



## NETWORK PROTOCOLS & SECURITY 23EC2210 R/A/E

Topic:

# Data Encryption Standard (DES)

Session – 37

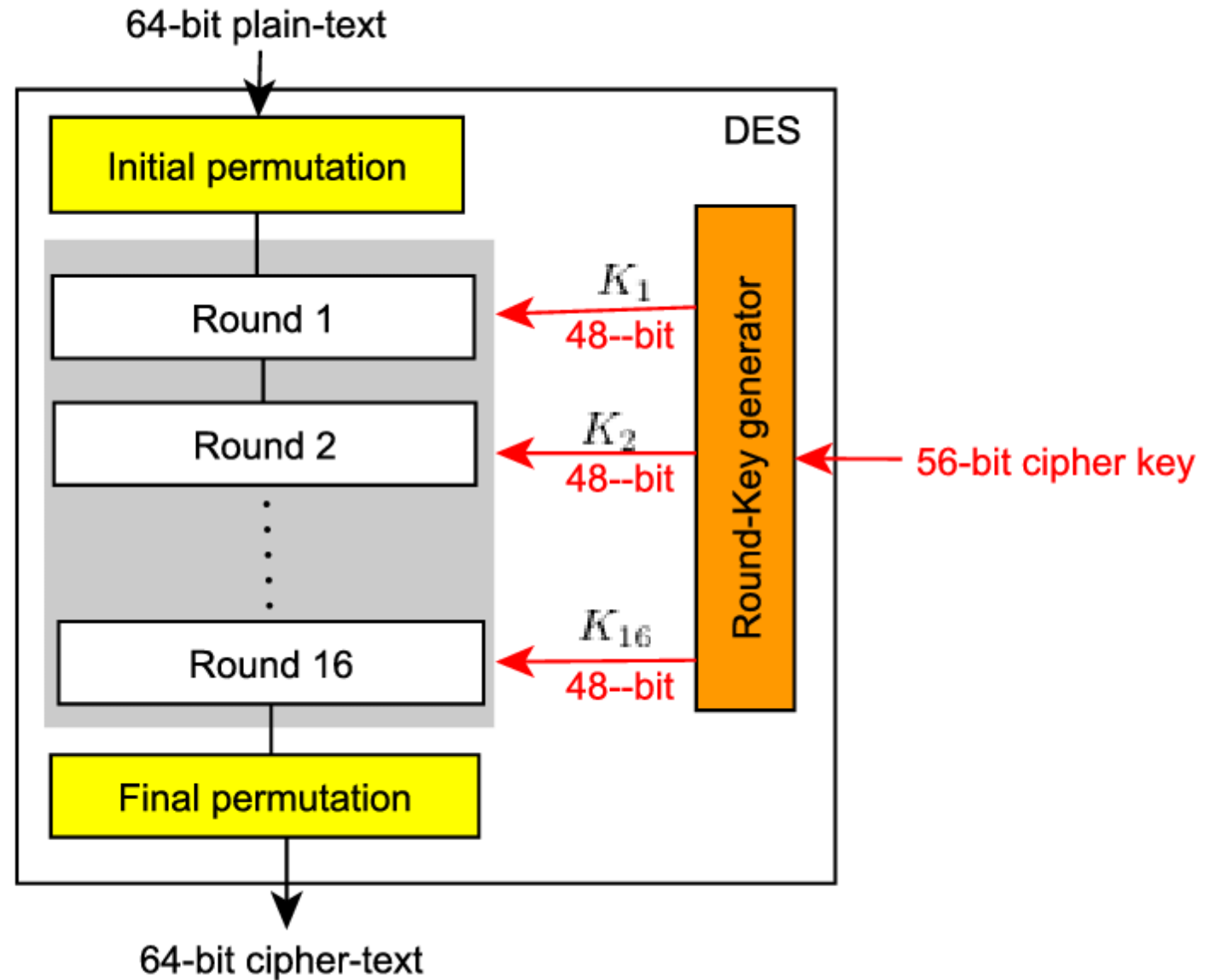
# Data Encryption Standard(DES)

