When differentiation based on access type (e.g. code vs. data) is supported the requests also carry an access-type (**AT**) indicator. The resource controllers may be configured with separate capacity and/or bandwidth allocations for each supported access-type. The CBQRI extension defines a 3-bit **AT** field, encoded as specified in [Table 1](#), in the register interface to configure differentiated resource allocation and monitoring for each **AT**.

*Table 1. Encodings of **AT** field*

Value	Name	Description
0	Data	Requests to access data.
1	Code	Requests for code execution.
2-5	Reserved	Reserved for future standard use.
6-7	Custom	Designated for custom use.

If a request is received with an unsupported **AT** value then the resource controller behaves as if the **AT** value was 0.

Chapter 3. Capacity-controller QoS Register Interface

Controllers, such as cache controllers, that support capacity allocation and capacity usage monitoring provide a memory-mapped register programming interface.

Table 2. Capacity-Controller QoS register layout

Offset	Name	Size	Description	Optional?
0	<code>cc_capabilities</code>	8	Capabilities	No
8	<code>cc_mon_ctl</code>	8	Usage monitoring control	Yes
16	<code>cc_mon_ctr_val</code>	8	Monitoring counter value	Yes
24	<code>cc_alloc_ctl</code>	8	Capacity allocation control	Yes
32	<code>cc_block_mask</code>	$M * 8$	Capacity block mask	Yes

The size of the `cc_block_mask` register is determined by the `NCBLKS` field of the `cc_capabilities` register.

The reset value is 0 for the following registers fields.

- `cc_mon_ctl` - `BUSY` and `STATUS` fields
- `cc_alloc_ctl` - `BUSY` and `STATUS` fields

The reset value is `UNSPECIFIED` for all other registers and/or fields.

The capacity controllers at reset must allocate all available capacity to `RCID` value of 0. When the capacity controller supports capacity allocation per access-type, then all available capacity is allocated to each supported access-type for `RCID=0`. The capacity allocation for all other `RCID` values is `UNSPECIFIED`. The behavior of handling a request with a non-zero `RCID` value before configuring the capacity controller with capacity allocation for that `RCID` is `UNSPECIFIED`.

3.1. Capabilities (`cc_capabilities`)

The `cc_capabilities` register is a read-only register that holds the capacity-controller capabilities.

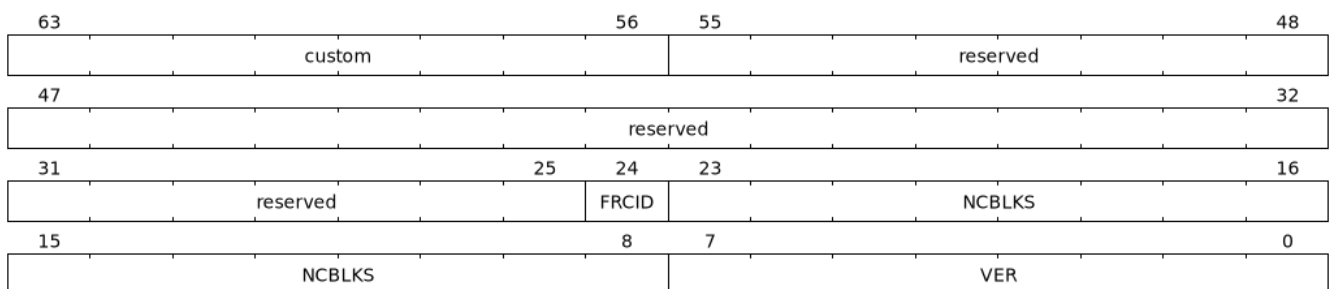


Figure 2. Capabilities register fields

The `VER` field holds the version of the specification implemented by the capacity controller. The low nibble is used to hold the minor version of the specification and the upper nibble is used to hold the

major version of the specification. For example, an implementation that supports version 1.0 of the specification reports 0x10.

The **NCBLKS** field holds the total number of allocatable capacity blocks in the controller. The capacity represented by an allocatable capacity block is **UNSPECIFIED**. The capacity controllers support allocating capacity in multiples of an allocatable capacity block.



For example, a cache controller that supports capacity allocation by ways may report the number of ways as the number of allocatable capacity blocks.

If **FRCID** is 1, the controller supports an operation to flush and deallocate the capacity blocks occupied by an **RCID**.

3.2. Usage monitoring control (**cc_mon_ctl**)

The **cc_mon_ctl** register is used to control monitoring of capacity usage by a **MCID**. When the controller does not support capacity usage monitoring the **cc_mon_ctl** register is read-only zero.

