

## Syntax

### Types

$$\begin{aligned}
P &::= \top \mid p : \tau \langle \overline{x : \tau} \rangle \mid P \wedge P \\
\mathbb{T}(\mathbb{B}) &::= \alpha^{\mathbb{B}} \mid x : \mathbb{T}(\mathbb{B}) \rightarrow \mathbb{T}(\mathbb{B}) \mid \text{TyCon}^{\mathbb{B}} \overline{\mathbb{T}(\mathbb{B})} \mathbb{B} \mid \text{Class } \overline{\mathbb{T}(\mathbb{B})} \\
\mathbb{P}(\mathbb{B}) &::= \forall P. \mathbb{P}(\mathbb{B}) \mid \mathbb{T}(\mathbb{B}) \\
\mathbb{S}(\mathbb{B}) &::= \forall \alpha. \mathbb{S}(\mathbb{B}) \mid \mathbb{P}(\mathbb{B}) \\
T_P &::= \mathbb{S}(P) \\
\hat{T} &::= \mathbb{S}(Q) \\
T &::= \mathbb{S}(E) \\
\tau &::= \mathbb{S}(\top) \\
e &::= x \mid C \mid e \mid \lambda x. e \\
&\quad \mid e @ \mid \hat{\Lambda} p. e \mid e[\tau] \mid \Lambda \alpha. e \\
&\quad \mid \text{case } x \text{ of } \mid_i C_i \overline{x_i} \rightarrow e_i \\
&\quad \mid \text{let } x = e \text{ in } e \mid \text{letrec } x = e \text{ in } e
\end{aligned}$$

### Rules

$$\boxed{\Gamma \vdash_{P,Q} e : T}$$

$$\begin{aligned}
&\frac{\Gamma \vdash_{P,Q} e : T_1 \quad \Gamma \vdash T_1 <: T_2 \quad \Gamma \vdash T_2}{\Gamma \vdash_{P,Q} e : T_2} \\
&\frac{\Gamma(x) = \{v : T \mid e\}}{\Gamma \vdash_{P,Q} x : \{v : T \mid e \wedge v = x\}} \\
&\frac{}{\Gamma \vdash_{P,Q} C : ty(c)} \\
&\frac{\Gamma \vdash_{P,Q} e_1 : x : T_x \rightarrow T \quad \Gamma \vdash_{P,Q} e_2 : T}{\Gamma \vdash_{P,Q} e_1 e_2 : T[x \mapsto e_2]} \\
&\frac{\Gamma, x : \hat{T}_x \vdash_{P,Q} e : \hat{T} \quad \Gamma \vdash x : \hat{T}_x \rightarrow \hat{T}}{\Gamma \vdash_{P,Q} \lambda x. e : x : \hat{T}_x \rightarrow \hat{T}} \\
&\frac{\Gamma \vdash_{P,Q} e : \forall p. T \quad e_p = \text{fresh } (p) \quad \Gamma, p \vdash e_p}{\Gamma \vdash_{P,Q} e @ : T[p \mapsto e_p]}
\end{aligned}$$

define : freshP, subP  
do I need the following rule?

$$\begin{aligned}
&\frac{\Gamma \vdash_{P,Q} e : T}{\Gamma \vdash_{P,Q} \hat{\Lambda} p. e : \forall p. T} \\
&\frac{\Gamma \vdash_{P,Q} e : \forall \alpha. T \quad \hat{T} = \text{fresh } (\tau)}{\Gamma \vdash_{P,Q} e[\tau] : T[\alpha \mapsto \hat{T}]}
\end{aligned}$$

define subT, freshT

$$\frac{\Gamma \vdash_{P,Q} e : T \quad \alpha \notin \Gamma}{\Gamma \vdash_{P,Q} \Lambda \alpha. e : \forall \alpha. T}$$

