

# Lab #3. ROP & Challenges

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# General Information

## ■ Check "Lab #3" in *Assignment* tab of *Cyber Campus*

- Skeleton code (Lab3.tgz) is attached in the post
- Deadline: **11/15 Friday 23:59**
- Submission will be accepted in that post, too
- Late submission due: **11/17 Sunday 23:59 (-20% penalty)**
- Delay penalty is applied uniformly (**not problem by problem**)

## ■ Please read the instructions in this slide carefully

- This slide is a step-by-step tutorial for the lab
- It also contains important submission guidelines
  - If you do not follow the guidelines, **you will get penalty**

# Remind: Cheating Policy

- **Cheating (code copy) is strictly forbidden in this course**
  - Read the orientation slide once more
- **Don't ask for solutions in the online community**
  - TA will regularly monitor the communities
- **Sharing your code with others is as bad as copying**
  - Your cooperation is needed to manage this course successfully
- **You must submit a report as well**
  - More instructions are provided at the end of this slide

# Skeleton Code Structure

- Copy Lab3.tgz into CSPRO server and decompress it
  - You must connect to [csproN.sogang.ac.kr](http://csproN.sogang.ac.kr) (N = 2, 3, or 7)
- Skeleton code has similar structure to the previous lab
  - 3-1/ ... 3-4/ : Problems that you have to solve
  - 3-5/ : *Bonus problem* for practice (**not included in grading**)
    - But this one can be important when preparing the lab exam
  - **check.py**, **config**: Files for self-grading
- This slide will provide a guide on writing ROP exploit
  - It also provides a detailed tutorial for solving 3-1

```
jschoi@cspro2:~$ tar -xzf Lab3.tgz
jschoi@cspro2:~$ ls Lab3
3-1  3-2  3-3  3-4  3-5  check.py  config
```

# Reading ~~secret~~.txt

*secret*

- In the lecture slide, we talked about `execve()` function
  - But there are other variants of like `execv()`, `execl()`, ...
- In this lab, you must run the following code\* with ROP
  - There are other ways to read `secret.txt`, but **don't use them**
    - Ex) Using `system()` instead, or spawning a shell with `execv()`
  - They may not work and you can even **get 0 point** in such cases
    - For instance, `system()` does not work properly with SUID

```
// You can run "cat secret.txt" with execv() as follow.
```

```
char *argv[3];
```

```
argv[0] = "/bin/cat";
```

```
argv[1] = "secret.txt";
```

```
argv[2] = NULL;
```

```
execv(argv[0], argv);
```

↓ 이 코드로 하기 (무조건)

: cat program 실행해서 txt file 출력.



\* Except for 3-3 (this one can be solved without control hijack)

# Example: Problem 3-1

- Target program (twice.c / twice.bin) is given

```
void run_cat(char *filepath) {  
    char *argv[3];  
    argv[0] = "/bin/cat";  
    argv[1] = filepath;  
    argv[2] = NULL;  
    execv(argv[0], argv);  
}
```

Your goal is to execute this function with "secret.txt"

→ 백조 함수

...

```
void vuln(void) {  
    char buf[20];  
    printf("Input your message in stack buffer: ");  
    read(0, buf, 64);  
}
```

You can see that BOF occurs here

# Finding ROP Gadgets

- In principle, you must disassemble all the addresses in the code section, which contains assembly instructions
- Pwntools offers `ROP()` API that does this automatically
  - `print(rop.rdi)`: Print gadgets that can affect `%rdi` register
- Tip: You can use `p64()` function to write concise code

- FYI, `u64()` <sup>(정수 → bytes type으로 return)</sup> function performs conversion in opposite direction <sup>반대방향으로 conversion (bytes type → 정수로 return)</sup> ⇒ 3-2번 문제

```
p = process("./twice.bin")
rop = ROP("./twice.bin") ⇒ Rop 객체 생성
# You can print the gadget information as follow. ⇒ Rop gadget 출력
print(rop.rdi)

# The following two lines have the same meaning.
rdi_gadget = b"\xb3\x12\x40\x00\x00\x00\x00\x00"
rdi_gadget = p64(0x4012b3) # More concise ⇒ 8byte
pop rdi gadget 주소
```

