

# Sorting Techniques

Sorting is one of the fundamental operations in computer science.

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## 1. Bubble Sort

**Overview:** Repeatedly compares adjacent elements and swaps them if they are in the wrong order. It "bubbles" the largest element to the end of the list.

- **Time Complexity:**
  - Best case:  $O(n)$  (if already sorted).
  - Worst case:  $O(n^2)$ .
- **Space Complexity:**  $O(1)$  (in-place).
- **Best for:** Educational purposes or extremely small datasets.
- **Why:**
  - Easy to implement and understand.
- **Limitations:**
  - Time complexity:  $O(n^2)$  for all cases.
- Highly inefficient for large datasets.

## Python Implementation

```
def bubble_sort(arr):  
    n = len(arr)  
    for i in range(n):  
        swapped = False  
        for j in range(0, n - i - 1): # Last i elements are already sorted  
            if arr[j] > arr[j + 1]:  
                arr[j], arr[j + 1] = arr[j + 1], arr[j]  
                swapped = True  
        if not swapped: # Break if no swaps were made (array is sorted)  
            break  
    return arr
```

```
print(bubble_sort([5, 3, 8, 6, 2]))
```

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## 2. Selection Sort

**Overview:** Repeatedly finds the smallest (or largest) element in the unsorted part of the list and places it in its correct position.

- **Time Complexity:**  $O(n^2)$  (always, as comparisons don't depend on order).
- **Space Complexity:**  $O(1)$  (in-place).

### Python Implementation

```
def selection_sort(arr):  
    n = len(arr)  
    for i in range(n):  
        min_index = i  
        for j in range(i + 1, n):  
            if arr[j] < arr[min_index]:  
                min_index = j  
        arr[i], arr[min_index] = arr[min_index], arr[i]  
    return arr  
  
print(selection_sort([5, 3, 8, 6, 2]))
```

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## 3. Insertion Sort

**Overview:** Builds the sorted array one element at a time by picking elements and inserting them into their correct position.

- **Time Complexity:**
  - Best case:  $O(n)$  (if already sorted).
  - Worst case:  $O(n^2)$ .
- **Space Complexity:**  $O(1)$  (in-place).
- **Best for:** Small or nearly sorted datasets.
- **Why:**

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- Simple and efficient for small input sizes ( $O(n)$  for nearly sorted data).
- In-place and stable.
- **Limitations:**
  - Time complexity:  $O(n^2)$  for large, unsorted datasets.

## Python Implementation

```
def insertion_sort(arr):  
    for i in range(1, len(arr)):  
        key = arr[i]  
        j = i - 1  
        while j >= 0 and key < arr[j]:  
            arr[j + 1] = arr[j]  
            j -= 1  
        arr[j + 1] = key  
    return arr  
  
print(insertion_sort([5, 3, 8, 6, 2]))
```

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## 4. Merge Sort

**Overview:** A divide-and-conquer algorithm that splits the array into halves, sorts each half recursively, and then merges them.

- **Time Complexity:**  $O(n \log n)$  (always).
- **Space Complexity:**  $O(n)$  (due to temporary arrays).
- **Best for:** Large datasets requiring stable sorting or when memory is not a concern.
- **Why:**
  - Always  $O(n \log n)$  time complexity.
  - Stable (maintains relative order of equal elements).
- **Limitations:**
  - Requires extra memory ( $O(n)$ ) for temporary arrays.

## Python Implementation

```
def merge_sort(arr):  
    if len(arr) <= 1:  
        return arr  
  
    mid = len(arr) // 2  
    left = merge_sort(arr[:mid])  
    right = merge_sort(arr[mid:])  
  
    return merge(left, right)  
  
def merge(left, right):  
    result = []  
    i = j = 0  
    while i < len(left) and j < len(right):  
        if left[i] < right[j]:  
            result.append(left[i])  
            i += 1  
        else:  
            result.append(right[j])  
            j += 1  
    result.extend(left[i:])  
    result.extend(right[j:])  
    return result  
  
print(merge_sort([5, 3, 8, 6, 2]))
```

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## 5. Quick Sort

**Overview:** A divide-and-conquer algorithm that selects a "pivot" and partitions the array such that all smaller elements are on one side and larger elements on the other.

