# Tabular Database Systems

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Database-External Data in Parquet Files 👪

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# 1 Columnar Compressed Data Storage Outside the DBMS: Parquet

Fragility, space requirements, and parsing effort render CSV files problematic for applications that read/write GBs or TBs of data.

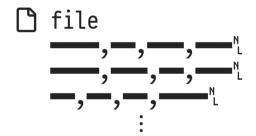
Alternative database-external file formats have been developed to address these issues. Among these, Parquet is widely used.

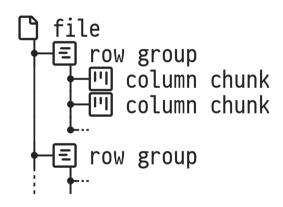
- Developed in the open since 2013 (initiated by Twitter et al).
- Stores columns of typed data.
- Built-in compression (on file and column levels).
- Incorporates rich metadata that supports projection and selection pushdown.
- Supported by libraries for a wide range of programming languages. Directly readable/writable by DuckDB .

<sup>&</sup>lt;sup>1</sup> Apache Parquet ▶ (sponsored by the Apache Software Foundation): open source, column-oriented data file format designed for efficient data storage and retrieval.

# CSV vs. Parquet

CSV	Parquet ##
monolithic	split into horizontal row groups column-based (chuṇk-by-chunk)
row-oriented (lines separated by \bar{\tau}) plain text, uncompressed	sunnorts compression:
prain cext, and any coolea	<pre>supports compression:     file-level (gzip, zstd,)     column-level (dictionary, RLE,)</pre>
untyped (requires parsing)	• column-level (dictionary, RLE,)
uncyped (requires parsing)	typed: • scalar (INT32/64, FLOAT, BYTE ARRAY,)
	• scalar (INT32/64, FLOAT, BYTE_ARRAY,)  [• nested records]  metadata for file/row groups/columns:  • file format version
no metadata (optional header row)	metadata for file/row groups/columns:
	• min/max/count statistics
	<ul> <li>min/max/count statistics</li> <li>cardinality, (un)compressed byte sizes</li> </ul>





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## A Sketch of Parquet's Storage Format

## Parquet File Magic Number "PAR1" Row Group **≡** Metadata Column 1: [1], Chunk 1 Column 2: [1], Chunk 1 Row Group 1 Column $c: \square$ , Chunk 1 more row groups Row Group **■** Metadata Column 1: 🗓 , Chunk N Column 2: 🗓 , Chunk N Row Group N Column $c: \square$ , Chunk NFile [ Metadata Length of Metadata (4 bytes) Magic Number "PAR1"

#### **□ Column Chunk Metadata**

```
Type
Compression Scheme
Codec (gzip, zstd, ...)
# of Values
Compressed Size
Uncompressed Size
Offset to First Value
Statistics (min/max/...)
Bloom Filter
:
```

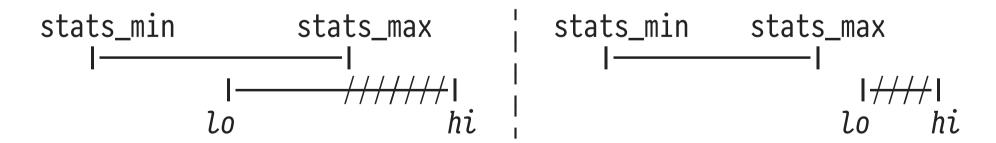
#### **■** Row Group Metadata

Column Info [][][][][]
Total Compressed Size
Total Uncompressed Size
# of Rows
(Sorting Columns [][][])
:

# 2 | Pushing Projection and Filtering Down into Parquet Reading

Parquet's file structure + metadata enables readers (like DuckDB's PARQUET\_SCAN operator) to **only access relevant data subsets.** 

- **Projection pushdown:** Exploit **column-based layout.** In each row group, use data\_page\_offset to navigate the file to only read required column chunk(s).
- Filter pushdown: Exploit statistics metadata. Entirely skip row group if stats\_min/stats\_max² for column c indicate that filter predicates like  $lo \le c \le hi$  will always fail  $(\frac{}{///})$ .



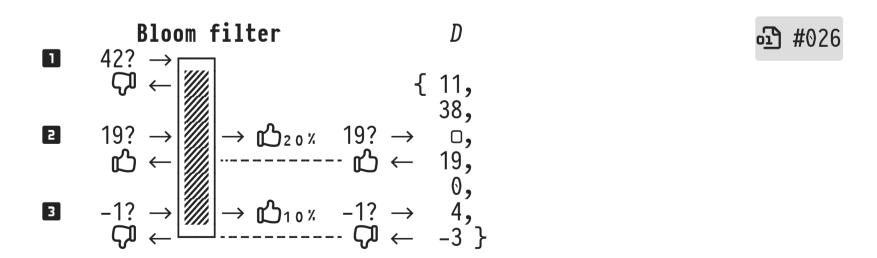
<sup>&</sup>lt;sup>2</sup> In the DB research literature, the (stats\_min, stats\_max) pairs in all row groups are sometimes collectively referred to as **zone map** (for column c).



## 3 Optional in Parquet Files: Bloom Filters

If column values are randomly shuffled, all values may occur in all row groups and the effectiveness of zone maps is largely lost.

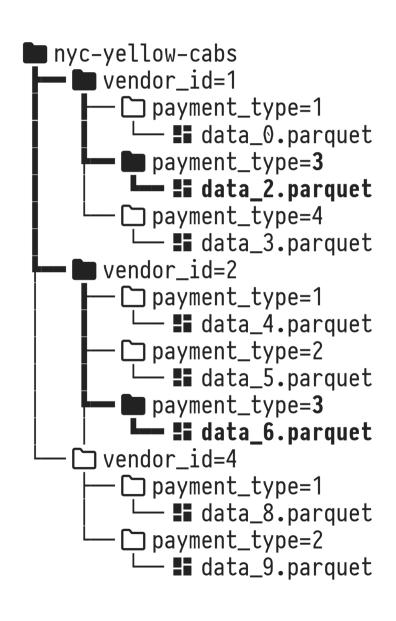
- **Bloom filters** are *compact* data structures that *over-approximate* the set *D* of distinct values<sup>3</sup> in a row group.
- Given the question  $x \in D$ ?, Bloom filters either respond with
  - ∘ definitely no (♥) or
  - o probably yes  $(\mathcal{D}_p)$ , where p is a false positive rate).



<sup>&</sup>lt;sup>3</sup> Set *D* is the *active* domain of the row group. Example: if the values in a column have type int32 with domain  $[-2^{31},2^{31})$ , we typically have  $D \in [-2^{31},2^{31})$  (read  $\in$  as "significantly smaller").

# 4 | Partitioned Files and Filter Pushdown (Hive Partitioning)

If data volumes get truly large, it may make sense to **partition rows** into a family of files. Which partition will a row belong to?



- Choose n ≥ 1 columns c<sub>1</sub>,...,c<sub>n</sub> as partition criteria
   (left: n=2, vendor\_id + payment\_type).
- Arrange partitions in tree hierarchy:
  - $\circ$  Depth of tree: n+1.
  - o Width of tree ≤ product of the size of active domains of columns  $c_i$ .
- Inner nodes: ☐ OS directories, leaf nodes: ■ data (Parquet or CSV).
- Filter query: traverse relevant subtrees/read relevant data files only (left: payment\_type = 3).
- Particularly useful when I/O is slow.