

# Design & Analysis of Algorithms

## Session – 34

# NP-HARD GRAPH PROBLEMS

The strategy to show that a problem  $L_2$  is NP-hard is

- (i) Pick a problem  $L_1$  already known to be NP-hard.
- (ii) Show how to obtain an instance  $I^1$  of  $L_2$  from any instance  $I$  of  $L_1$  such that from the solution of  $I^1$ 
  - We can determine (in polynomial deterministic time) the solution to instance  $I$  of  $L_1$ .
- (iii) Conclude from (ii) that  $L_1 \alpha L_2$ .
- (iv) Conclude from (i),(ii), and the transitivity of  $\alpha$  that satisfiability  $\alpha L_1$   
 $L_1 \alpha L_2$

$\therefore$  Satisfiability  $\alpha L_2$

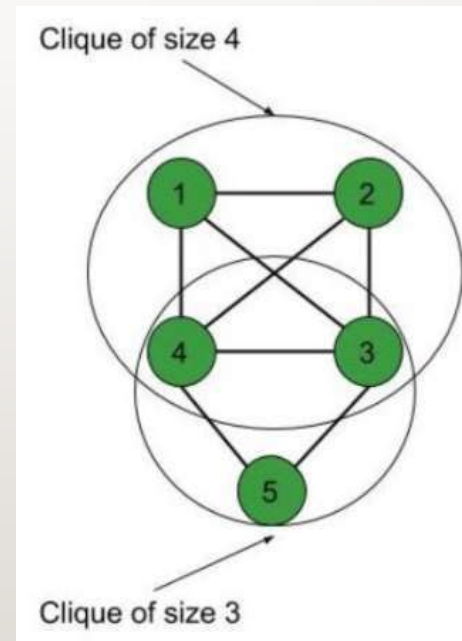
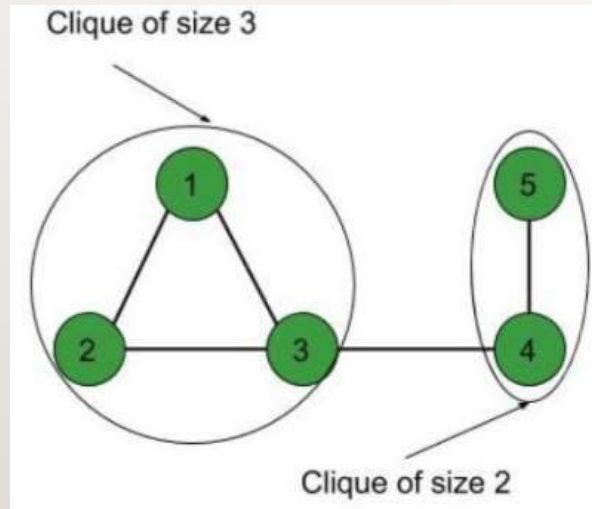
$\therefore L_2$  is NP-hard

# Clique Decision Problem(CDP) :

## Clique:

Clique is a maximal complete sub graph of a graph  $G = (V,E)$

– Size of a clique is the number of vertices in it



# Theorem: CNF- satisfiability $\alpha$ Clique Decision Problem

Proof:

Let  $F = \bigwedge_{1 \leq i \leq k} C_i$  be a propositional formula in CNF. Let  $x_i, 1 \leq i \leq n$  be the variables in  $F$ .

We shall show how to construct from  $F$  a graph  $G = (V, E)$  such that  $G$  will have a clique of size at least  $k$  if  $F$  is satisfiable.

If the length of  $F$  is  $m$ , then  $G$  will be obtainable from  $F$  in  $O(m)$  time.

Hence, if we have a polynomial time algorithm for CDP, then we can obtain a polynomial time algorithm for CNF-satisfiability using this construction.

For any  $F$ ,  $G = (V, E)$  is defined as follows:  $V = \{ \langle \sigma, i \rangle \mid \sigma \text{ is a literal in clause } C_i \}$ ;  $E = \{ (\langle \sigma, i \rangle, \langle \delta, j \rangle) \mid i \neq j \text{ and } \sigma \neq \delta \}$ .

A sample construction is given in Example.

## Example:

Consider  $F = (x_1 \vee x_2 \vee x_3) \wedge (\bar{x}_1 \vee \bar{x}_2 \vee \bar{x}_3)$ .

The construction of Theorem yields the graph:

This graph contains six cliques of size two.

Consider the clique with vertices  $\{ \langle x_1, 1 \rangle, \langle \bar{x}_2, 2 \rangle \}$ .

By setting  $x_1 = \text{true}$  and  $\bar{x}_2 = \text{true}$  (i.e.  $x_2 = \text{false}$ )

$F$  is satisfied.

$x_3$  may be set either to true or false.

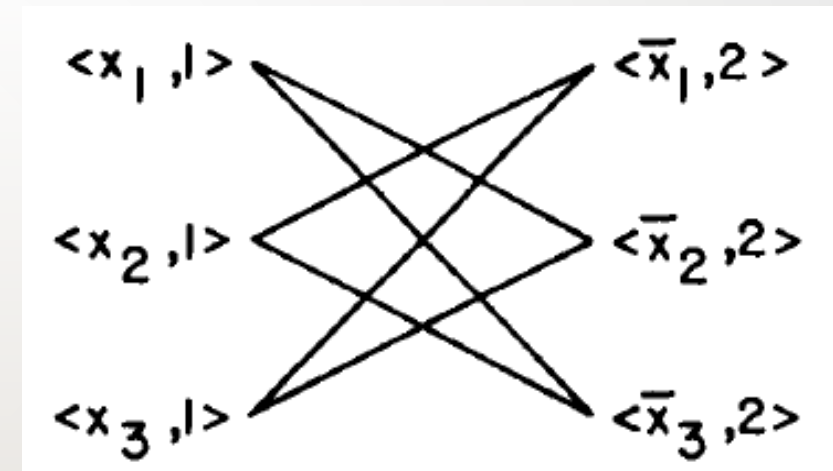


Figure: A sample graph and satisfiability

## Questions:

1. Discuss in detail about Clique Decision Problem
2. Reduce CNF-Satisfiability problem into CDP and Solve

# THANK YOU