

CS341: Computer Architecture Lab

# Lab Assignment 4

## Report

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## Abstract

Summarize the objective of the lab, what experiments you have conducted, what were the results that you have obtained in a clear and concise manner. Numbers matter, not just words only, for ex. *very high, slow* etc.

## Part 0: Getting Things Ready

### Install Intel VTune Profiler

Installed successfully using the stand-alone app using offline installer script present in this link.

It was pretty easy to install VTune using the script.

While installing it showed that I didn't have XCB and DRM packages installed. Upon checking, I confirmed that they were already present.

Even though it failed prerequisites, there was a next option. I didn't face any issues for the rest of installation process.

From start to end, it took around 10-12 minutes to have the application installed, followed by 3-5 minutes for tutorial.

### Install Docker

Had docker setup from other projects.

Version: 20.10.9

### Pull ChampSim Image

Pulled 0xd3ba/champsim-lab:latest

# Part 1: Profiling with VTune

## 1.1 bfs.cpp

### Performance Snapshot

- IPC: 1.830
- Logical Core Utilization: 8.2% (0.979 out of 12)
- Physical Core Utilization: 16.2% (0.973 out of 6)
- Memory bound: 32.0% of Pipeline slots

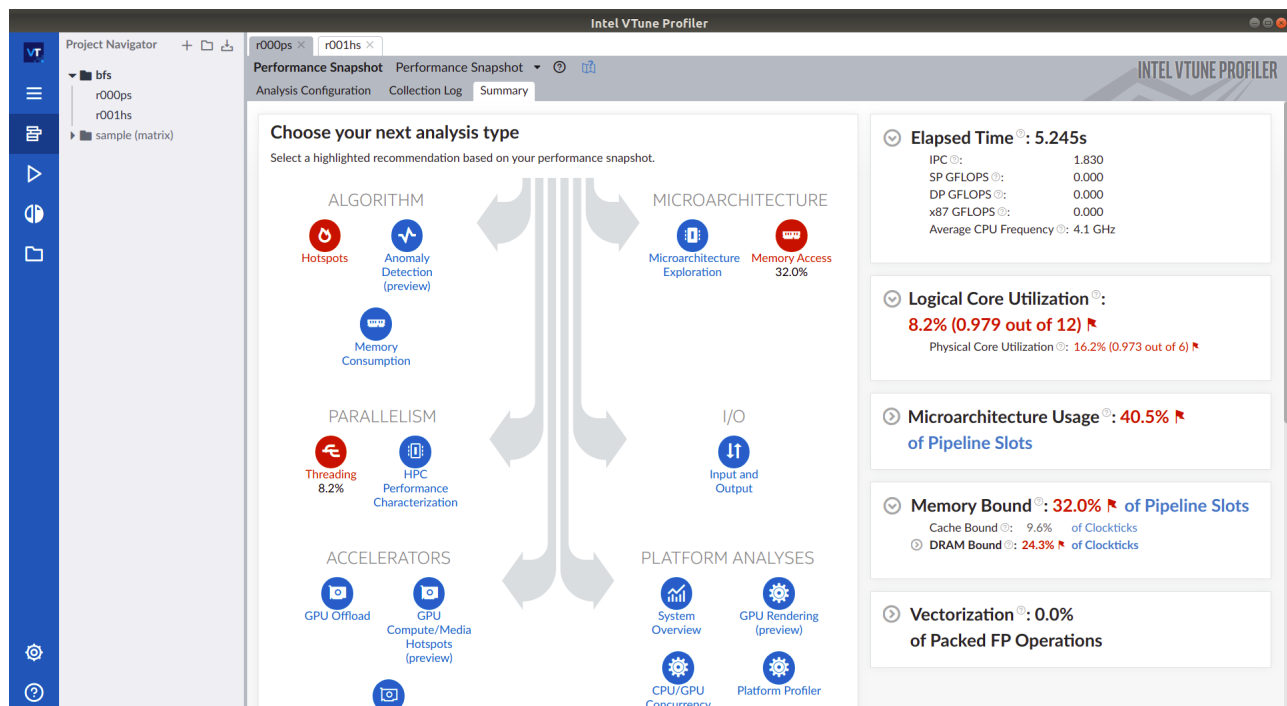


Figure 1.1: Performance Snapshot for bfs.cpp

## Top 5 Functions by CPU Time

Function	Module	CPU Time
bfs	bfs.o	2.621s
main	bfs.o	1.156s
_int_free	libc-2.27.so	0.236s
_int_malloc	libc-2.27.so	0.154s
__gnu_cxx::new_allocator<Node*>::construct<Node*, Node* const&>	bfs.o	0.124s

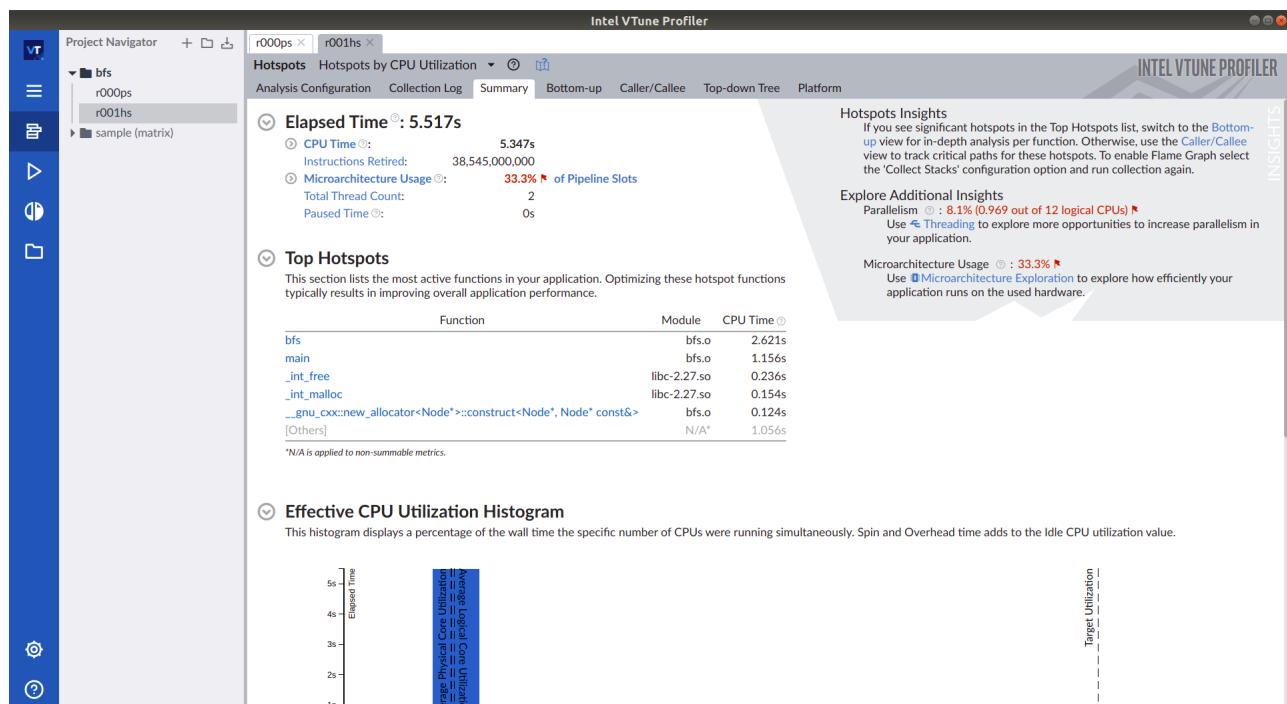


Figure 1.2: Top Functions by CPU Time for bfs.cpp

## Top 5 Source lines by CPU Utilization

Source	Function	CPU Utilization
if (left_child) node_Q.push(left_child);	inline void bfs(Node *root)	22.7%
bfs(root);	int main()	21.6%
right_child = curr_node->right;	inline void bfs(Node *root)	17.2%
for (int i = 0; i < q_size; i++) {	inline void bfs(Node *root)	4.8%
left_child = curr_node->left;	inline void bfs(Node *root)	3.2%

