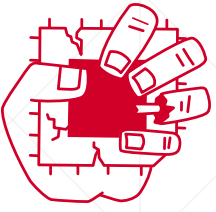


ZombieLoad

Cross-Privilege-Boundary Data Sampling

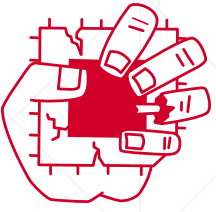
Michael Schwarz, Moritz Lipp, **Daniel Moghimi**, Jo Van Bulck, Julian Stecklina,
Thomas Prescher, Daniel Gruss



whoami

- Daniel Moghimi (@danielmgmi)
- Computer Security Enthusiast *since 2010*
 - Reverse Engineering
 - Application Security
- PhD Student *since 2017*
 - Microarchitectural Security
 - Side Channels
 - Breaking Cryptographic Implementations

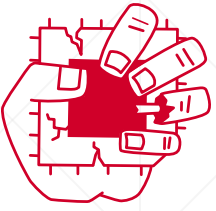




2018: Meltdown Attack?

```
char secret = *(char *) 0xffffffff81a0123;  
printf("%c\n", secret);
```

- Segmentation **Fault!** Inaccessible Address.

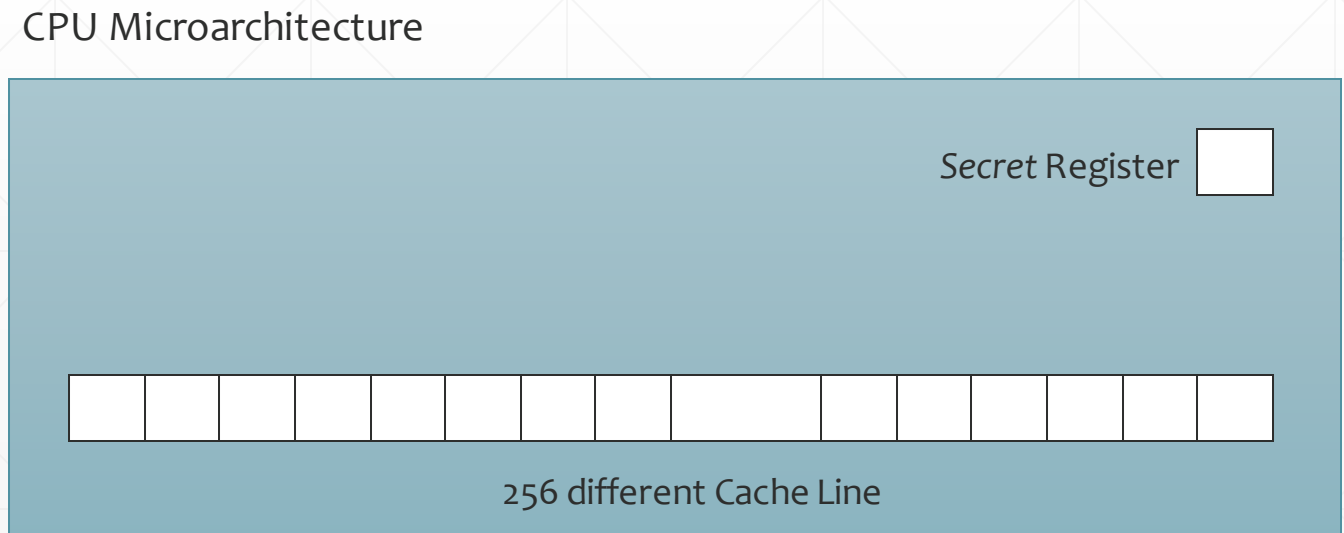
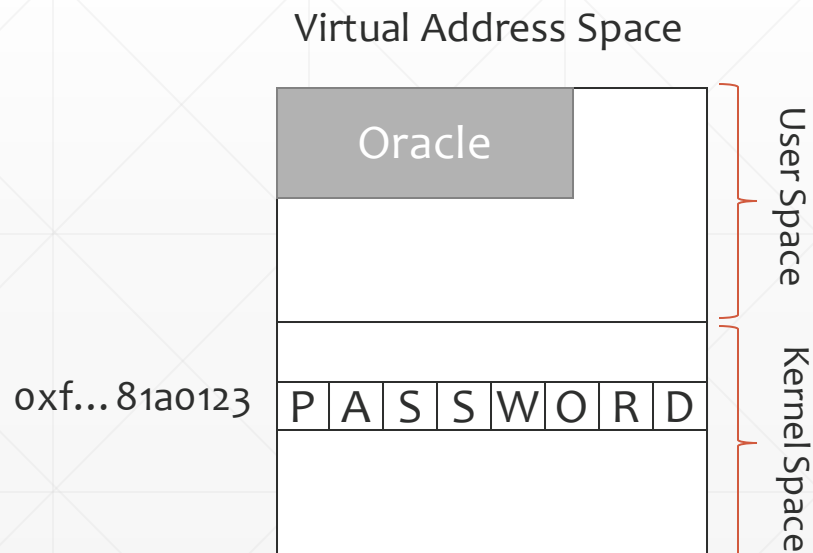


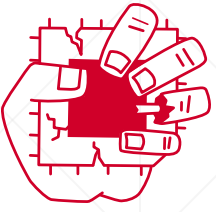
2018: Meltdown Attack?



```
char oracle[4096 * 256];
char secret = *(char *) 0xffffffff81a0123;
char x = oracle[secret * 4096];

for(char secret = 0; i < 256; i++){
    if(flush_reload(oracle[i * 4096]) < threshold)
        printf("%c\n", i);
}
```



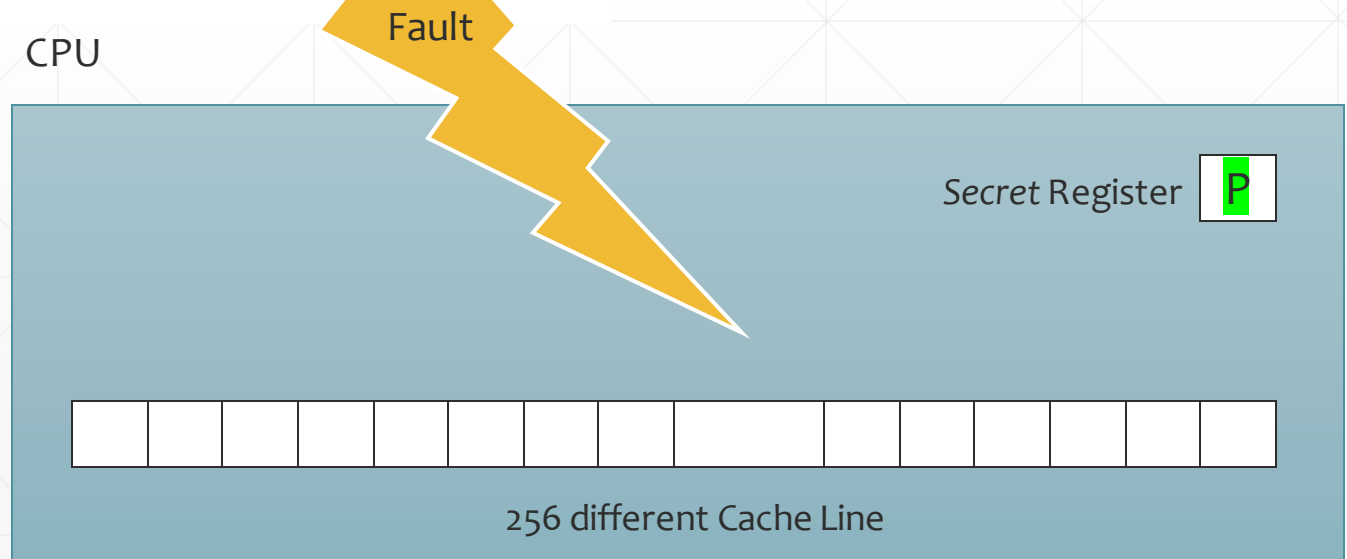
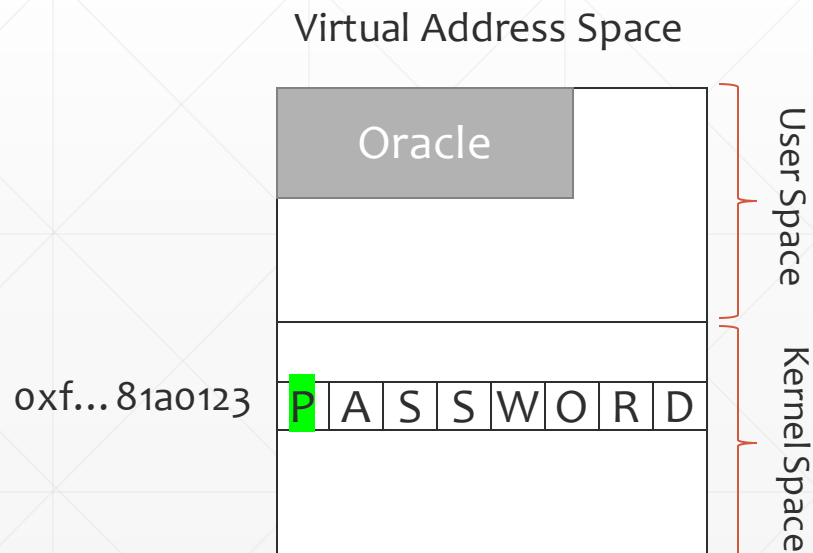


