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Date	Student Name

Experiment Title: Analysis of Algorithm based on Arrays

Aim/Objective: To understand the concept and implementation of Basic programs on arrays.

Description: The students will understand and able to implement programs on Arrays.

Pre-Requisites:

Knowledge: Arrays

Tools: Code Blocks/Eclipse IDE

Pre-Lab:

Calculate the time complexity of the following:

1. Algorithm linear_search(arr, target):

for i in range(len(arr)):

if arr[i] == target:

return i

return -1

• Procedure/Program: `#include <studio.h>`
`int linear_search (int arr[], int size, int target)`
`{`
`for (int i = 0; i < size; i++) {` $- (n+1)$
`if (arr[i] == target)` $-$
`{`
`return i;` $-$
`}` $-$
`}` $-$
 $(n+1) + n + 1$
 $= 2n + 1$
 $O(n)$

• Data and Results: let arr size is (n) the
Best case: - the element is the first element in array
Time complexity $O(1)$
worst case: The element is not present in the array
Time complexity $O(n)$ $- O(n)$

• Analysis and Inferences: The linear search of an algorithm in each element in the array until it finds the target element or reaches
Best case: $O(1)$; - worst case: $O(n)$

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2. Algorithm function(arr):

n = len(arr)

for i in range(n): $\rightarrow (n+1)$

for j in range(0, n-i-1): $n(n+1)$

if arr[j] > arr[j+1]: $n(n)$

arr[j], arr[j+1] = arr[j+1], arr[j] \cap

return arr;

- Procedure/Program: no. of bubble-sort(int arr[], info)
for(int i=0; i<n-1; i++) $\cap (n+1)$
for(int j=0; j<n-i-1; j++) $\cap (n+1)$
if(arr[j] > arr[j+1]) $\cap (n)$
int temp = arr[j]; $\cap (n)$
arr[j] = arr[j+1]; $\cap (n)$
arr[j+1] = temp; $\cap (n)$

- Data and Results: let input array [13, 25, 12, 22, 11]
Step-1: [25, 34, 12, 11, 22]
Step-2: [25, 11, 34, 1, 22]
Step-3: [12, 25, 11, 34, 22]
Step-4: [12, 11, 20, 22, 34]
Step-5: [11, 12, 22, 25, 34]

• Analysis and Inferences:

Best case - The array is already sorted make a Passes but no swap performed $O(n)$

Worst case - it takes $n-1$ comparisons

this result is fine completely $O(n^2)$

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3. The median of a list of numbers is essentially its middle element after sorting. The same number of elements occur after it as before. Given a list of numbers with an odd number of elements, find the median? Also find its time complexity.

Example arr=[5,3,1,2,4]

The sorted array arr=[1,2,3,4,5]. The middle element and the median is 3.

Description: Write an algorithm findMedian with the following parameter(s):

arr[n]: an unsorted array of integers

Returns the median of the array

Sample Input

7

0 1 2 4 6 5 3

Sample Output

3

• Procedure/Program:

```

int compare(const void*a, const void*b)
{
    return (*(int*)a) - (*(int*)b);
}

int findMedian(int arr[], int n) {
    qsort(arr, n, sizeof(int), compare);
    int mid = n/2;
    return arr[mid];
}

```

• Data and Results:

let array [1, 2, 3, 4, 5, 6, 7]
Median = 4

• Analysis and Inferences: The most time-complexity of an algorithm. The time complexity = $O(n \log n)$

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In-Lab:

- Given an array of strings, sort them in alphabetical order. Find the

• Proc
Data

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{

ment after sorting. The same number
with an odd number of elements, find

the median is 3.
wing parameter(s):

```
const void* a)
fa) * (int*) b);
(int n) {
(int), compare)
);
```

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In-Lab:

- Given an array of strings arr[], the task is to sort the array of strings according to frequency of each string, in ascending order. If two elements have same frequency, then they are to be sorted in alphabetical order.

Input: arr[] = {"Ramesh", "Mahesh", "Mahesh", "Ramesh"}

Output: {"Mahesh", "Ramesh"}

Explanation: As both the strings have the same frequency, the array is sorted in alphabetical order.

Find the time and space complexity for the procedure/Program.