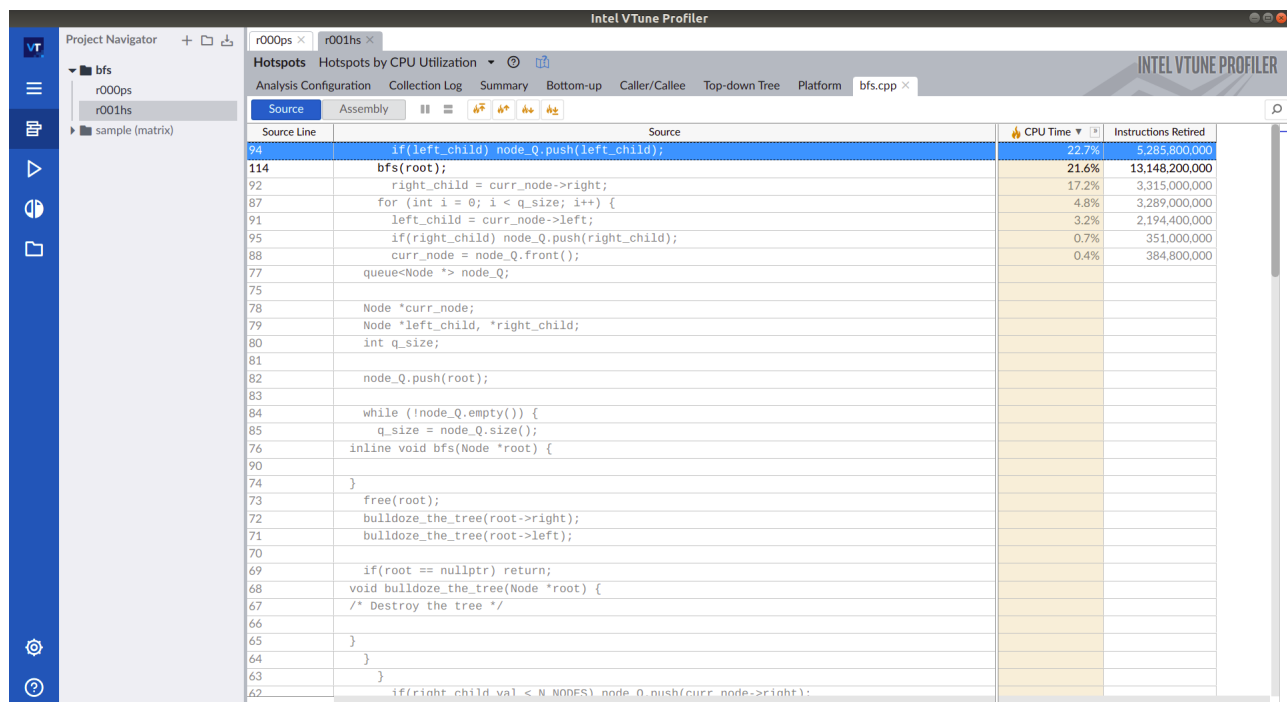


Figure 1.3: Top Source lines by CPU Utilization for bfs.cpp

## Inference

We see that majority of the time goes in function `bfs`.

The time consuming line in `main` is calling `bfs(root)`. This would be due to the need to set the function stack and as it is done for `N_LOOPS` times, it climbs above `plant_a_tree` and `bulldoze_the_tree` stack setup.

Every line in `bfs` is called repeatedly due to 3 level of loops (2 level in `bfs` and 1 level in `main`).

These lines combined consume the majority of CPU time (50%).

## 1.2 matrix\_multi.cpp

### Performance Snapshot

- IPC: 0.874
- Logical Core Utilization: 8.2% (0.982 out of 12)
- Physical Core Utilization: 16.3% (0.976 of 6)
- Memory bound: 59.1% of Pipeline slots

