
CS 471/571 (Fall 2023):
Introduction to Artificial Intelligence

Lecture 4: Constraint Satisfaction Problems
(Part 1)

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Source: <http://ai.berkeley.edu/home.html>



Reminder

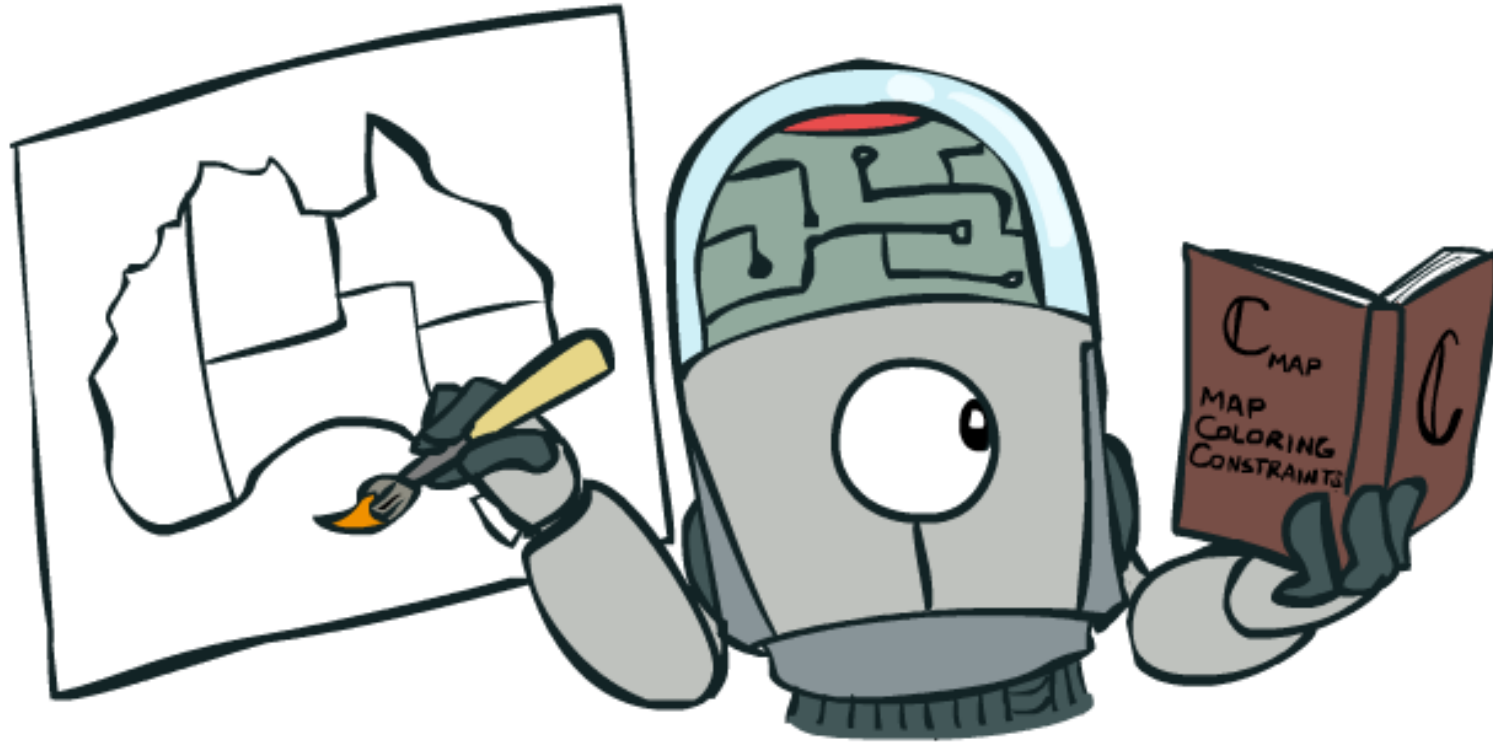
- Written assignment 1: Search
 - Deadline: Oct 11th, 2023
- Project 1: Search
 - Deadline: Oct 16th, 2023

Today

- Constraint Satisfaction Problems
 - Backtracking Search
 - Filtering
 - Ordering

