

SUM OF SUBSETS PROBLEM BY BACKTRACKING











SUM OF SUBSETS

Problem Definition: Given n distinct positive numbers wi, and m, find all subsets whose sums are m.

- Explicit Constraints :
- Xi = { j / j is an integer and 1 <= j <= n }
- Implicit Constraints :
- 1. No two xi's are same
- 2. Sum of the chosen weights must be equal to m.
- 3. To avoid generation of multiple instances of the same subsets)

We can formulate this problem using Fixed tuple or Variable tuples.











VARIABLE- SIZED TUPLE

Ex:- n=4, (w1, w2, w3, w4)= (11,13, 24, 7), m=31.

- Solutions are (11, 13, 7) and (24, 7)
- Rather than representing the solution by wi's, we can represent by giving the indices of these wi
- Now the solutions are (1, 2, 4) and (3, 4).
- Different solutions may have different-sized tuples.
- We use the following condition to avoid generating multiple instances of the same subset (e.g., (1,2,4) and (1,4,2))

$$x_i < x_i + 1$$

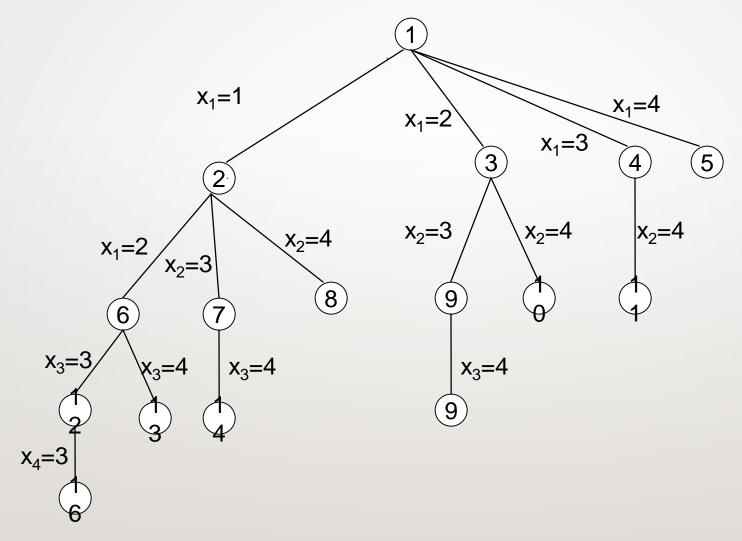








Subset Tree for n=4 (variable- sized tuple)



Nodes are numbered as in Breadth first search.











FIXED- SIZED TUPLE

- In this method, each solution subset is represented by an n-tuple (x1,x2,....x n).
- x i = 0 if wi is not chosen and xi = 1 if wi is chosen

Ex:- n=4, (w1, w2, w3, w4)= (11,13, 24, 7), m=31.

Solutions are (11, 13, 7) and (24, 7)









• Solutions are:

• X=(1,1,0,1) and X=(0,0,1,1)