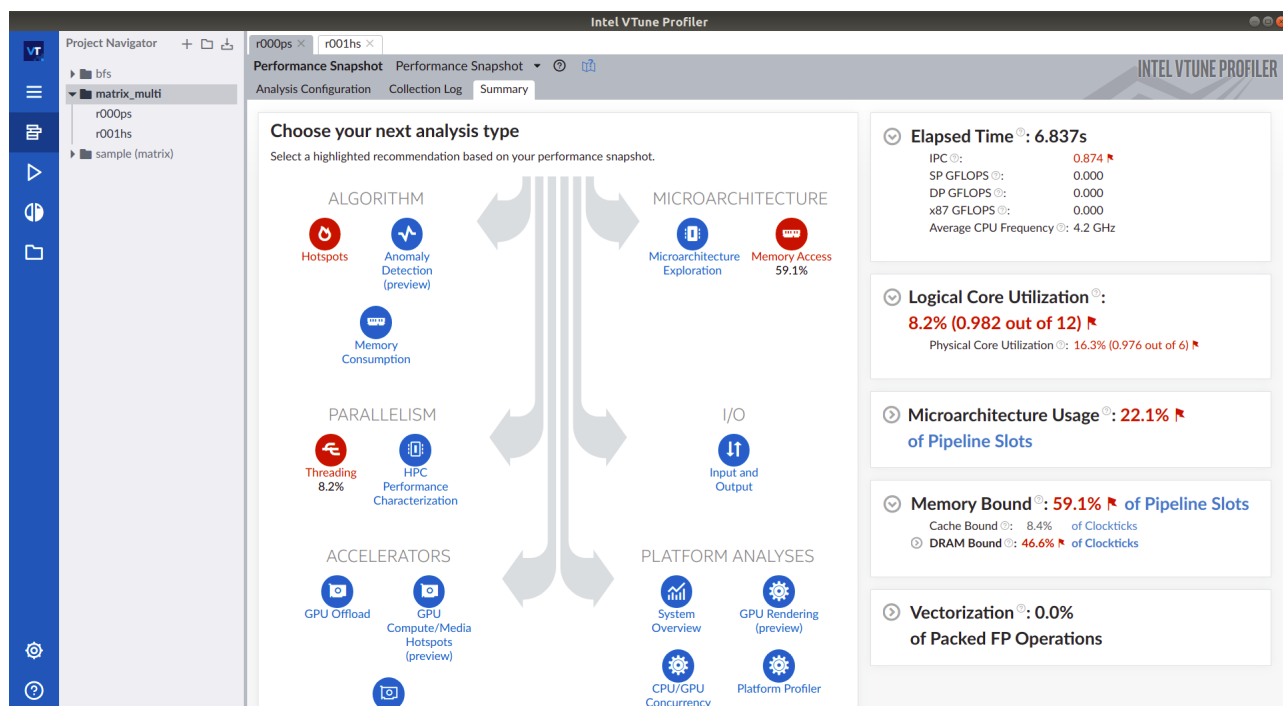


Figure 1.4: Performance Snapshot for `matrix_multi.cpp`



## Top Functions by CPU Time

Function	Module	CPU Time
matrix_product	matrix_multi.o	6.597s

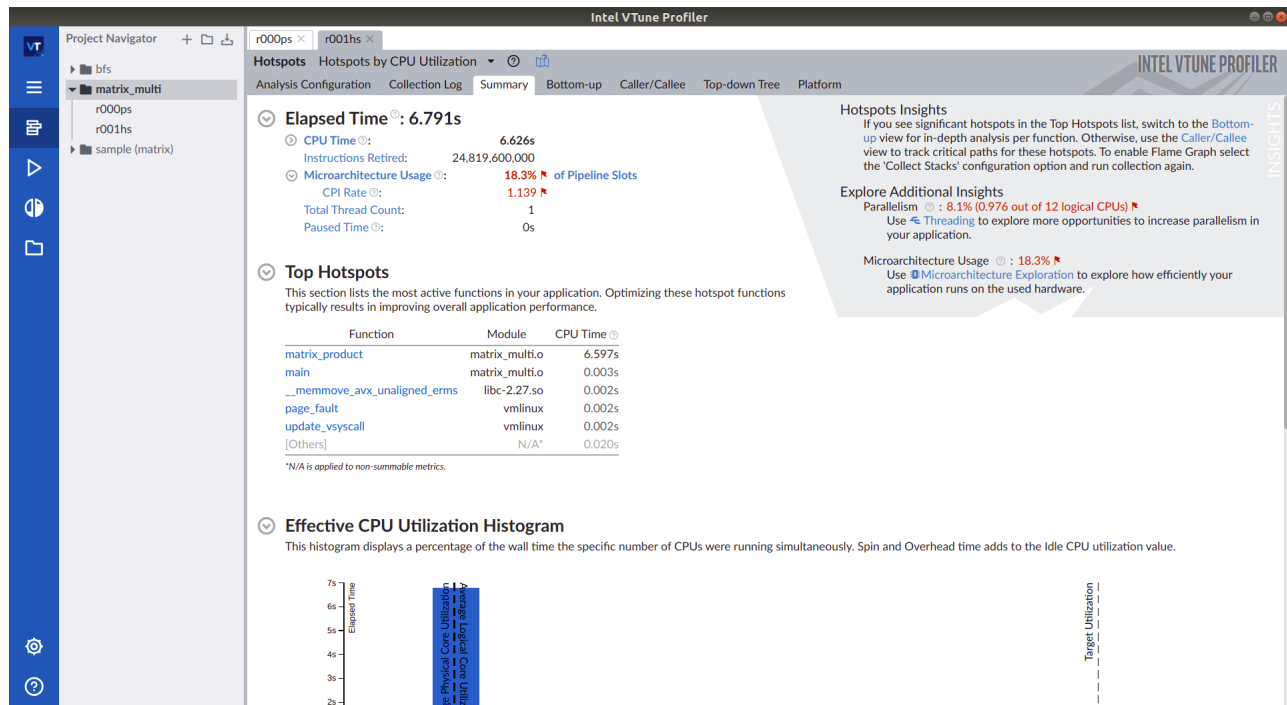


Figure 1.5: Top Functions by CPU Time for matrix\_multi.cpp

## Top 2 Source lines by CPU Utilization

Source	Function	CPU Utilization
<code>C[i][j] += A[i][k] * B[k][j];</code>	<code>void matrix_product()</code>	90.9%
<code>for (int k = 0; k &lt; N_DIMS; k++) {</code>	<code>void matrix_product()</code>	8.6%

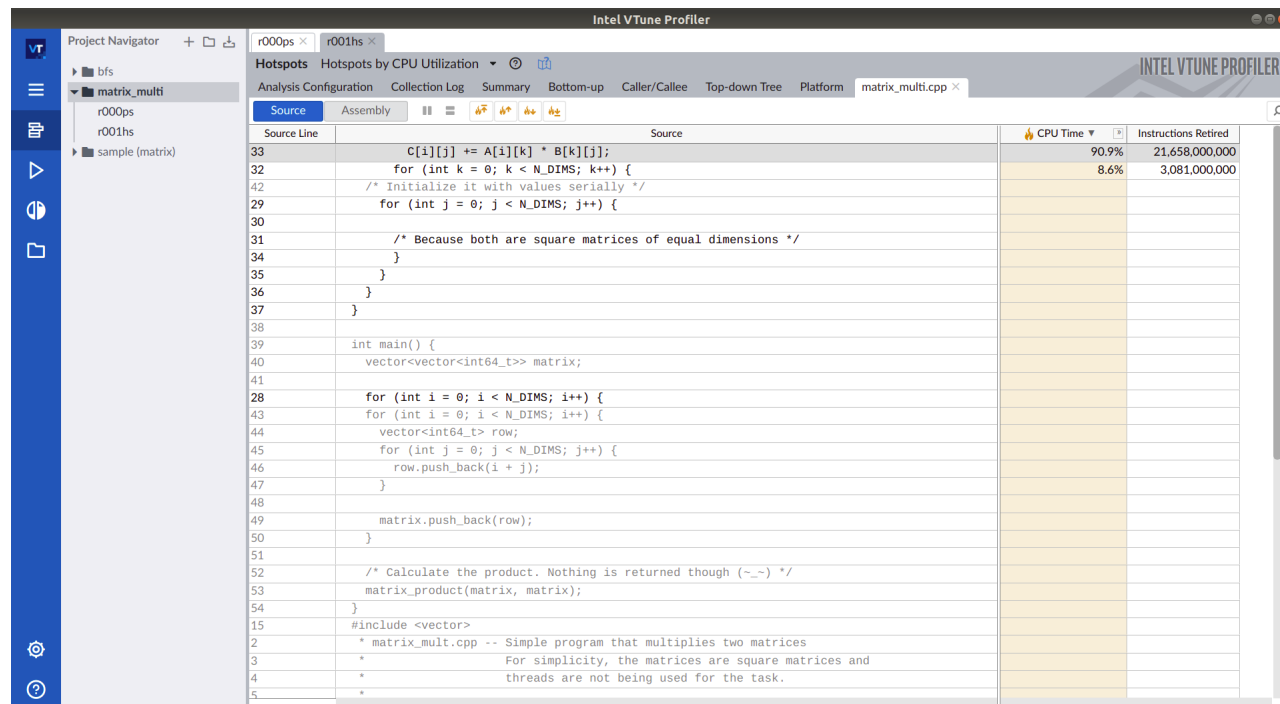


Figure 1.6: Top Source lines by CPU Utilization for `matrix_multi.cpp`

## Inference

We see that majority of the time goes in function `matrix_product`.

As discussed in lectures, `ijk` is not the optimal loop order for memory access. We see that the inner loop takes most of the time.

And we access and modify `k` at every iteration of the inner loop, it is the line to take second most CPU time.