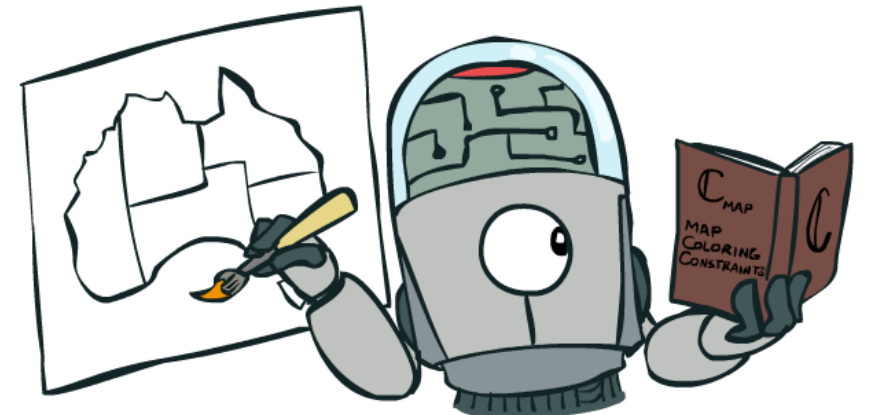
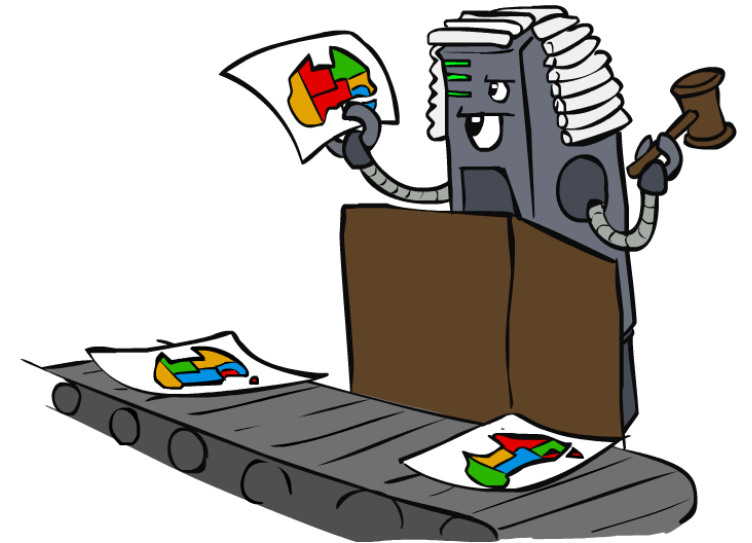