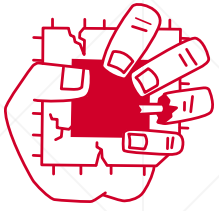
2018: Meltdown Attack?



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char secret = *(char *) 0xffffffff81a0123;
char x = oracle[secret * 4096];

for(char secret = 0; i < 256; i++){
    if(flush_reload(oracle[i * 4096]) < threshold)
        printf("%c\n", i);
}
```



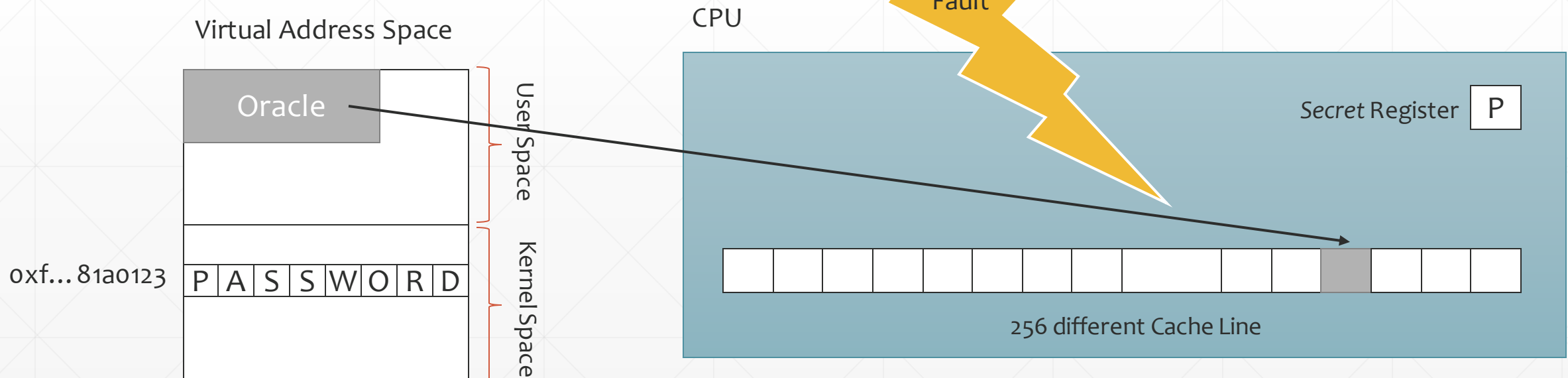


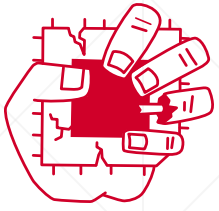
2018: Meltdown Attack?



→
`char oracle[4096 * 256];
char secret = *(char *) 0xffffffff81a0123;
char x = oracle[secret * 4096];`

```
for(char secret = 0; i < 256; i++){  
    if(flush_reload(oracle[i * 4096]) < threshold)  
        printf("%c\n", i);  
}
```





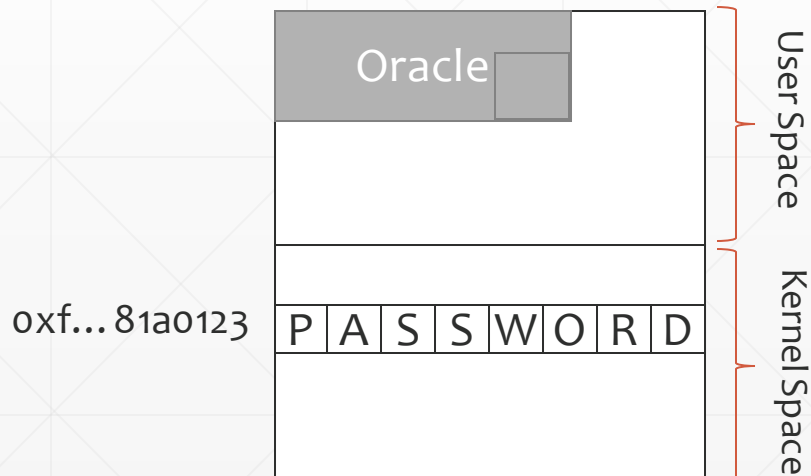
2018: Meltdown Attack?

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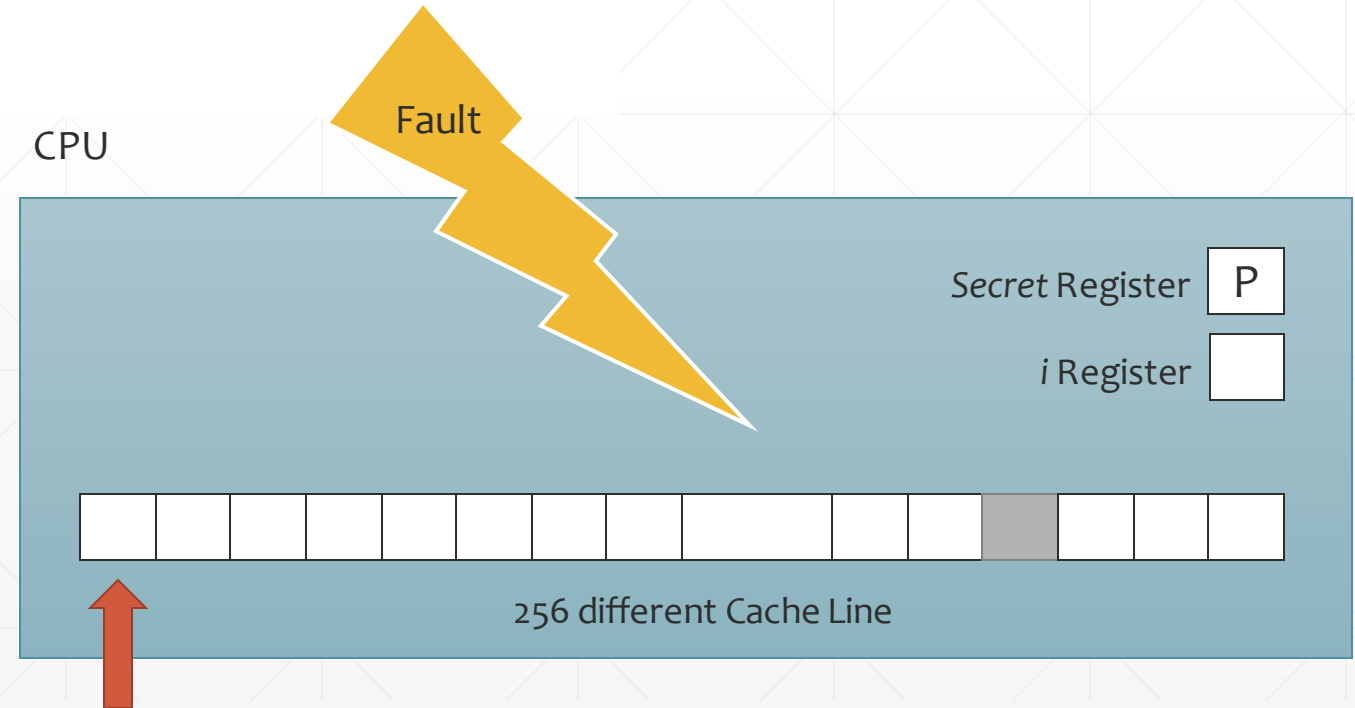


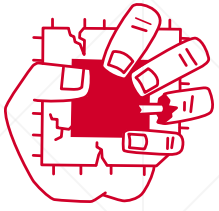
```
for(char secret = 0; i < 256; i++){  
    if(flush_reload(oracle[i * 4096]) < threshold)  
        printf("%c\n", i);  
}
```

Virtual Address Space



CPU





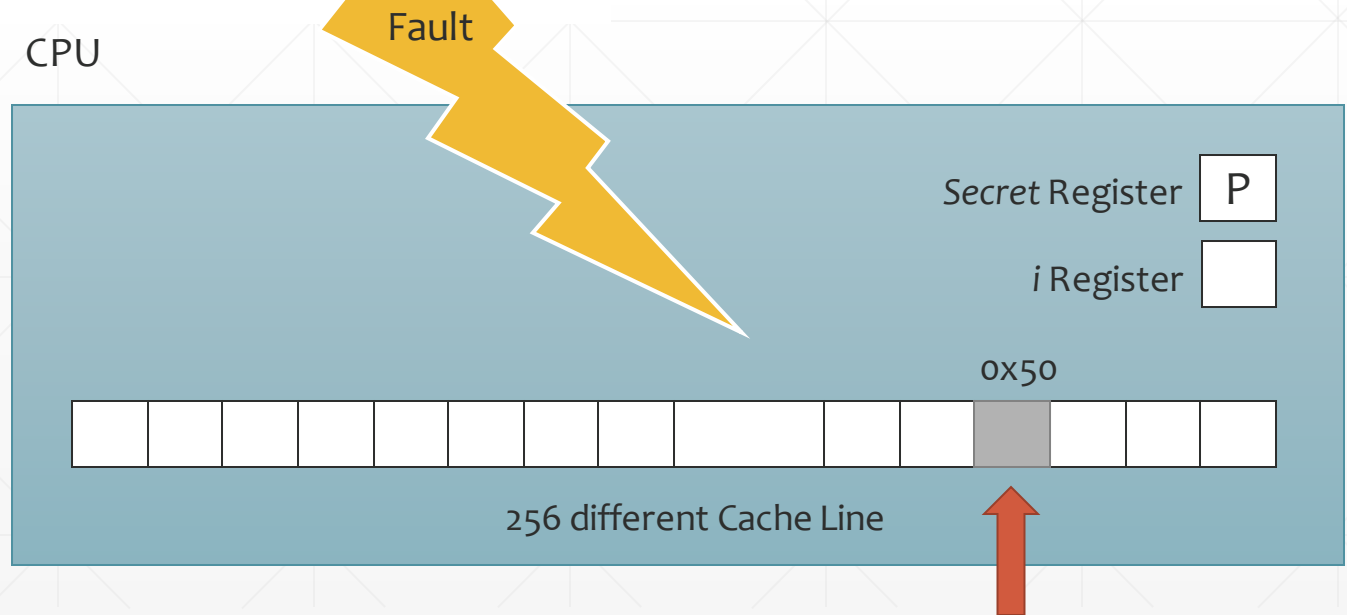
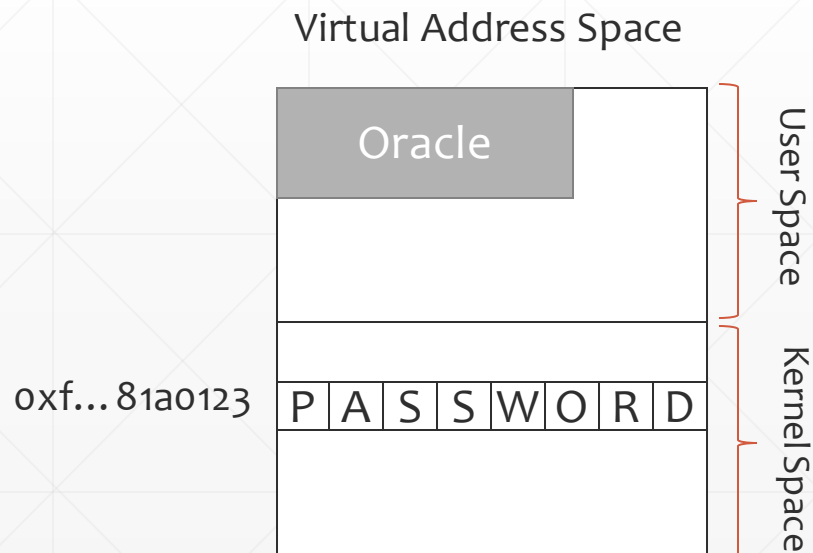
2018: Meltdown Attack?

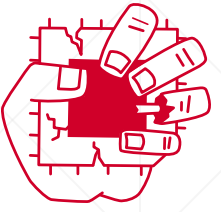
```
char oracle[4096 * 256];  
char secret = *(char *) 0xffffffff81a0123;  
char x = oracle[secret * 4096];
```



➔

```
for(char secret = 0; i < 256; i++){  
    if(flush_reload(oracle[i * 4096]) < threshold)  
        printf("%c\n", i);  
}
```

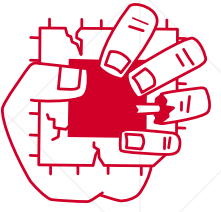


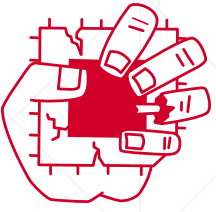


Meltdown-style Attacks !!!

- What if we dereference addresses that causes other faults/assists?
- What if the Microarchitecture is in a specific state during the fault?
- Where does this data leak from?!

ZombieLoad !?!





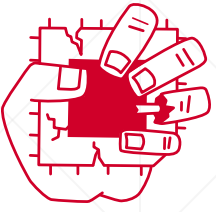
ZombieLoad !?!

