

## 1. Natural Language Description of Algorithm

The algorithm starts by accepting an input that defines the DTM's configuration, including symbols, states, and transition functions. It then begins to simulate that DTM, processing input strings and applying transition rules to modify the tape and machine state as defined. The simulator will attempt to solve specific computational problems, such as addition, subtraction, and multiplication. The algorithm will output detailed transition states.

## 2. Pseudocode

```

SIMULATEDTM(InputString)
1  Tape = InputString + BlankSymbol
2  HeadPosition = 0
3  CurrentState = 'q0'
4  TransitionCount = 0
5  Delta = DTM Operating mode instructions based on type
6  while CurrentState not in {qY, qN}
7      ReadSymbol = Tape[HeadPosition]
8      if TransitionFunction exists for (CurrentState, ReadSymbol)
9          (NextState, WriteSymbol, MoveDirection) = TransitionFunction(CurrentState, ReadSymbol)
10         Tape[HeadPosition] = WriteSymbol
11         if MoveDirection = 'R'
12             HeadPosition = HeadPosition + 1
13         elseif MoveDirection = 'L'
14             HeadPosition = HeadPosition - 1
15         CurrentState = NextState
16         TransitionCount = TransitionCount + 1
17         if TransitionCount > 30
18             Write Tape to File
19     else
20         Break
21 return Tape, CurrentState
    
```

## 3. Analytically derived Asymptotic family for algorithm

SIMULATEDTM( <i>InputString</i> )		
1	Tape = InputString + BlankSymbol	$c_1$ 1
2	HeadPosition = 0	$c_2$ 1
3	CurrentState = $q_0$	$c_3$ 1
4	TransitionCount = 0	$c_4$ 1
5	Delta = DTM Operating mode instructions based on type	$c_5$ 1
6	<b>while</b> CurrentState not in { $q_A$ , $q_N$ }	$c_5$ $T(n)$
7	ReadSymbol = Tape[HeadPosition]	$c_6$ 1
8	<b>if</b> TransitionFunction exists for (CurrentState, ReadSymbol)	$c_7$ 1
9	(NextState, WriteSymbol, MoveDirection) = delta(CurrentState, ReadSymbol)	$c_8$ 1
10	Tape[HeadPosition] = WriteSymbol	$c_9$ 1
11	<b>if</b> MoveDirection = 'R'	$c_{10}$ 1
12	HeadPosition = HeadPosition + 1	$c_{11}$ 1
13	<b>Else If</b> MoveDirection = 'L'	$c_{12}$ 1
14	HeadPosition = HeadPosition - 1	$c_{13}$ 1
15	CurrentState = NextState	$c_{14}$ 1
16	TransitionCount = TransitionCount + 1	$c_{15}$ 1
17	<b>Else</b>	$c_{16}$ 1
18	Break	$c_{17}$ 1
19	<b>return</b> Tape, CurrentState	$c_{18}$ 1

As we can see in the above analysis there is only 1 real variable to the DTM which I consider  $T(n)$ . As the length of input does not actually affect the speed at which the while loop increases I had to use a different name than  $n$ .  $T(n)$  can be described as the function of the Turing Machine and its overall time complexity would be  $O(T(n))$ . As  $n$  only really affects the function  $T$ , not directly the time complexity.

#### 4. Analysis

As I determined in my Asymptotic analysis of the pseudocode  $O(T(n))$  seemed to be the best representation of how it ran. This can be seen with a few black box testings I performed:

Base DTM:

Input: 1110, Times While Loop ran: 7

Input: 11110, Times While Loop ran: 7

Input: 11101, Times While Loop ran: 10

DTM Addition:

Input: 100+11, Times While Loop ran: 25

Input: 111+1011, Times While Loop ran: 67

Input: 101100101+101101101, Times While Loop ran: 232

DTM Subtraction:

Input: 100-10, Times While Loop ran: 25

Input: 101-10, Times While Loop ran: 25

Input: 101111-101, Times While Loop ran: 43

DTM Multiplication:

Input: 1\*1, Times While Loop ran: 295

Input: 10\*10, Times While Loop ran: 557

Input: 1110\*10, Times While Loop ran: 825

As we can see, while input size can make a difference to the DTM Runtime, what truly makes a difference is what instructions it is using and what is input, neither of which can easily be graphed in a simple  $n$  format.