### 1.3 matrix\_multi\_2.cpp

#### Performance Snapshot

- IPC: 1.339
- Logical Core Utilization: 8.2% (0.981 out of 12)
- Physical Core Utilization: 16.2% (0.973 of 6)
- Memory bound: 38.0% of Pipeline slots

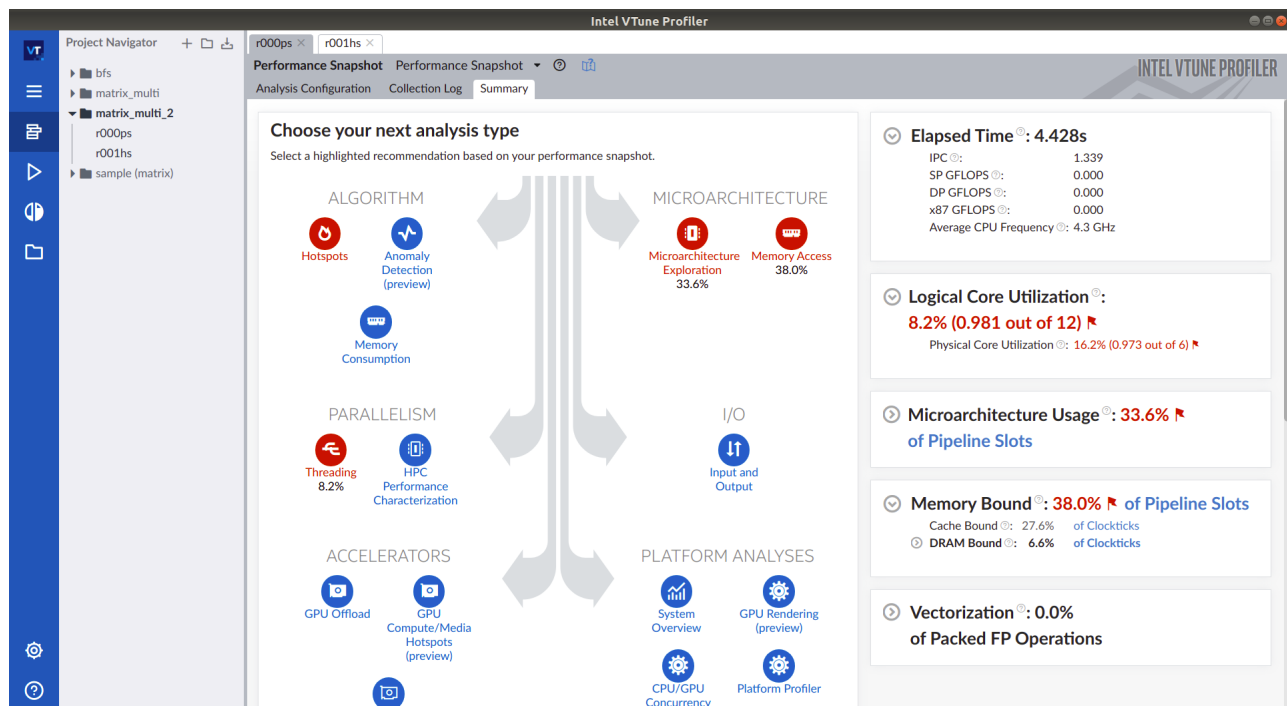


Figure 1.7: Performance Snapshot for `matrix_multi_2.cpp`

## Top Functions by CPU Time

Function	Module	CPU Time
matrix_product	matrix_multi_2.o	4.492s

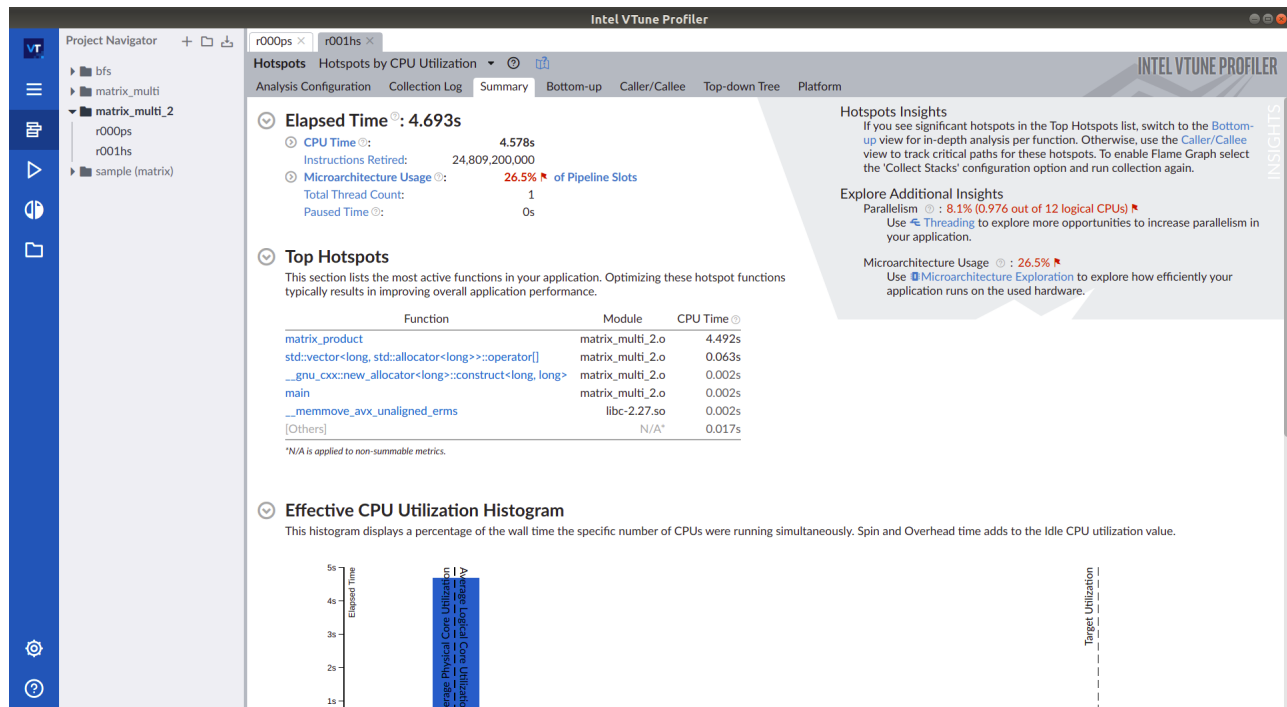


Figure 1.8: Top Functions by CPU Time for matrix\_multi\_2.cpp

## Top 2 Source lines by CPU Utilization

Source	Function	CPU Utilization
<code>C[i][j] += A[i][k] * B[k][j];</code>	<code>void matrix_product()</code>	84.4%
<code>for (int k = 0; k &lt; N_DIMS; k++) {</code>	<code>void matrix_product()</code>	13.7%