63												56			
		WPRI													
55												48			
		WPRI													
47												40			
		WPRI													
39		38												32	
BUSY (RO)		STATUS (RO)													
31		29				28		27				24			
WPRI				ATV (WARL)		EVT_ID (WARL)									
23		20				19				16					
EVT_ID (WARL)						MCID (WARL)									
15												8			
		MCID (WARL)													
7		5				4		0							
AT (WARL)				OP (WARL)											

Figure 3. Usage monitoring control (**cc_mon_ctl**)

Capacity controllers that support capacity usage monitoring implement a usage monitoring counter for each supported **MCID**. The usage monitoring counter may be configured to count a monitoring event. When an event matching the event configured for the **MCID** occurs then the monitoring counter is updated. The event matching may optionally be filtered by the access-type.

The **OP** field is used to instruct the controller to perform an operation listed in Table 3.

The **EVT_ID** field is used to program the identifier of the event to count in the monitoring counter selected by **MCID**. The **AT** field is used to program the access-type to count.

When the **EVT_ID** for a **MCID** is selected as 0, the counter retains its current value but stops counting. When **EVT_ID** is programmed with a non-zero and legal value, the counter resets to 0 and starts counting the programmed events for requests with matching **MCID** and **AT**(if **ATV** is 1). When **ATV** is 0, the counter counts requests with all access-types and the **AT** value is ignored.

A controller that does not support monitoring by access-type may hardwire the **ATV** and **AT** fields to 0.

Table 3. Usage monitoring operations (OP)

Operation	Encoding	Description
—	0	Reserved for future standard use.
CONFIG_EVENT	1	Configure the counter selected by MCID to count the event selected by EVT_ID , AT , and ATV . The EVT_ID encodings are listed in Table 4.
READ_COUNTER	2	Snapshot the value of the counter selected by MCID into cc_mon_ctr_val register. The EVT_ID , AT and ATV fields are not used by this operation.
—	3-23	Reserved for future standard use.
—	24-31	Designated for custom use.

Table 4. Usage monitoring event ID (EVT_ID)

Event ID	Encoding	Description
None	0	Counter does not count and retains its value.
Occupancy	1	Counter is incremented by 1 when a request with a matching MCID and AT allocates a unit of capacity. The counter is decremented by 1 when a unit of capacity is de-allocated.
—	2-127	Reserved for future standard use.
—	128-256	Designated for custom use.

When the **cc_mon_ctl** is written, the controller may need to perform several actions that may not complete synchronously with the write. A write to the **cc_mon_ctl** sets the **BUSY** bit to 1 indicating the controller is performing the requested operation. When the **BUSY** bit reads 0 the operation is complete and the **STATUS** field provides a status value (Table 5) of the requested operation.

Table 5. **cc_mon_ctl.STATUS** field encodings

STATUS	Description
0	Reserved
1	Operation was successfully completed.
2	Invalid operation (OP) requested.
3	Operation requested for an invalid MCID .
4	Operation requested for an invalid EVT_ID .
5	Operation requested for an invalid AT .
6-63	Reserved for future standard use.
64-127	Designated for custom use.

Behavior of writes to the **cc_mon_ctl** when **BUSY** is 1 is **UNSPECIFIED**. Some implementations may ignore the second write and others may perform the operation determined by the second write. Software must verify that **BUSY** is 0 before writing **cc_mon_ctl**.

3.3. Monitoring counter value (`cc_mon_ctr_val`)

The `cc_mon_ctr_val` is a read-only register that holds a snapshot of the counter requested by `READ_COUNTER` operation.

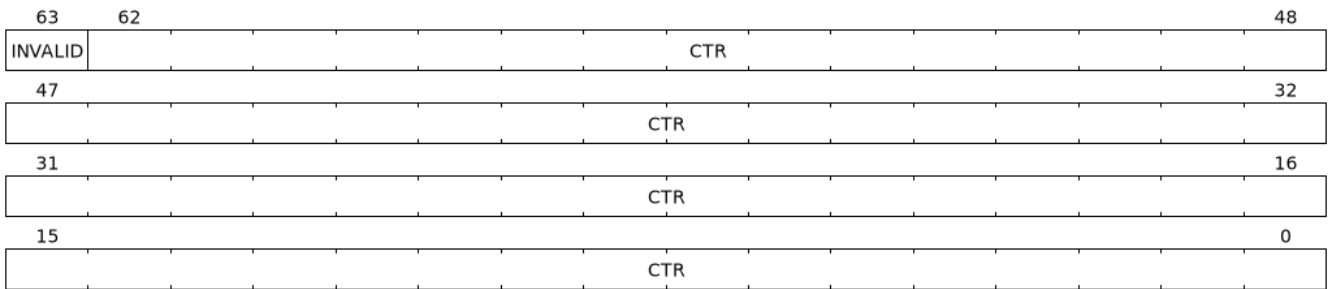


Figure 4. Usage monitoring counter value (`cc_mon_ctr_val`)

The counter is valid if the `INVALID` field is 0. The counter may be marked `INVALID` if it underflows or the controller for `UNSPECIFIED` reasons determine the count to be not valid. The counters marked `INVALID` may become valid in future.