# Attaching GDB to Process

## ■ Assume that you wrote the exploit code below

- It uses ROP gadget to change the value of %rdi into 0x4142

## ■ Let's use gdb to check if this works as expected

- Previously, we launched **gdb** and started a process from there
- This time, let's run the script and attach to the **running process**

```
p = process("./twice.bin")
```

```
# You can use this line to pause the script for a while.
```

```
input("Attach GDB now and press enter to continue: ")
```

```
...
```

→ 4바이트로부터 입력값을 읽어들이는 (target program과의 상호작용을 위해 필요)

```
print(p.recvuntil(b"stack buffer: "))
```

```
rdi_gadget = p64(0x4012b3)
```

```
p.send(b"a" * 0x28 + rdi_gadget + p64(0x4142))
```

(4 bytes) → saved return address 0x4142로 바꾼다

```
input("Done, but let me wait for a while...")
```



# Attaching GDB to Process

- You must open **two terminals** and switch between them
  - When launching `gdb`, specify the **process id (pid)** to attach

## Step 1. Start the exploit script (1<sup>st</sup> terminal)

```
jschoi@cspiro2:~Lab3/3-1$ ./exploit-twice.py  
[+] Starting local process './twice.bin': pid 6936  
Attach GDB now and press enter to continue:
```

(안녕하세요)

## Step 2. Attach and set breakpoints (2<sup>nd</sup> terminal)

```
jschoi@cspiro2:~Lab3/3-1$ gdb -q ./twice.bin 6936
```

```
Reading symbols from ./twice.bin...
```

```
...
```

```
(gdb) b * 0x4011f0
```

```
Breakpoint 1 at 0x4011f0
```

```
(gdb) c
```

```
Continuing.
```

break point

target  
program

⇒ 6936에 attach 하겠다.

argument

# Attaching GDB to Process

## ■ You must open **two terminals** and switch between them

- In the 2<sup>nd</sup> terminal, you can use the **gdb** commands to debug

### Step 3. Resume the exploit script (1<sup>st</sup> terminal)

```
jschoi@cspro2:~Lab3/3-1$ ./exploit-twice.py
[+] Starting local process './twice.bin': pid 6936
Attach GDB now and press enter to continue: Let's go!
b'Input your message in global buffer: '
b'Input your message in stack buffer: '
Done, but let me wait for a while...
```

You must type in something like this

(Enter키만 눌러도 되네.)

### Step 4. Now the breakpoint is hit (2<sup>nd</sup> terminal)

...

Breakpoint 1, 0x00000000004011f0 in vuln ()

(gdb) x/2xg \$rsp

0x7ffe17aa3248: 0x00000000004012b3      0x0000000000004142

⇒ stack의 상황 조사 또는 변화 확인.

# Demonstration

# Obtaining Function Offset

- For problem 3-2, you will have to obtain the offset of a function within the `libc` library
  - Recall that you need this information to figure out the address of `execv()` function, using memory disclosure
- You can do this easily by using the `pwntools` API
  - Then you don't have to hard-code constants in your script

```
# You can investigate the offset of libc functions as follow.  
libc = ELF("/lib/x86_64-linux-gnu/libc.so.6")  
read_offset = libc.symbols['read']  
execv_offset = libc.symbols['execv']  
print("Offset of read() within library: 0x%x" % read_offset)  
print("Offset of execv() within library: 0x%x" % execv_offset)
```

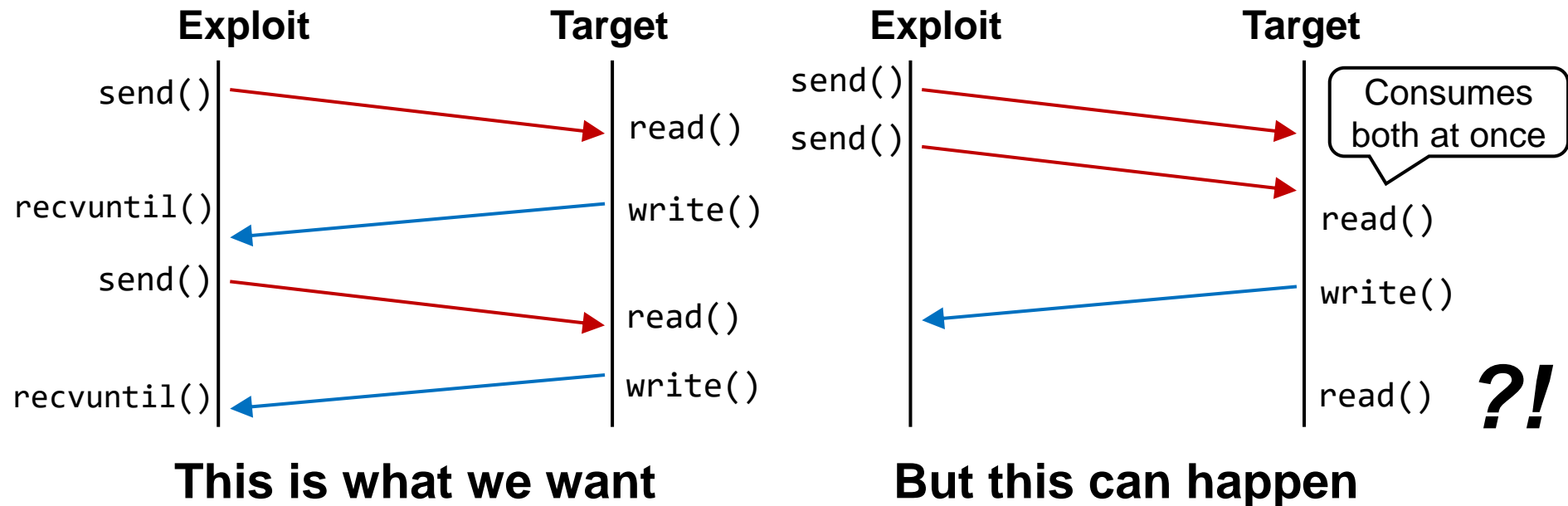
⇒ 好簡単

# Hints

- In 3-1, be careful in handling newline (`\n`) and null (`\0`)
  - Recall that `pwntools`' `sendline()` implicitly appends `'\n'`  
*⇒ gets 함수가 null로 replace한다.*
- In 3-2, you must leak the addresses of `libc` functions
  - Try to disclose the library addresses stored in **GOT** *⇒ ROP5랑 39p 참조*
  - A function's **GOT** entry is filled in when it's called for the first time  
*⇒ puts나 read 함수가 호출될때 GOT table이 업데이트가 된다. ⇒ 한 번이라도 dump한 적이 없는 곳에서 GOT를 찾아낸다.*
- In 3-3, you will have to exploit a **format string bug** to disclose the memory content of an arbitrary address
- In 3-4, you must exploit a **use-after-free** vulnerability
  - First, examine the behavior of memory allocator, by writing a simple program with `malloc()` and `free()` sequence  
*⇒ malloc과 free 순서까지 따져 작동하는 법 알아보기.*
  - In other words, think about how to make the allocator return the freed block that you want

# Caution: Reliability of Exploit

- In this lab, your script has to be especially careful in interacting with the target program carefully
  - Make sure that you send and receive message **step by step**
  - If not, your exploit code may not work reliably (if it doesn't work during the actual grading for this reason, I will deduct point)