• Procedure/Program:

```
Data struct {
    char str[100];
    int freq;
    string freq;
}

int compare (const void* a, const void* b)
{
    string freq *strA = (string freq*) a;
    string freq *strB = (string freq*) b;
    if (strA->freq == strB->freq)
        return strcmp(strA->str, strB->str);
    return strA->freq - strB->freq;
}
```

• Data and Results:

```
{ Ramesh, 2; Mahesh, 2}
{ Mahesh, 2; Ramesh, 2}
sort array (Mahesh, Ramesh)
```

```
void sortByFrequency (char*
                        a[], n)
{
    string frequency[100];
    int freq count = 0;
    int found;
    for (int i = 0; i < n; i++)
    {
        found = 0;
        for (int j = 0; j < freq count; j++)
        {
            if (strcmp(freq[j], strA[i]) == 0)
            {
                freq[j] = freq[j] + strA[i];
                found = 1;
                break;
            }
        }
    }
}
```

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• Analysis and Inferences:

2. Given an array of strings words [] and the sequential order of alphabets, our task is to sort the array according to the order given. Assume that the dictionary and the words only contain lowercase alphabets.

Input: words = {"word", "world", "row"},

Order= "worldabceefghijklmnpqstuvwxyz"

Output: {"world", "word", "row"}

Explanation: According to the given order 'l' occurs before 'd' hence the words "world" will be kept first.

Find the time and space complexity for the procedure/Program.

• Procedure/Program:

```
int compare(const void *a, const void *b,
            void *order - map) {
    string order *strA = (string order *) a;
    string order *strB = (string order *) b;
    char *word1 = strA -> str;
    char *order = S((int *) order - map);
    int len1 = strlen(word1);
    int len2 = strlen(word2);
```

```
for(int i=0; i<len1; i++) {
    if(word1[i] != word2[i])
        return order[(int)word1[i]] -
            order[(int)word2[i]];
}
return len1 - len2;
}
```

• Data and Results:

words = {"word", "world", "row"}
 order: "worldabceefghijklmnpqstuvwxyz"
 {w, o, d, l, a, b, c, e, f, g, h, i, j, k, l, m, n, o, p, q, r, s, t, u, v, w, x, y, z}
 Fig -->
 sorted word = {"world", "word", "row"}

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• Analysis and Inferences:

creating ordermap $O(16) = O(1)$ there are 26 characters in the alph
 The complexity is $O(m \log n)$ where m is the number of words
 space of complexity $O(16) = O(1)$
 $O(n)$

Post-Lab:

Evaluate the time complexity of pre-lab using master's theorem.

• Procedure/Program:

$$T(n) = aT(n/b) + f(n) \quad T = 3T(n/2) + n$$

where

$$a = 3$$

$$n/b = n/2$$

n = size of an input $f(n) = n$

a = number of each supportable All sub Problems
 are assumed to have the same size

$f(n)$ = cost of the work done outside the recursive call which including the cost of dividing the Problem

$a > 1$ and $b > 1$ are constants

and $f(n)$ is as asymptotically for

• Data and Results:

$$\text{If } f(n) = O(n \log b a^{-s}) \text{ then} \\ T(n) = O(n \log b a)$$

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$$E \quad f(n) = O(n \log b) \quad \text{the } T(n) = O(n \log a \cdot \log n)$$

$$f(n) = a(n \log b \text{ atc}) \quad \text{the } T(n) = O(f(n))$$

• Analysis and Inferences:

n = size of an input
 a = number of each subproblem
 note $a=1$ and $b=1$ are constants and
 $f(n)$ is a significantly

• Sample VIVA-VOCE Questions (In-Lab):

1. What is the significance of analyzing both time complexity and space complexity while evaluating algorithms?
2. How does the choice of data structures impact both time and space complexity in algorithm implementations? Give examples.
3. When comparing two algorithms with different time and space complexities, how do you decide which one is more suitable for a particular problem or application?
4. What is the difference between best-case, average-case, and worst-case time complexity? Provide examples.
5. In the context of sorting algorithms, compare the time and space complexities of algorithms like Quick Sort and Merge Sort. Which one is more time-efficient, and which one is more space-efficient?

Evaluator Remark (if Any):	Marks Secured: <u>42</u> out of 50
	Signature of the Evaluator: <u>R. Chiranjeev</u> With Date

Evaluator MUST ask Viva-voce prior to signing and posting marks for each experiment.

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$$T(n) = O(n \log a \cdot \log b)$$
$$T(n) = O(F(n) \cdot g(n))$$

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constants and

complexity and space complexity when

time and space complexity in algorithms

and space complexities, how do you

m or application?

and worst-case time complexity?

space complexities of algorithms

efficient, and which one is more

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each experiment.

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- ① Time complexity of an algorithm is the amount of time taken for an algorithm and function of time and input.
- ② Using the hash table can reduce the average time complexity of search operation of complexity $O(1)$ (or) $O(n)$.
- ③ If two algorithms with the different time and space complexity the using memory and input.
- ④ Best case:- minimum steps on input data of an element.
 - Worst case:- the function which performs the minimum number of steps input has size.
 - Average case:- The steps perform the average number of steps of an element.

Time complexity is $3(n \log n)$ of a quick sort of an algorithm

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