A counter may underflow when capacity is de-allocated following a reset of the counter to 0. This may be due to the `MCID` being reallocated to a new workload while the capacity controller still holds capacity allocated by the workload to which the `MCID` was previously allocated. The counter value should typically stabilize to reflect the usage of the new workload after the workload has executed for a short duration following the counter reset.



Some implementations may not store the `MCID` of the request that caused the capacity to be allocated with every unit of capacity in the controller to optimize on the storage overheads. Such controllers may in turn rely on statistical sampling to report the capacity usage by tagging only a subset of the capacity units.

Techniques such as set-sampling in caches have been shown to provide statistically accurate cache occupancy information with a relatively small sample size such as 10% [3].

When the controller has not observed enough samples to provide an accurate value in the monitoring counter the controller may report the counter as being `INVALID` till more accurate measurements are available.

3.4. Capacity allocation control (`cc_alloc_ctl`)

The `cc_alloc_ctl` register is used to control allocation of capacity to a `RCID` per `AT`. If a controller does not support capacity allocation then the register is read-only zero. If the controller does not support capacity allocation per access-type then the `AT` field is read-only zero.

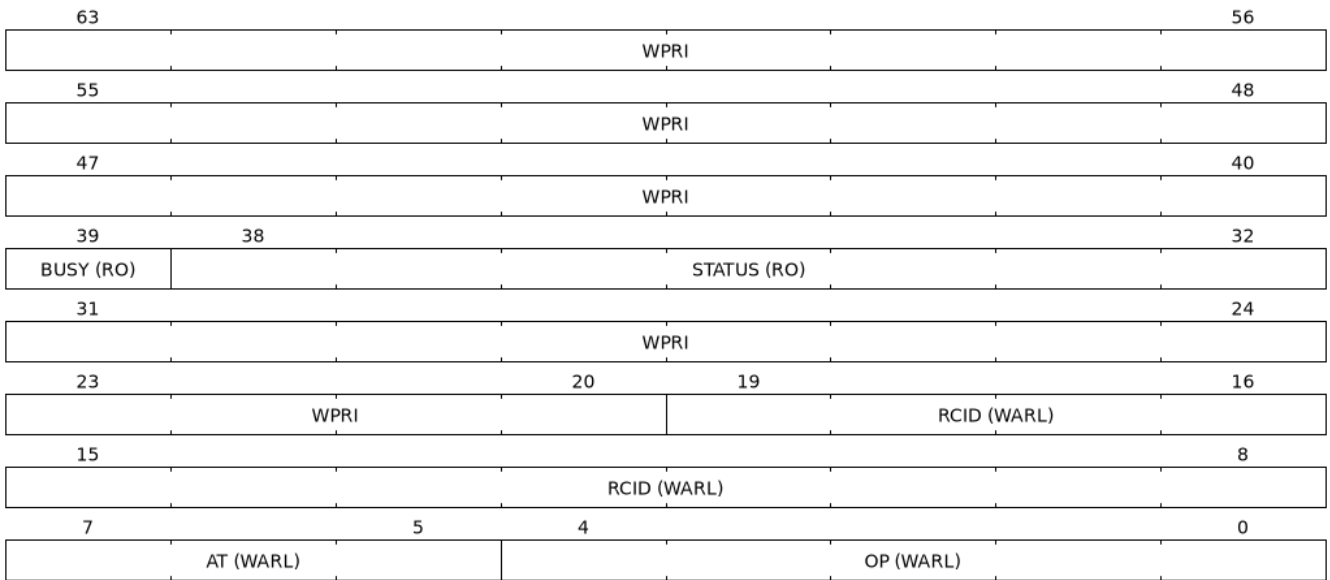


Figure 5. Capacity allocation control (`cc_alloc_ctl`)

The `OP` field is used to instruct the capacity controller to perform an operation listed in Table 6. The `cc_alloc_ctl` register is used in conjunction with the `cc_block_mask` register to perform capacity allocation operations. When the requested operation uses the operands configured in `cc_block_mask`, software must first program the `cc_block_mask` register with the operands for the operation before requesting the operation.

Table 6. Capacity allocation operations (`OP`)

Operation	Encoding	Description
—	0	Reserved for future standard use.
<code>CONFIG_LIMIT</code>	1	The <code>CONFIG_LIMIT</code> operation is used to establish a limit for capacity allocation for requests by <code>RCID</code> and of access-type <code>AT</code> . The capacity allocation is specified in <code>cc_block_mask</code> register.
<code>READ_LIMIT</code>	2	The <code>READ_LIMIT</code> operation is used to read back the previously configured allocation limits for requests by <code>RCID</code> and of type <code>AT</code> . The current configured allocation limit is written to <code>cc_block_mask</code> register on completion of the operation.
<code>FLUSH_RCID</code>	3	The <code>FLUSH_RCID</code> operation requests the controller to deallocate the capacity allocated for use by the specified <code>RCID</code> for access-type specified by <code>AT</code> . The <code>cc_block_mask</code> register is not used for this operation.
—	4-23	Reserved for future standard use.
—	24-31	Designated for custom use.

