

Zusatz 12. Übung

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1 a)

```
h    LOAD 1; LOAD 2; LEQ; JMC h.1;  
      LOAD 1; STORE 1; WRITE 1; JMP 0;  
h.1  LOAD 2; LOAD 1; STORE 2; STORE 1;  
      JMP h;
```

1 b)

```
int main(){
    int x1, x2, x3;
    scanf("%d", &x1);    scanf("%d", &x2);
    while(x2>0){
        if(x1>=x2)x1=x1-1;
        else{
            x3=x1;
            x1=x2;
            x2=x3;}
        printf("%d", x1);
        return 0;
    }
```

$$\begin{aligned}
& \left\{ \begin{pmatrix} \delta(\gamma(x_3), \gamma(\gamma(\alpha)), \sigma(\gamma(x_2), x_1)) \\ \delta(\gamma(x_3), \gamma(x_3), \sigma(\gamma(\alpha), \gamma(\gamma(x_2)))) \end{pmatrix} \right\} \\
& \xRightarrow{Dek.} \left\{ \begin{pmatrix} \gamma(x_3) \\ \gamma(x_3) \end{pmatrix}, \begin{pmatrix} \gamma(\gamma(\alpha)) \\ \gamma(x_3) \end{pmatrix}, \begin{pmatrix} \sigma(\gamma(x_2), x_1) \\ \sigma(\gamma(\alpha), \gamma(\gamma(x_2))) \end{pmatrix} \right\} \\
& \xRightarrow{Dek.} \left\{ \begin{pmatrix} x_3 \\ x_3 \end{pmatrix}, \begin{pmatrix} \gamma(\alpha) \\ x_3 \end{pmatrix}, \begin{pmatrix} \gamma(x_2) \\ \gamma(\alpha) \end{pmatrix}, \begin{pmatrix} x_1 \\ \gamma(\gamma(x_2)) \end{pmatrix} \right\} \\
& \xRightarrow{Dek.} \left\{ \begin{pmatrix} x_3 \\ x_3 \end{pmatrix}, \begin{pmatrix} \gamma(\alpha) \\ x_3 \end{pmatrix}, \begin{pmatrix} x_2 \\ \alpha \end{pmatrix}, \begin{pmatrix} x_1 \\ \gamma(\gamma(x_2)) \end{pmatrix} \right\}
\end{aligned}$$

$$\xRightarrow{Elim.} \left\{ \begin{pmatrix} \gamma(\alpha) \\ x_3 \end{pmatrix}, \begin{pmatrix} x_2 \\ \alpha \end{pmatrix}, \begin{pmatrix} x_1 \\ \gamma(\gamma(x_2)) \end{pmatrix} \right\}$$

$$\xRightarrow{Vert.} \left\{ \begin{pmatrix} x_3 \\ \gamma(\alpha) \end{pmatrix}, \begin{pmatrix} x_2 \\ \alpha \end{pmatrix}, \begin{pmatrix} x_1 \\ \gamma(\gamma(x_2)) \end{pmatrix} \right\}$$

$$\xRightarrow{Sub.} \left\{ \begin{pmatrix} x_3 \\ \gamma(\alpha) \end{pmatrix}, \begin{pmatrix} x_2 \\ \alpha \end{pmatrix}, \begin{pmatrix} x_1 \\ \gamma(\gamma(\alpha)) \end{pmatrix} \right\}$$

3 a)

Sei a ein Datentyp beliebig aber fest.

Induktionsanfang:

Sei $xs :: [a]$ mit $xs = []$. Dann gilt

$$\text{breadth}(\text{toTree} xs) = \text{breadth}(\text{toTree } [])$$

$$\stackrel{Z12}{=} \text{breadth}(\text{Leaf})$$

$$\stackrel{Z4}{=} 1$$

$$\stackrel{Z16}{=} \text{pow}(0)$$

$$\stackrel{Z8}{=} \text{pow}(\text{length } []) = \text{pow}(\text{length } xs)$$

3 b)

Induktionsvoraussetzung:

Sei $xs :: [a]$ so, dass gilt:

IV: $\text{breadth}(\text{toTree } xs) = \text{pow}(\text{length } xs)$

Es sei $x :: a$ beliebig aber fest. Behauptung: Dann gilt

$\text{breadth}(\text{toTree}(x : xs)) = \text{pow}(\text{length}(x : xs))$

3 b)

$$\begin{aligned} \text{breadth}(\text{toTree}(x : xs)) &\stackrel{Z13}{=} \text{breadth}(\text{Branch } x (\text{toTreexs})(\text{toTreexs})) \\ &\stackrel{Z5}{=} \text{breadth}(\text{toTree } xs) + \text{breadth}(\text{toTree } xs) \\ &\stackrel{IV}{=} 2 * \text{pow}(\text{length } xs) \\ &\stackrel{Z17}{=} \text{pow}(1 + \text{length } xs) \\ &\stackrel{Z9}{=} \text{pow}(\text{length}(x : xs)) \end{aligned}$$