```
mov 0x401234, %rsi
```

```
mov (%rsi), %rax
```

Virtual Address

0x000401	234
----------	-----



ZombieLoad !?!

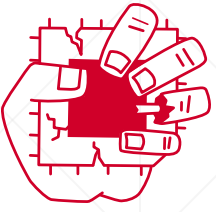
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Virtual Address

0x000401	234
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ZombieLoad !?!

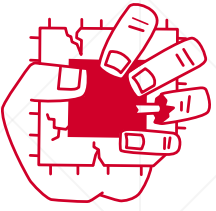
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Virtual Address

0x000401	234
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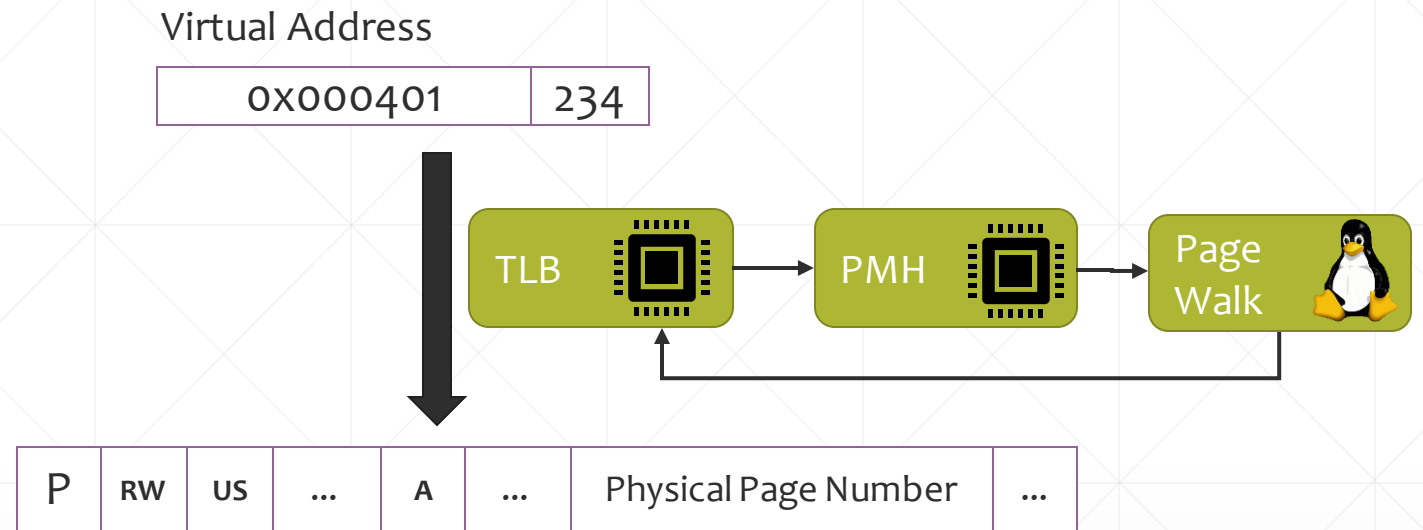


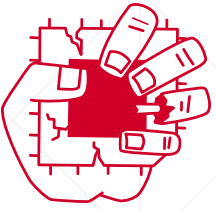


ZombieLoad !?!

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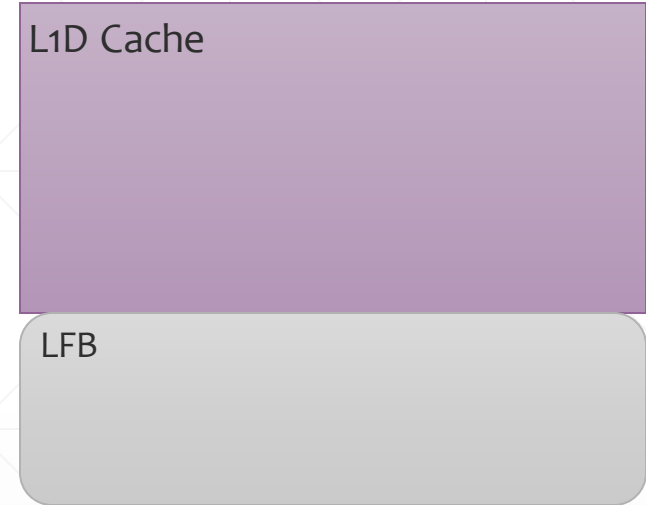
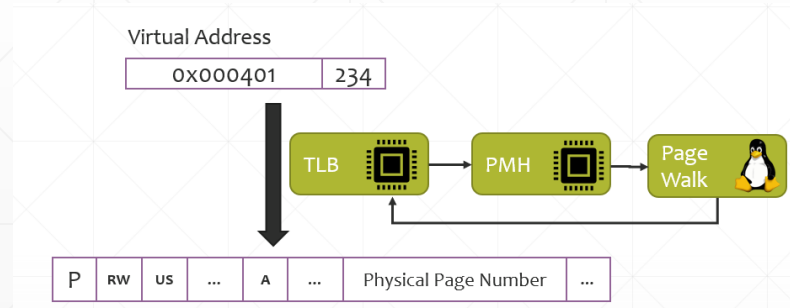


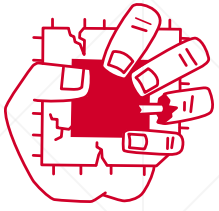


ZombieLoad !?!

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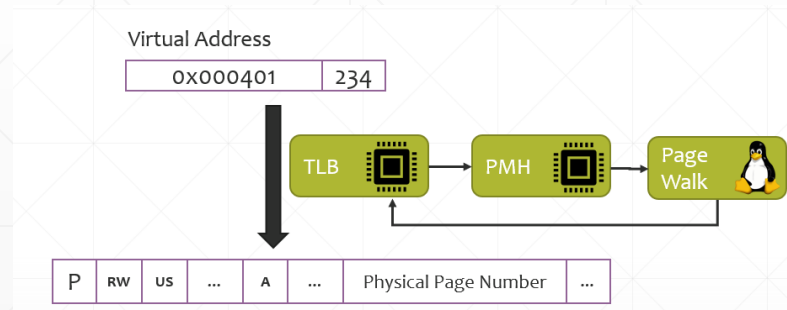
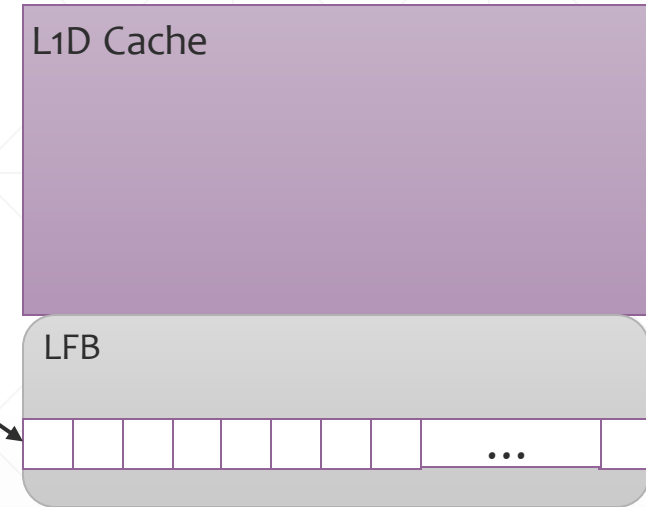


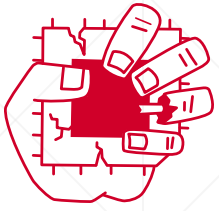


