

Experiment #13		Student ID	
Date		Student Name	

Experiment Title: Implementation of Programs/ Algorithms on CDP, NCDP and AOG problems.

Aim/Objective: To understand the concept and implementation of Basic program on NCDP as NP Hard problem

Description: The students will understand and able to implement programs on NCDP as NP Hard problem.

Pre-Requisites:

Knowledge: NCDP as NP Hard problem in C/C++/Python Tools: Code Blocks/Eclipse IDE

Pre-Lab:

1) Vijval has an array A of length N.

In one operation, Vijval can choose any two **distinct** indices i,j ($1 \leq i,j \leq N, i \neq j$) and **either** change A_i to A_j **or** change A_j to A_i .

Find the **minimum** number of operations required to make all the elements of the array **equal**.

Input Format

- o First line will contain T, number of test cases. Then the test cases follow.
- o First line of each test case consists of an integer N - denoting the size of array A.
- o Second line of each test case consists of N space-separated integers $1, 2, \dots, A_1, A_2, \dots, A_N$ - denoting the array A.

Output Format

For each test case, output the minimum number of operations required to make all the elements equal. Constraints

$$1 \leq T \leq 100$$

$$2 \leq N \leq 1000$$

$1 \leq A_i \leq 1000$ Sample 1:

Input4

4

4 5 6 7

3

6 6 6

4

Course Title	Design and Analysis of Algorithms	ACADEMIC YEAR: 2024-25
Course Code(s)	23CS2205R	Page 87 of 93

Experiment #13		
Date	Student ID	Student Name

3 3 2 2

3

2 2 3

Output3

0

2

1

Explanation:

Test Case 1: You can make all the elements equal in 22 operations. In the first operation, you can choose indices 4,5 and convert A1 to A2. So the array becomes [4,4,6,7]. Now you can choose indices 4,6 and convert A6 to A4, etc., so the final array becomes [4,4,4,4].

Test Case 2: Since all the elements are already equal there is no need to perform any operation.

- **Procedure/Program:**

```
#include <stdio.h>
int main()
{
    int T;
```

```
    scanf ("%d", &T);
```

```
    while (T--) {
```

```
        int N;
```

```
        scanf ("%d", &N);
```

```
        int A[N];
```

```
        int count[1001] = {0};
```

- **Data and Results:**

```
        for (int i=0; i<N; i++) {
```

Course Title	Design and Analysis of Algorithms	ACADEMIC YEAR: 2024-25
Course Code(s)	23CS2205R	Page 88 of 93

```
scanf("%d", &A[i]);
```

```
count[A[i]]++;
```

}

```
int maxCount = 0;
```

```
for (int i=1; i<=1000; i++) {
```

```
    if (count[i] > maxCount) {
```

```
        maxCount = count[i];
```

}

}

```
int minOperations = N - maxCount;
```

```
printf("%d\n", minOperations);
```

}

```
return 0;
```

}

Data

=

Test cases illustrate operations needed

for equalizing array elements

Result

= operations required for equal elements in

arrays are calculated accurately

- **Analysis and Inferences:**

Analysis

Minimum operations depend on frequency of most common array element
inferences

In-Lab: Higher frequency elements reduce required operations

1. The Clique Decision Problem belongs to NP-Hard. Prove that the Boolean Satisfiability problem reduces to the Clique Decision Problem

- **Procedure/Program:**

1) Input: A CNF formula ϕ with n variables and m clauses

2) graph construction:

- create a graph G with vertices representing literals (both x and $\neg x$)

- connect vertices with edges if:

- They belong to different clauses

- They are not contradictory

(i.e., no edges between x and $\neg x$)

- **Data and Results:** 3) clique size: set $k = n$

Course Title	Design and Analysis of Algorithms	ACADEMIC YEAR: 2024-25
Course Code(s)	23CS2205R	Page 89 of 93

Conclusion

=

- A satisfying assignment for ϕ corresponds to a clique of size n in G .
- Therefore, SAT reduces to another clique decision problem, providing that clique is NP-hard.

DATA and Result

Σ Σ \leftarrow

SAT reduces to clique decision problem, proving NP-hardness effectively.

Experiment #13		Student ID	
Date		Student Name	

- **Analysis and Inferences:**

clique Decision Problem's NP-hardness reflects complexity of Boolean satisfiability.

Post-Lab

1) You are given two integer A and B.

You need to compute and output the Greatest common divisor and Least common multiple of these 2 numbers and store them in the variables GCD and LCM.

Input Format

The first line of input will contain a single integer T, denoting the number of test cases. Each test case consists of a single line of input - the integer A and B.

Output Format

For each test case, output the GCD and LCM on one line

Sample 1:

Input

2

4 9

24 32

Output

1 36

8 96

- **Procedure/Program:**

```
#include <cmath>
int gcd(int a, int b) {
    while (b != 0) {
        int temp = b;
        b = a % b;
        a = temp;
    }
}
```

Course Title	Design and Analysis of Algorithms	ACADEMIC YEAR: 2024-25
Course Code(s)	23CS2205R	Page 90 of 93

```
3
return a;
3
int lcm (int a, int b, int gcdValue)
return (a * b) / gcdValue;
3
int main()
int T;
scanf ("%d", &T);
while (T--) {
int A, B;
scanf ("%d %d", &A, &B);
int gcdValue = gcd (A, B);
int lcmValue = lcm (A, B, gcdValue);
printf ("%d %d\n", gcdValue, lcmValue);
}
return 0;
3
```

Experiment #13		Student ID	
Date		Student Name	

• Data and Results:

Data Result
 2 GCD and LCM values
 4 9 computed successfully for all
 24 32 test cases

• Analysis and Inferences:

Analysis
 GCD helps simplify fractions; LCM aids in finding common multiples

• Sample VIVA-VOCE Questions:

- 1) Differentiate between CDP and NCDP? GCD and LCM are essential for number applications and computations
- 2) List the NP-Hard Graph problems?
- 3) What is Reducibility? theory
- 4) Identify one difference between Satisfiability and Reducibility?
- 5) What is AOG?

Evaluator Remark (if Any):	Marks Secured: ___ out of 50
	Signature of the Evaluator with Date

Evaluator MUST ask Viva-voce prior to signing and posting marks for each experiment.

Course Title	Design and Analysis of Algorithms	ACADEMIC YEAR: 2024-25
Course Code(s)	23CS2205R	Page 91 of 93

1) CDP is deterministic, NCDP is non-deterministic

Polynomial-time

2) Hamiltonian cycle, clique, vertex cover,
graph coloring

3) process of converting one problem
to another efficiently

4) satisfiability checks solutions;
reducibility simplifies problem complexity

5) bipartite with no both "and" and "or"
nodes