

## Exercise 9 Answer Sheet — Logic Theory (2), 80424

June 5, 2025



## Question 1

### Part a

We will show that the set,

$$H = \{e \in \mathbb{N} \mid (e, 0) \in \text{dom } U\}$$

is recursively-enumerable but not recursive.

*Proof.*  $H$  is  $\Sigma_1^0$  by the definition of  $U$ , implies that it is recursively-enumerable, then it is sufficient to show that it is not recursive. Let  $f : \mathbb{N} \rightarrow \mathbb{N}$  be the function defined by,

$$f(x) = x.$$

This is indeed a recursive function (even primitive-recursive), and for every  $e \in D$ , when  $D = \{e \in \mathbb{N} \mid (e, e) \in \text{dom } U\}$ , we get  $f(e) = e \in H$ , as  $e$  witnessing bound larger than 0. It follows that  $f$  is a recursive reduction of  $D$  to  $H$  and implies that  $H$  is not recursive.  $\square$

### Part b

We will conclude that  $H^C = \mathbb{N} \setminus H$  is not recursively-enumerable.

*Proof.*  $H$  is a relation, meaning that if it were to be  $\Pi_1^0$  it would be  $\Delta_1^0$  and recursive. If  $H^C$  is recursively-enumerable then  $H^C = \{e \in \mathbb{N} \mid (e, 0) \in \text{dom}(\neg U)\}$  is  $\Sigma_1^0$ , and implying by negation that  $H = \neg H^C$  (in the sense of relations) is  $\Pi_1^0$ . But  $H$  is recursively-enumerable and  $\Sigma_1^0$ , meaning that it is indeed  $\Delta_1^0$ , in contradiction to its being not recursive.  $\square$

### Part c

Let us define the set,

$$T = \{e \in \mathbb{N} \mid \{x \mid (e, x) \in \text{dom } U\} = \mathbb{N}\}$$

We will show that  $T$  is not recursively-enumerable.

*Proof.* We assume in contradiction that  $T$  is recursively-enumerable.

By the equivalence to recursively-enumerable functions proposition we can assume that there is a function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that  $f$  is primitive-recursive and  $T = \{f(x) \mid x \in \mathbb{N}\}$ . Let us define  $g : \mathbb{N}^2 \rightarrow \mathbb{N}$  by,

$$g(n, x) = U(f(n), x) + 1$$

$g$  is indeed total recursive function as  $U$  is recursive and  $f$  is total such that  $U(f(n), \cdot)$  is total recursive function. Then by Kleene's recursion theorem there is  $e \in \mathbb{N}$  such that,

$$U(e, x) = g(e, x) = U(f(e), x) + 1.$$

$\square$