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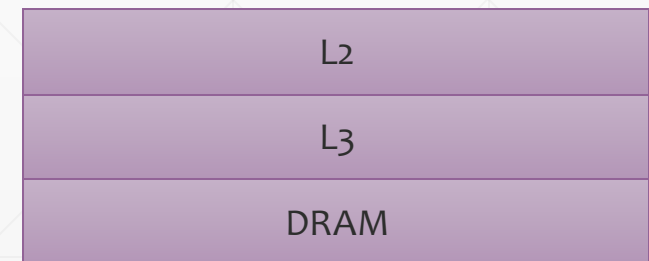
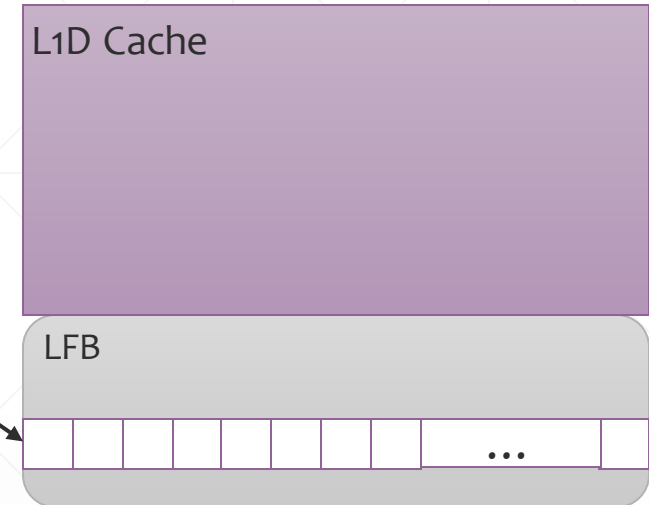
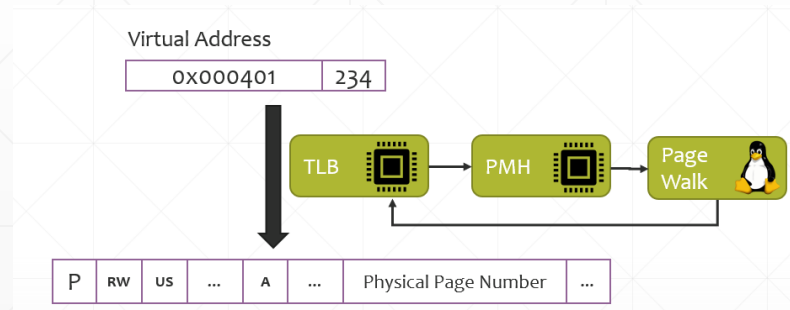


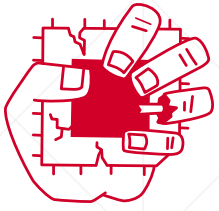


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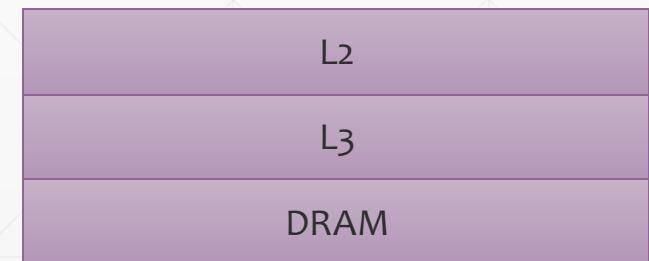
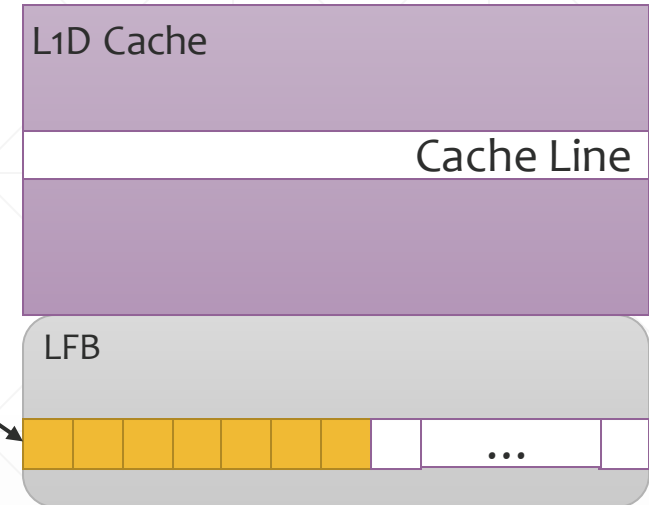
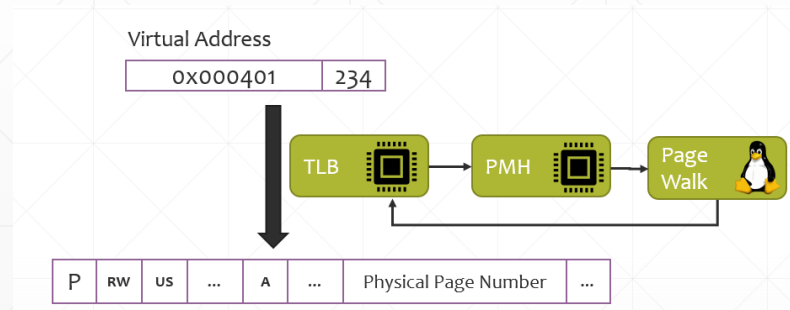


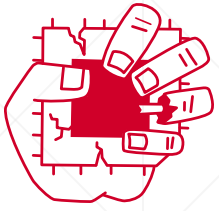


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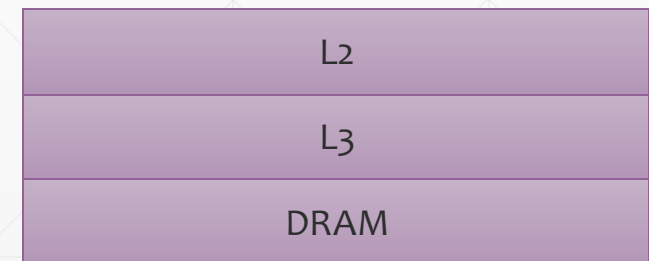
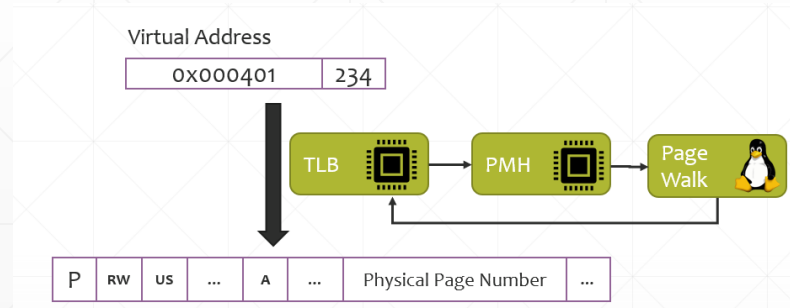
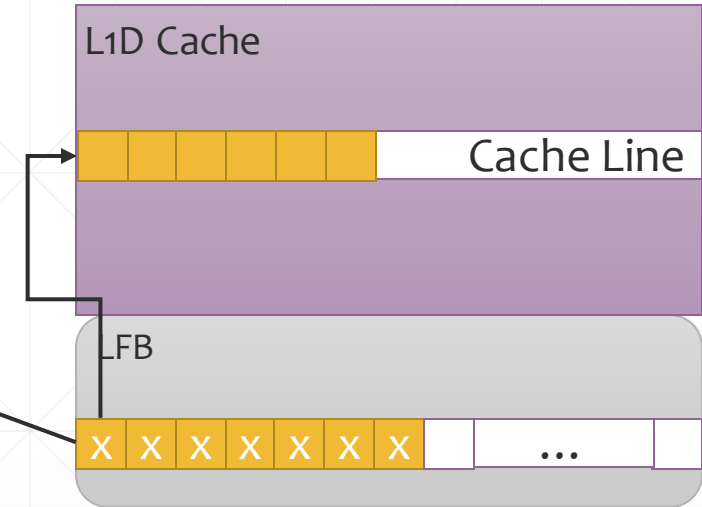


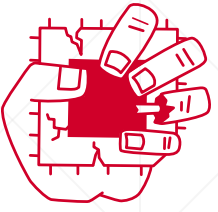


ZombieLoad !?!

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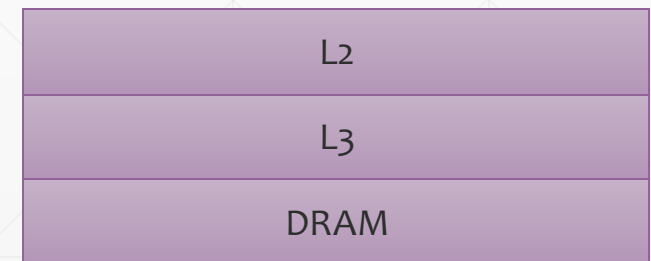
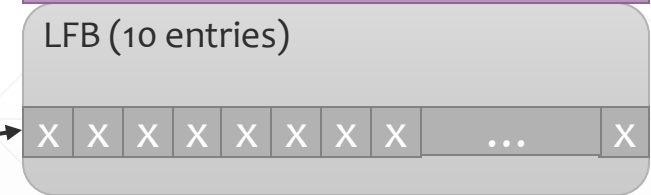
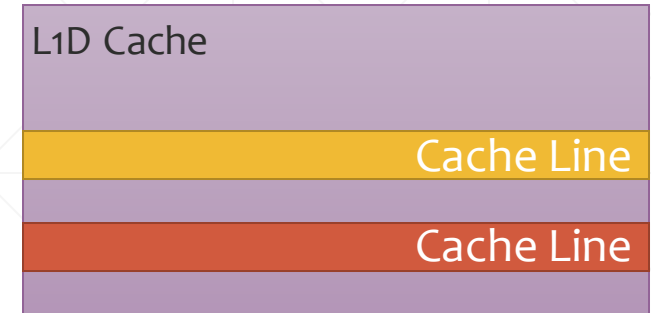
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mov (%rsi), %rax
```