Capacity controllers enumerate the allocatable capacity blocks in `NCBLKS` field of the `cc_capabilities` register. The `cc_block_mask` register is programmed with a bit-mask where each bit represents a capacity block to be allocated.

A capacity allocation must be configured for each supported access-type by the controller. Identical limits may be configured for two or more access-types if different capacity allocation per access-type is not required. The behavior is `UNSPECIFIED` if capacity is not allocated for each access-type

supported by the controller.



A cache controller that supports capacity allocation indicates the number of allocatable capacity blocks in `cc_capabilities.NCBLKS` field. For example, consider a cache with `NCBLKS=8`. In this example, the `RCID=5` has been allocated capacity blocks numbered 0 and 1 for requests with access-type `AT=0` and has been allocated capacity blocks numbered 2 for requests with access-type `AT=1`. The `RCID=3` in this example has been allocated capacity blocks numbered 3 and 4 for both `AT=0` and `AT=1` access-types as separate capacity allocation by access-type is not required for this workload. Further in this example, the `RCID=6` has been configured with the same capacity block allocations as `RCID=3`. This implies that they share a common capacity allocation in this cache but have been associated with different `RCID` to allow differentiated treatment in another capacity and/or bandwidth controller.

	7	6	5	4	3	2	1	0
<code>RCID=3, AT=0</code>	0	0	0	1	1	0	0	0
<code>RCID=3, AT=1</code>	0	0	0	1	1	0	0	0
<code>RCID=5, AT=0</code>	0	0	0	0	0	0	1	1
<code>RCID=5, AT=1</code>	0	0	0	0	0	1	0	0
<code>RCID=6, AT=0</code>	0	0	0	1	1	0	0	0
<code>RCID=6, AT=1</code>	0	0	0	1	1	0	0	0

The `FLUSH_RCID` operation may incur long latency to complete. New requests to the controller by the `RCID` being flushed are allowed. The controller is allowed to deallocate capacity that was allocated after the operation was initiated.



For cache controllers, the `FLUSH_RCID` operation may perform an operation similar to that performed by the RISC-V `CBO.FLUSH` instruction on each cache block that is part of the allocation configured for the `RCID`.

The `FLUSH_RCID` operation may be used as part of reclaiming a previously allocated `RCID` and associating it with a new workload. When such a reallocation is performed the capacity controllers may have capacity allocated by the old workload and thus for a short warmup duration the capacity controller may be enforcing capacity limits that reflect the usage by the old workload. Such warmup durations are typically not statistically significant but if that is not desired then the `FLUSH_RCID` operation may be used to flush and evict capacity allocated to the old workload.

When the `cc_alloc_ctl` is written, the controller may need to perform several actions that may not complete synchronously with the write. A write to the `cc_alloc_ctl` sets the `BUSY` bit to 1 indicating the controller is performing the requested operation. When the `BUSY` bit reads 0 the operation is complete and the `STATUS` field provides a status value (Table 7) of the requested operation.

Table 7. `cc_alloc_ctl.STATUS` field encodings

STATUS	Description
0	Reserved
1	Operation was successfully completed.
2	Invalid operation (OP) requested.
3	Operation requested for an invalid RCID.
4	Operation requested for an invalid AT.
5	Invalid capacity block mask specified.
6-63	Reserved for future standard use.
64-127	Designated for custom use.

Behavior of writes to the `cc_alloc_ctl` or `cc_block_mask` when `BUSY` is 1 is `UNSPECIFIED`. Some implementations may ignore the second write and others may perform the operation determined by the second write. Software must verify that `BUSY` is 0 before writing `cc_alloc_ctl` or `cc_block_mask`.

3.5. Capacity block mask (`cc_block_mask`)

The `cc_block_mask` is a WARL register. When the controller does not support capacity allocation i.e. `NCBLKS` is 0, then this register is read-only 0.

The register has `NCBLKS` Section 3.1 bits each corresponding to one allocatable capacity block in the controller. The width of this register is variable but always a multiple of 64 bits. The width in bits is determined as

$$BMW = \frac{NCBLKS + 63}{64} \times 64$$

Bits `NCBLKS-1:0` are read-write and the bits `BMW-1:NCBLKS` are read-only zero.

To configure capacity allocation limit for an `RCID` and `AT`, the `cc_block_mask` is first programmed with a bit-mask identifying the capacity blocks to be allocated and then the `cc_alloc_ctl` register written to request a `CONFIG_LIMIT` operation for the `RCID` and `AT`. Once a capacity allocation limit has been established, a request may be allocated capacity in a capacity block if the capacity block mask programmed for `RCID` and `AT` associated the request has the corresponding bit set to 1 in the capacity block mask. At least one capacity block must be allocated using `cc_block_mask` when allocating capacity. Allocating overlapping capacity block masks among `RCID` and `AT` is allowed.