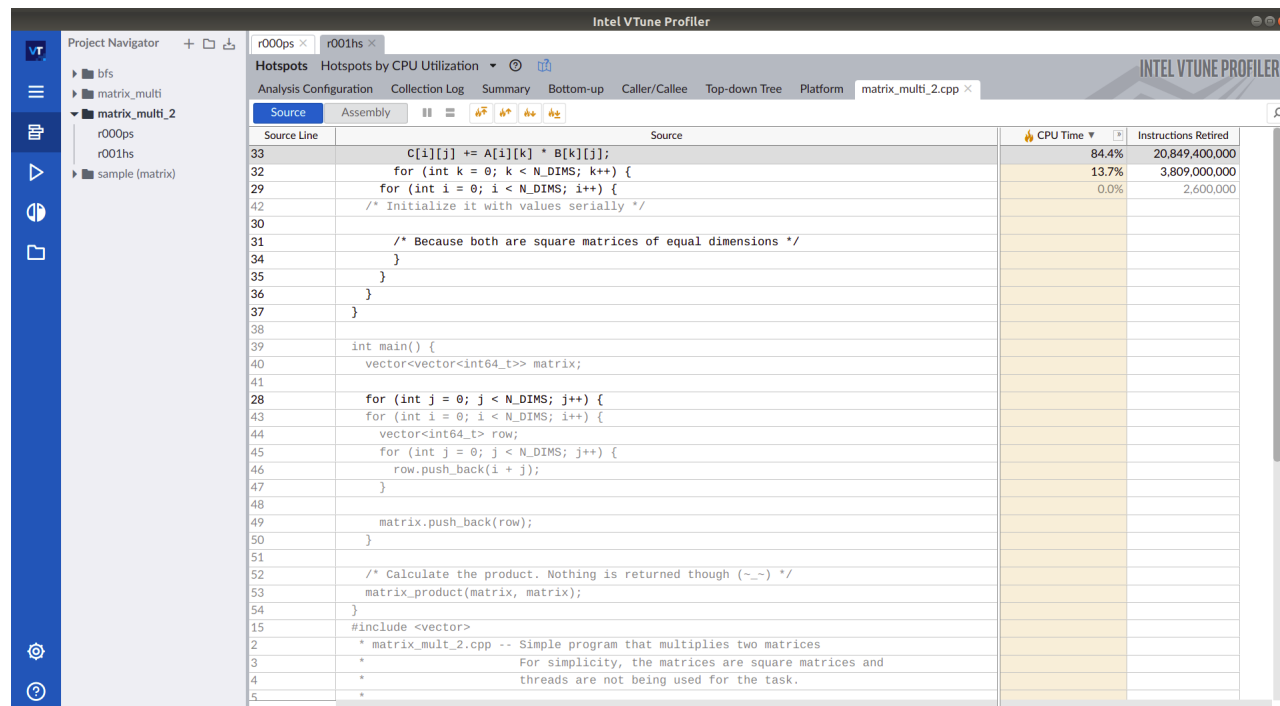


Figure 1.9: Top Source lines by CPU Utilization for `matrix_multi_2.cpp`

## Inference

We see that majority of the time goes in function `matrix_product`.

As discussed in lectures, `jik` (same as `ijk`) is not the optimal loop order for memory access. We see that the inner loop takes most of the time.

And we access and modify `k` at every iteration of the inner loop, it is the line to take second most CPU time.

## 1.4 quicksort.cpp

### Performance Snapshot

- IPC: 0.748
- Logical Core Utilization: 8.0% (0.966 out of 12)
- Physical Core Utilization: 15.7% (0.941 of 6)
- Memory bound: 23.0% of Pipeline slots

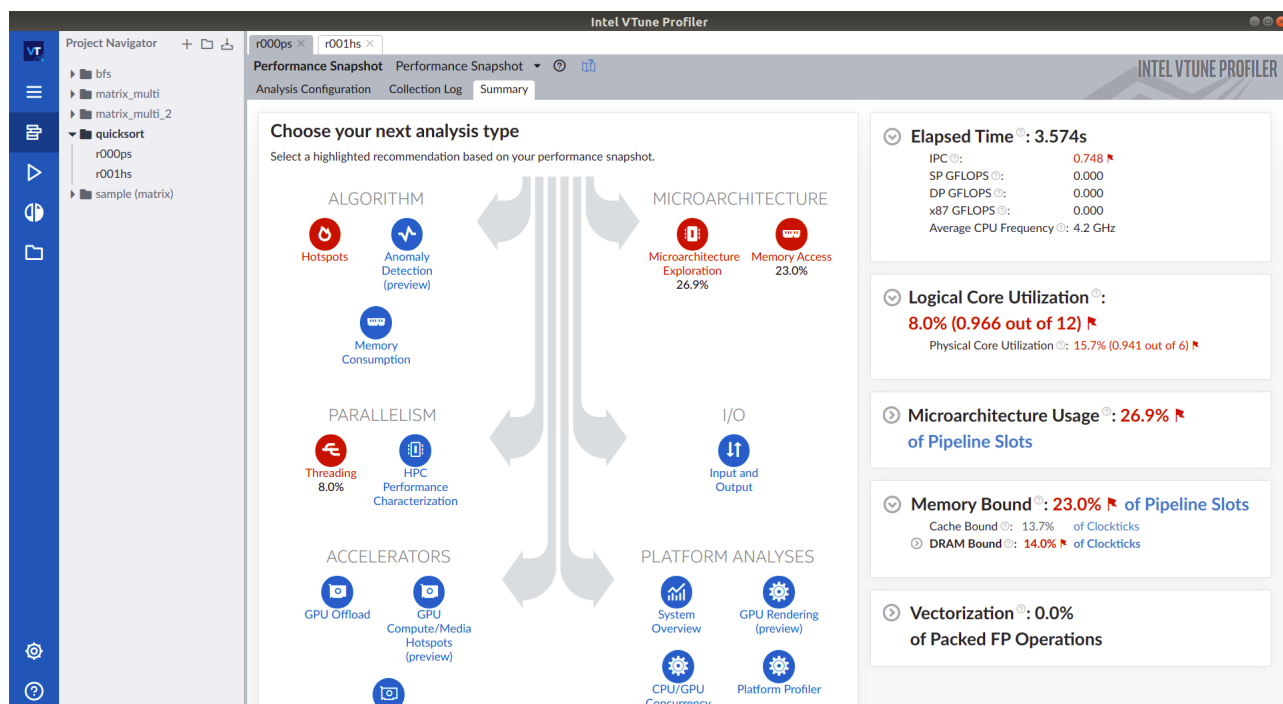


Figure 1.10: Performance Snapshot for quicksort.cpp

## Top Functions by CPU Time

Function	Module	CPU Time
__memmove_avx_unaligned_erms	libc-2.27.so	0.875s
page_fault	vmlinux	0.511s
clear_page_erms	vmlinux	0.228s
prepare_exit_to_usermode	vmlinux	0.226s
perf_iterate_ctx	vmlinux	0.155s
Others	N/A	1.545s