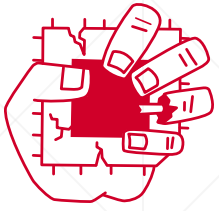




ZombieLoad Attack !?!

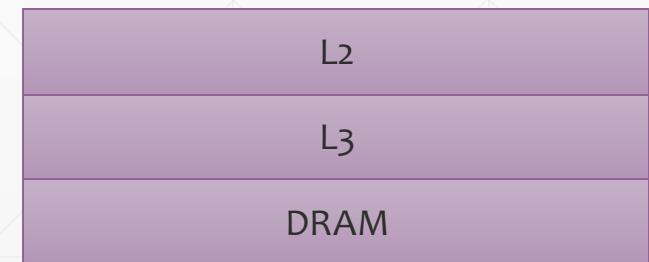
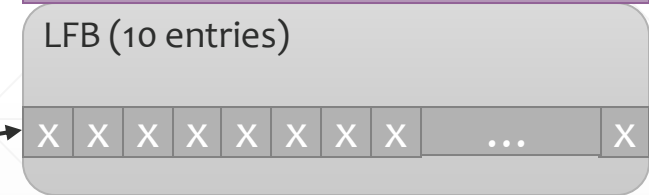
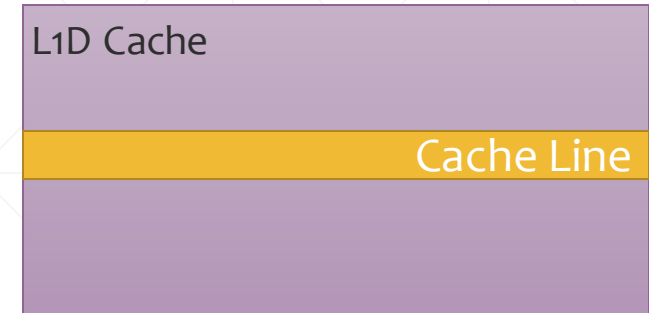
P	RW	US	...	A	...	Physical Page Number	...
---	----	----	-----	---	-----	----------------------	-----

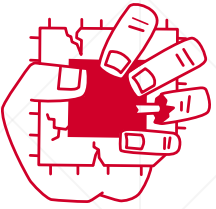




ZombieLoad Attack !?!

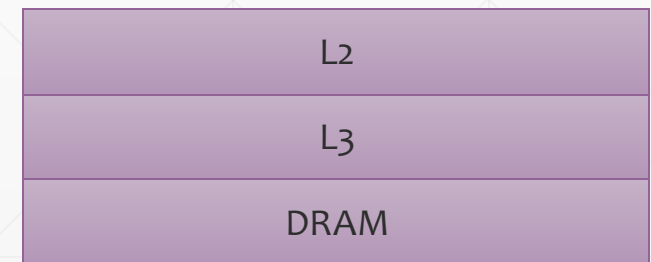
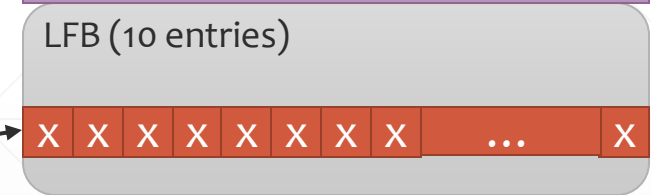
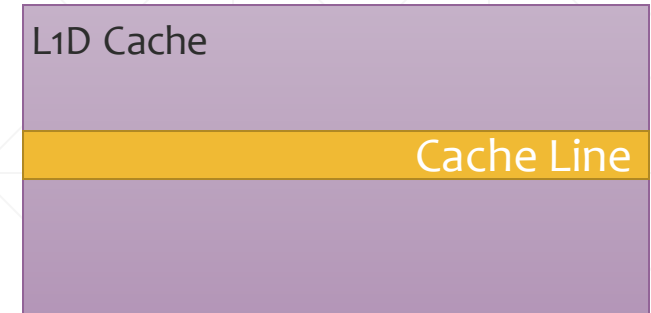
P	RW	US	...	A	...	Physical Page Number	...
---	----	----	-----	---	-----	----------------------	-----

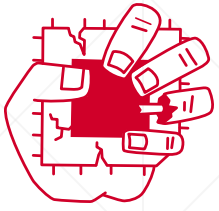




ZombieLoad Attack !?!

P	RW	US	...	A	...	Physical Page Number	...
---	----	----	-----	---	-----	----------------------	-----





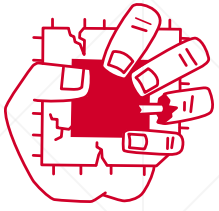
ZombieLoad Attack !?!

P	RW	US	...	A	...	Physical Page Number	...
---	----	----	-----	---	-----	----------------------	-----



X X X X



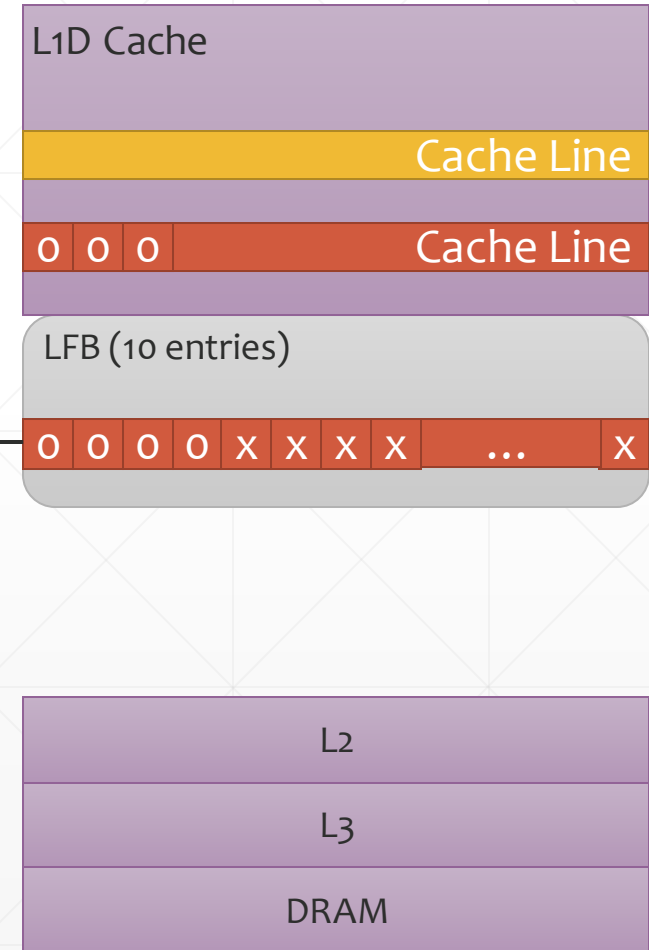


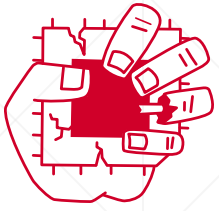
ZombieLoad Attack !?!

P	RW	US	...	A	...	Physical Page Number	...
---	----	----	-----	---	-----	----------------------	-----



X X X X



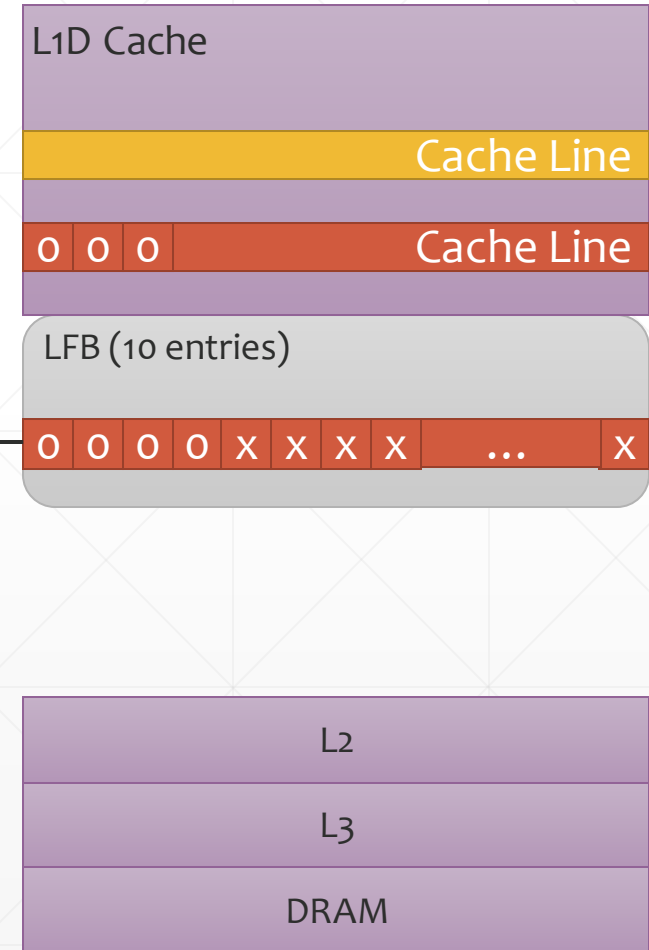


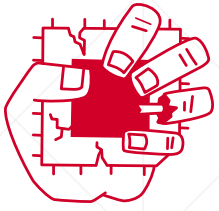
ZombieLoad Attack !?!