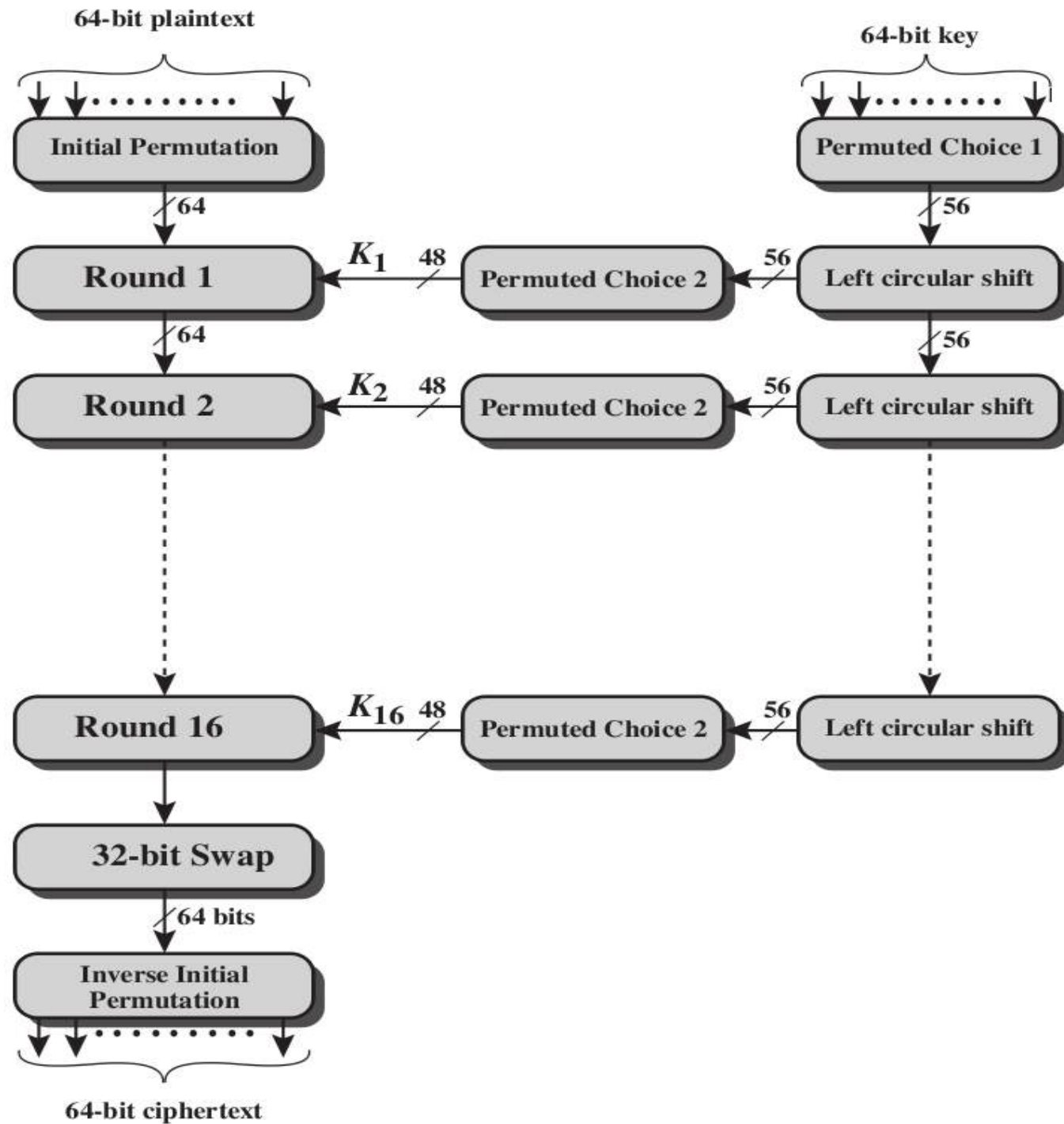
## Block Cipher

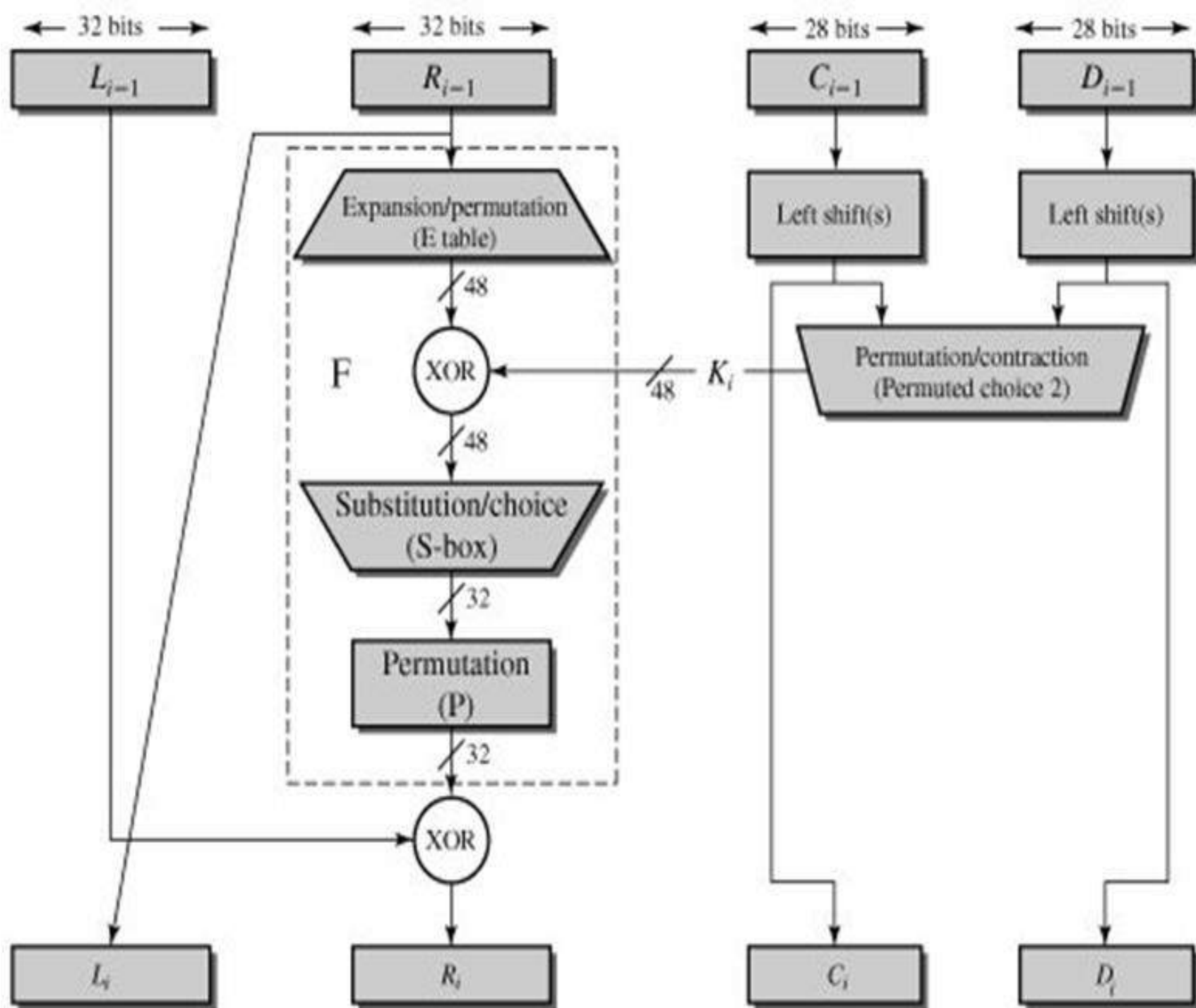
# Data Encryption Standard(DES)

- **Data Encryption Standard (DES)** is the **symmetric block** cipher which encrypts a **64-bit plain text** in a **64-bit ciphertext**.
- Introduced by the National Institute of Standard and Technology (NIST) in the 1970s.
- Initially, DES was only used in financial applications but later it was accepted as the cryptographic algorithm by other organizations too.

