

Pass-2 Program :

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

struct symbol {
    int index;
    char label[20];
    int address;
} symtab[50];

int symcount = 0;

int get_address_by_index(int idx) {
    for (int i = 0; i < symcount; i++) {
        if (symtab[i].index == idx)
            return symtab[i].address;
    }
    return -1;
}

void load_symtab() {
    FILE *fp = fopen("symtab.txt", "r");
    while (fscanf(fp, "%d %s %d", &symtab[symcount].index, symtab[symcount].label,
    &symtab[symcount].address) != EOF) {
        symcount++;
    }
}
```

```
}

fclose(fp);

}

int main() {

FILE *ic = fopen("intermediate.txt", "r");

FILE *mc = fopen("machine_code.txt", "w");

load_symtab();

char line[100];

int lc = 100;

while (fgets(line, sizeof(line), ic)) {

int opcode, reg, symidx;

char type[5];

if (sscanf(line, "(%[^,],%d) (R,%d) (S,%d)", type, &opcode, &reg, &symidx) == 4) {

int addr = get_address_by_index(symidx);

fprintf(mc, "%d\t%02d\t%d\t%03d\n", lc, opcode, reg, addr);

lc++;

} else if (sscanf(line, "(%[^,],%d)", type, &opcode) == 2 && strcmp(type, "IS") == 0) {

fprintf(mc, "%d\t%02d\t0\t000\n", lc, opcode);

lc++;

} else if (strstr(line, "(DL,02)")) {

int constant;

sscanf(line, "(DL,02) (C,%d)", &constant);
```

```

        fprintf(mc, "%d\t00\t0\t%03d\n", lc, constant);

        lc++;

    } else if (strstr(line, "(DL,01)")) {

        int size;

        sscanf(line, "(DL,01) (C,%d)", &size);

        for (int i = 0; i < size; i++) {

            fprintf(mc, "%d\t00\t0\t000\n", lc);

            lc++;

        }

    }

}

fclose(ic);

fclose(mc);

printf("Pass-II complete.\n");

return 0;

}

```

Output :

machinecode.txt :

100	00	0	000
101	04	1	102
102	00	0	005
103	05	1	102

Algorithm of Pass 2

- **Initialize:**
 - Load the symbol table and literal table created during Pass 1.
 - Prepare for object code
 - **Process Each Line of Intermediate Code:**
 - For each line in the intermediate representation
 - a. **Parse the line** into label, mnemonic, and operand.
 - b. **Handle Pseudo-Ops** (e.g., START, END, RESW, RESB):
 - Process directives appropriately (e.g., START specifies the starting address).
 - c. **Process Mnemonics:**
 - Search the opcode table for the mnemonic.
 - If found, generate the corresponding machine code.
 - If the operand is a symbol, look it up in the symbol table to get its address.
 - If the operand is a literal, handle it accordingly.
 - d. **Handle Literals:**
 - Assign addresses to literals and include them in the object code.
 - e. **Handle Forward References:**
 - If an operand refers to a label not yet defined, leave a placeholder and backpatch it later.
 - f. **Generate Object Code:**
 - Assemble the object code instruction.
 - If the object code will not fit into the current text record, write the current text record to the object program and initialize a new one.
 - Add the object code to the text record.
- **Finalize:**
 - Write the last text record to the object program.
 - Write the end record to the object program.
 - Write the last listing line.

▀ **%02d:** Prints a signed decimal integer with a minimum width of 2 characters, padding with leading zeros if necessary.

▀ **%03d**: Prints a signed decimal integer with a minimum width of 3 characters, padding with leading zeros if necessary.

Breakdown of (%[^,], %d)

1. **(**: Matches the literal opening parenthesis `(` in the input string.
2. **%[^,]**: This is a scanset conversion specifier. It reads characters into a string until it encounters a comma `,`. The `^` inside the brackets negates the set, so it matches any character except a comma.
3. **,**: Matches the literal comma `,` in the input string.
4. **%d**: Reads an integer (signed decimal) from the input string.
5. **)**: Matches the literal closing parenthesis `)` in the input string.