$$\frac{\begin{array}{c} \Gamma \vdash_{P,Q} e : \mathbf{C}^{e_c} \bar{T} \bar{e} \\ \forall i. \{ \Gamma \vdash_{P,Q} C_i : \forall \bar{\alpha}. \forall \bar{p}. y_1 : T_1 \rightarrow \dots y_j : T_j \rightarrow \dots y_n : T_n \rightarrow \mathbf{C}^{e'_c} \bar{T}' \bar{e}' \\ \Gamma, x_{ij} : T_j [\bar{\alpha} \mapsto \bar{T}] [\bar{p} \mapsto \bar{e}] \vdash_{P,Q} e_i : T \} \end{array}}{\Gamma \vdash_{P,Q} \text{case } e \text{ of } \mid_i C_i \bar{x}_i \rightarrow e_i : T}$$

define unifyTypes

$$\frac{\Gamma \vdash_{P,Q} e_1 : T_x \quad \langle P(x), T_x \rangle \models T'_x \quad \Gamma, x : T'_x \vdash_{P,Q} e_2 : T}{\Gamma \vdash_{P,Q} \text{let } x = e_1 \text{ in } e_2 : T}$$

$$\frac{\Gamma \vdash \hat{T}_x \quad \Gamma \vdash_{P,Q} e_1 : \hat{T}_x \quad \langle P(x), \hat{T}_x \rangle \models \hat{T}'_x \quad \Gamma, x : \hat{T}'_x \vdash_{P,Q} e_2 : \hat{T}}{\Gamma \vdash_{P,Q} \text{letrec } x = e_1 \text{ in } e_2 : \hat{T}}$$

$$\boxed{\Gamma \vdash T_1 <: T_2}$$

$$\frac{[\Gamma] \wedge [e_1] \Rightarrow [e_2]}{\Gamma \vdash \alpha^{e_1} <: \alpha^{e_2}}$$

$$\frac{\Gamma \vdash T_{x_2} <: T_{x_1} \quad \Gamma, x_2 : T_{x_2} \vdash T_1 [x_1 \mapsto x_2] <: T_2}{\Gamma \vdash x_1 : T_{x_1} \rightarrow T_1 <: x_2 : T_{x_2} \rightarrow T_2}$$

$$\frac{\Gamma \vdash T_1 <: T_2}{\Gamma \vdash \forall \alpha. T_1 <: \forall \alpha. T_2}$$

$$\frac{[\Gamma] \wedge [e_1] \Rightarrow [e_2] \quad \forall i. \Gamma \vdash T_{1_i} <: T_{2_i} \quad \forall j. [\Gamma] \wedge [e_{1_j}] \Rightarrow [e_{2_j}]}{\Gamma \vdash \mathbf{C}^{e_1} \bar{T}_1 \bar{e}_1 <: \mathbf{C}^{e_2} \bar{T}_2 \bar{e}_2}$$

no subtype for tforallPr!

no subtype for tclass!

$$\boxed{\Gamma \vdash T}$$

$$\frac{\Gamma, v : \text{shape}(T) \vdash e : \text{bool}}{\Gamma \vdash \alpha^e}$$

$$\frac{\Gamma \vdash T_x \quad \Gamma, x : T_x \vdash T}{\Gamma \vdash x : T_x \rightarrow T}$$

$$\frac{\Gamma \vdash T}{\Gamma \vdash \forall \alpha. T}$$

define tyConP

$$\frac{\Gamma, v : \text{shape}(T) \vdash e : \text{bool} \quad \forall i. \Gamma \vdash T_i \quad \bar{p} = \text{predicates}(\mathbf{C}) \quad \forall j. \Gamma, p_j \vdash e_j}{\Gamma \vdash \mathbf{C}^e \bar{T} \bar{e}}$$

no subtype for tforallPr!

no subtype for tclass!

$$\boxed{\Gamma, p \vdash e}$$

$$\frac{\Gamma, v : \tau, \overline{x : \tau} \vdash e : \text{bool}}{\Gamma, p : \tau \langle \overline{x : \tau} \rangle \vdash e}$$