- DES is a block cipher.
- Encrypts blocks of 64-bits.
- Uses a key of 56-bit.
- Same key is used in encryption and decryption process of DES.
- No. of Rounds: 16







# S-DES (Simplified DES)

- Used for better understanding of DES.
- Plaintext size: 8-bit
- Key : 10-bit
- No. of rounds: 2
- Ciphertext : 8-bit

# S-DES example

S-DES example



## S-DES (Simplified DES)

Input-Plaintext  $\rightarrow$  8-bit  
Key  $\rightarrow$  10-bit

Output-Ciphertext  $\rightarrow$  8-bit  
No. of Rounds  $\rightarrow$  2

### Encryption

8-bit Plaintext

Initial Permutation

$f_k$

$K_1$  (8-bit)

Swap

$f_k$

$K_2$  (8-bit)

Inverse Initial Permutation

8-bit Ciphertext

### Key-Generation

10-bit key

Permutation( $P_{10}$ )

Left Circular Shift

10 bits

Permutation( $P_8$ )

Left Circular Shift

10-bit

Permutation( $P_8$ )

## S-DES Example:

Given PT = 1 0 0 1 0 1 1 1  
Key = 1 0 1 0 0 0 0 1 0  
P<sub>10</sub> = 3 5 2 7 4 10 1 9 8 6  
P<sub>8</sub> = 6 3 7 4 8 5 10 9  
I/P = 2 6 3 1 4 8 5 7  
E/P = 4 1 2 3 2 3 4 1