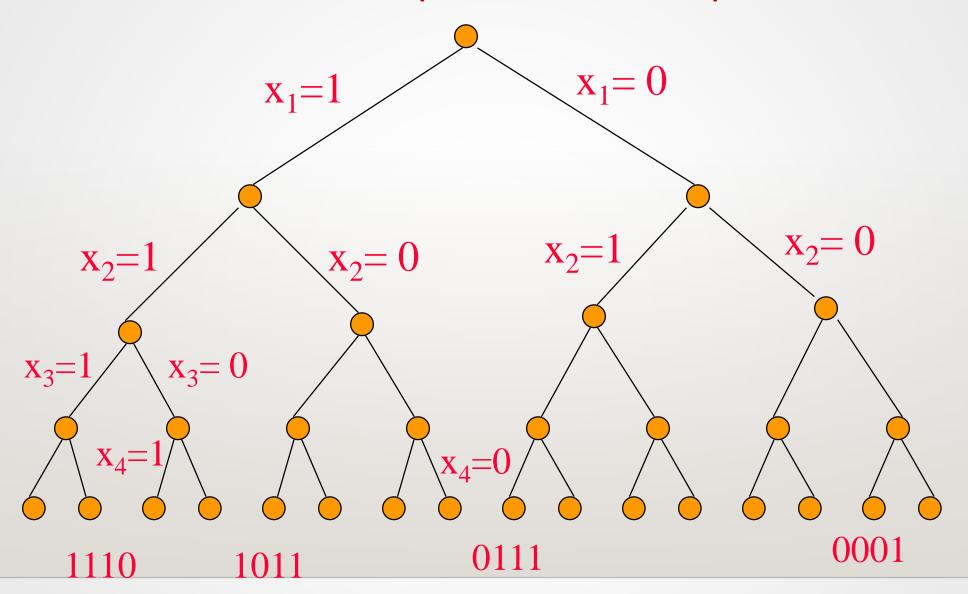








SUBSET TREE FOR N = 4 (FIXED - SIZED TUPLE)





BOUNDARY CONDITIONS OF SUM OF SUBSETS

Simple choice for the bounding Function is Bk (X1 ... Xk) = true iff

$$\sum_{i=1}^{k} w_i x_i + \sum_{i=k+1}^{n} w_i \ge m$$

Clearly x1 ...xk can not lead to an answer node if this condition is not satisfied.

$$\sum_{i=1}^{k} w_{i} x_{i} + w_{k+1} > m$$

Assuming wi's in non decreasing order, (x1... xk) cannot lead to an answer node if So, the bounding functions we use are therefore

$$B_{k}(x_{1},...,x_{k}) = true \ iff \sum_{i=1}^{k} w_{i}x_{i} + \sum_{i=k+1}^{n} w_{i} \ge m$$

$$and \sum_{i=1}^{k} w_{i}x_{i} + w_{k+1} \le m$$









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Algorithm SumOfSub(s, k, r)
// Find all subsets of w[1:n] that sum to m
// It is assumed that w[1] \le m and \sum wi \ge m
           x[k]=1; // left child
           if(s + w[k] = m) then write(x[1:k]); // Subset found
        else if (s + s [k] + s [k+1] \le m)
               then SumOfSub(s+w[k], k+1, r-w[k])
           // Generate right child and evaluate Bk
     if ((s+r-w[k] \ge m)) and (s+w[k+1] \le m)) then
                      x[k] = 0;
                      SumOfSub(s, k+1, r-w[k]);
```





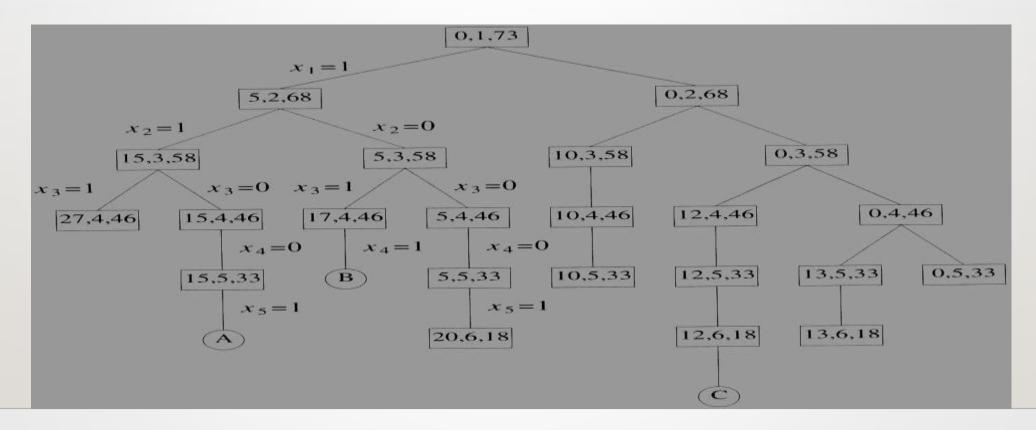






EX:- N=6, M=30, W [1:6]= { 5,10,12,13,15,18 }
PORTION OF STATE SPACE TREE GENERATED BY SUMOFSUB.
CIRCULAR NODES INDICATE SUBSETS WITH SUMS EQUAL TO M.

- Example 7.6 (Figure 7.10)
- n=6, w[1:6]={5,10,12,13,15,18}, m=30













SAMPLE QUESTIONS

- What is the Sum of Subsets problem
- Describe the problem statement of the Sum of Subsets. What are the inputs and outputs of the problem
- Explain about state space tree in sum of subsets problem
- Solve sum of subsets by using Input: set[] = {4, 16, 5, 23, 12}, sum = 9