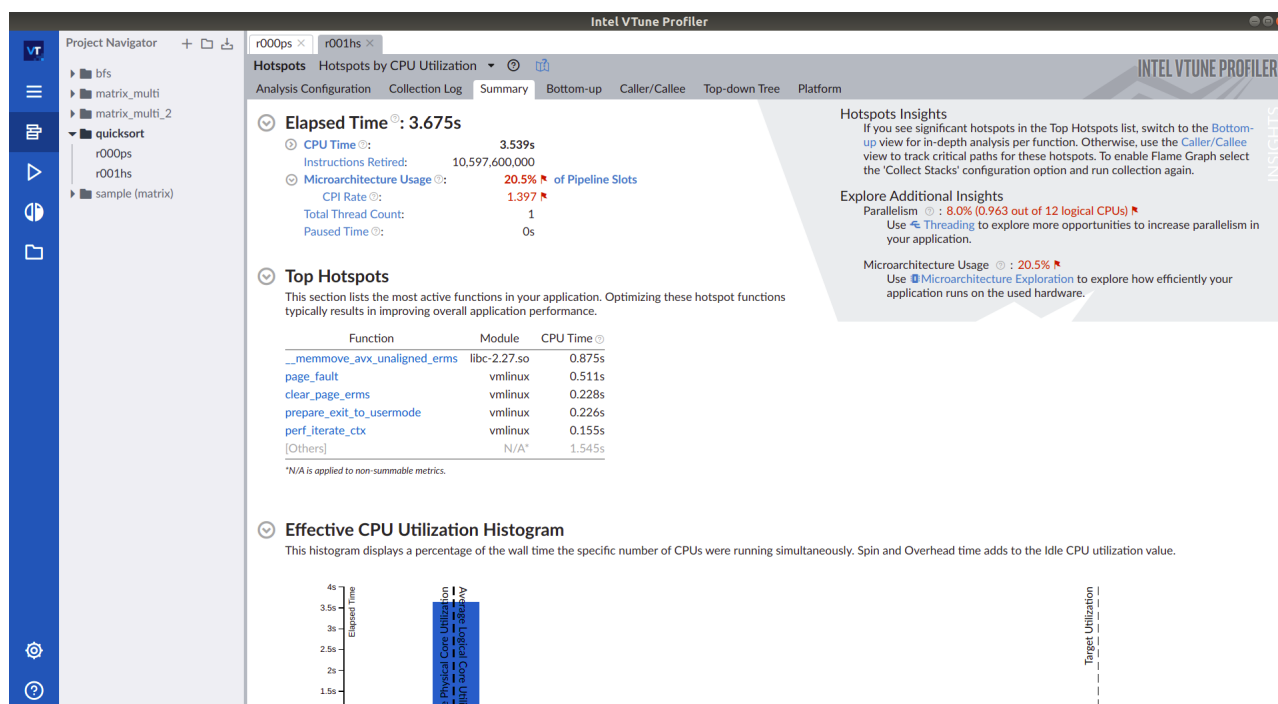


Figure 1.11: Top Functions by CPU Time for quicksort.cpp

## Top 5 Source lines by CPU Utilization

Source	Function	CPU Utilization
b = c;	void swap()	2.1%
if (nums[i] < pivot) {	long partition()	1.9%
slow_ptr++;	long partition()	1.7%
for (long i = lo; i < hi; i++) {	long partition()	0.6%
a = b;	void swap()	0.2%

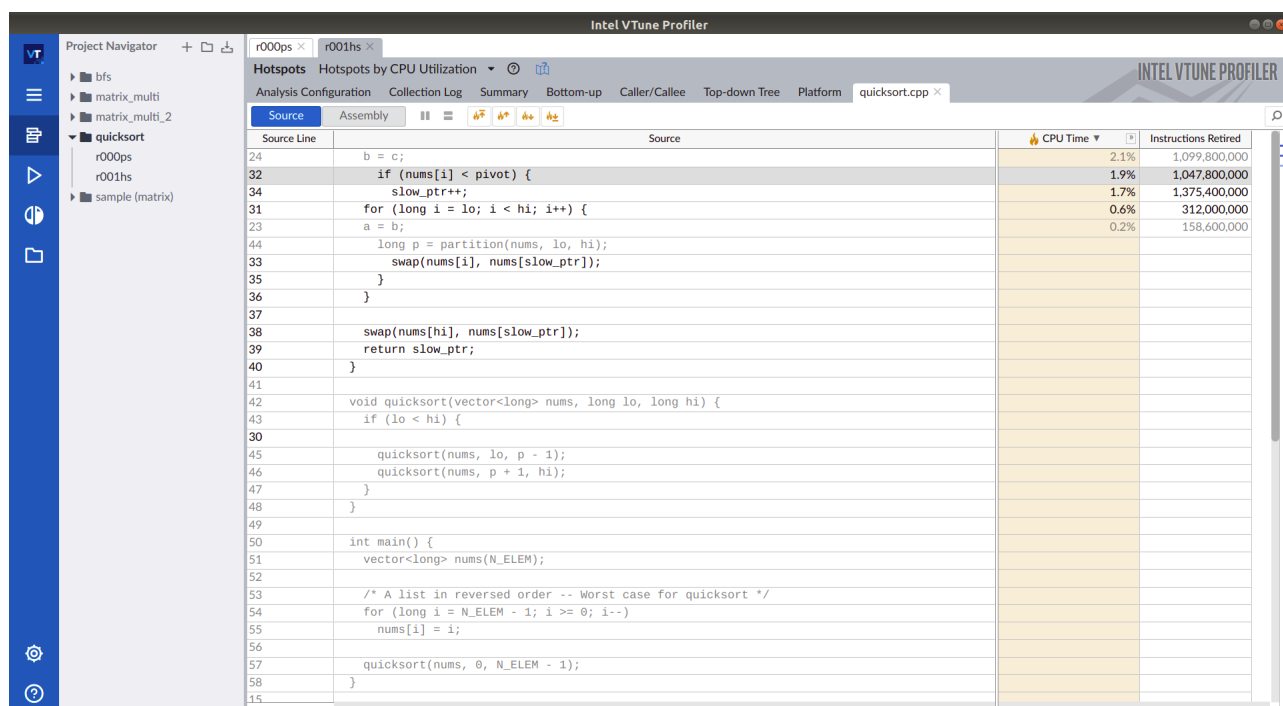


Figure 1.12: Top Source lines by CPU Utilization for quicksort.cpp

## Inference

We see that majority of the time goes in handling page faults and `memmove`. Possible explanation is that because it crosses my limit of RAM and overflows in swap memory, we might be getting page faults and page needs to be loaded back from the swap memory. (We discussed this in OS course)

The lines (present in `quicksort.cpp`) consuming the majority of time is mostly because of the number of times it is executed.

Quicksort algorithm is mostly partitioning and swapping, so those two functions take the majority of the time.



## Part 2: Simulating with ChampSim

### 2.1 Prepare traces

Generate tracer for champsim:

```
cd /champsim/ChampSim/tracer; ./make_tracer.sh;
```

Used pin to generate traces:

```
pin -t obj-intel64/champsim_tracer.so -o <program>.trace -t 21000000 - <executable>;
```

Used xz to compress the traces:

```
xz -vz <program>.trace -threads=0;
```

Program	Parameters	Execution time	Trace size
bfs.o	N_NODES (1«15); N_LOOPS 1000;	3.4 s	2092 KB
matrix_multi.o	N_DIMS 700;	4.5 s	2172 KB
matrix_multi_2.o	N_DIMS 700;	4.2 s	2160 KB
quicksort.o	N_ELEM (1«14);	3.1 s	4172 KB

### 2.2 Setup Configurations

To speed up the task and to avoid having to reset to default values again and again, I prepared 7 copies of ChampSim in the docker container.

This helped me run the 28 simulations in parallel and the entire experiment took 5-6 minutes to finish.

To setup each of the Configurations, I had to modify `./inc/cache.h`.

I have used CACTI to compute the latency updates for changes in cache size.

Theoretically, there would be change in latency when we change associativity as well but as it wasn't present in PS I have ignored it.

As hinted by professor, latency of 12-way cache is computed by taking mean of latency of 8-way and 16-way caches keeping sets as constant.

For L1I and L1D caches, the change in access time was small, so the latency remains same across the configurations.

Also, I have rounded off the latency to nearest integers to avoid issues with ChampSim.

(For example, I got L1I\_LATENCY to be 3.8 and 4.2 for half and double size cache respectively. I have consider both of them as 4 - same as baseline)

Updates in ./inc/cache.h

Line	Parameter	Baseline	Direct Mapped	Fully Associative	Reduced Size	Doubled Size	Reduced MSHR	Doubled MSHR
46	L1I_SET	64	64*8	1	32	128	64	64
47	L1I_WAY	8	1	8*64	8	8	8	8
51	L1I_MSHR_SIZE	8	8	8	8	8	4	16
52	L1I_LATENCY	4	4	4	4	4	4	4
55	L1D_SET	64	64*12	1	32	128	64	64
56	L1D_WAY	12	1	12*64	12	12	12	12
60	L1D_MSHR_SIZE	16	16	16	16	16	8	32
61	L1D_LATENCY	5	5	5	5	5	5	5
64	L2C_SET	1024	1024*8	1	512	2048	1024	1024
65	L2C_WAY	8	1	8*1024	8	8	8	8
69	L2C_MSHR_SIZE	32	32	32	32	32	16	64
70	L2C_LATENCY	10	10	10	9	13	10	10
73	LLC_SET	2048	2048*16	1	1024	4096	2048	2048
74	LLC_WAY	16	1	16*2048	16	16	16	16
78	LLC_MSHR_SIZE	64	64	64	64	64	32	128
79	LLC_LATENCY	20	20	20	17	24	20	20

## 2.3 Run simulations

Build champsim:

```
./build_champsim.sh bimodal no no no no lru 1;
```

Run simulation for traces:

```
./run_champsim.sh bimodal-no-no-no-no-lru-1core 10 10 bfs.trace.xz &;
./run_champsim.sh bimodal-no-no-no-no-lru-1core 10 10 matrix_multi.trace.xz &;
./run_champsim.sh bimodal-no-no-no-no-lru-1core 10 10 matrix_multi_2.trace.xz &;
./run_champsim.sh bimodal-no-no-no-no-lru-1core 10 10 quicksort.trace.xz &;
```

We can run all 28 simulations in the background and watch the processes using `watch -n 0.5 ps;`.