P	RW	US	...	A	...	Physical Page Number	...
---	----	----	-----	---	-----	----------------------	-----



X X X X



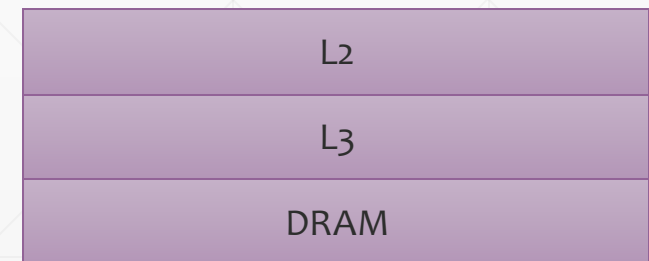
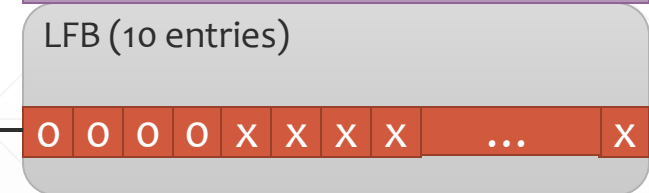
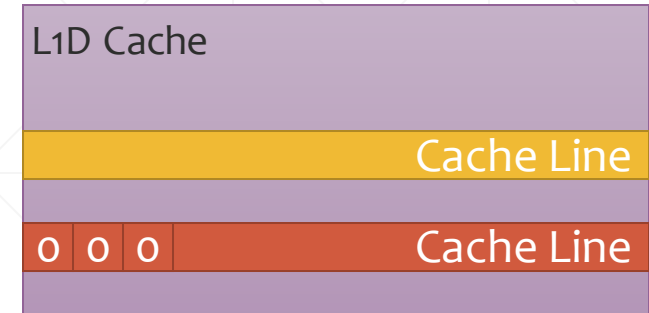


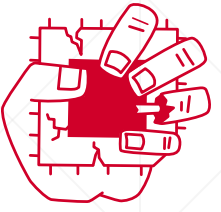
ZombieLoad Attack !?!



Variant 1

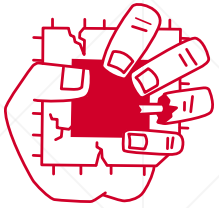
Variant 3





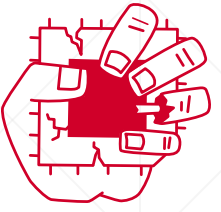
Microcode Assist on 'A' Bit

- The CPU tells the OS if a page has been accessed or not by setting the Access Bit
- The OS can clear the bit and causes a microcode assist
- The microcode assist flushes the pipeline while setting the access bit



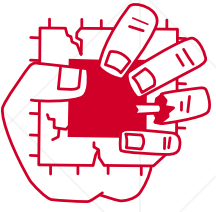
ZombieLoad VS. other Meltdown-Style Attacks

	Page Number			Page Offset	
Meltdown	51	Physical	12	11	0
	47	Virtual	12		
Foreshadow	51	Physical	12	11	0
	47	Virtual	12		
Fallout	51	Physical	12	11	0
	47	Virtual	12		
ZombieLoad	51	Physical	12	11	6
	47	Virtual	12		



What can we do with this data leakage?

- Attack across Process Context Switches
- Attack across Simultaneous Multithreading (SMT) AKA. Intel Hyperthreading

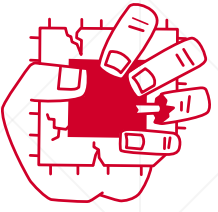


Data Sampling - Domino Attack

- We may leak bytes of data from other unimportant fill buffer entries
- Leak domino bytes to perform error correction

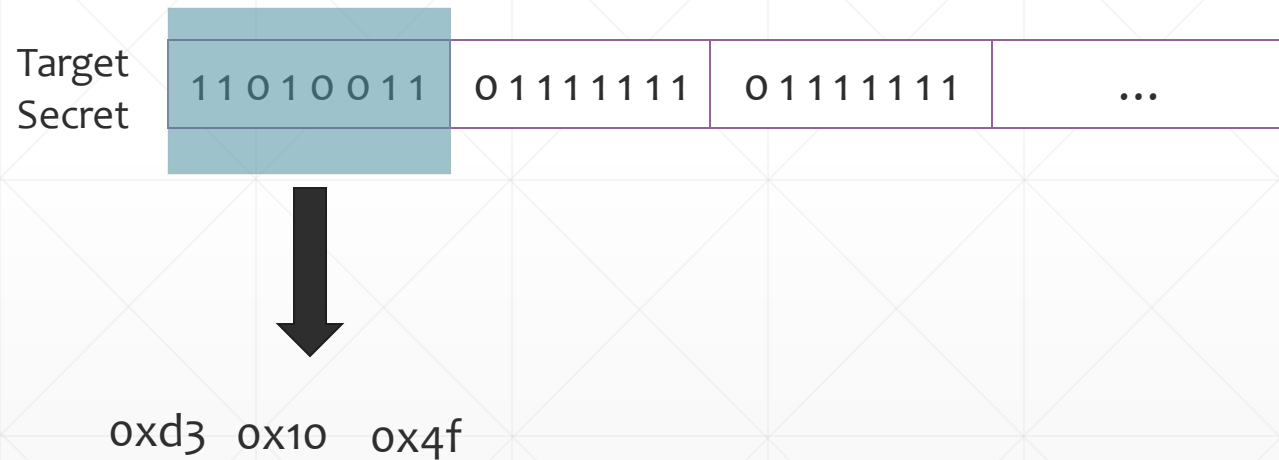
Target
Secret

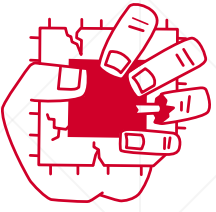
11010011	01111111	01111111	...
----------	----------	----------	-----



Data Sampling - Domino Attack

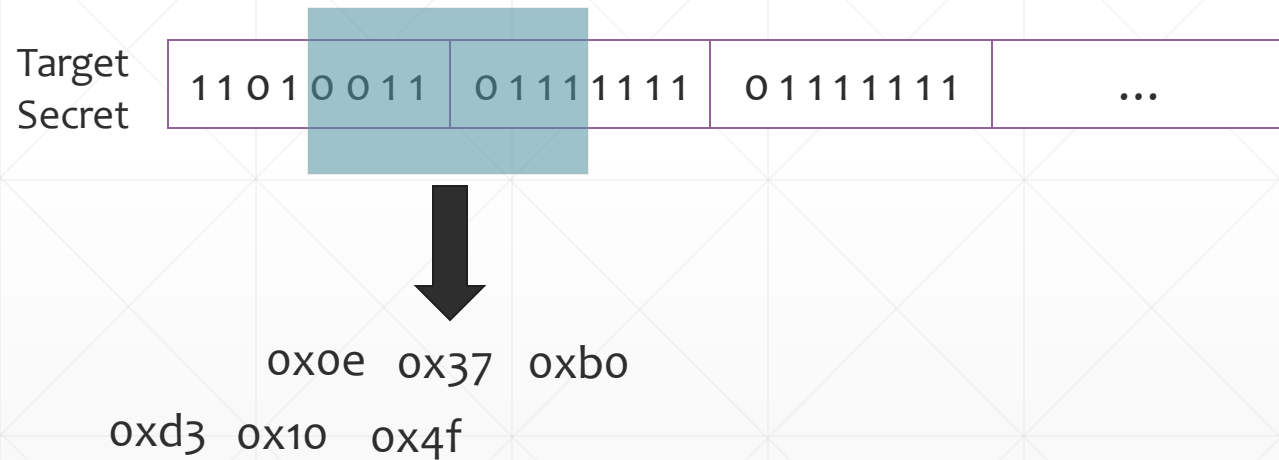
- We may leak bytes of data from other unimportant fill buffer entries
- Leak domino bytes to perform error correction

