Mathematisch-Naturwissenschaftliche Fakultät Wilhelm-Schickard-Institut für Informatik Datenbanksysteme · Prof. Dr. Grust





# **Functional Programming**

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# Assignment #7

Submission Deadline: Thu, 12.12.2019

Wichtig: Im Zeitraum vom 6.12.2019 bis zum 19.12.2019 (16 Uhr) findet die Lehrevaluation zu dieser Veranstaltung statt. Sie wurden dazu per E-Mail an Ihre studentische Adresse eingeladen. Bitte nehmen Sie daran teil, das Feedback ist uns sehr wichtig!

# Exercise 1: Deep Embedding of a While-Language

(10 Points)

In the lecture we have implemented a deeply embedded, typed expression language over Integers and Booleans. To ensure that only well-typed expressions can be built, the module makes use of Haskell's *Generalized Algebraic Data Types* (GADTs). GADT syntax allows us to check the typability of our embedded language at compile time. In this exercise we are going to extend this DSL to a simple *While-Language*.

Inspect your repository to find the module <code>DeepWhile</code> which contains a slightly modified version of the lecture's code to start with.

The given module exports the AST data type Expr a, together with its constructors and an observer function eval :: Expr a -> a which evaluates DSL expressions.

1. First, extend the expression language with a relational operator Lt which takes two expressions  $e_1$ ,  $e_2$  of type Expr Int and — when evaluated — returns whether  $e_1 < e_2$  or not. Complete the Show instance and eval observer accordingly.

### Example:

```
eval (Lt (VarI 41) (VarI 42)) \equiv True
```

- 2. Extend the expression language with support for a *variable environment* (see type Env b in DeepWhile<sup>1</sup>) which binds values of a particular type b to variable names (type String) that can be used in the DSL expression:
  - (a) Equip the data type Expr with a second type variable b that determines the type of environment values (Expr a b: evaluates to a value of type a, based on an environment of variables of type b).
  - (b) Adapt eval to additionally take an environment of type Env b:

<sup>&</sup>lt;sup>1</sup>The environment is implemented as a dictionary using the module Data.Map. See https://hackage.haskell.org/package/containers-0.6.2.1/docs/Data-Map-Lazy.html for further information.

```
eval :: Env b \rightarrow Expr a b \rightarrow a
```

(c) Provide an additional constructor Var which takes a variable name (String). Var represents a lookup of the variable name in the environment on evaluation. If the variable is not bound in the environment an error "Variable not bound!" is thrown at runtime. Complete Show and eval accordingly.

#### **Example:**

```
eval (M.singleton "a" 42) (Var "a") \equiv 42
```

- 3. Define a new data type Stmt b for While-Language statements operating on an environment of type b. Use GADT syntax to define the following constructors:
  - (a) Assign takes a variable name (String) and an expression (Expr) that evaluates to a value of type b. It represents the assignment of the expression result to a variable with the given name.
  - (b) While takes a Boolean expression and a statement (Stmt b) which is to be repeated while the condition holds. Both, the expression and the loop statement are evaluated based on an environment of type b.
  - (c) Seq takes two statements (Stmt b) and represents their sequential execution.

Define a Show instance for (Stmt b).

4. Write a function

```
runS :: Env b -> Stmt b -> Env b
```

which runs a given statement based on a given environment. The – possibly modified – environment is returned. Whenever statements are run sequentially (or looped) the modifications in the environment of the previous statement are visible in the execution of the subsequent.

## Example:

5. Finally, define and export a function

```
run :: String -> Stmt b -> Maybe b
```

which takes a variable name v and a statement s. This function should run s with an empty environment and returns the value bound to v after evaluation (or Nothing if v was not bound by the program).

### Example:

```
\texttt{run} \ \texttt{p} \ \equiv \ 42
```

Solve the following puzzle:

For each number n between 0 and 20, find at least one mathematical expression which evaluates to n and is an arbitrary combination of exactly four numbers 4 using the following arithmetic operations: addition (a+b), subtraction (a-b), multiplication (a\*b), division (a/b), exponentiation  $(a^b)$ , square root  $(\sqrt{a})$  and factorial of numbers (4!).

#### Example:

$$0 = \frac{4}{4} * 4 - 4$$
  $1 = \left(\frac{4}{4}\right)^{4^4}$  ...

- 1. Define a data type ExprTree to represent such mathematical expressions.
- 2. Write a function eval :: ExprTree -> Maybe Float which evaluates an expression to a number if possible; if the expression is not evaluable due to division by zero return Nothing, instead.
- 3. Write a function trees :: [ExprTree] -> [ExprTree] which takes a list of leaf nodes and returns a list of all expression trees that can be built in combination of these leaf nodes. Assume that a sequential application of the unary square-root-operation  $(e.g. \sqrt{\sqrt{a}})$  is not allowed, while factorial is only applied to leaf nodes (4!), but not to other expressions (e.g. (a+b)!).

**Example:** (Note: This example uses a human-digestable notation for **ExprTree** values.)

```
trees [4,4,4,4] => [(4+4+4+4), ..., (4*4+4+4), ..., (4*(4+4)+4), ..., (4*(4+4)+4), ..., ((4/4)^4)^4, ..., ((4!+4)/4+4), ..., [
```

**Hint:** It might be useful to write a helper function splits :: [a] -> [([a], [a])] which returns all combinations of a list split in two parts.

```
Example: splits [1..4] \Rightarrow [([1],[2,3,4]),([1,2],[3,4]),([1,2,3],[4])]
```

4. Finally, write a function solution :: Int -> Maybe ExprTree that returns one arbitrary "expression of four 4s" which evaluates to a given number  $n \in \{1, ..., 20\}$ . Return Nothing if no such expression was found.