The entire experiment ends in around 10 minutes.

### 2.3.1 bfs.trace.xz

Parameter		Baseline	Direct Mapped	Fully Associative	Reduced Size	Doubled Size	Reduced MSHR	Doubled MSHR
cycles IPC		11505060 0.869183	11548164 0.865938	11510528 0.86877	11646705 0.858612	11410653 0.876374	11505205 0.869172	11504968 0.86919
Cache Misses	L1D	21942	30732	21942	21947	21931	21942	21972
	L1I	1	201163	1	9	1	1	1
	L2C	21464	23645	21727	21881	14648	21464	21464
	LLC	13437	14032	13437	14920	13436	13437	13437
MPKI	L1D	2.1942	3.0732	2.1942	2.1947	2.1931	2.1942	2.1972
	L1I	0.0001	20.1163	0.0001	0.0009	0.0001	0.0001	0.0001
	L2C	2.1464	2.3645	2.1727	2.1881	1.4648	2.1464	2.1464
	LLC	1.3437	1.4032	1.3437	1.4920	1.3436	1.3437	1.3437
Avg Miss Latency	L1D	91.8093	71.4266	91.6041	130.174	85.5957	91.7903	91.844
	L1I	215	14.1125	215	57.6667	229	215	215
	L2C	78.5156	73.5786	77.3577	116.525	107.193	78.4964	78.5508
	LLC	77.5087	79.0059	76.5858	132.767	76.5338	77.4779	77.5649

### 2.3.2 matrix\_multi.trace.xz

Parameter		Baseline	Direct Mapped	Fully Associative	Reduced Size	Doubled Size	Reduced MSHR	Doubled MSHR
cycles IPC		17624045 0.567407	17696897 0.565071	17624045 0.567407	17625492 0.56736	17624129 0.567404	17624452 0.567394	17623986 0.567409
Cache Misses	L1D	7392	20961	7393	9130	7381	7392	7392
	L1I	0	8396	0	174	0	0	0
	L2C	7269	7707	7269	7481	7268	7269	7269
	LLC	7268	7548	7268	7268	7268	7268	7268
MPKI	L1D	0.7392	2.0961	0.7393	0.9130	0.7381	0.7392	0.7392
	L1I	0.0000	0.8396	0.0000	0.0174	0.0000	0.0000	0.0000
	L2C	0.7269	0.7707	0.7269	0.7481	0.7268	0.7269	0.7269
	LLC	0.7268	0.7548	0.7268	0.7268	0.7268	0.7268	0.7268
Avg Miss Latency	L1D	117.04	52.756	117.026	101.301	130.94	116.558	118.607
	L1I	-	14.2587	-	20.1954	-	-	-
	L2C	103.766	103.328	103.766	106.711	114.666	103.277	105.36
	LLC	73.7643	76.0135	73.7643	83.4732	77.6564	73.275	75.3581

### 2.3.3 matrix\_multi\_2.trace.xz

Parameter		Baseline	Direct Mapped	Fully Associative	Reduced Size	Doubled Size	Reduced MSHR	Doubled MSHR
cycles IPC		17622240 0.567465	17694653 0.565143	17622240 0.567465	17623493 0.567424	17622347 0.567461	17622647 0.567452	17622268 0.567464
Cache Misses	L1D	7392	21323	7393	9140	7381	7392	7392
	L1I	0	6867	0	174	0	0	0
	L2C	7269	7857	7269	7483	7268	7269	7269
	LLC	7268	7465	7268	7268	7268	7268	7268
MPKI	L1D	0.7392	2.1323	0.7393	0.9140	0.7381	0.7392	0.7392
	L1I	0.0000	0.6867	0.0000	0.0174	0.0000	0.0000	0.0000
	L2C	0.7269	0.7857	0.7269	0.7483	0.7268	0.7269	0.7269
	LLC	0.7268	0.7645	0.7268	0.7268	0.7268	0.7268	0.7268
Avg Miss Latency	L1D	116.464	53.2717	116.329	101.704	130.203	116.168	117.712
	L1I	-	14.8274	-	19.4425	-	-	-
	L2C	103.181	105.285	103.057	107.273	113.948	102.88	104.45
	LLC	73.1791	81.0347	73.0553	84.0922	76.9375	72.8782	74.4483

### 2.3.4 quicksort.trace.xz

Parameter		Baseline	Direct Mapped	Fully Associative	Reduced Size	Doubled Size	Reduced MSHR	Doubled MSHR
cycles IPC		23473142 0.426019	32062866 0.311887	23473142 0.426019	23492281 0.425672	23585074 0.423997	23473142 0.426019	23473142 0.426019
Cache Misses	L1D	62136	1191142	62136	62136	62136	62136	62136
	L1I	0	102	0	0	0	0	0
	L2C	12300	16609	12300	14630	12297	12300	12300
	LLC	12296	18305	12293	12297	12290	12296	12296
MPKI	L1D	6.2136	119.1142	6.2136	6.2136	6.2136	6.2136	6.2136
	L1I	0.0000	0.0102	0.0000	0.0000	0.0000	0.0000	0.0000
	L2C	1.2300	1.6609	1.2300	1.4630	1.2297	1.2300	1.2300
	LLC	1.2296	1.8305	1.2293	1.2297	1.2290	1.2296	1.2296
Avg Miss Latency	L1D	37.1162	11.9622	37.2176	52.8647	40.495	37.1162	37.1162
	L1I	-	90.6176	-	-	-	-	-
	L2C	111.698	119.162	112.211	165.043	113.641	111.698	111.698
	LLC	81.725	81.2332	82.2574	165.422	76.6843	81.725	81.725

## 2.4 Results

### 2.4.1 IPC

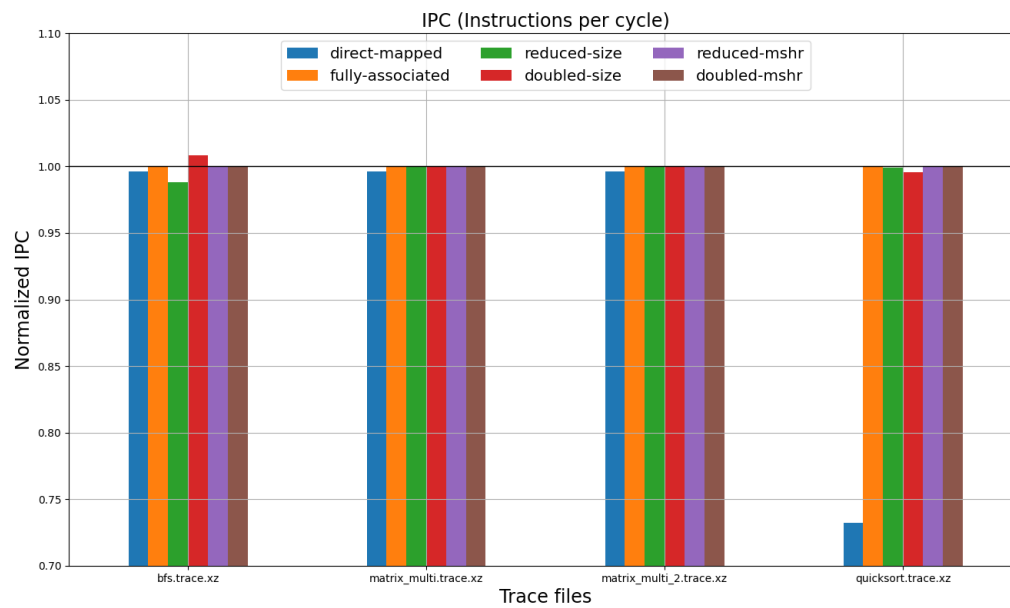


Figure 2.1: IPC (Instructions per cycle)

## 2.4.2 MPKI for L1 caches (L1D + L1I)

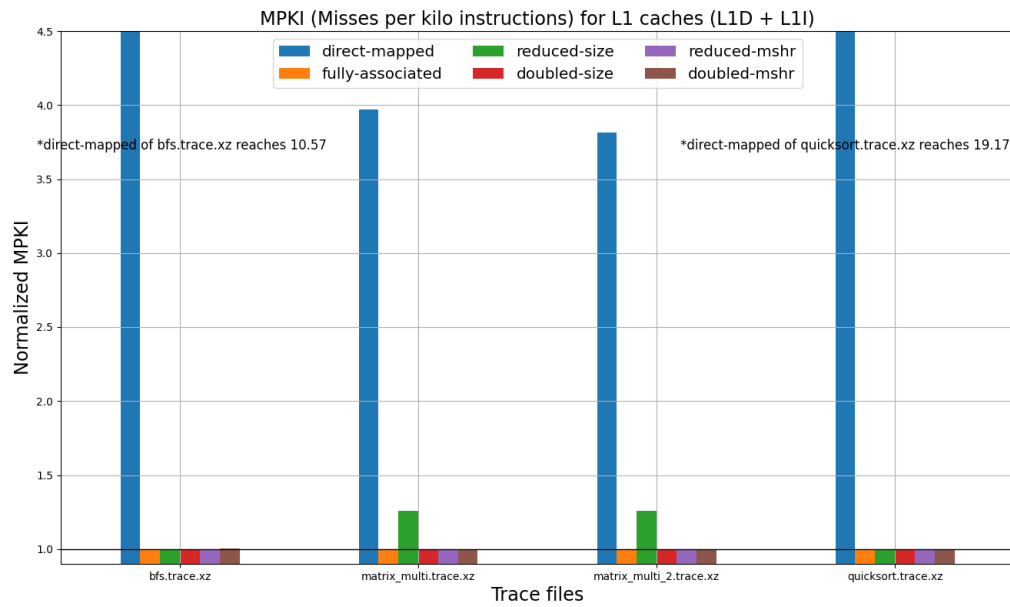


Figure 2.2: MPKI (Misses per kilo instructions) for L1 caches (L1D + L1I)

### 2.4.3 MPKI for L2 cache (L2C)

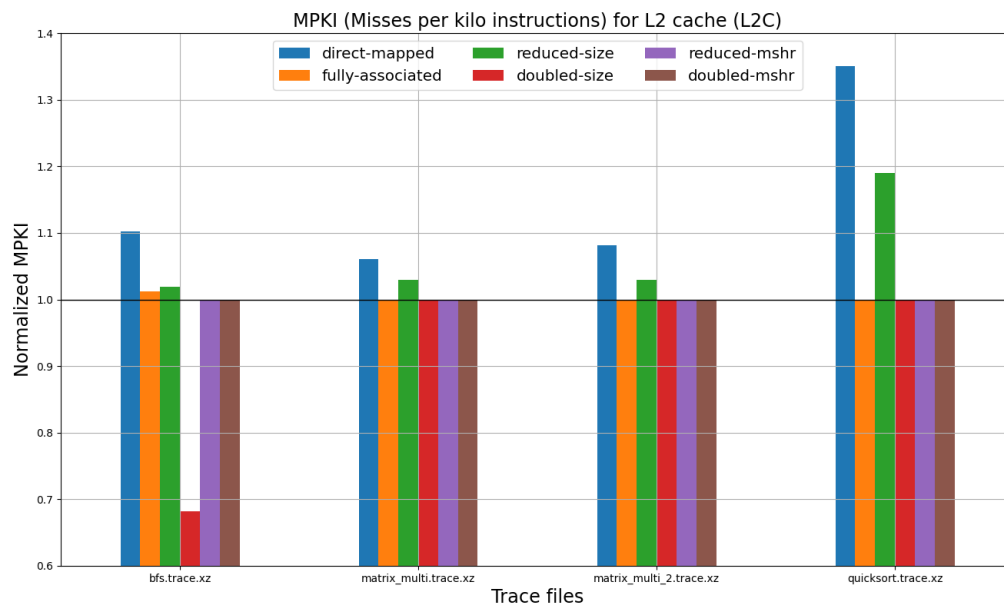


Figure 2.3: MPKI (Misses per kilo instructions) for L2 cache (L2C)

#### 2.4.4 MPKI for L3 cache (LLC)

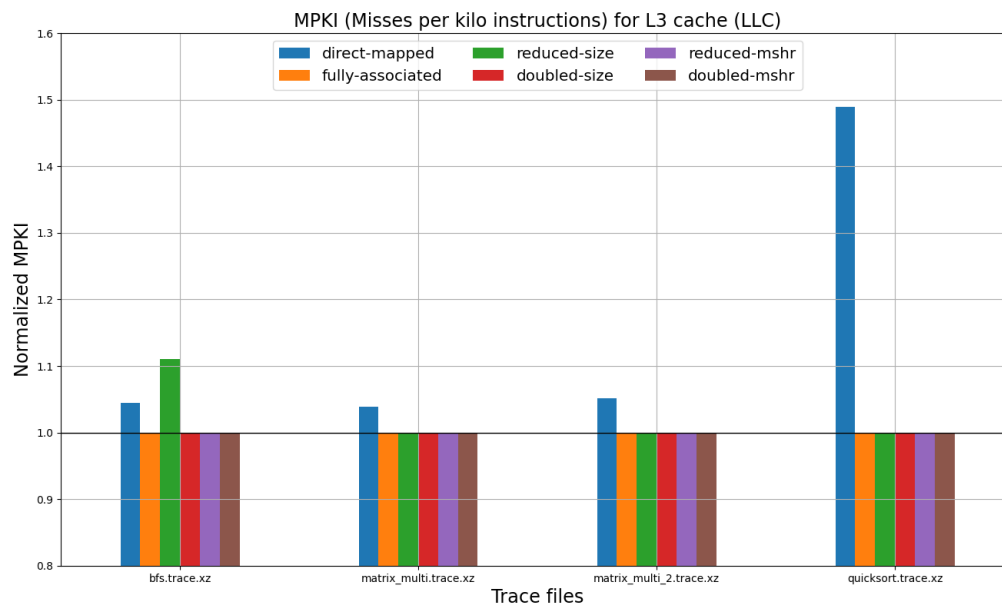


Figure 2.4: MPKI (Misses per kilo instructions) for L3 cache (LLC)

#### 2.4.5 Average Miss Latency