- **Time Complexity:**
  - Best/Average case:  $O(n \log n)$ .
  - Worst case:  $O(n^2)$  (if pivot is poorly chosen, e.g., smallest or largest element).
- **Space Complexity:**  $O(\log n)$  (due to recursion).
- **Best for:** Large datasets in general-purpose applications.
- **Why:**
  - Average time complexity:  $O(n \log n)$ .
  - In-place (requires little additional memory).
  - Efficient for random datasets.
- **Limitations:**
  - Worst-case time complexity:  $O(n^2)$  (can be mitigated with randomized pivot selection).

### Python Implementation

```
def quick_sort(arr):  
    if len(arr) <= 1:  
        return arr  
  
    pivot = arr[len(arr) // 2]  
    left = [x for x in arr if x < pivot]  
    middle = [x for x in arr if x == pivot]  
    right = [x for x in arr if x > pivot]  
  
    return quick_sort(left) + middle + quick_sort(right)  
  
print(quick_sort([5, 3, 8, 6, 2]))
```

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## 6. Heap Sort

**Overview:** Builds a binary heap and repeatedly extracts the maximum (or minimum) element.

- **Time Complexity:**  $O(n \log n)$  (always).
- **Space Complexity:**  $O(1)$  (in-place).
- **Best for:** Memory-constrained systems where in-place sorting is required.
- **Why:**
  - Time complexity:  $O(n \log n)$  (worst-case, best-case, and average-case).
  - In-place and does not require additional memory.
- **Limitations:**
  - Not stable (relative order of equal elements is not preserved).

Slower than Quick Sort in practice.

### Python Implementation

```
def heapify(arr, n, i):
```

```
    largest = i
```

```
    left = 2 * i + 1
```

```
    right = 2 * i + 2
```

```
    if left < n and arr[left] > arr[largest]:
```

```
        largest = left
```

```
    if right < n and arr[right] > arr[largest]:
```

```
        largest = right
```

```
    if largest != i:
```

```
        arr[i], arr[largest] = arr[largest], arr[i]
```

```
        heapify(arr, n, largest)
```

```
def heap_sort(arr):
```

```
    n = len(arr)
```

```
    for i in range(n // 2 - 1, -1, -1):
```

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```
heapify(arr, n, i)
```

```
for i in range(n - 1, 0, -1):
```

```
    arr[i], arr[0] = arr[0], arr[i]
```

```
    heapify(arr, i, 0)
```

```
return arr
```

```
print(heap_sort([5, 3, 8, 6, 2]))
```

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## 7. Counting Sort

- **Best for:** Sorting integers or datasets with a small range of values.
  - **Why:**
    - Time complexity:  $O(n+k)$ , where  $k$  is the range of input.
    - Very fast for small ranges.
  - **Limitations:**
    - Requires additional memory ( $O(k)$ ).
    - Only works for discrete keys (e.g., integers).
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## 8. Radix Sort

- **Best for:** Large datasets with integers or strings.
  - **Why:**
    - Time complexity:  $O(nk)$ , where  $k$  is the number of digits or characters.
    - Can be faster than  $O(n \log n)$  for specific datasets.
  - **Limitations:**
    - Requires extra memory for intermediate buckets.
    - Not comparison-based.
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### Comparison of Sorting Techniques

Algorithm	Time Complexity (Best)	Time Complexity (Worst)	Space Complexity	Stable?
Bubble Sort	$O(n)$	$O(n^2)$	$O(1)$	Yes
Selection Sort	$O(n^2)$	$O(n^2)$	$O(1)$	No
Insertion Sort	$O(n)$	$O(n^2)$	$O(1)$	Yes
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n)$	Yes
Quick Sort	$O(n \log n)$	$O(n^2)$	$O(\log n)$	No
Heap Sort	$O(n \log n)$	$O(n \log n)$	$O(1)O(1)$	No

The "best" sorting technique depends on the context, including factors like the size of the dataset, memory constraints, and whether the data is already partially sorted.

### Best Sorting Algorithm by Use Case

Use Case	Recommended Algorithm
Small datasets	Insertion Sort
Large, general-purpose datasets	Quick Sort
Stable sorting required	Merge Sort
Memory-constrained systems	Heap Sort
Integer datasets with small range	Counting Sort
Sorting integers/strings (large range)	Radix Sort
Nearly sorted datasets	Insertion Sort
Educating beginners	Bubble Sort