

# **Advanced Algorithms & Data Structures**











## Complex



Hands-on Technology

Brainstorming

Department of CSE

ADVANCED ALGORITHMS AND DATA STRUCTURES 23CS03HF

**Topic:** 

**Single Source Shortest Path Problem** 

Session - 23



Writing (Minute Paper)



Think-Pair-Share

**Groups Evaluations** 

Informal Groups

Self-assessment

Pause for reflection

Large Group Discussion

Case Studies

Triad Groups

Peer Review









### AIM OF THE SESSION



To familiarize students with the concept of Single Source Shortest Path Problem

## INSTRUCTIONAL OBJECTIVES



This Session is designed to:

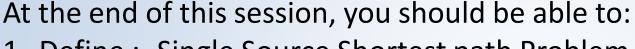
1.Demonstrate :- Single Source Shortest Path Problem.

===

2.Describe :- solve single-source shortest path problem using Dijkstra's

Algorithm.

### **LEARNING OUTCOMES**





- 1. Define :- Single Source Shortest path Problem.
- 2. Describe :- solve single-source shortest path problem using Dijkstra's Algorithm
- 3. Summarize:- Finding the shortest paths from a single source vertex to all other vertices in a weighted graph.



## **Shortest Path Problem**

- Shortest path problem is a problem of finding the shortest path(s) between vertices of a given graph.
- Shortest path between two vertices is a path that has the least cost as compared to all other existing paths.

## **Applications**

- Google Maps
- Road Networks
- Logistics Research



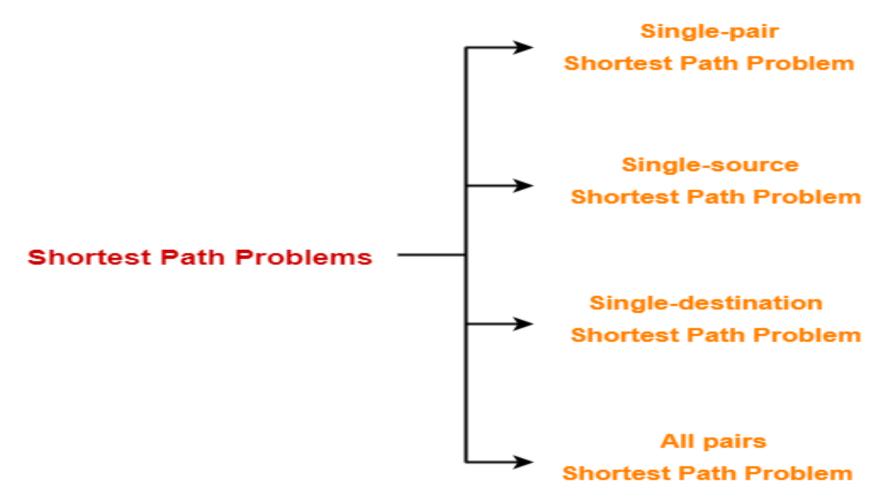








## **Types of Shortest Path Problem-**













## Single Source Shprtest Path Problem

- It is a shortest path problem where the shortest path from a given source vertex to all other remaining vertices is computed.
- Dijkstra's Algorithm and Bellman Ford Algorithm are the famous algorithms used for solving single-source shortest path problem.









## Dijkstra Algorithm

- Dijkstra Algorithm is a very famous greedy algorithm.
- It is used for solving the single source shortest path problem.
- It computes the shortest path from one particular source node to all other remaining nodes of the graph.









#### **Conditions**

It is important to note the following points regarding Dijkstra Algorithm-

- Dijkstra algorithm works only for connected graphs.
- Dijkstra algorithm works only for those graphs that do not contain any negative weight edge.
- The actual Dijkstra algorithm does not output the shortest paths.
- It only provides the value or cost of the shortest paths.
- Dijkstra algorithm works for directed as well as undirected graphs.