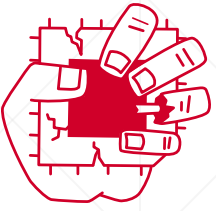




Data Sampling - Domino Attack

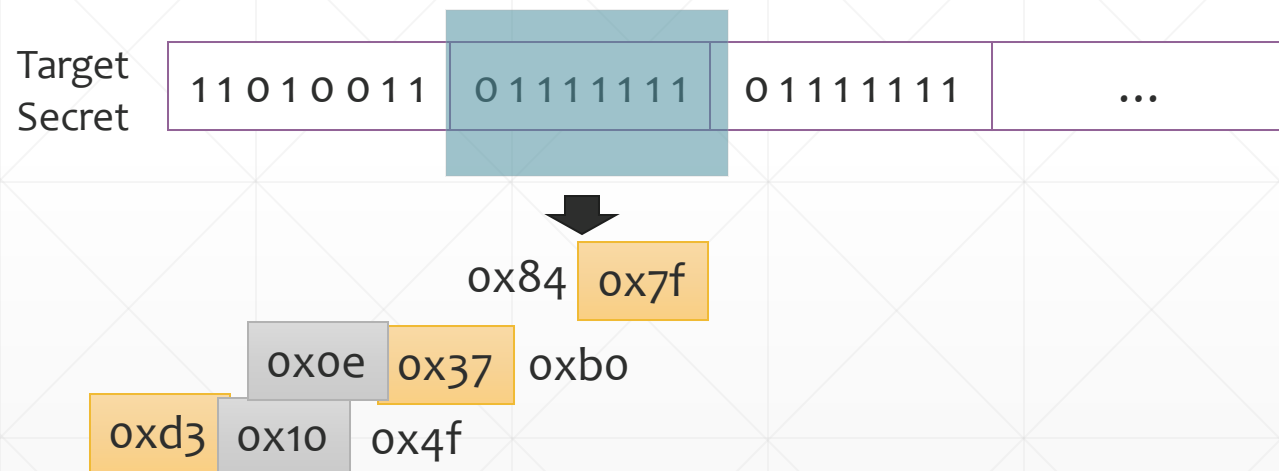
- We may leak bytes of data from other unimportant fill buffer entries
- Leak domino bytes to perform error correction

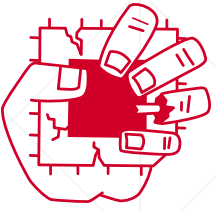




Data Sampling - Domino Attack

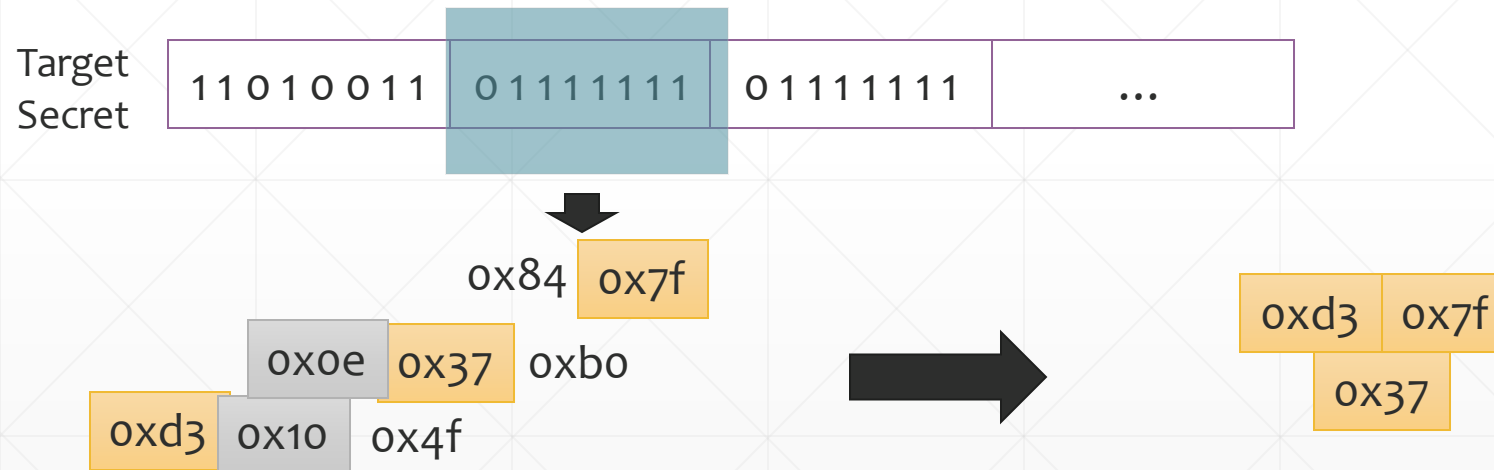
- We may leak bytes of data from other unimportant fill buffer entries
- Leak domino bytes to perform error correction

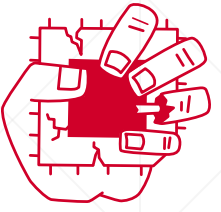




Data Sampling - Domino Attack

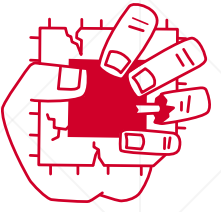
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Recovering Intel SGX Sealing Key

- Intel SGX allow developers to have hardware support for TEE
- Malicious OS is part of the threat model
- We can read register values of a trusted enclave with help of a malicious OS

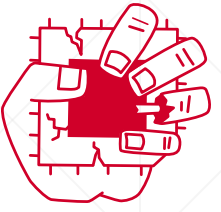


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sgx-step

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mov  
add  
xor  
mov 0x4142434445464748, %rax  
call  
nop  
jmp
```

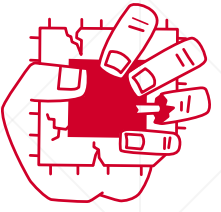


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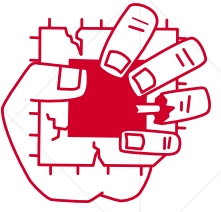


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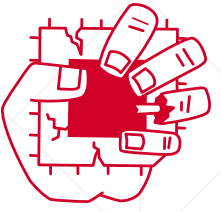


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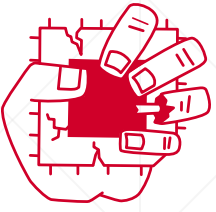
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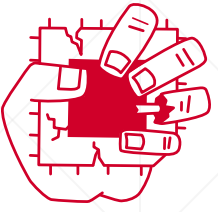
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Try to Execute



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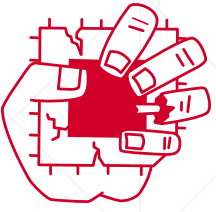
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Try to Execute



Exception



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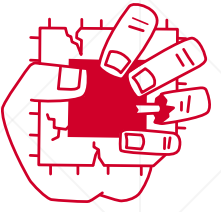
Mark Non-Executable

Try to Execute

Handle Exception

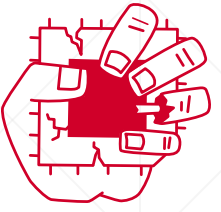
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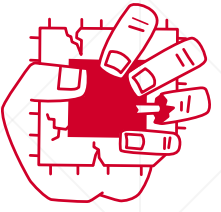


Recovering Intel SGX Sealing Key

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- Repeated Context Switch in the transient domain w/ the same register values



ZombieLoad Demo! Recovering Linux Shadow



Is there any Mitigation?

- Intel suggested an instruction sequence to fill all the buffers across context switch
- Disable hyperthreading

Questions?!



<https://zombieloadattack.com/>