A set-associative cache controller that supports capacity allocation by ways may advertize `NCBLKS` to be the number of ways per set in the cache. Allocating capacity in such a cache for an `RCID` and `AT` involves selecting a subset of ways and programming the mask of the selected ways in `cc_block_mask` when requesting the `CONFIG_LIMIT` operation.

To read out the capacity allocation limit for an `RCID` and `AT`, the `READ_LIMIT` operation is requested using `cc_alloc_ctl`. On successful completion of the operation the `cc_block_mask` holds the

configured capacity allocation limit.

Chapter 4. Bandwidth-controller QoS Register Interface

Controllers, such as memory controllers, that support bandwidth allocation and bandwidth usage monitoring provide a memory-mapped register programming interface.

Table 8. Bandwidth-Controller QoS register layout

Offset	Name	Size	Description	Optional?
0	<code>bc_capabilities</code>	8	Capabilities	No
8	<code>bc_mon_ctl</code>	8	Usage monitoring control	Yes
16	<code>bc_mon_ctr_val</code>	8	Monitoring counter value	Yes
24	<code>bc_alloc_ctl</code>	8	Bandwidth allocation control	Yes
32	<code>bc_bw_alloc</code>	8	Bandwidth allocation	Yes

The reset value is 0 for the following registers fields.

- `bc_mon_ctl` - `BUSY` and `STATUS` fields
- `bc_alloc_ctl` - `BUSY` and `STATUS` fields

The reset value is `UNSPECIFIED` for all other registers and/or fields.

The bandwidth controllers at reset allocate all available bandwidth to `RCID` value of 0. When the bandwidth controller supports bandwidth allocation per access-type, one of the supported access-types for `RCID=0` is allocated all available bandwidth and all other access-types share the bandwidth allocation with that access-type. The bandwidth allocation for all other `RCID` values is `UNSPECIFIED`.

4.1. Capabilities (`bc_capabilities`)

The `bc_capabilities` register is a read-only register that holds the bandwidth-controller capabilities.

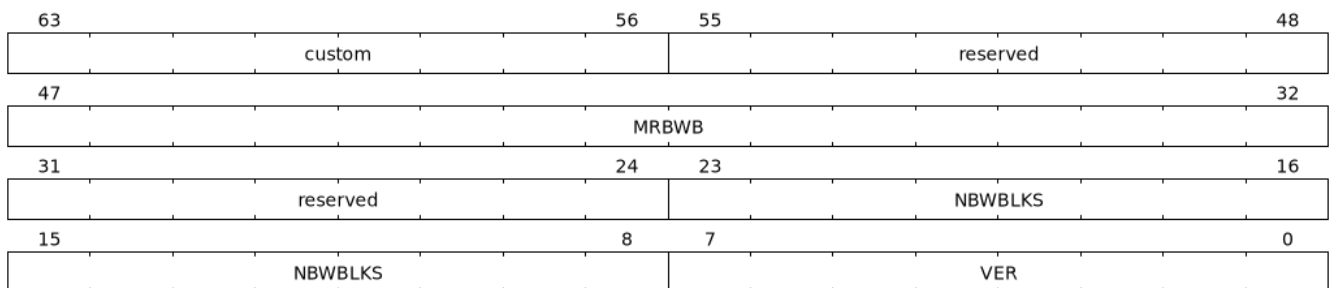


Figure 6. Capabilities register fields

The `VER` field holds the version of the specification implemented by the bandwidth controller. The low nibble is used to hold the minor version of the specification and the upper nibble is used to hold the major version of the specification. For example, an implementation that supports version 1.0 of the specification reports 0x10.

The `NBWBLKS` field holds the total number of available bandwidth blocks in the controller. The

bandwidth represented by an available bandwidth block is **UNSPECIFIED**. The bandwidth controller supports reserving bandwidth in multiples of a bandwidth block and thereby support proportional allocation of bandwidth.

Bandwidth controllers may limit the maximum bandwidth that may be reserved to be smaller than **NBWLKS**. The **MRBWB** field reports the maximum bandwidth blocks that may be reserved.



The bandwidth controller needs to meter the bandwidth usage by a workload to determine if its exceeding its allocations and if required invoke necessary measures to throttle the bandwidth usage by the workload. The instantaneous bandwidth use by a workload may thus undershoot or overshoot the configured allocation. The QoS capabilities are statistical in nature and are typically designed to enforce the configured bandwidth over larger time windows. Not allowing all available bandwidth blocks to be reserved for allocation may allow the bandwidth controller with handle such transient inaccuracies.

4.2. Usage monitoring control (**bc_mon_ctl**)

The **bc_mon_ctl** register is used to control monitoring of bandwidth usage by a **MCID**. When the controller does not support bandwidth usage monitoring the **bc_mon_ctl** register is read-only zero.