Constraint Satisfaction Problems

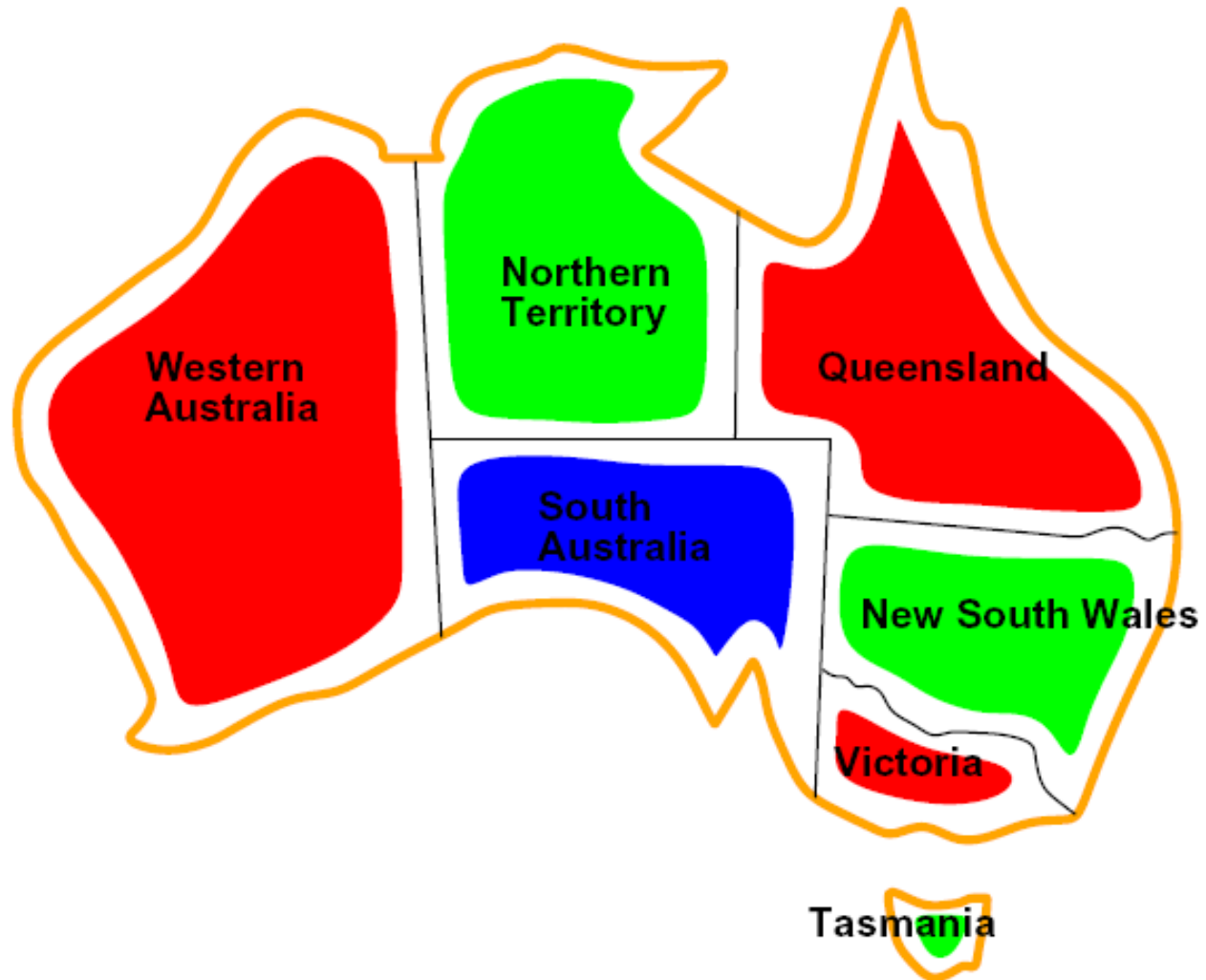


Constraint Satisfaction Problems

- Standard search problems:
 - State is a “black box”: arbitrary data structure
 - Goal test can be any function over states
 - Successor function can also be anything
- Constraint satisfaction problems (CSPs):
 - A special subset of search problems
 - State is defined by **variables X_i** with values from a **domain D** (sometimes D depends on i)
 - Goal test is a **set of constraints** specifying allowable combinations of values for subsets of variables
- Allows useful general-purpose algorithms with more power than standard search algorithms



CSP Examples



Example: Map Coloring

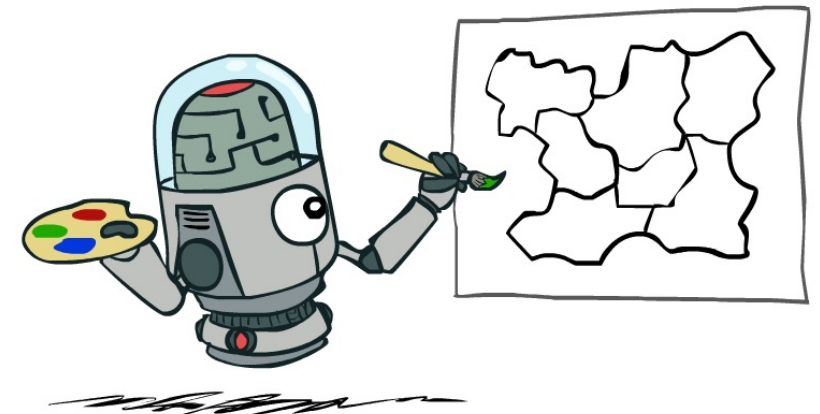
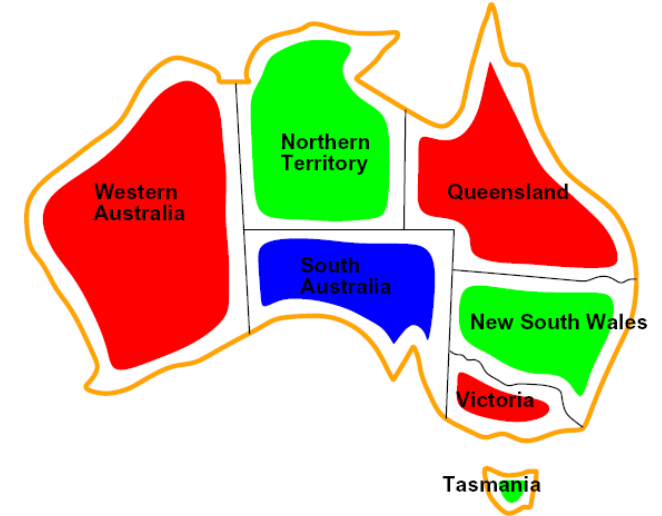
- Variables: WA, NT, Q, NSW, V, SA, T
- Domains: $D = \{\text{red, green, blue}\}$
- Constraints: adjacent regions must have different colors

