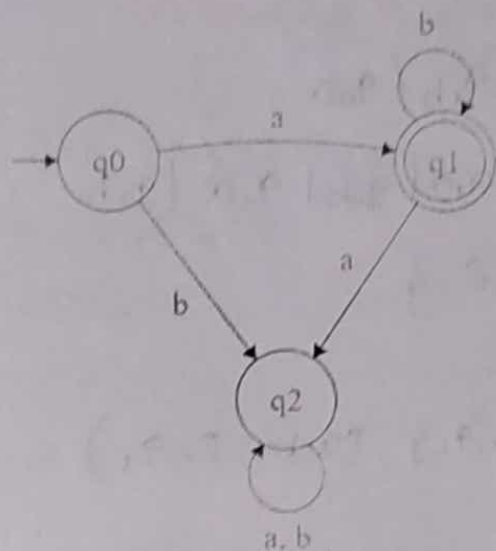


Pre-Tutorial (To be completed by student before attending tutorial session)

1. Consider the following automaton that accepts the language ab^*



Generate a Right Linear Grammar for the automaton.

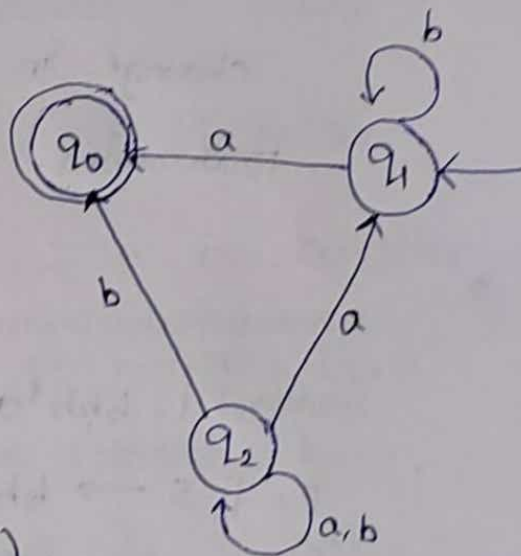
Solution:

$$P: \{ q_1 \rightarrow bq_1 / aq_0$$

$$q_2 \rightarrow aq_1 / bq_0 / aq_2 / bq_2$$

$$q_0 \rightarrow \epsilon \}$$

$$G = (\{q_0, q_1, q_2\}, \{a, b\}, P, q_1)$$



2. Write the Left Linear Grammar for the above Finite Automaton. What did you observe? Explain.

Solution:

$$P = \{ q_1 \rightarrow q_1 b / q_0 a$$

$$q_2 \rightarrow q_1 a / q_0 b / q_1 a / q_2 b$$

$$q_0 \rightarrow \epsilon \}$$

$$G = (\{q_0, q_1, q_2\}, \{a, b\}, P, q_1)$$

Here from right linear grammar, the change in left grammar is only in productions.

3. Write Left-Linear Grammar that generates the language $L(bbb^*aaaa^*)$

Solution: $L = bbb^*aaaa^*$

$$P = \{ S \rightarrow b b A a a a B$$

$$A \rightarrow b A / \epsilon$$

$$B \rightarrow a B / \epsilon \}$$

$$\{ S \rightarrow \overset{aaa}{A} \overset{bbb}{B} \overset{aaa}{A} \overset{bbb}{B} A$$

$$A \rightarrow \overset{aaa}{A} b A / \epsilon$$

$$B \rightarrow a B / \epsilon \}$$

Left linear Grammar:

$$P' = \{ S \rightarrow B a a a A b b$$

$$A \rightarrow A b / \epsilon$$

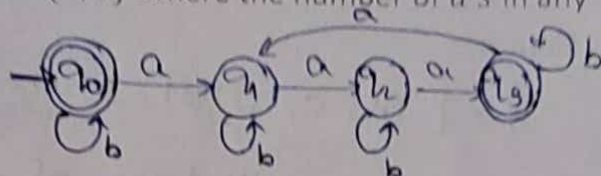
$$B \rightarrow B a / \epsilon \}$$

IN-TUTORIAL (To be carried out in presence of faculty in classroom)

1. Find a regular grammar for the language over $\{a, b\}$ where the number of a 's in any string is divisible by 3.

Solution:

$$L = (aaa)^*$$



$$G = (\{q_0, q_1, q_2, q_3\}, \{a, b\}, P, q_0)$$

$$P = \{ q_0 \rightarrow bq_0 \mid aq_1 \mid \epsilon$$

$$q_1 \rightarrow bq_1 \mid aq_2$$

$$q_2 \rightarrow bq_2 \mid aq_3$$

$$q_3 \rightarrow bq_3 \mid aq_0 \mid \epsilon \}$$

2. Let L_1 and L_2 be two languages with $L_1 \subseteq L_2$. In what condition L_2 must be a regular language?

Solution:

for $L_1 = L_2$ both L_1 and L_2 are regular

languages. For $L_1 \subset L_2$, there is no condition

for L_2 to be a regular or non-regular

Because here there is no condition on L_1 .

so, $L_1 \subseteq L_2$ is regular and L_2 is also

regular language.