63								56
		WPRI						
55								48
		WPRI						
47								40
		WPRI						
39		38						32
BUSY (RO)	STATUS (RO)							
31		29		28		27		24
		WPRI		ATV (WARL)		EVT_ID (WARL)		
23				20		19		16
		EVT_ID (WARL)			MCID (WARL)			
15								8
		MCID (WARL)						
7		5		4		0		
		AT (WARL)		OP (WARL)				

Figure 7. Usage monitoring control (**bc_mon_ctl**)

Bandwidth controllers that support bandwidth usage monitoring implement a usage monitoring counter for each supported **MCID**. The usage monitoring counter may be configured to count a monitoring event. When an event matching the event configured for the **MCID** occurs then the monitoring counter is updated. The event matching may optionally be filtered by the access-type. The monitoring counter for bandwidth usage counts the number of bytes transferred by requests matching the monitoring event as the requests go past the monitoring point.

The **OP** field is used to instruct the controller to perform an operation listed in [Table 9](#).

The **EVT_ID** field and is to program the identifier of the event to count in the counter selected by **MCID**. The **AT** field is used to program the access-type to count.

When the **EVT_ID** for a **MCID** is selected as 0, the counter retains its current value but stops counting.

When **EVT_ID** is programmed with a non-zero and legal value, the counter resets to 0, the **OVF** bit is cleared to 0, and the counter starts counting the programmed events for requests with **MCID** and **AT** (if **ATV** is 1). When **ATV** is 0, the counter counts requests with all access-types and the **AT** value is ignored.

A controller that does not support monitoring by access-type may hardwire the **ATV** and **AT** fields to 0.

Table 9. Usage monitoring operations (OP)

Operation	Encoding	Description
—	0	Reserved for future standard use.
CONFIG_EVENT	1	Configure the counter selected by MCID to count the event selected by EVT_ID , AT , and ATV . The EVT_ID encodings are listed in Table 10.
READ_COUNTER	2	Snapshot the value of the counter selected by MCID into bc_mon_ctr_val register. The EVT_ID , AT and ATV fields are not used by this operation.
—	3-23	Reserved for future standard use.
—	24-31	Designated for custom use.

Table 10. Bandwidth monitoring event ID (EVT_ID)

Event ID	Encoding	Description
None	0	Counter does not count and retains its value.
Total Read and Write byte count	1	Counter is incremented by the number of bytes transferred by a matching read or write request as the requests go past the monitor.
Total Read byte count	2	Counter is incremented by the number of bytes transferred by a matching read request as the requests go past the monitor
Total Write byte count	3	Counter is incremented by the number of bytes transferred by a matching write request as the requests go past the monitor
—	4-127	Reserved for future standard use.
—	128-256	Designated for custom use.

When the **bc_mon_ctl** is written, the controller may need to perform several actions that may not complete synchronously with the write. A write to the **bc_mon_ctl** sets the **BUSY** bit to 1 indicating the controller is performing the requested operation. When the **BUSY** bit reads 0 the operation is complete and the **STATUS** field provides a status value (Table 11) of the requested operation.

Table 11. **bc_mon_ctl.STATUS** field encodings

STATUS	Description
0	Reserved

STATUS	Description
1	Operation was successfully completed.
2	Invalid operation (OP) requested.
3	Operation requested for invalid MCID.
4	Operation requested for invalid EVT_ID.
5	Operation requested for invalid AT.
6-63	Reserved for future standard use.
64-127	Designated for custom use.

Behavior of writes to the `bc_mon_ctl` when `BUSY` is 1 is `UNSPECIFIED`. Some implementations may ignore the second write and others may perform the operation determined by the second write. Software must verify that `BUSY` is 0 before writing `bc_mon_ctl`.

4.3. Monitoring counter value (`bc_mon_ctr_val`)

The `bc_mon_ctr_val` is a read-only register that holds a snapshot of the counter requested by `READ_COUNTER` operation.

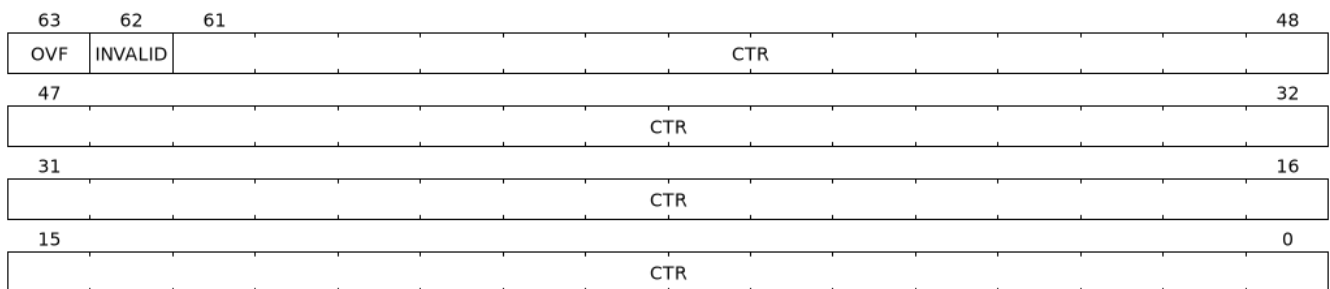


Figure 8. Usage monitoring counter value (`bc_mon_ctr_val`)

The counter is valid if the `INVALID` field is 0. The counter may be marked `INVALID` if the controller for `UNSPECIFIED` reasons determine the count to be not valid. The counters marked `INVALID` may become valid in future. If an unsigned integer overflow of the counter occurs then the `OVF` bit is set to 1.



