

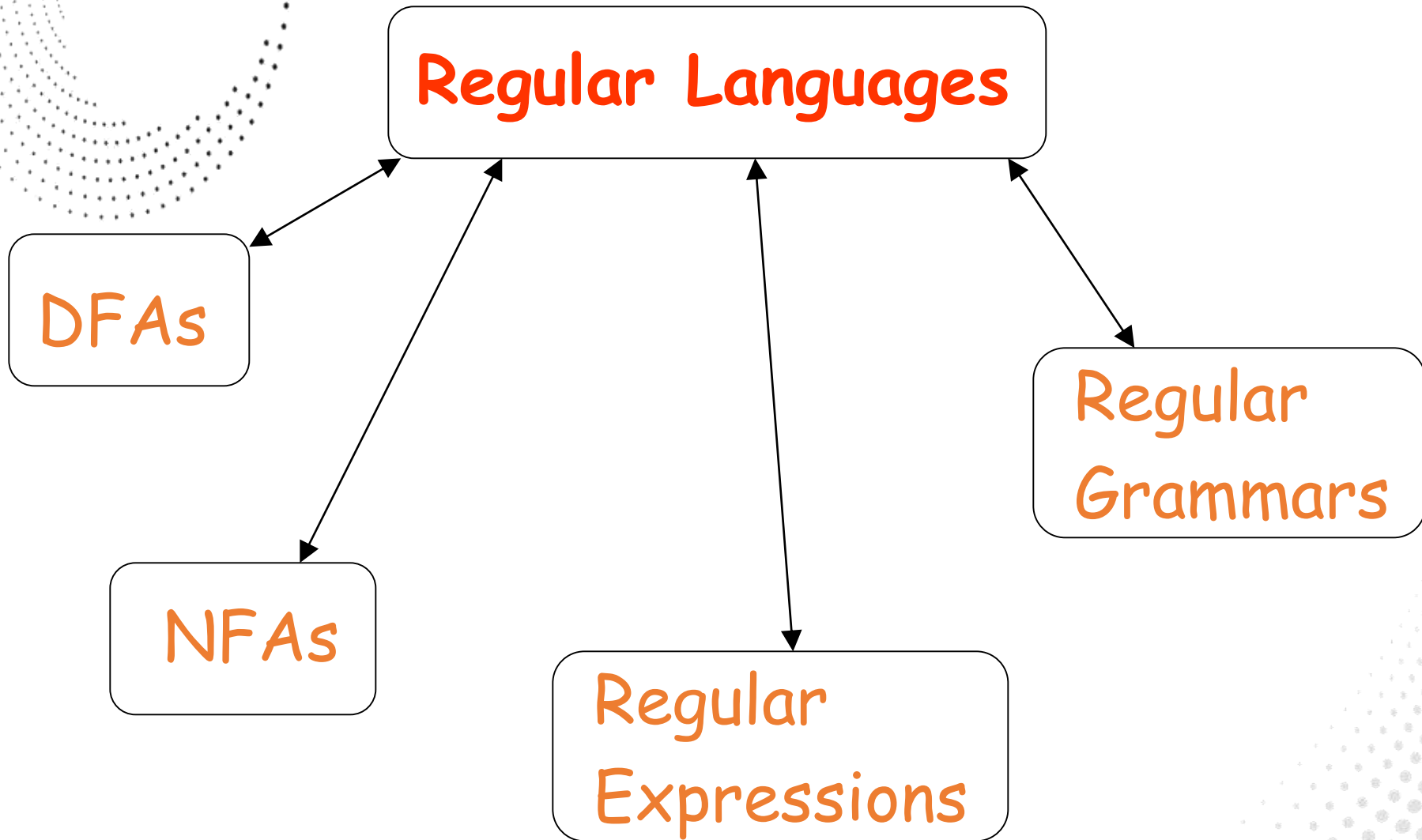
23MT2014

## THEORY OF COMPUTATION

Topic:

# PROPERTIES OF REGULAR LANGUAGE AND PUMPING LEMMA

Session – 9-a



When we say: We are given  
a Regular Language  $L$

We mean: Language  $L$  is in a standard  
representation

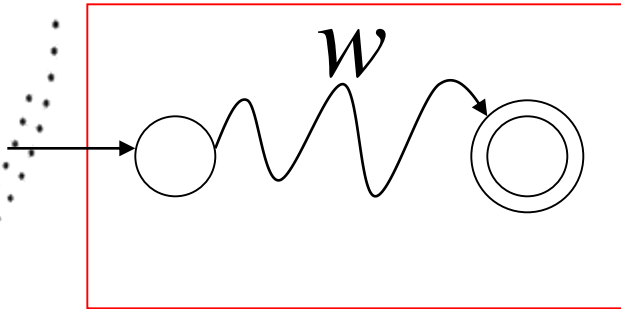
# Elementary Questions about Regular Languages

# Membership Question

**Question:** Given regular language  $L$   
and string  $w$   
how can we check if  $w \in L$ ?

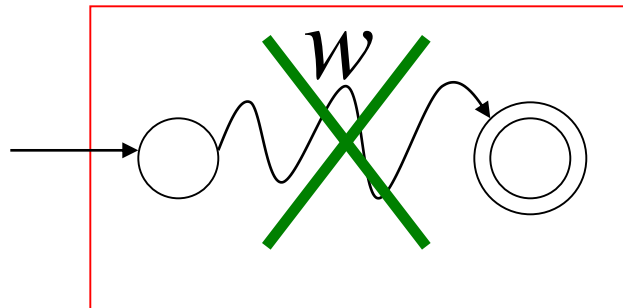
**Answer:** Take the DFA that accepts  $L$   
and check if  $w$  is accepted

DFA



$w \in L$

DFA



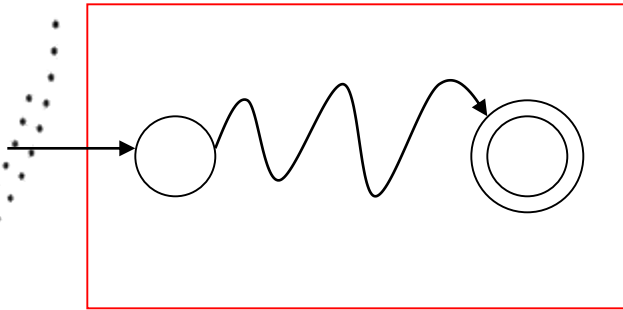
$w \notin L$

**Question:** Given regular language  $L$   
how can we check  
if  $L$  is empty:  $(L = \emptyset)$  ?

**Answer:** Take the DFA that accepts  $L$

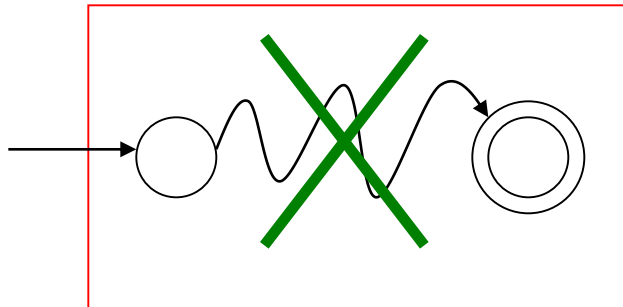
Check if there is any path from  
the initial state to a final state

DFA



$$L \neq \emptyset$$

DFA



$$L = \emptyset$$

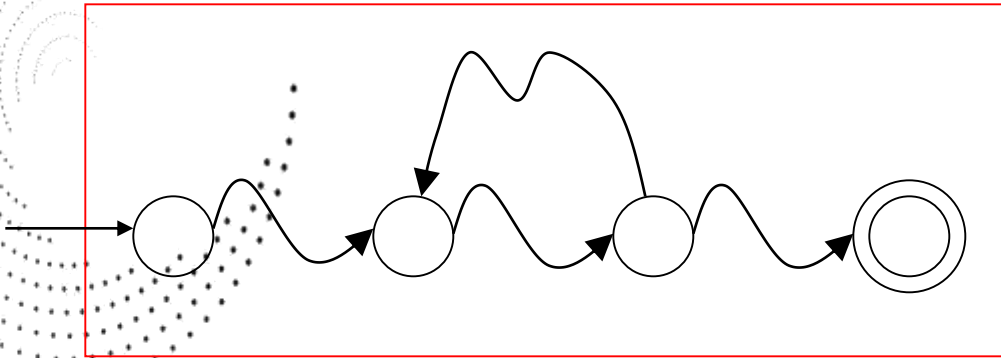


**Question:** Given regular language  $L$   
how can we check  
if  $L$  is finite?

**Answer:** Take the DFA that accepts  $L$

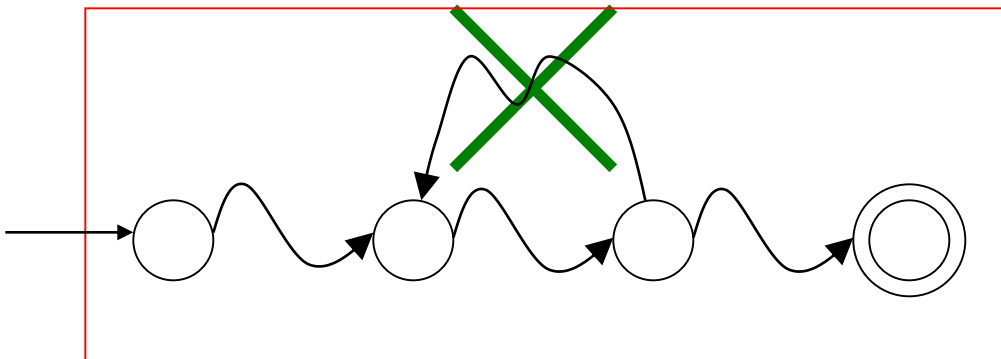
Check if there is a walk with cycle  
from the initial state to a final state

DFA



$L$  is infinite

DFA

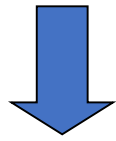


$L$  is finite

**Question:** Given regular languages  $L_1$  and  $L_2$   
how can we check if  $L_1 = L_2$  ?

**Answer:** Find if  $(L_1 \cap \overline{L_2}) \cup (\overline{L_1} \cap L_2) = \emptyset$

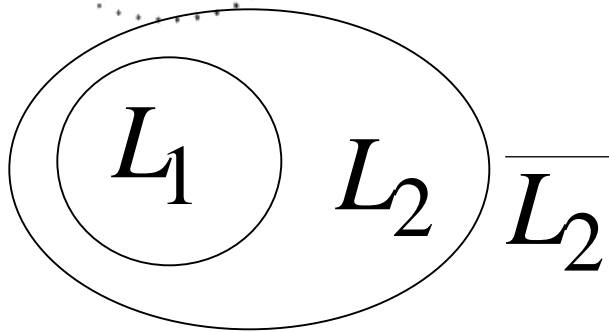
$$(L_1 \cap \overline{L_2}) \cup (\overline{L_1} \cap L_2) = \emptyset$$



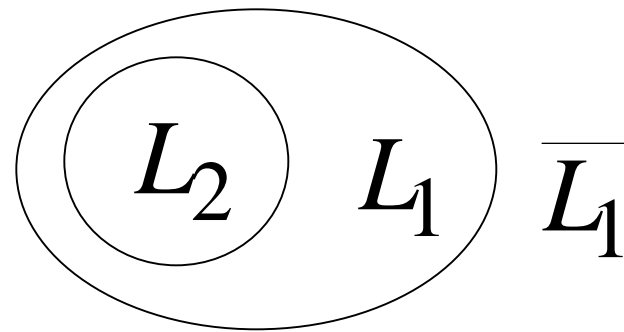
$$L_1 \cap \overline{L_2} = \emptyset$$

and

$$\overline{L_1} \cap L_2 = \emptyset$$



$$L_1 \subseteq L_2$$



$$L_2 \subseteq L_1$$



$$L_1 = L_2$$

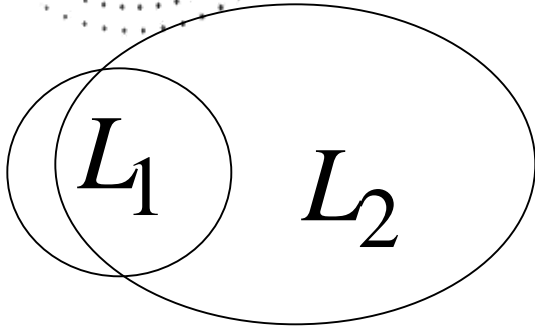
$$(L_1 \cap \overline{L_2}) \cup (\overline{L_1} \cap L_2) \neq \emptyset$$



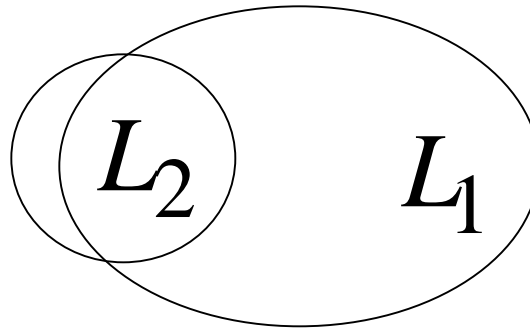
$$L_1 \cap \overline{L_2} \neq \emptyset$$

or

$$\overline{L_1} \cap L_2 \neq \emptyset$$



$$L_1 \not\subseteq L_2$$



$$L_2 \not\subseteq L_1$$



$$L_1 \neq L_2$$

# Non-regular languages

## Non-regular languages

$$\{a^n b^n : n \geq 0\}$$

$$\{vv^R : v \in \{a,b\}^*\}$$

## Regular languages

$$a^*b$$

$$b^*c + a$$

$$b + c(a + b)^*$$

*etc...*

How can we prove that a language  $L$  is not regular?

Prove that there is no DFA that accepts  $L$

**Problem:** this is not easy to prove

**Solution:** the Pumping Lemma !!!



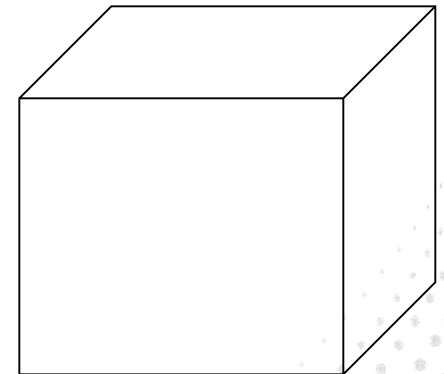
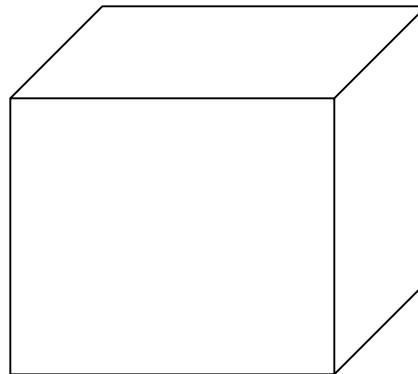
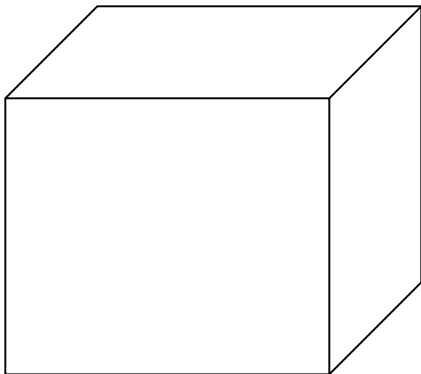


# The Pigeonhole Principle

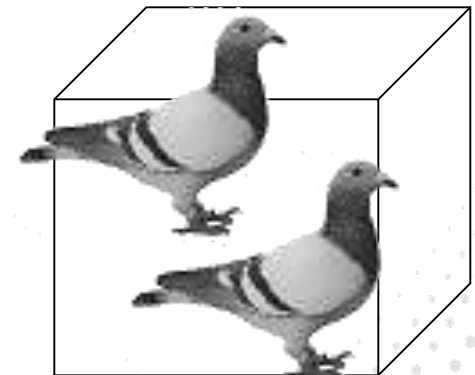
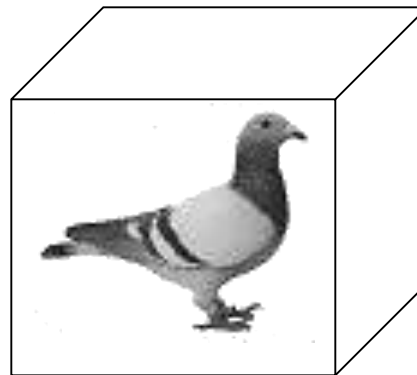
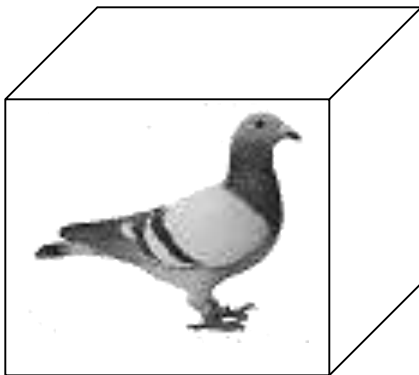
4 pigeons



3 pigeonholes



A pigeonhole must contain at least two pigeons

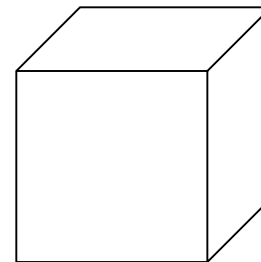
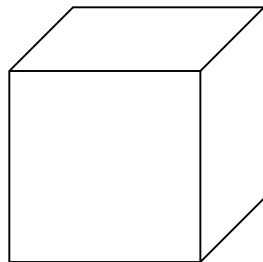
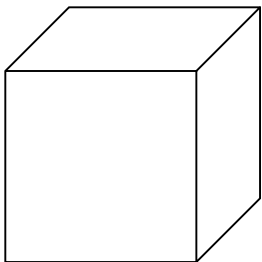


$n$  pigeons



$m$  pigeonholes

$n > m$



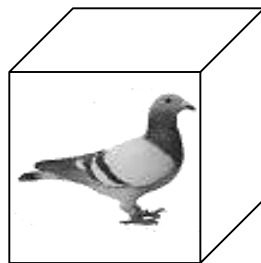
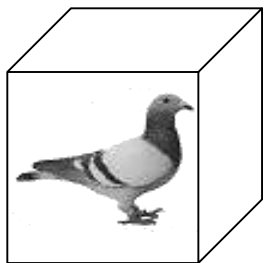
# The Pigeonhole Principle

$n$  pigeons

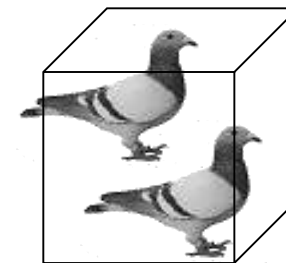
$m$  pigeonholes

$$n > m$$

There is a pigeonhole with at least 2 pigeons

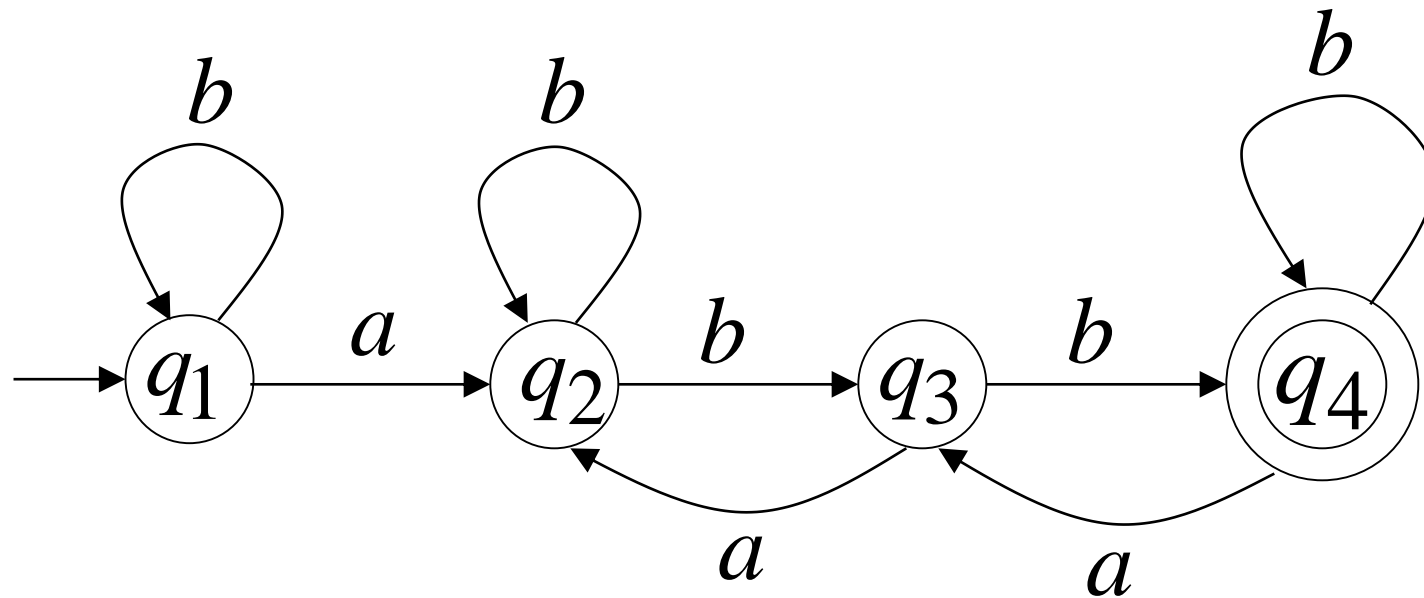


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# The Pigeonhole Principle and DFAs

# DFA with 4 states



In walks of strings:

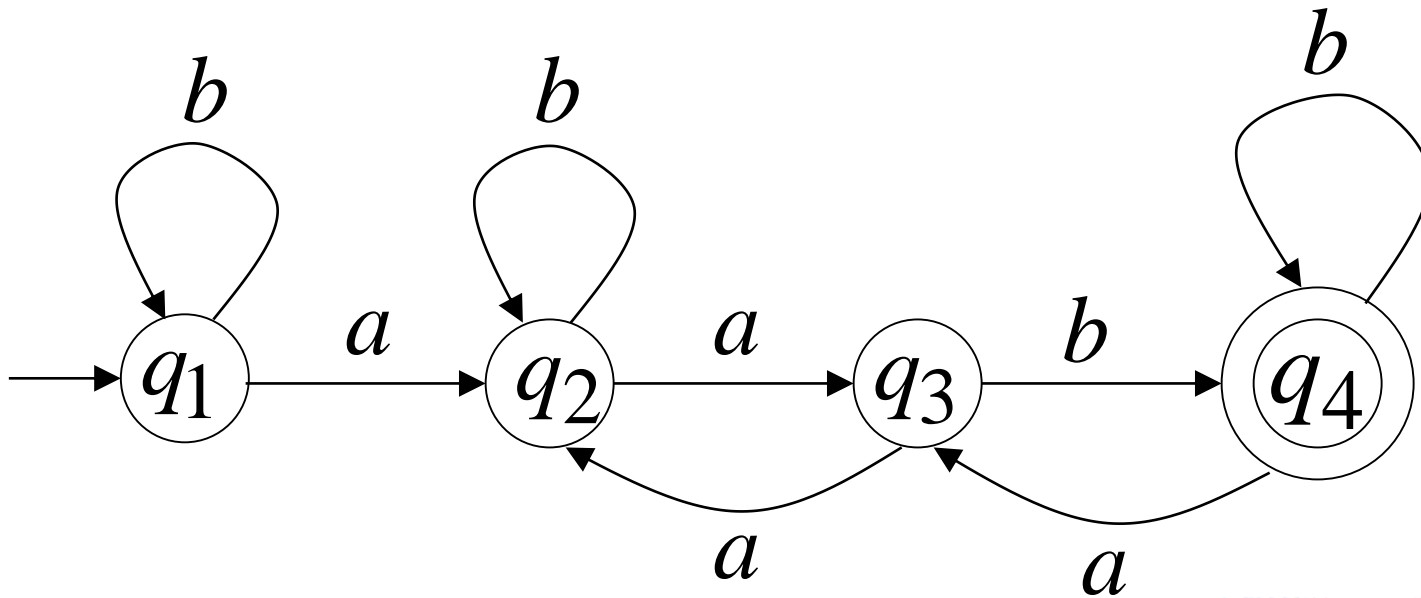
$a$

$aa$

$aab$

no state

is repeated





In walks of strings:

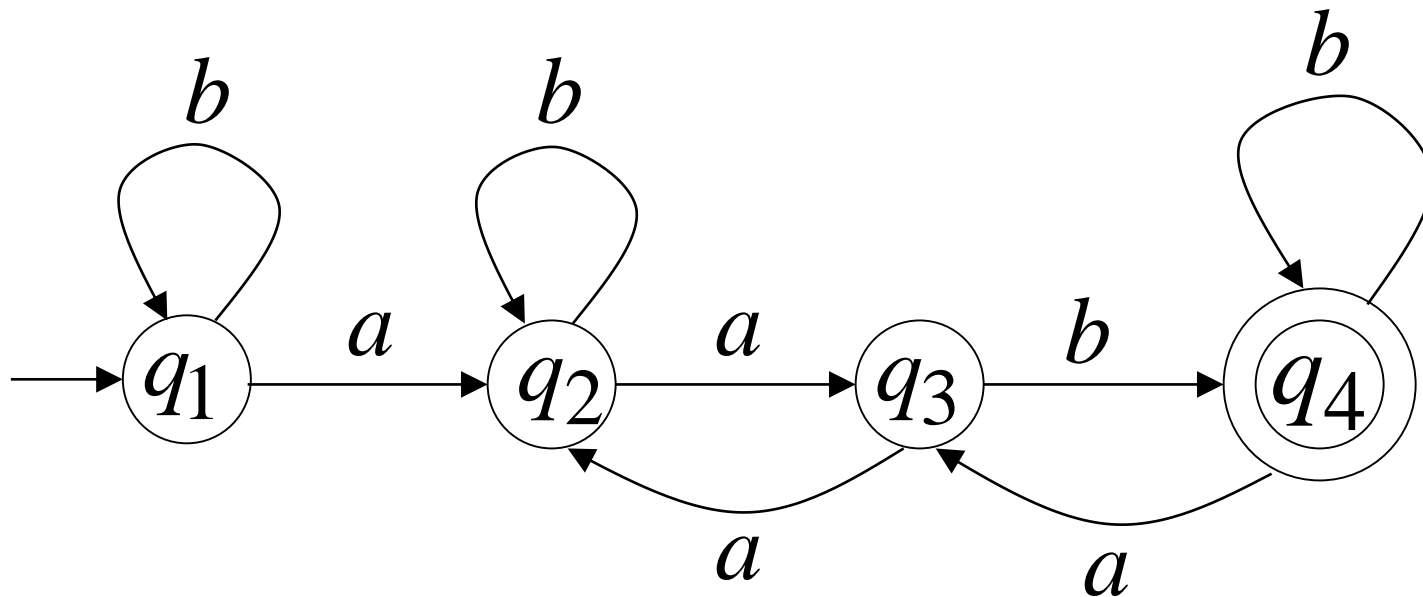
$aabb$

$bbaa$

$abbabb$

$abbbabbabb...$

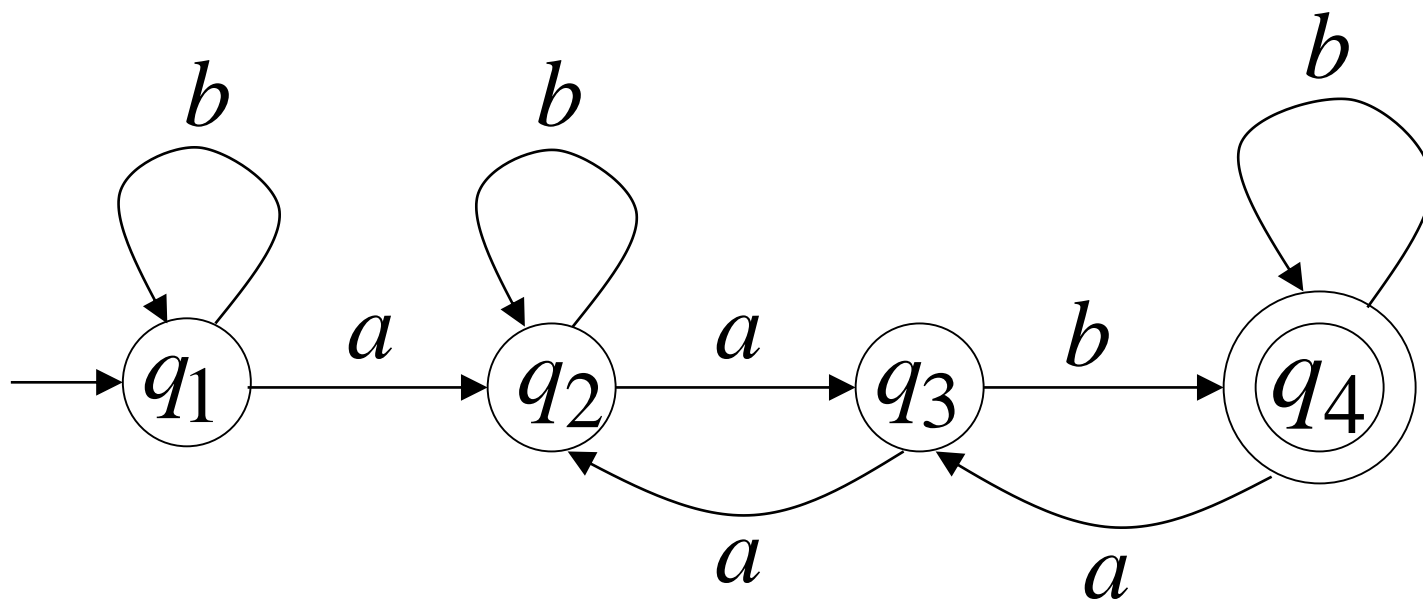
a state  
is repeated



If string  $w$  has length  $|w| \geq 4$ :

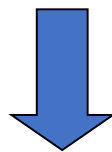
Then the transitions of string  $w$   
are more than the states of the DFA

Thus, a state must be repeated

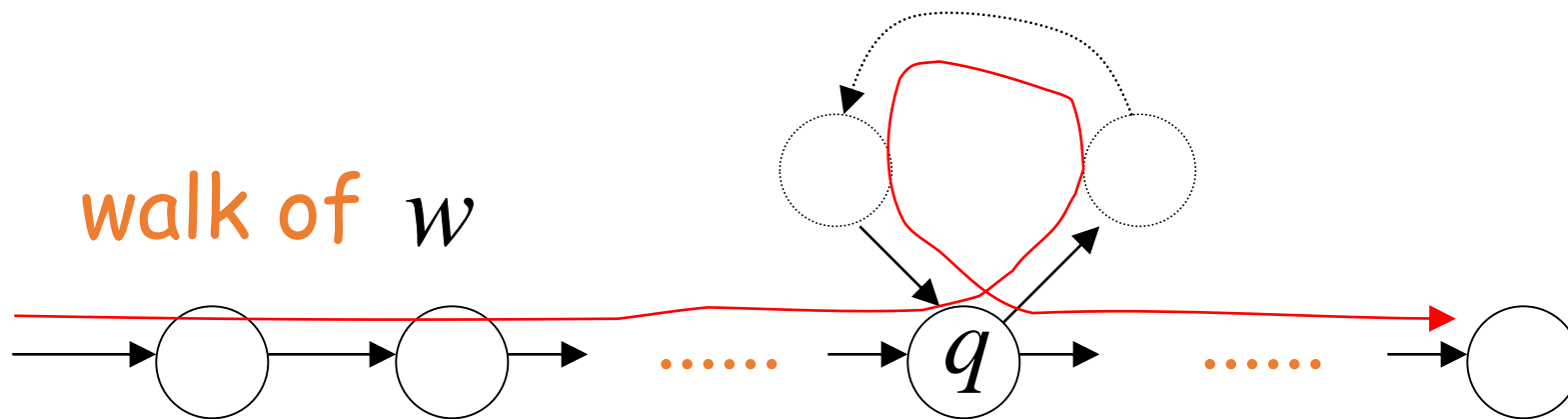


 In general, for any DFA:

String  $w$  has length  $\geq$  number of states

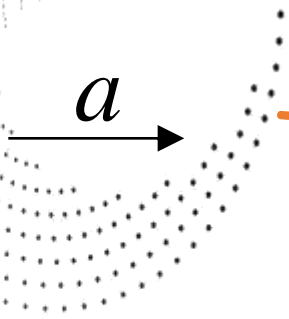


A state  $q$  must be repeated in the walk of  $w$



Repeated state

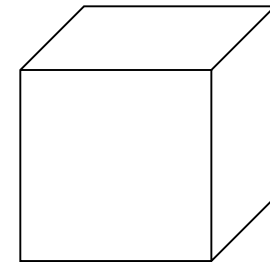
In other words for a string  $w$ :



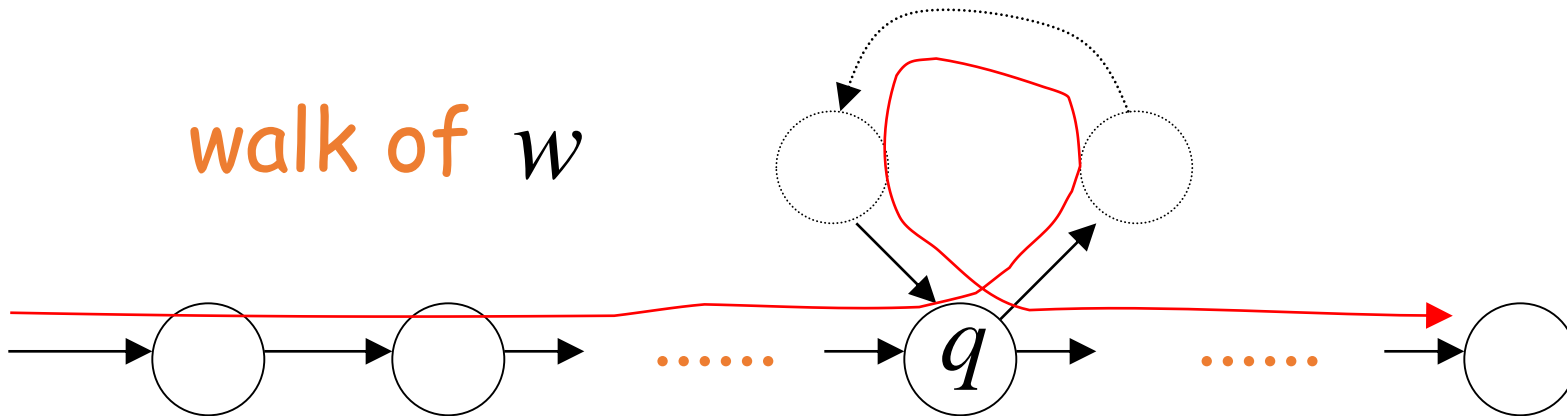
transitions are pigeons



states are pigeonholes



walk of  $w$



Repeated state

# The Pumping Lemma:

- Given a infinite regular language  $L$
- there exists an integer  $m$
- for any string  $w \in L$  with length  $|w| \geq m$
- we can write  $w = x y z$
- with  $|x y| \leq m$  and  $|y| \geq 1$
- such that:  $x y^i z \in L \quad i = 0, 1, 2, \dots$



**THANK YOU**



**Team – TOC**