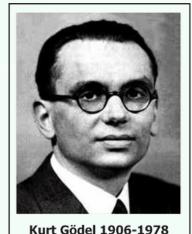
串 Karp-Rabin算法: 串即是数

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All things are numbers. - Pythagoras

凡物皆数: Gödel Numbering

- ❖ 逻辑系统的符号、表达式、公式、命题、定理、公理等,均可表示为自然数
- ❖ 每个有限维的自然数向量(包括字符串),都唯一对应于某个自然数
- ❖ 素数序列: p(k) = 第k个素数 = 2, 3, 5, 7, 11, 13, 17, 19, ...



$$\langle a_1, a_2, \dots, a_n \rangle \rightleftharpoons p(1)^{1+a_1} \times p(2)^{1+a_2} \times \dots \times p(n)^{1+a_n}$$

 $\langle \, | \, 3 | \, \, | \, 1 | \, \, | \, 4 | \, \, | \, 1 | \, \, | \, 5 | \, \, | \, 9 | \, \, |$

$$2^{4} \times 3^{2} \times 5^{5} \times 7^{2} \times 11^{6} \times 13^{10} \times 17^{3} \times 19^{7}$$

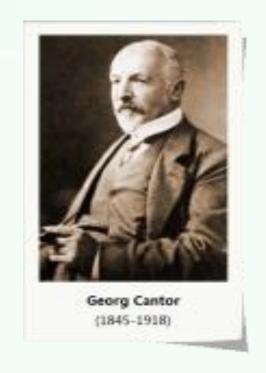
"godel" =
$$2^{1+7} \times 3^{1+15} \times 5^{1+4} \times 7^{1+5} \times 11^{1+12} = 139,869,560,310,664,817,087,943,919,200,000$$

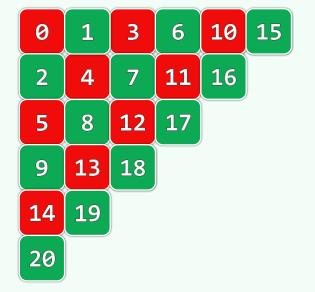
❖ 若果真如RAM模型所假设的字长无限,则只需一个寄存器即可...

凡物皆数: Cantor Numbering

$$cantor_2(i,j) = [(i+j)^2 + 3 \cdot i + j]/2$$

 $cantor_2(2,3) = [(2+3)^2 + 3 \cdot 2 + 3]/2 = 17$
 $cantor_2(3,2) = [(3+2)^2 + 3 \cdot 3 + 2]/2 = 18$





$$cantor_{n+1}(a_1, \ldots; a_{n-1}, a_n, a_{n+1}) =$$

$$cantor_n(a_1, \ldots; a_{n-1}, cantor_2(a_n, a_{n+1}))$$

凡物皆数: Cantor Numbering

❖ 长度有限的字符串,都可视作d=1+ |∑|进制的自然数

"decade" =
$$453145_{(10)}$$

$$//d = 1 + ('i'-'a') = 10$$

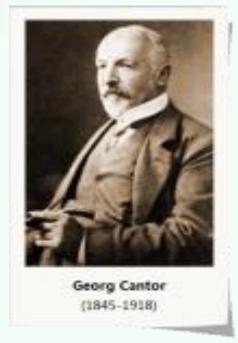


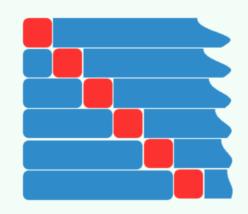
"bgahbhahbhdei..." = 0.2718281828459...

"fade" = 0.6145 = 0.6144999... = "faddiii..."

� Cantor's Diagonal: $\aleph_0 < \aleph_1$







串亦为数

- ◆ 十进制串,可直接视作自然数 //指纹 (fingerprint),等效于多项式法P = "82818"T = 271 82818 284590452353602874713527
- → 一般地,随意对字符编号{ 0, 1, 2, ..., d 1 } //设d = |∑|
 于是,每个字符串都对应于一个 d 进制自然数 //尽管不是单射
 "CAT" = 2 0 19 (26) = 1371(10) //∑ = { A, B, C, ..., Z }
 "ABBA" = 0 1 1 0 (26) = 702(10)
- ❖ P在T中出现 仅当 T中某一子串与P相等 //为什么不是 "当" ?
- **❖ 这,不已经就是一个算法了吗?!** //具体如何实现?
- ❖ 问题似乎解决得很顺利,果真如此简单吗? //复杂度?