Implicit: $WA \neq NT$

Explicit: $(WA, NT) \in \{(\text{red, green}), (\text{red, blue}), \dots\}$

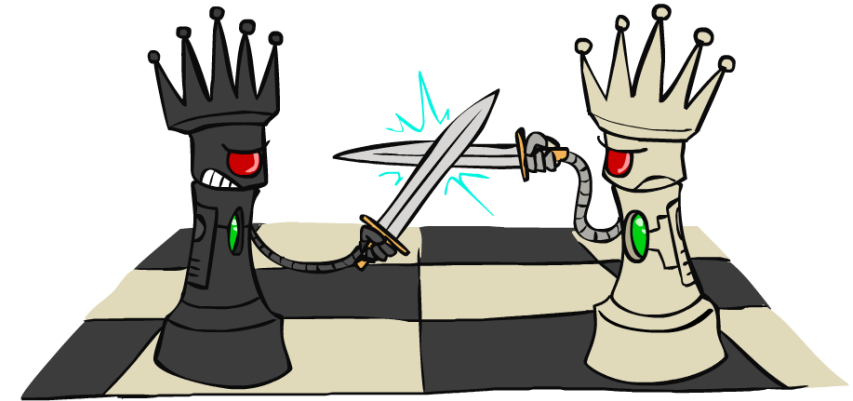
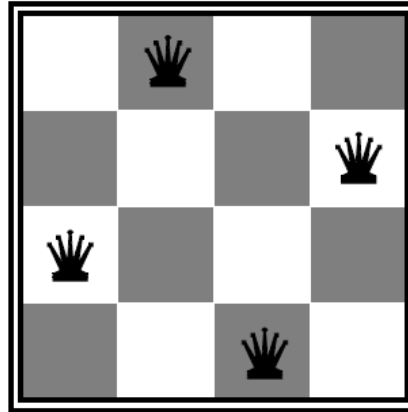
- Solutions are assignments satisfying all constraints, e.g.:

$\{WA=\text{red}, NT=\text{green}, Q=\text{red}, NSW=\text{green}, V=\text{red}, SA=\text{blue}, T=\text{green}\}$



Example: N-Queens

- Formulation 1:
 - Variables: X_{ij}
 - Domains: $\{0, 1\}$
 - Constraints



$$\forall i, j, k \quad (X_{ij}, X_{ik}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\forall i, j, k \quad (X_{ij}, X_{kj}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\forall i, j, k \quad (X_{ij}, X_{i+k, j+k}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\forall i, j, k \quad (X_{ij}, X_{i+k, j-k}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\sum_{i,j} X_{ij} = N$$



Example: N-Queens

- Formulation 2:

- Variables: Q_k

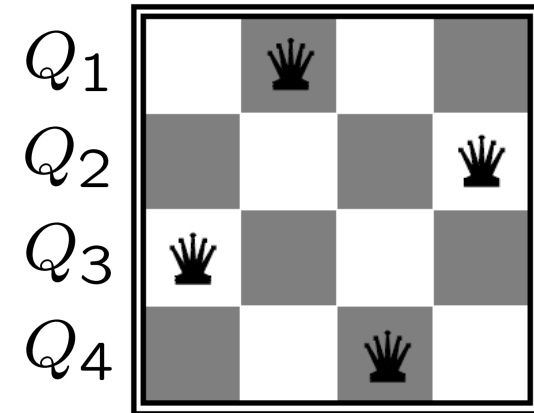
- Domains: $\{1, 2, 3, \dots, N\}$

- Constraints:

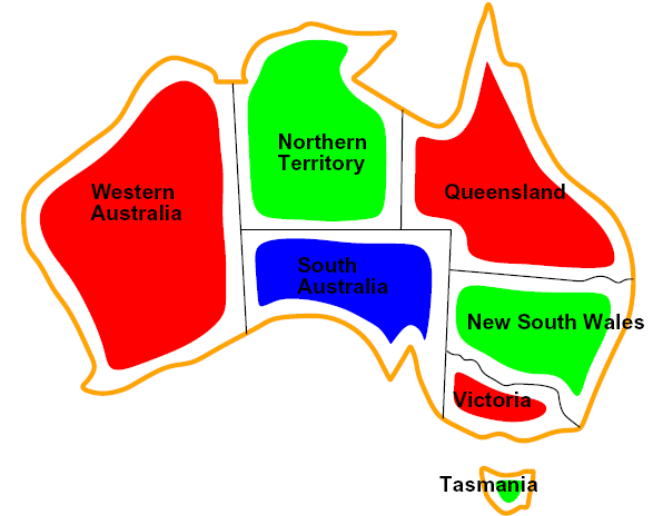
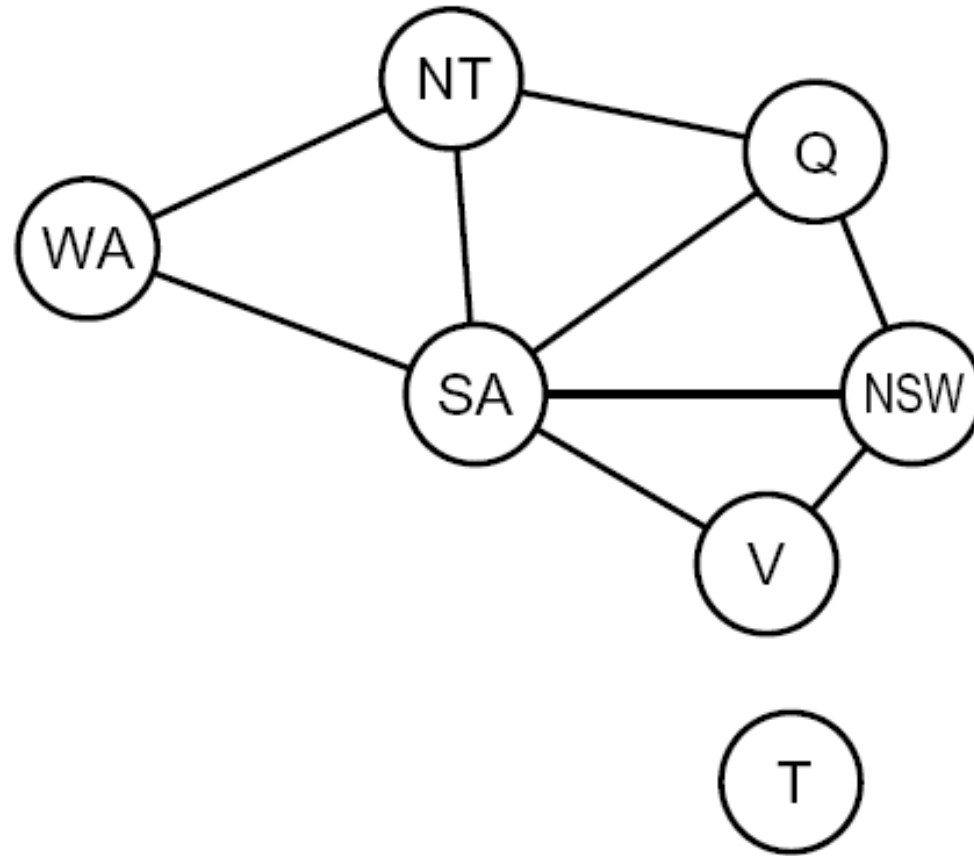
Implicit: $\forall i, j \text{ non-threatening}(Q_i, Q_j)$

Explicit: $(Q_1, Q_2) \in \{(1, 3), (1, 4), \dots\}$

...

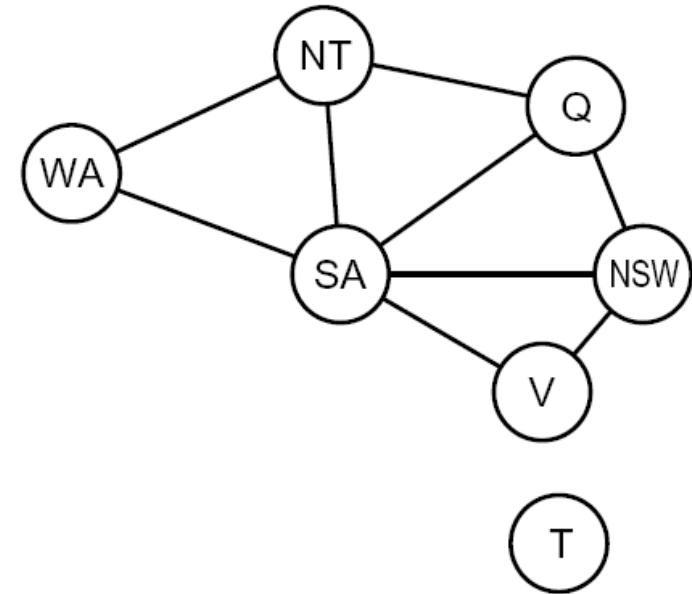


Constraint Graphs

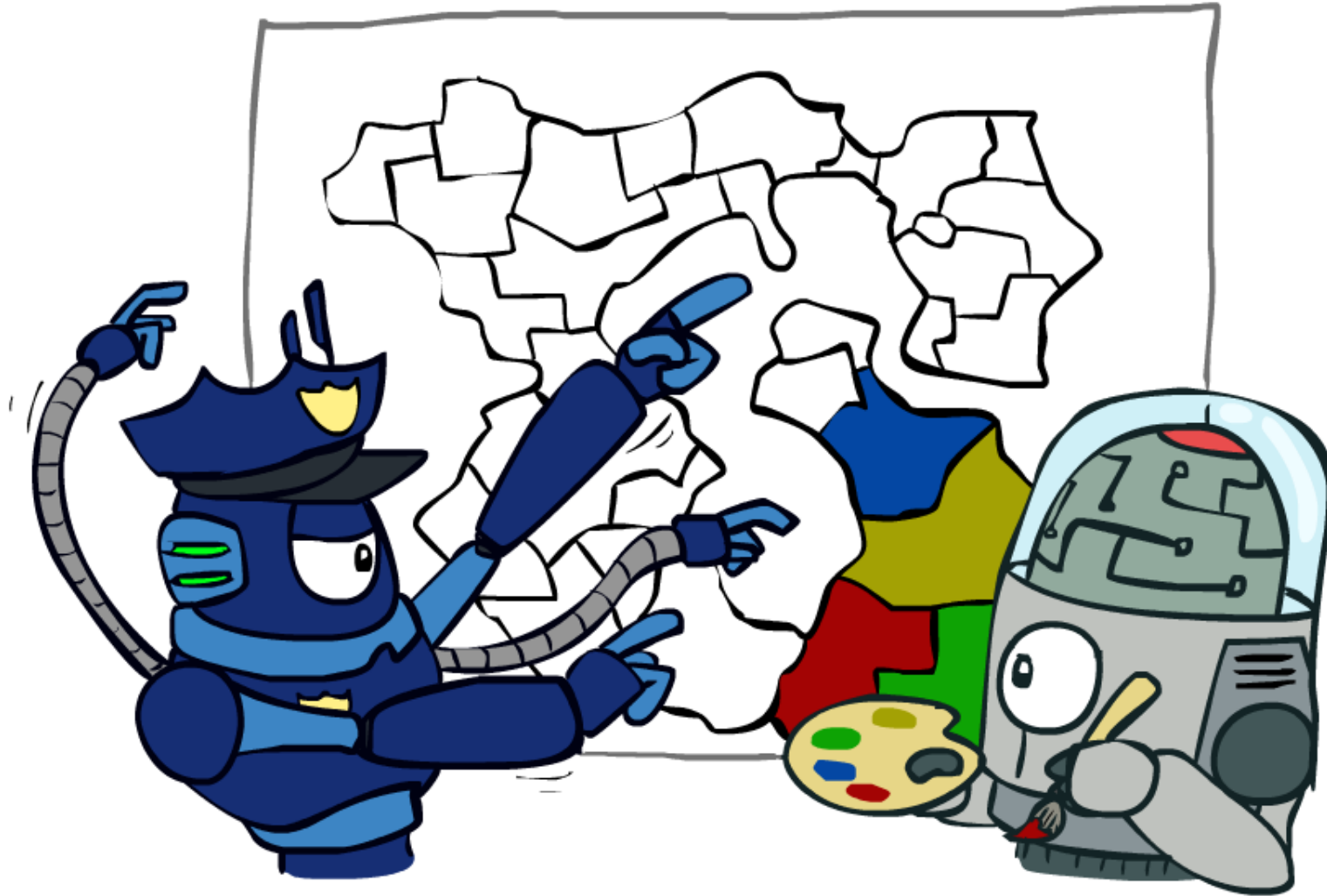


Constraint Graphs

- Binary CSP: each constraint relates (at most) two variables
- Binary constraint graph: nodes are variables, arcs show constraints
- General-purpose CSP algorithms use the graph structure to speed up search. E.g., Tasmania is an independent subproblem!



Varieties of CSPs and Constraints