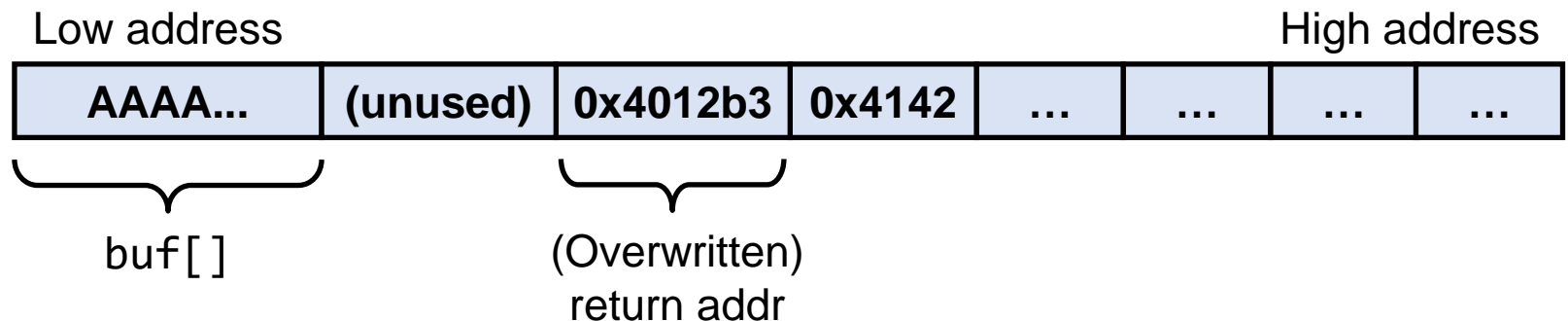


# Report Guideline

- Write report for **3-2 and 3-3** (not required for 3-1 and 3-4)
  - The role of report is to prove that you solved them on your own
  - If you couldn't solve a problem, don't have to write its report
  - Report will not give you point; it is only used to deduct point
- This time, I will provide a template for each problem
  - Make sure that your report contains the requested content
- If you used ChatGPT to write your exploit code, clearly describe it in your report (review the orientation slide)

# Report Template for 3-2

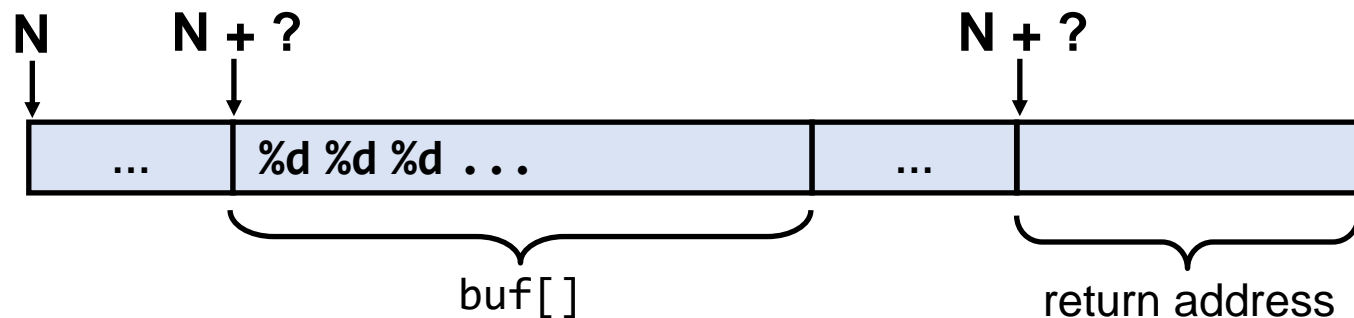
- Draw the state of stack frame after your input overflows the buffer (see the example below)
  - Draw it just as I did in the lecture slide of the ROP chapter
  - Explain the **meaning (role) of each memory** block in the figure
  - Ex) If it's a gadget address, explain **what that gadget does**
  - Ex) If it's an address of a function, explain **what arguments you are trying to pass**, and **why you are doing that**
    - If you are passing a pointer (memory address) as a function argument, explain what is stored in that address





# Report Template for 3-3

- Draw the state of `main()`'s stack frame immediately before `printf()` is called
  - Clearly indicate the positions of `buf[]` and `saved return address` in the stack frame
  - `N` must be the value of `%rsp` at address `0x4011e9`
  - You will be entering format specifiers as input ("`%d%d%d...`"); so explain `which stack position is consumed` by each specifier
  - Justify `why you repeat` each format specifier `for certain number of times`



# Problem Information

- There are four problems you have to solve (25 pt. each)
  - Problem 3-1: `twice.bin` (★) → Pop Gadget 0/0
  - Problem 3-2: `substr.bin` (★★★★)
  - Problem 3-3: `fsb.bin` (★★☆)
  - Problem 3-4: `item.bin` (★★☆)
- You'll get the point for each problem if the exploit works
  - No partial point for non-working exploit
- If the report does not clearly explain how you analyzed and solved the problem, you will **lose points**
  - You can write the report in Korean or English
  - Due to the limited time, I will randomly select a problem to grade

# Submission Guideline

## ■ You should submit four exploit scripts and report

- Problem 3-1: `exploit-twice.py`
- Problem 3-2: `exploit-substr.py`
- Problem 3-3: `exploit-fsb.py`
- Problem 3-4: `exploit-item.py`
- **Don't forget the report: `report.pdf`**
- 3-5 is a bonus problem, so you don't have to submit it

## ■ Submission format

- Upload these files directly to *Cyber Campus* (**do not zip them**)
- **Do not change the file name** (e.g., adding any prefix or suffix)
- If your submission format is wrong, you will get **-20% penalty**