Foundation Session - 1

Agenda

- 1. Asymptotic Analysis
- 2. Rules to follow
- 3. Big-O
- 4. Big-Theta
- 5. Big-Omega
- 6. Best-Case v/s Average-Case v/s Worst-Case Complexity
- 7. Auxiliary Space v/s Space Complexity
- 8. String Sorting Analysis
- 9. Recursive Runtimes

2 rules:

- 1. Drop-off the constants
- 2. Drop-off the irrelevant terms

```
arr => matrix of strings
For (i = 0; i < n; i++) {
      For (int j = 0; j < m; j++) {
             Cout << arr[i][j] << endl;
      }
}
TC => Not of O(n*m) => O(n*m*max_len_of_string)
Aux Space \Rightarrow O(1)
SC =>
Printing an Integer \Rightarrow O(1)
32-bit integer => [0, 0, 1, 1 .....]
Printing a String => O(length of String)
Question:
You have an array of Strings.
Array size = n
Length of all the strings = m
TODO: Sort each of the strings in the array individually and then sort the
complete array.
Input = ["abd", "dfg", "bca", "ddb"]
Step-1: ["abc", "dfg", "abc", "bdd"]
Step-2: ["abc", "abc", "bdd", "dfg"]
```

Prerequisites:

- Sorting an array of integer => O(n*log(n))
- Comparing two strings => O(m)

Solution:

Part-1: Sort each string of size m

- Sorting 1 string = O(m * log(m))
- Sorting n strings = O(n*m*log(m))

Part-2: Sort the complete array lexicographically

- Sorting an integer array of size n = O(n*log(n))
- Time taken to compare 2 ints = O(1)
- Sorting does O(n*log(n)) number of comparisons
- Time taken to compare 2 strings = O(m)
- Number of comparisons = O(n*log(n))
- Sorting an string array of size n = O(m*n*log(n))

```
Overall TC = O(n*m*log(m)) + O(m*n*log(n)) = O(n*m*(log(m) + log(n)))
```

Recursive Analysis:

```
Int fib(int n) {
            If (n <= 1) return n;
            Return fib(n-1) + fib(n-2);
}

TC: O(2^n)
AS: O(n)

Int fun(int n) {
            If (n <= 1) return n;
            For (int i = 0; i < n; i++) { ... do something O(1) ...}</pre>
```