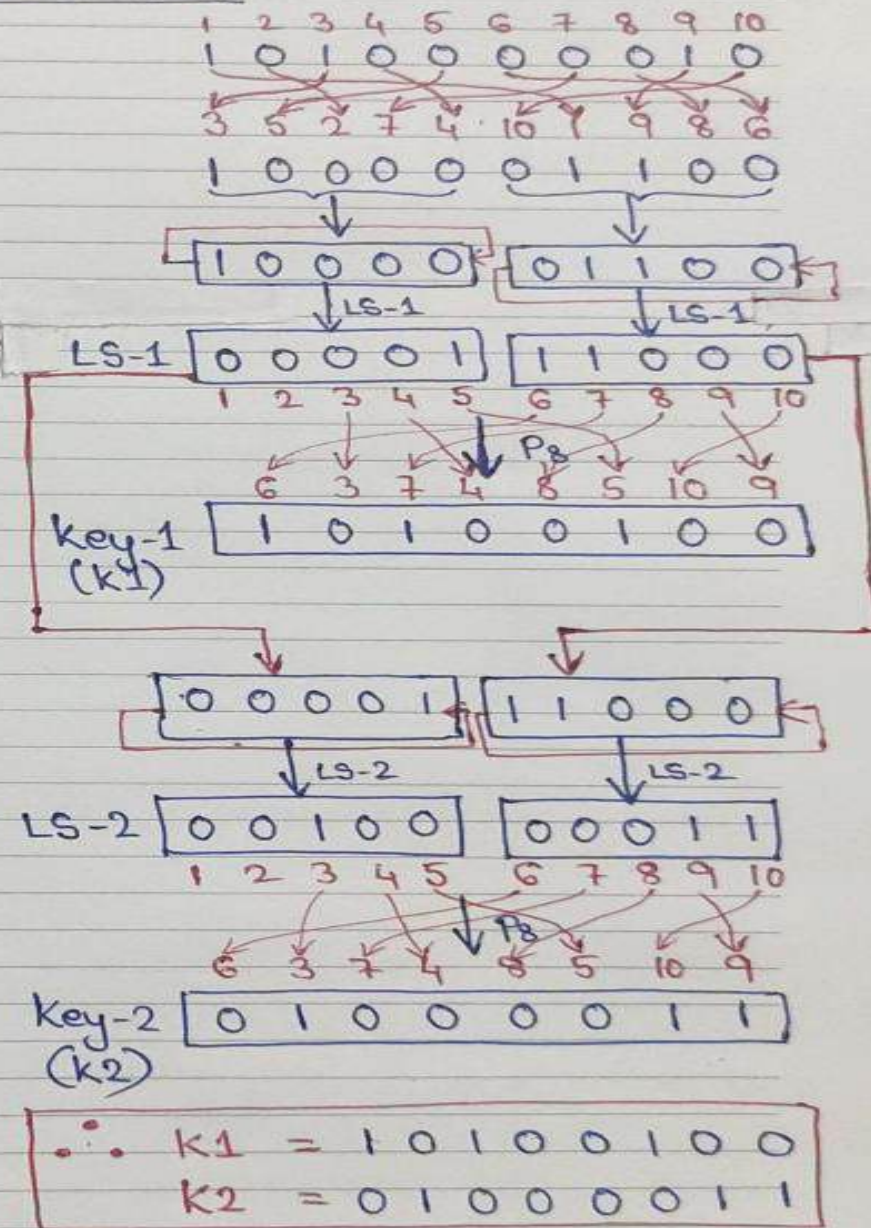
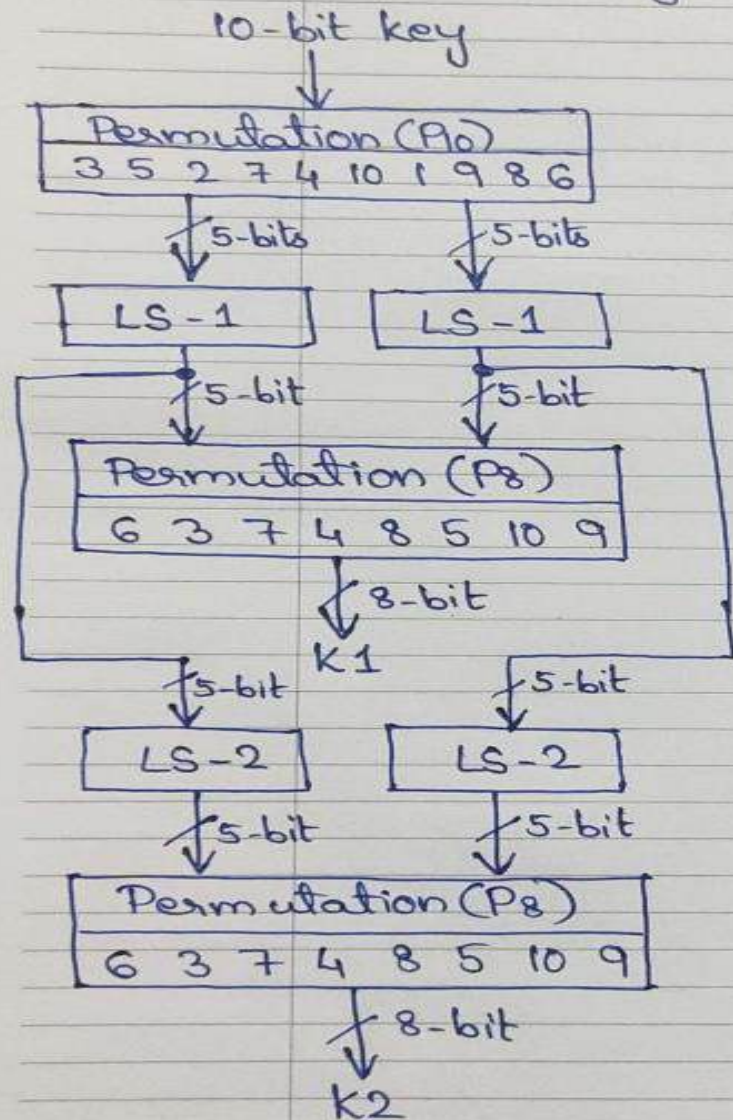
$$S_0 = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 1 & 0 & 3 & 2 \\ 3 & 2 & 1 & 0 \\ 0 & 2 & 1 & 3 \\ 3 & 1 & 3 & 2 \end{bmatrix} \end{matrix} \quad S_1 = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 1 & 2 & 3 \\ 2 & 0 & 1 & 3 \\ 3 & 0 & 1 & 0 \\ 2 & 1 & 0 & 3 \end{bmatrix} \end{matrix}$$

$$P_4 = 2 \ 4 \ 3 \ 1$$

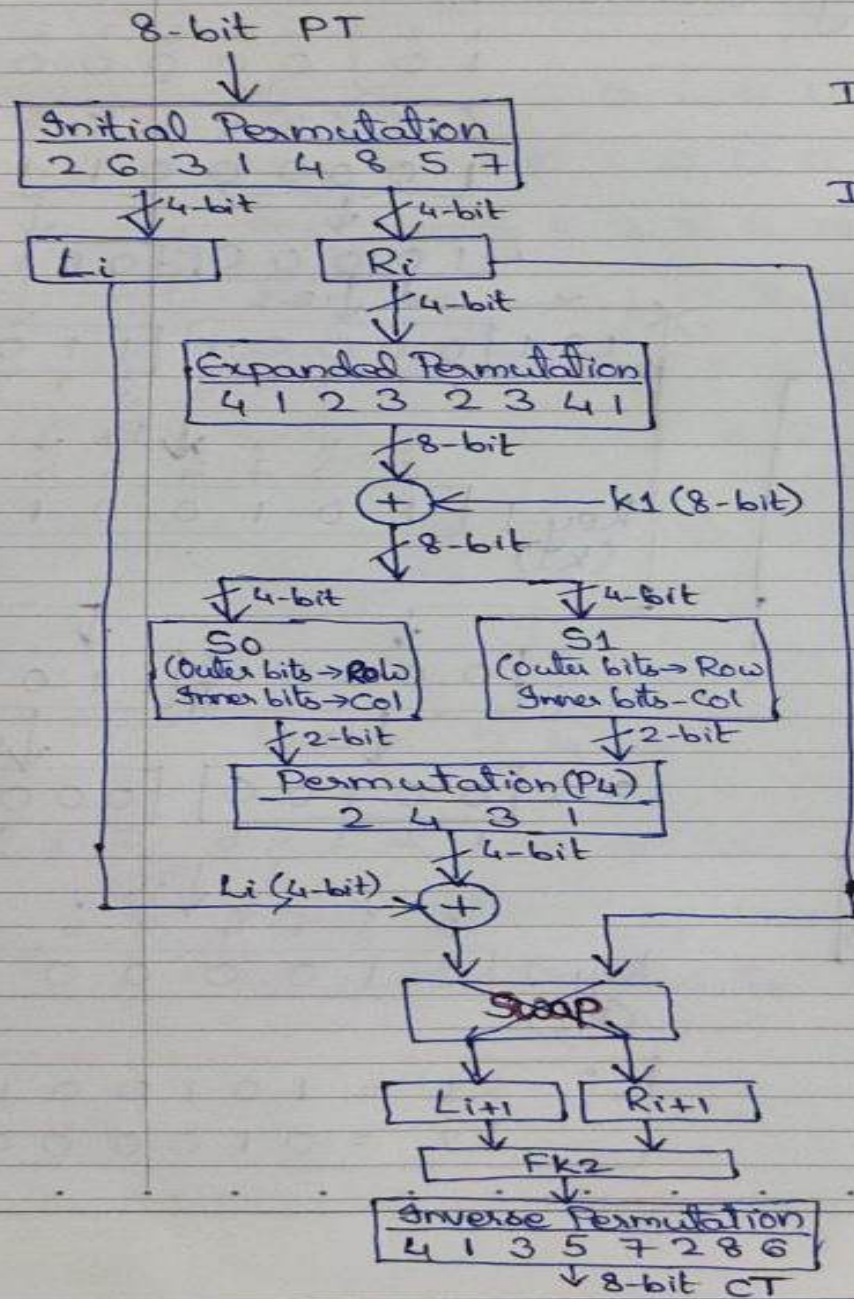
Note : S-boxes S<sub>0</sub> & S<sub>1</sub> are  
4x4 matrix of range (0-3)  
To convert 8-bits to 4-bits