A counter may be marked as `INVALID` if the controller has not been able to establish an accurate counter value for the monitored event.

The counter provides the byte transferred by requests matching the `EVT_ID` as the requests go past the monitoring point. A bandwidth value may be determined by reading the byte count value at two instances of time `T1` and `T2`. If the value of the counter at time `T1` was `B1` and at time `T2` is `B2` then the bandwidth is as follows. The frequency of the time source is T_{freq} .

$$Bandwidth = T_{freq} \times \frac{B2 - B1}{T2 - T1}$$

The width of the counter is `UNSPECIFIED`.



The width of the counter is `UNSPECIFIED` but is recommended to be wide enough to not cause more than one overflow per sample when sampled at a frequency of 1

Hz.

If an overflow was detected then software may discard that sample and reset the counter and overflow indication by reprogramming the event using **CONFIG_EVENT** operation.

4.4. Bandwidth Allocation control (**bc_alloc_ctl**)

The **bc_alloc_ctl** register is used to control allocation of bandwidth to a **RCID** per **AT**. If a controller does not support bandwidth allocation then the register is read-only zero. If the controller does not support bandwidth allocation per access-type then the **AT** field is read-only zero.

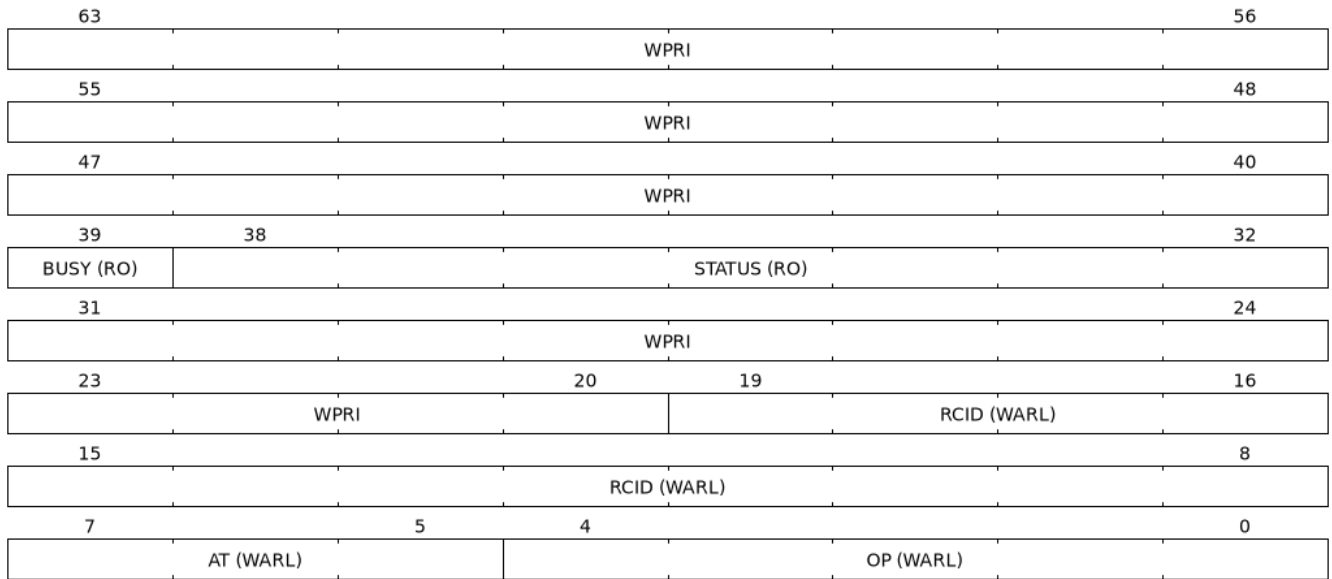


Figure 9. Bandwidth allocation control (**bc_alloc_ctl**)

The **OP** field is used to instruct the bandwidth controller to perform an operation listed in [Table 12](#). The **bc_alloc_ctl** register is used in conjunction with the **bc_bw_alloc** register to perform bandwidth allocation operations. When the requested operation uses the operands configured in **bc_bw_alloc**, software must first program the **bc_bw_alloc** register with the operands for the operation before requesting the operation.