3. Let $L = \{w \in \{0,1\}^* : n_0(w) = n_1(w)\}$. Is L regular? If you find it to be not, prove it. Here $n_0(w)$ = count of zeros in w , and $n_1(w)$ = count of 1's in w .

Solution:

$$L = \{01, 0101, 1100, 001011, \dots\}$$

By using applications of pumping lemma,

Step 1:

let, constant $n = 4$

Step 2: $z = uv^i w$, $|z| \geq n$

where, $|uv| \leq n$

$$|v| \geq 1$$

$$z = \frac{1100}{uv \quad w}$$

$$4 \geq 4$$

$$|uv| \leq n \Rightarrow 2 \leq 4 \checkmark$$

$$|v| \geq 1 \Rightarrow 1 \geq 1 \checkmark$$

$$uv^i w \in L$$

$$i = 0, 1, 2, 3, \dots$$

$$i = 0 \rightarrow 1100 = 100 \notin L$$

L is not regular.

Tutorial #

07

Student ID

2100030642

Date

12/9/24

Student Name

S. Arditi

4. Is the language $L = \{a^n b^n : n \geq 0\}$ a regular language? If you find it to be not a regular language prove it. Write each step.

Solution:

$$L = \{ \epsilon, ab, aabb, aaabbb, \dots \}$$

1. let $n = 4$

2. $z = uvw$

$$z = \underline{aabb}$$

$$\quad \quad \quad \underline{uv} \quad \underline{w}$$

$$|z| \geq n \Rightarrow |z| \geq 4$$

$$|uv| \leq n \Rightarrow 2 \leq 4$$

$$|v| \geq 1 \Rightarrow 1 \geq 1$$

3. $uv^i w$

$$i = 0, 1, 2, 3, \dots$$

$$i = 0 \Rightarrow a(a^0)b^4 = aabb \notin L$$

$\therefore L$ is not regular.

5. Consider the grammar: $S \rightarrow aS \mid bA$, $A \rightarrow bA \mid \epsilon$. Is it a regular grammar? it to be a regular grammar determine whether it is right linear or left linear language generated by the given grammar.

Solution:

$$P = \{ S \rightarrow aS \mid bA \\ A \rightarrow bA \mid \epsilon \}$$

$$G = (\{S, A\}, \{a, b\}, P, S)$$

It is ~~regular~~ regular grammar.

Because it consists,

it is in the form $T^*V \mid T^*$.

Given production is in the form of T^*V .

So, it is Right regular grammar.

Post-Tutorial (To be carried out by student after attending tutorial session)

Consider the grammar: $S \rightarrow aS \mid bS \mid \epsilon$. Is it a regular grammar? If you find it to be regular, determine its type. Write the language generated by the grammar.

ution: It is in the form of $T^*V \mid T^*$. So, it is regular grammar.

$$S \rightarrow aS \mid bS \mid \epsilon$$

$$\therefore L = (a^*b^*)^*$$

Let $L_1 = \{a^n : n > 0\}$ and $L_2 = \{b^n : n > 0\}$ be two regular languages. Consider another language L_3 where $L_3 = L_1 \cup L_2$. Is L_3 a regular language? Explain.

Solution:

If L_1 ~~and~~ is a regular language.

L_2 is a regular language.

If L_1 and L_2 are regular languages, then $L_3 = L_1 \cup L_2$ is regular language.

$$L_3 = \{a^n, b^n : n > 0\}$$

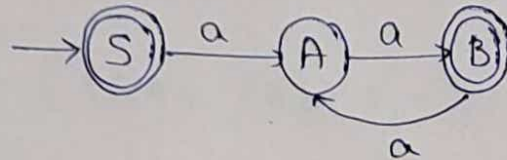
$\therefore L_3$ is regular language.

3. Consider the language $L = \{a^n; n \text{ is even}\}$. Is L a regular language? P_1

Solution:

Solution:

$$L = \{ \epsilon, aa, aaaa, \dots \} = (aa)^*$$



$$P = \{ S \rightarrow QA \mid \epsilon \}$$

$$A \rightarrow QB$$

$$B \rightarrow \{aA | a \in \Sigma\}$$

$$G = (\{S, A, B\}, \{a\}, P, S)$$

\therefore It is regular language (T^*v/T^*)

(For Evaluator's use only)

Comment of the Evaluator (if Any)	Evaluator's Observation
	Marks Secured: _____ out of <u>50</u>
	Full Name of the Evaluator:
	Signature of the Evaluator Date of
	Evaluation: