Assignment 4: CS 754

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Question:

We have studied two greedy algorithms for compressive recovery in class - MP and OMP.

Find out a research paper that proposes a greedy algorithm for CS recovery that is different from OMP and MP. Write down the algorithm in your report, state the key theorem and explain the meaning of the terms involved.

Answer:

There are various greedy algorithms that focus on solving the central problem in CS Recovery, The Basis Pursuit problem: $\min \|x\|_{\ell_1}$ subject to $y = \Phi x$. Here, we'll describe one such algorithm.

The Algorithm: Subspace Pursuit

Paper Title: Subspace Pursuit for Compressive Sensing Signal Reconstruction

Paper by: Wei Dai and Olgica Milenkovic, Department of Electrical and Computer Engineering, UIUC

Link to the paper: http://arXiv.org/abs/0803.0811v3

Introduction

The main contribution of this paper is a new algorithm, termed the subspace pursuit (SP) algorithm. It has provable reconstruction capability comparable to that of LP methods, and exhibits the low reconstruction complexity of matching pursuit techniques for very sparse signals. The algorithm can operate both in the noiseless and noisy regime, allowing for exact and approximate signal recovery, respectively. The basic idea behind the SP algorithm is borrowed from coding theory, more precisely, the A* order-statistic algorithm for additive white Gaussian noise channels.

Some Definitions

(Note: Φ^* denotes the transpose of the real valued matrix Φ)

1. Truncation and span:

Let
$$\mathbf{\Phi} \in \mathbb{R}^{m \times N}$$
, $\mathbf{x} \in \mathbb{R}^N$ and $I \subset \{1, 2, ..., N\}$

 Φ_I denotes the matrix consisting of the columns Φ of with indices $i \in I$.

 x_I is composed of the entries of x indexed by $i \in I$.

 $span(\Phi_I)$ denotes the space spanned by the columns of the matrix Φ_I .

2. Projection and Residue:

Let $\mathbf{y} \in \mathbb{R}^m$ and $\mathbf{\Phi} \in \mathbb{R}^{m \times n}$.

Suppose $\Phi^*\Phi$ is invertible.

Then, the projection of y onto $span(\Phi)$ is defined as:

$$oldsymbol{y}_{n}=proj(oldsymbol{y},oldsymbol{\Phi})=oldsymbol{\Phi}oldsymbol{\Phi}^{\dagger}oldsymbol{y}$$

Here, \dagger denotes psuedo-inverse: $\Phi^{\dagger} = (\Phi^*\Phi)^{-1}\Phi^*$

The residue vector of the projection is:

$$y_r = resid(y, \Phi) = y - y_p$$

The Psuedo-Code of the SP Algorithm

Input: K, Φ, y

Initialisation:

- 1. $T^0 = \{ K \text{ indices corresponding to the largest magnitude entries in the vector } \mathbf{\Phi}^* \mathbf{y} \}$
- 2. $\mathbf{y_r^0} = resid(\mathbf{y}, \mathbf{\Phi}_{T^0})$

Iteration: At the i^th iteration, do:

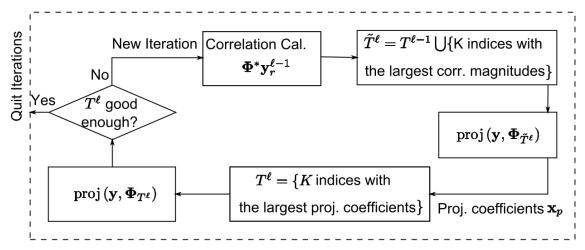
- 1. $\tilde{T}^l = T^{l-1} \cup \{ K \text{ indices corresponding to the largest magnitude entries in the vector } \Phi^* y_r^{i-1} \}$
- 2. Set $\boldsymbol{x}_p = \boldsymbol{\Phi}_{T^i}^{\dagger}$
- 3. $T^i = \{ K \text{ indices corresponding to the largest magnitude entries in the vector } \boldsymbol{x}_p \}$
- 4. $\mathbf{y_r^i} = resid(\mathbf{y}, \mathbf{\Phi}_{T^i})$
- 5. If $\|\boldsymbol{y_r^i}\|_2 > \boldsymbol{y_r^{i-1}}$, let $T^i = T^{i-1}$ and quit the iteration

Output:

Let $x_{T^i} = \Phi_{T^i}^{\dagger}$, and set all other entries to 0. Output this estimated signal x.

Flow Chart

The following Flow Chart can assist in better understanding of the algorithm:



(b) Iterations in the proposed Subspace Pursuit Algorithm: a list of K candidates, which is allowed to be updated during the iterations, is maintained.

It is directly taken from the paper, but it does help a great deal in understanding of the algorithm.

Performace Bound Theorems

Non-Noisy Case: Theorem 1

Let $x \in \mathbb{R}^k$ be a K-sparse signal, and let its corresponding measurement be $y = \Phi x \in \mathbb{R}^m$.

If the sampling matrix Φ satisfies the RIP with constant $\delta_{3K} < 0.165$, then the SP algorithm is guaranteed to **exactly** recover x from y via a finite number of iterations.

Since this is an exact recovery, "bounds" are not needed.

Noisy Case: Theorem 9

Let $x \in \mathbb{R}^k$ be a K-sparse signal, and let its corresponding measurement be $y = \Phi x + e \in \mathbb{R}^m$, where e denotes the noise vector.

Suppose that the sampling matrix satisfies the RIP with parameter $\delta_{3K} < 0.083$.

Then the reconstruction distortion of the SP algorithm satisfies

$$\|\boldsymbol{x} - \hat{\boldsymbol{x}}\|_{2} \le c_{K} \|\boldsymbol{e}\|_{2}$$

Here, x is the original signal, and \hat{x} is the estimated signal.

Also, c_K is a constant independent of \boldsymbol{x} , equal to $\frac{1 + \delta_{3K} + \delta_{3K}^2}{\delta_{3K}(1 - \delta_{3K})}$

Approximate Case: Theorem 9, Corollary 1

Let $x \in \mathbb{R}^k$ be an approximately K-sparse signal, and let its corresponding measurement be $y = \Phi x + e \in \mathbb{R}^m$, where e denotes the noise vector.

Suppose that the sampling matrix satisfies the RIP with parameter $\delta_{6K} < 0.083$.

Then the reconstruction distortion of the SP algorithm satisfies

$$\|\boldsymbol{x} - \hat{\boldsymbol{x}}\|_{2} \le c_{2K}(\|\boldsymbol{e}\|_{2} + c_{K}^{'}) \|\boldsymbol{x} - \boldsymbol{x}_{K}\|_{1}$$

Repeating notations are same. x_K is the the vector obtained from x by maintaining the K entries with largest magnitude and setting all other entries in the vector to zero.

$$c_K^{'}$$
 is a constant independent of \boldsymbol{x} , equal to $\sqrt{\frac{1+\delta_{6K}}{K}}$

Part (1)

Question:

A, a $m \times n$ matrix of rank greater than r, is known apriori.

We seek to minimize $J(\mathbf{Q})$:

$$J(\boldsymbol{Q}) = \|\boldsymbol{A} - \boldsymbol{Q}\|_F^2$$
, where \boldsymbol{Q} is a rank- r matrix, with $r < m, \ r < n$

Answer:

The arg min of J(Q) is $A_r = U \Sigma_r V^T$

where $U\Sigma V^T$ is the singular value decomposition of A

 Σ_r is same as the matrix Σ , with all but the top r Singular values in Σ taken as 0.

Justification:

A proof of the result above is here: https://link.springer.com/article/10.1007%2FBF02288367. Here, we provide another argument.

For a matrix M let $\sigma_i(M)$ denote the i^{th} largest singular value. WLOG assume $n \geq m$. Then, we have the following:

Lemma: For $m \times n$ matrices X, Y with $q \leq i, j \leq n$, we have

$$\sigma_{i+j-1}(X+Y) \le \sigma_i(X) + \sigma_j(Y)$$

It is one of Weyl's inequalities.

A proof can be found here: https://qchu.wordpress.com/2017/03/13/singular-value-decomposition/, under the section "Additive perturbation (Weyl)".

In this inequality, substitute X = A - B, Y = B, j = r + 1 to get:

$$\sigma_{i+r}(A) \le \sigma_i(A-B) + \sigma_{r+1}(B)$$

Thus, we have,

$$||A - B||_F^2 = \sum_{i=1}^n \sigma_i^2 (A - B) \ge \sum_{i=1}^{n-r} \sigma_i^2 (A - B)$$
$$\ge \sum_{i=r+1}^n \sigma_i^2 (A)$$

And equality is attained when we set $B = A_r$ as defined earlier.

An Intuitive Argument:

Suppose A has rank-p.

Let $U\Sigma V^T = \sum_{i=1}^p \sigma_i u_i v_i^T$ be the singular value decomposition of A.

For N, a $n \times d$ matrix: think of the rows of N as n points in d-dimensional space. The Frobenius norm of N is the square root of the sum of the squared distance of the points to the origin. We will use this interpretation in this problem.

Claim 4.1.1: Interpretation of A_r

The rows of A_r are the projections of the rows of A onto the subspace V_r spanned by the first r right singular vectors of A.

Proof: Let a be an arbitrary row vector.

Since the v_i are orthonormal, the projection of the vector a onto V_r is given by $\sum_{i=1}^r (a \cdot v_i) v_i^T$

Thus, the matrix whose rows are the projections of the rows of A onto V_r is given by $\sum_{i=1}^r Av_i v_i^T$

This last expression simplifies to
$$\sum_{i=1}^r A v_i v_i^T = \sum_{i=1}^r \sigma_i u_i v_i^T = A_r$$

Thus the claim is proved.

Claim 4.1.2: For any matrix B of rank at most r, $||A - A_r||_F^2 \le ||A - B||_F^2$

Let **B** minimize $||A - B||_F^2$ among all rank r or less matrices.

Let V be the space spanned by the rows of B. The dimension of V is at most k.

Since B minimizes $||A - B||_F^2$, it must be that each row of B is the projection of the corresponding row of A onto V, otherwise replacing the row of B with the projection of the corresponding row of A onto V does not change V and hence the rank of B but would reduce $||A - B||_F^2$

Since each row of B is the projection of the corresponding row of A, it follows that $||A - B||_F^2$ is the sum of squared distances of rows of A to V.

Since A_r minimizes the sum of squared distance of rows of A to any r-dimensional subspace, the claim follows.

Usage:

The best rank-r approximation has an interesting application to image compression.

This is based on the fact that, digital images are in fact, stored as matrices; and it is empirically observed that these are inherently low rank.

Thus, we store the image matrix as its low(enough) rank approximation instead.

Suppose \mathcal{I} is an image with values in $\mathbb{R}^{m \times n}$. Then the space needed to store \mathcal{I} is $\mathcal{O}(m * n)$. We store info about the rank-r approximation of \mathcal{I} instead, where $r \ll \min\{m, n\}$.

We perform the SVD of M, and only store the r largest singular values and vectors:

$$\mathcal{I}_{m imes n} pprox oldsymbol{U_r} oldsymbol{\Sigma_r} oldsymbol{V_r^T} = \mathcal{I}_r$$

We store U_r, Σ_r, V_r instead of \mathcal{I} . The cost is:

Cost of
$$U_r$$
 + Cost of Σ_r + Cost of V_r = Total cost $mk + k + nk = k(m + n + 1)$

This is one magnitude smaller than $\mathcal{O}(m*n)$ when $r \ll \min\{m, n\}$.

The above compression algorithm is otherwise arbitrary. The reason why it works with so promising results, is because of the particular solution to the optimization problem we solved before.

Part (2)

Question:

Matrices \boldsymbol{A} and \boldsymbol{B} are both known.

We seek to minimize $J(\mathbf{R})$ where \mathbf{R} is constrained to be orthonormal, and:

$$J(\mathbf{R}) = \|\mathbf{A} - \mathbf{R}\mathbf{B}\|_F^2$$
, where $\mathbf{A} \in \mathbb{R}^{n \times m}$, $\mathbf{B} \in \mathbb{R}^{n \times m}$, $\mathbf{R} \in \mathbb{R}^{n \times n}$, $m > n$

Answer:

The arg min of $J(\mathbf{R})$ with constraint $\mathbf{R}^T \mathbf{R} = \mathbf{I}$, is $\mathbf{R} = \mathbf{V} \mathbf{U}^T$ where $\mathbf{U} \mathbf{S} \mathbf{V}^T$ is the singular value decomposition of $\mathbf{B} \mathbf{A}^T$

Justification:

First, we simplify $J(\mathbf{R})$ as follows:

$$J(\mathbf{R}) = \|\mathbf{A} - \mathbf{R}\mathbf{B}\|_{F}^{2}$$

$$= trace((\mathbf{A} - \mathbf{R}\mathbf{B})^{T}(\mathbf{A} - \mathbf{R}\mathbf{B}))$$

$$= trace(\mathbf{A}^{T}\mathbf{A} + \mathbf{B}^{T}\mathbf{B} - \mathbf{A}^{T}\mathbf{R}\mathbf{B} - (\mathbf{A}^{T}\mathbf{R}\mathbf{B})^{T}) \quad \dots \text{(As } \mathbf{R}^{T}\mathbf{R} = \mathbf{I} \text{)}$$

$$= trace(\mathbf{A}^{T}\mathbf{A} + \mathbf{B}^{T}\mathbf{B}) - 2trace(\mathbf{A}^{T}\mathbf{R}\mathbf{B}) \quad \dots \text{(As } trace(\mathbf{K}^{T}) = trace(\mathbf{K}))$$

Now, $trace(\mathbf{A}^T \mathbf{A} + \mathbf{B}^T \mathbf{B})$ is a constant.

Thus to minimize $J(\mathbf{R})$ is same as to maximize $trace(\mathbf{A}^T\mathbf{R}\mathbf{B})$ under same constraint.

Further we'll make use of the property that trace(AB) = trace(BA)

Thus, firstly, we have $trace(\mathbf{A}^T\mathbf{R}\mathbf{B}) = trace(\mathbf{R}\mathbf{B}\mathbf{A}^T)$

Now suppose the singular value decomposition of BA^T is USV^T . Then,

$$trace(\boldsymbol{RBA^T}) = trace(\boldsymbol{RUSV^T}) = trace(\boldsymbol{V^TRUS}) = trace(\boldsymbol{XS}) = \sum_{i=1}^{n} X_{ii}S_{ii}$$

where, $X = V^T RU$ is an orthonormal matrix due to which $|X_{ii}| \leq 1$

The singular values S_{ii} are all non-negative, thus the maximum will be obtained when for all $i, X_{ii} = 1$, i.e. X = I.

Thus we need to have $V^T R U = I$, that is $R = V U^T$

Usage:

The given optimization problem is an important and fundamental problem in various fields, like Computer Vision, Computer Graphics, Medical Imaging (especially in a sub-branch called as 'statistical shape analysis') and many more.

In image processing specifically, the solution to above optimization problem is required in quite a few places, of which we'll explain one here.

We'll describe the instance where it is encountered in one of the algorithms in dictionary learning. We use it in the method known as "Union of Orthonormal bases"

We represent the signal as: $X = AS + \epsilon$, where $X \in \mathbb{R}^{d \times N}$ is the signal, $A \in \mathbb{R}^{d \times Md}$ is the over-complete dictionary that is a union of ortho-normal bases, $S \in \mathbb{R}^{Md \times N}$ is the sparse co-efficients vector.

A is a the row concatenation of ortho-normal bases A_i for $i \in \{1, 2, ..., M\}$

X is a known signal. Assuming we have fixed bases stored in A, the coefficients in X can be estimated using block coordinate descent (BCR).

Given the coefficients, we next want to update the dictionary.

For all m, we do the following:

- 1. Get the residual vector: $\boldsymbol{X_m} = \boldsymbol{X} \sum_{j \neq m} \boldsymbol{A_j} \boldsymbol{S_j}$
- 2. Solve for A_m as:

$$A_m = \operatorname{arg\,min}_A \|X_m - AS_m\|^2 \text{ s.t. } AA^T = A^TA = I$$

Here, $\boldsymbol{X_m} \in \mathbb{R}^{d \times N}, \, \boldsymbol{S_m} \in \mathbb{R}^{d \times N} \, \text{ and } \, \boldsymbol{A} \in \mathbb{R}^{d \times d}$

It is in this step where we use the minimization problem.

Part (a)

Question:

What is hyperspectral unmixing?

You may use an equation to support your answer with symbol meanings carefully explained.

Answer:

Meaning of the Terms

a. Spectral Imaging:

Spectral imaging is imaging that uses multiple bands across the electromagnetic spectrum. While an ordinary camera captures light across three wavelength bands in the visible spectrum, red, green, and blue (RGB), spectral imaging encompasses a wide variety of techniques that go beyond RGB. Spectral imaging may use the infrared, the visible spectrum, the ultraviolet, x-rays, or some combination of the above. It may include the acquisition of image data in visible and non-visible bands simultaneously, illumination from outside the visible range, or the use of optical filters to capture a specific spectral range. It is also possible to capture hundreds of wavelength bands for each pixel in an image.

b. Hyperspectral Imaging:

Hyperspectral imaging, like other spectral imaging, collects and processes information from across the electromagnetic spectrum.[1] The goal of hyperspectral imaging is to obtain the spectrum for each pixel in the image of a scene, with the purpose of finding objects, identifying materials, or detecting processes. In hyperspectral imaging, a complete spectrum or some spectral information (such as the Doppler shift or Zeeman splitting of a spectral line) is collected at every pixel in an image plane. A hyperspectral camera uses special hardware to capture hundreds of wavelength bands for each pixel, which can be interpreted as a complete spectrum. In other words, the camera has a high spectral resolution.

c. Hyperspectral Unmixing:

Hyperspectral unmixing is a procedure that decomposes the measured pixel spectrum of hyperspectral data into a collection of constituent spectral signatures (or end-members) and a set of corresponding fractional abundances.

Mathematical Formulation

We focus on a relatively simplistic but very representative model, known as the Linear Mixing Model(LMM).

Despite the fact that the LMM is not always true, especially under certain scenarios that exhibit strong non-linearity, it is generally recognized as an acceptable model for many real-world scenarios.

We assume a macroscopic mixing scale in which the incident light interacts with only one material before reflecting off.

Then, the LMM is described as follows:

Let $y_m[n]$ denote the hyper-spectral camera's measurement at spectral band m and at pixel n.

Suppose M is the number of spectral bands, and L is the number of pixels.

Consider the signal $y[n] = [y_1[n], y_2[n], \cdots, y_M[n]]^T \in \mathbb{R}^M$

The LMM is now given by:

$$y[n] = \sum_{i=1}^{n} a_i s_i[n] + v[n] = As[n] + v[n].$$

for $n = 1, 2, \ldots, L$, where:

- 1. Each $a_i \in \mathbb{R}^M$ for i = 1, ..., N is called an endmember signature vector which contains the spectral components of a specific material (indexed by i) in the scene
- 2. N is the number of endmembers, or materials, in the scene.
- 3. $A = [a_1 \dots a_N] \in \mathbb{R}^{M \times N}$ is called the endmember matrix.
- 4. $s_i[n]$ describes the contribution of material i at pixel n. $s[n] = [s_1[n] \dots s_N[n]] \in \mathbb{R}^N$ is called the abundance vector at pixel n.
- 5. $v[n] \in \mathbb{R}^M$ simply denotes the noise vector at pixel n.

Some characteristics of the formulation are:

- M is often large—typically more than 200
- The mixing process described by the LMM Equation is a consequence of limited spatial resolution of hyper-spectral cameras. Specifically, one pixel may not be spatially fine enough to contain one material only.
- By nature, the abundance vectors s[n] should satisfy, for every $n, s_i[n] \ge 0$ and $\sum_{i=1}^n s_i[n] = 1$

Part (b)

Question:

In equation 40 of the paper, explain how non-negative matrix factorization is used for hyperspectral unmixing.

Answer:

Treat each y[n] for n = 1, 2, ..., L as a column vector.

Then these are column concatened to form $Y = [y[1] \ y[2] \ \dots \ y[L]] \in \mathbb{R}^{M \times L}$.

A similar process is carried out on the RHS also, we can express the linear model as:

$$Y = AS + V$$

Here, $\boldsymbol{Y} \in \mathbb{R}^{M \times L}$ is as above, $\boldsymbol{A} \in \mathbb{R}^{M \times N}$ is the endmember matrix, $\boldsymbol{S} \in \mathbb{R}^{N \times L}$ is the abundance matrix, and $\boldsymbol{V} \in \mathbb{R}^{M \times L}$ is the noise matrix.

This formulation has special properties:

First, we note that since the end-member matrix \boldsymbol{A} contains the spectral components of specific materials, it is non-negative.

Also, the abundance matrix S has every element in each vector non-negative as explained in the previous part- thus this matrix is non-negative as well.

Thirdly, N, the number of endmembers, or materials, in the scene- is small, N < M, N < L.

Moreover the vectors s[n] are empirically (In geoscience and remote sensing, a tremendous amount of effort has been spent on measuring and recording spectral samples of many different materials) observed to be sparse vectors.

Furthermore, the optimization problem we solve here due to the nature of the LMM, is

$$\min_{\boldsymbol{A}\succeq \boldsymbol{0}, \boldsymbol{S}\succeq \boldsymbol{0}}\|\boldsymbol{Y}-\boldsymbol{A}\boldsymbol{S}\|_F^2$$

This formulation is, in a way, equivalent to the NMF problem, which was:

Minimize
$$E(W, H) = ||Y - WH||$$
 such that $W \succeq 0, H \succeq 0$

Suppose we obtain factors A and S by non-negative matrix factorization (NMF) of Y.

Then, we can use A as an estimate of the end-members; and S as the corresponding abundances.

This is what we wanted to estimate: to decompose the measured pixel spectrum of hyperspectral data into a collection of constituent spectral signatures (or end-members) and a set of corresponding fractional abundances.

As we have accomplished this, we have accomplished hyperspectral unmixing. Thus, non-negative matrix factorization can be used for hyperspectral unmixing in the way as described.

Part (c)

Question:

Explain the improvement to non-negative matrix factorization proposed in equation 41 of the paper. (You may explain any two forms each for g and h.)

Answer:

First of all, we impose all the constraints on S that are required:

Let $S = \{S \mid \forall n : s[n] \succeq \mathbf{0}; \mathbf{1}^T s[n] = 1\}$. It is clear that S must belong to S.

Let Q be the set : $Q = \{(A, S) | A \succeq 0, S \in S\}$ Our optimization is now over Q.

Now we investigate problems with NMF.

First, NMF is NP-hard in general. For this reason, optimization schemes we see in the current NMF-based blind HU developments are rather pragmatic.

Second, NMF may not guarantee solution uniqueness. This is a serious issue to the blind HU application, since it means that an NMF solution may not necessarily be the true endmembers and abundances, even in the noiseless case.

In blind HU, NMF is modified to fit the problem better. For this, we use regularizers on A and S.

A general formulation is:

$$\min_{O} \|\boldsymbol{Y} - \boldsymbol{A}\boldsymbol{S}\|_{F}^{2} + \lambda \cdot g(\boldsymbol{A}) + \mu \cdot h(\boldsymbol{S})$$

Here, g and h are regularizers, and $\lambda, \mu > 0$ are some constants. Some regularizers are:

1. Abundance regularizer:

In reality, a given spectral signature is usually composed of a limited number of materials in a hyperspectral scene, and hence the abundance regularization should be selected to be sparsity-prompting. We impose sparsity constraints on the ℓ_1 norm of the columns of S.

We modify the general equation as:

$$g(\mathbf{A}) = 0$$
 and $h(\mathbf{S}) = \sum_{i,j} |S_{i,j}|$

2. The $L_{1/2}$ regularizer:

The $L_{1/2}$ regularizer is an alternative to the L_1 counterpart. The $L_{1/2}$ regularizer is a sparsity-promoting function. Furthermore, the $L_{1/2}$ regularizer not only can provide sparse solutions close to those yielded when L_0 is used but is also computationally efficient.

We modify the general equation as:

$$g(\boldsymbol{A}) = 0$$
 and $h(\boldsymbol{S}) = \|S\|_{1/2} = \sum_{k,n=1}^{K,N} \boldsymbol{s}_n(k)^{1/2}$

3. Minimum Volume Regularizer:

Although we see many choices with the regularizers, the philosophies behind the choices follow a few core principles. For the endmember regularizer g, the principle can be traced back to Minimum Volume in Convex Geometry. A classical example is minimum volume constrained NMF. Here, we modify the general equation as:

$$g(\mathbf{A}) = (vol(B))^2$$
 and $h(\mathbf{S}) = 0$

Here, vol(B) is the simplex volume corresponding to A. (In geometry, a simplex is a generalization of the notion of a triangle or tetrahedron to arbitrary dimensions.)