Table 12. Bandwidth allocation operations (**OP**)

Operation	Encoding	Description
—	0	Reserved for future standard use.
CONFIG_LIMIT	1	The CONFIG_LIMIT operation is used to establish reserved bandwidth allocation for requests by RCID and of access-type AT . The bandwidth allocation is specified in bc_bw_alloc register.
READ_LIMIT	2	The READ_LIMIT operation is used to read back the previously configured bandwidth allocation for requests by RCID and of type AT . The current configured allocation is written to bc_bw_alloc register on completion of the operation.
—	3-23	Reserved for future standard use.

Operation	Encoding	Description
—	24-31	Designated for custom use.

A bandwidth allocation must be configured for each supported access-type by the controller. When differentiated bandwidth allocation based on access-type is not required, one of the access-types may be designated to hold a default bandwidth allocation and the other access-types configured to share the allocation with the default access-type. The behavior is **UNSPECIFIED** if bandwidth is not allocated for each access-type supported by the controller.

When the **bc_alloc_ctl** is written, the controller may need to perform several actions that may not complete synchronously with the write. A write to the **bc_alloc_ctl** sets the **BUSY** bit to 1 indicating the controller is performing the requested operation. When the **BUSY** bit reads 0 the operation is complete and the **STATUS** field provides a status value (Table 13) of the requested operation.

Table 13. **bc_alloc_ctl.STATUS** field encodings

STATUS	Description
0	Reserved
1	Operation was successfully completed.
2	Invalid operation (OP) requested.
3	Operation requested for an invalid RCID .
4	Operation requested for an invalid AT .
5	Invalid/unsupported reserved bandwidth blocks requested.
6-63	Reserved for future standard use.
64-127	Designated for custom use.

4.5. Bandwidth allocation (**bc_bw_alloc**)

The **bc_bw_alloc** is used to program reserved bandwidth blocks (**Rbwb**) for an **RCID** for requests of access-type **AT**. The bandwidth is allocated in multiples of bandwidth blocks and the value in **Rbwb** must be at least 1 and must not exceed **MRBWB** else the operation fails with **STATUS=5**. The sum of **Rbwb** allocated across all **RCID** must not exceed **MRBWB** else the operation fails with **STATUS=5**.

63						56
				WPRI		
55						48
				WPRI		
47						40
				WPRI		
39						32
				WPRI		
31	30	29	28	27		24
	WPRI	useDef (WARL)	isDef (WARL)		Mweight (WARL)	
23			20	19		16
		Mweight (WARL)			WPRI	
15						8
				Rbwb (WARL)		
7						0
				Rbwb (WARL)		

Figure 10. Bandwidth allocation (bc_bw_alloc)

Bandwidth allocation is typically enforced by the bandwidth controller over finite accounting windows. The process involves measuring the bandwidth consumption over an accounting window and using the measured bandwidth to determine if an **RCID** is exceeding its bandwidth allocations for each access-types. The specifics of how the accounting window is implemented is **UNSPECIFIED** but is expected to provide a statistically accurate control of the bandwidth usage over a few accounting intervals.

The **Rbwb** represents the bandwidth that is made available to a **RCID** for requests matching **AT** even when all other **RCID** are using their full allocation of bandwidth.

If there is non-reserved or unused bandwidth available in an accounting interval then additional bandwidth may be made available to **RCID** that contend for that bandwidth. The non-reserved or unused bandwidth is proportionately shared by the contending **RCIDs** using the configured **Mweight**. The **Mweight** parameter is a number between 0 and 255. A larger weight implies a greater fraction of the bandwidth. The sharing of non-reserved bandwidth is not differentiated by access-type. The **Mweight** parameter must be programmed identically for all access-types. If this parameter is programmed differently for each access-type then the controller may use the parameter configured for any of the access-types but the behavior is otherwise well defined. The share of unused bandwidth made available to **RCID=x** when it contends with another **RCID** is determined by the **Mweight** of **RCID=x** divided by the sum of **Mweight** of all other contending **RCID**. This ratio **P** is as follows:

$$P = \frac{Mweight_x}{\sum_{r=1}^n Mweight_r}$$



The bandwidth enforcement is typically work-conserving and allows unused bandwidth to be used by requestors even if they have consumed their **Rbwb**.

When contending for unused bandwidths the weighted share is typically computed among the **RCIDs** that are actively generating requests in that accounting interval.

When **isDef** is 1, the bandwidth allocation programmed for that access-type is considered the

default bandwidth allocation. If unique bandwidth allocation is not required for an access-type then the `useDef` may be used to indicate that requests of this access-type share bandwidth allocated for the default access-type. Only one access-type may be designated as the default access-type. The `isDef` and `useDef` fields are reserved if the controller does not support bandwidth allocation per access-type. If more than one access-types have been configured with `isDef` programmed to 1 then the implementation may associate an access-type programmed with `useDef` with any of those but the behavior is otherwise well defined. If an access-type is configured with `useDef` as 1 but no other access-type is configured to be the default access-type then the implementation may select any of the access-types as the default access-type but the behavior is otherwise well defined.