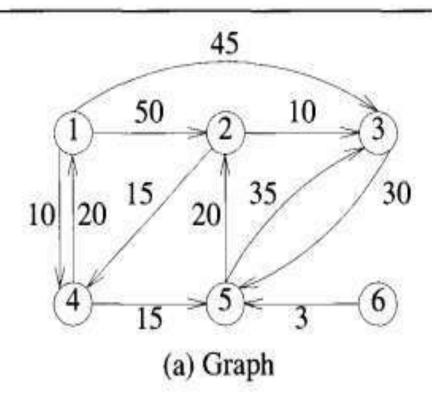












	Path	Length		
1)	1,4	10		
2)	1, 4, 5	25		
3)	1, 4, 5, 2	45		
	1, 3	45		
(b) SI	nortest path	s from 1		









## **Implementation**

Source Verki	intermediate Vertex	dist[2]	dic+[3]	dist[u]	dist[5]	Ca) taib	Path	Distance
1	-	50	45	10	<i>∞</i> 0	∞	1>4	10
	4	50	45	_	25	∞	1-4-5	25
		45 <b>®</b>	45	_	_	∞	1747572	45
	2	_	45	-	_	∞	1→3	45
	3	-		-		∞	No path to 6	









```
Algorithm ShortestPaths(v, cost, dist, n)
\frac{1}{2} \frac{3}{4} \frac{4}{5} \frac{6}{7} \frac{8}{9}
     // dist[j], 1 \le j \le n, is set to the length of the shortest
      // path from vertex v to vertex j in a digraph G with n
     // vertices. dist[v] is set to zero. G is represented by its
         cost adjacency matrix cost[1:n,1:n].
          for i := 1 to n do
           \{ // \text{ Initialize } S. 
                S[i] := false; dist[i] := cost[v, i];
10
          S[v] := \mathbf{true}; \ dist[v] := 0.0; // \ \mathrm{Put} \ v \ \mathrm{in} \ S.
11
          for num := 2 to n-1 do
12
13
14
                // Determine n-1 paths from v.
15
                Choose u from among those vertices not
                in S such that dist[u] is minimum;
16
17
                S[u] := \mathbf{true}; // \operatorname{Put} u \text{ in } S.
                for (each w adjacent to u with S[w] = false) do
18
19
                      // Update distances.
                     if (dist[w] > dist[u] + cost[u, w]) then
20
                                dist[w] := dist[u] + cost[u, w];
21
22
23
```



## **Time Complexity Analysis-**

- The given graph G is represented as an adjacency matrix.
- Priority queue Q is represented as an unordered list.

Here,

- A[i,j] stores the information about edge (i,j).
- Time taken for selecting i with the smallest dist is O(V).
- For each neighbor of i, time taken for updating dist[j] is O(1) and there will be maximum V neighbors.
- Time taken for each iteration of the loop is O(V) and one vertex is deleted from Q.
- Thus, total time complexity becomes O(V<sup>2</sup>).











#### **SUMMARY**

- •Objective: Find the shortest paths from a single source vertex to all other vertices in a weighted graph.
- •Approach: Initialize distances from the source to all vertices as infinity, update distances using relaxation to minimize path weights, and repeat until all shortest paths are found.
- •Algorithmic Solutions: Algorithms like Dijkstra's (for non-negative weights) used to solve this problem efficiently.











## **SELF-ASSESSMENT QUESTIONS**

Which of the following conditions must be true for Dijkstra's algorithm to work correctly?

- (a) The graph must be directed.
- (b) The graph must be undirected.
- (c) The graph must have non-negative edge weights.
- (d) The graph must have no cycles.

What is the main application of single source shortest path algorithms?

- (a) Finding Minimum Spanning Tree
- (b) Finding shortest paths from a single source node to all other nodes
- (c) Detecting cycles in a graph
- (d) Sorting nodes in topological order











## **TERMINAL QUESTIONS**

- 1. Can you explain the basic architecture of Dijkstra's algorithm along with an example?
- 2. What are some real-world scenarios where Dijkstra's algorithm can be applied effectively?









#### REFERENCES FOR FURTHER LEARNING OF THE SESSION

#### **Reference Books:**

- 1. Introduction to Algorithms, Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, and Clifford Stein., 3rd, 2009, The MIT Press.
- 2 Algorithm Design Manual, Steven S. Skiena., 2nd, 2008, Springer.
- 3 Data Structures and Algorithms in Python, Michael T. Goodrich, Roberto Tamassia, and Michael H. Goldwasser., 2nd, 2013, Wiley.
- 4 The Art of Computer Programming, Donald E. Knuth, 3rd, 1997, Addison-Wesley Professiona.

#### **MOOCS:**

- 1. https://www.coursera.org/specializations/algorithms?=
- 2.https://www.coursera.org/learn/dynamic-programming-greedy-algorithms#modules











# **THANK YOU**

