## Key - Generation



# S-DES Encryption



IP : 1 2 3 4 5 6 7 8  
2 6 3 1 4 8 5 7

IP<sup>-1</sup> : 1 2 3 4 5 6 7 8  
4 1 3 5 7 2 8 6



PT : 1 2 3 4 5 6 7 8  
 IP : 2 6 3 1 4 8 5 7

After IP : 0 1 0 1 1 1 0 1

Round-1  
 Li : 0 1 0 1 Ri : 1 1 0 1

Ri : 1 2 3 4

EP : 4 1 2 3 2 3 4 1

EP on Ri : 1 1 1 0 1 0 1 1

K1 : 1 0 1 0 0 1 0 0

EP  $\oplus$  K1 : 0 1 0 0 1 1 1 1

Row  
 S0 : 0 1 0 0      S1 : 1 1 1 1  
 Col  
 00  $\rightarrow$  0th Row  
 10  $\rightarrow$  2nd col      11  $\rightarrow$  3rd Row  
                          10  $\rightarrow$  3rd col

S0 o/p : 3 = 11      S1 o/p : 3 = 11

1 2 3 4  
 1 1 1 1

P4 : 2 4 3 1  $\Rightarrow$  1 1 1 1  
 Li : 0 1 0 1  
 P4  $\oplus$  Li : 1 0 1 0

Swap :

Li+1 : 1 1 0 1      Ri+1 : 1 0 1 0

To Round-2

S0 = 0 1 2 3  
 1 3 2 1 0  
 2 0 2 1 3  
 3 3 1 3 2

S0 o/p = 3

S1 = 0 1 2 3  
 1 2 0 1 3  
 2 3 0 1 2  
 3 2 1 0 3

S1 o/p = 3

## Round-2:

$L_{i+1} : 1\ 1\ 0\ 1$        $R_{i+1} : 1\ 0\ 1\ 0$   
 $R_{i+1} : 1\ 0\ 1\ 0$   
 $EP : 4\ 1\ 2\ 3\ 2\ 3\ 4\ 1$

$EP \text{ on } R_{i+1} : 0\ 1\ 0\ 1\ 0\ 1\ 0\ 1$

$K_2 : 0\ 1\ 0\ 0\ 0\ 0\ 1\ 1$

$EP \oplus K_2 : 0\ 0\ 0\ 1\ 0\ 1\ 1\ 0$

$S_0 : 0\ 0\ 0\ 1$

$S_1 : 0\ 1\ 1\ 0$

$01 \rightarrow 1^{\text{st}} \text{ Row}$   
 $00 \rightarrow 0^{\text{th}} \text{ Col}$

$00 \rightarrow 0^{\text{th}} \text{ Row}$   
 $11 \rightarrow 3^{\text{rd}} \text{ Col}$

$S_0 \text{ o/p} : 3 = 11$

$S_1 \text{ o/p} : 3 = 11$

$1\ 2\ 3\ 4$   
 $1\ 1\ 1\ 1$

$P_4(2, 4, 3, 1) : 1\ 1\ 1\ 1$

$L_{i+1} : 1\ 1\ 0\ 1$   
 $P_4 \oplus L_{i+1} : 0\ 0\ 1\ 0$   
 $1\ 2\ 3\ 4$

$1\ 0\ 1\ 0$   
 $5\ 6\ 7\ 8$

Swap not required after Last Round

$IP^{-1} : 4\ 1\ 3\ 5\ 7\ 2\ 8\ 6$

After  $IP^{-1} : 0\ 0\ 1\ 1\ 1\ 0\ 0\ 0$

$\therefore CT : 0\ 0\ 1\ 1\ 1\ 0\ 0\ 0$