

Homework 3

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Due: 2024/3/24 10:59pm

1. Please reinterpret the following story from the Bayesian perspective.

狼来了：从前有个放羊娃，每天都把羊群带到山上去吃草。山里有狼出没。第一天，放羊娃觉得无聊，想要作弄山下耕作的村民。他朝着山下大喊“狼来了！狼来了”，村民们信以为真，冲上山来准备帮助他，发现被欺骗了，大家很生气。第二天，放羊娃故技重施，村民们虽然有点迟疑，但还是冲上山来准备打狼，结果又一次发现被欺骗了，大家非常生气。第三天，狼真的来了，此时放羊娃慌了，哭着向山下大喊“狼来了！狼来了！”，请求村民的帮助。但这一次村民们认为他又在撒谎，无人相信他。最后他所有的羊都被狼吃掉了。



2. A fair die is rolled repeatedly, and a running total is kept (which is, at each time, the total of all the rolls up until that time). Let p_n be the probability that the running total is ever exactly n (assume the die will always be rolled enough times so that the running total will eventually exceed n , but it may or may not ever equal n).
- Write down a recursive equation for p_n (relating p_n to earlier terms p_k in a simple way). Your equation should be true for all positive integers n , so give a definition of p_0 and p_k for $k < 0$ so that the recursive equation is true for small values of n .
 - Find p_7 .
 - Give an intuitive explanation for the fact that $p_n \rightarrow 1/3.5 = 2/7$ as $n \rightarrow \infty$.
3. A sequence of $n \geq 1$ independent trials is performed, where each trial ends in “success” or “failure” (but not both). Let p_i be the probability of success in the i^{th} trial, $q_i = 1 - p_i$, and $b_i = q_i - 1/2$, for $i = 1, 2, \dots, n$. Let A_n be the event that the number of successful trials is even.
- Show that for $n = 2$, $P(A_2) = 1/2 + 2b_1b_2$.

(b) Show by induction that

$$P(A_n) = 1/2 + 2^{n-1}b_1b_2 \dots b_n$$

(This result is very useful in cryptography. Also, note that it implies that if n coins are flipped, then the probability of an even number of Heads is $1/2$ if and only if at least one of the coins is fair.) *Hint*: Group some trials into a super-trial.

(c) Check directly that the result of (b) is true in the following simple cases: $p_i = 1/2$ for some i ; $p_i = 0$ for all i ; $p_i = 1$ for all i .

4. A message is sent over a noisy channel. The message is a sequence x_1, x_2, \dots, x_n of n bits ($x_i \in \{0, 1\}$). Since the channel is noisy, there is a chance that any bit might be corrupted, resulting in an error (a 0 becomes a 1 or vice versa). Assume that the error events are independent. Let p be the probability that an individual bit has an error ($0 < p < 1/2$). Let y_1, y_2, \dots, y_n be the received message (so $y_i = x_i$ if there is no error in that bit, but $y_i = 1 - x_i$ if there is an error there).

To help detect errors, the n th bit is reserved for a parity check: x_n is defined to be 0 if $x_1 + x_2 + \dots + x_{n-1}$ is even, and 1 if $x_1 + x_2 + \dots + x_{n-1}$ is odd. When the message is received, the recipient checks whether y_n has the same parity as $y_1 + y_2 + \dots + y_{n-1}$. If the parity is wrong, the recipient knows that at least one error occurred; otherwise, the recipient assumes that there were no errors.

- (a) For $n = 5, p = 0.1$, what is the probability that the received message has errors which go undetected?
 - (b) For general n and p , write down an expression (as a sum) for the probability that the received message has errors which go undetected.
 - (c) Give a simplified expression, not involving a sum of a large number of terms, for the probability that the received message has errors which go undetected.
5. For x and y binary digits (0 or 1), let $x \oplus y$ be 0 if $x = y$ and 1 if $x \neq y$ (this operation is called exclusive or (often abbreviated to XOR), or addition mod 2).
- (a) Let $X \sim \text{Bern}(p)$ and $Y \sim \text{Bern}(1/2)$, independently. What is the distribution of $X \oplus Y$?
 - (b) With notation as in sub-problem (a), is $X \oplus Y$ independent of X ? Is $X \oplus Y$ independent of Y ? Be sure to consider both the case $p = 1/2$ and the case $p \neq 1/2$.
 - (c) Let X_1, \dots, X_n be i.i.d. (i.e., independent and identically distributed) $\text{Bern}(1/2)$ R.V.s. For each nonempty subset J of $\{1, 2, \dots, n\}$, let

$$Y_J = \bigoplus_{j \in J} X_j.$$

Show that $Y_J \sim \text{Bern}(1/2)$ and that these $2^n - 1$ R.V.s are pairwise independent, but not independent.

6. **(Optional Challenging Problem)** By LOTP for problems with recursive structure, we generate many difference equations.

(a) Solve the following difference equation:

$$p \cdot f_{i+1} - f_i + q \cdot f_{i-1} = -1, 1 \leq i \leq N - 1$$

where $0 < p < 1$, $q = 1 - p$, N is a constant, $f_0 = 0$, $f_N = 0$.

(b) Solve the following difference equation:

$$f_{i+1} = b \cdot f_i + a \cdot f_{i-1} + h, i \geq 1.$$

where h is a constant.

(c) Solve the following difference equation:

$$f_{i+1} = b \cdot f_i + a \cdot f_{i-1} + g(i), i \geq 1.$$

where $g(i)$ is a function of i .