Varieties of CSPs

- Discrete Variables

- Finite domains

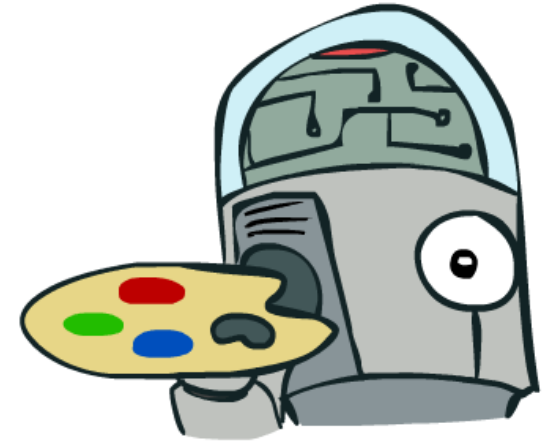
- E.g., Boolean CSPs, including Boolean satisfiability (NP-complete)

- Infinite domains (integers, strings, etc.)

- E.g., job scheduling, variables are start/end days for each job

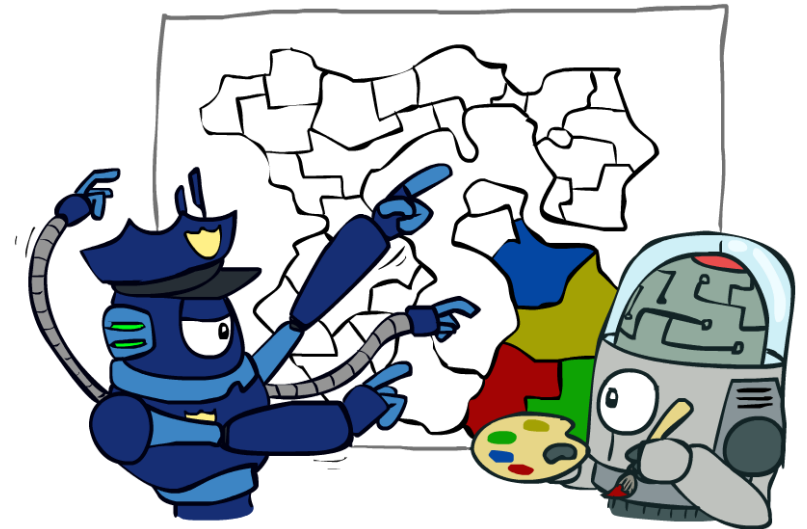
- Continuous variables

- E.g., start/end times for Hubble Telescope observations

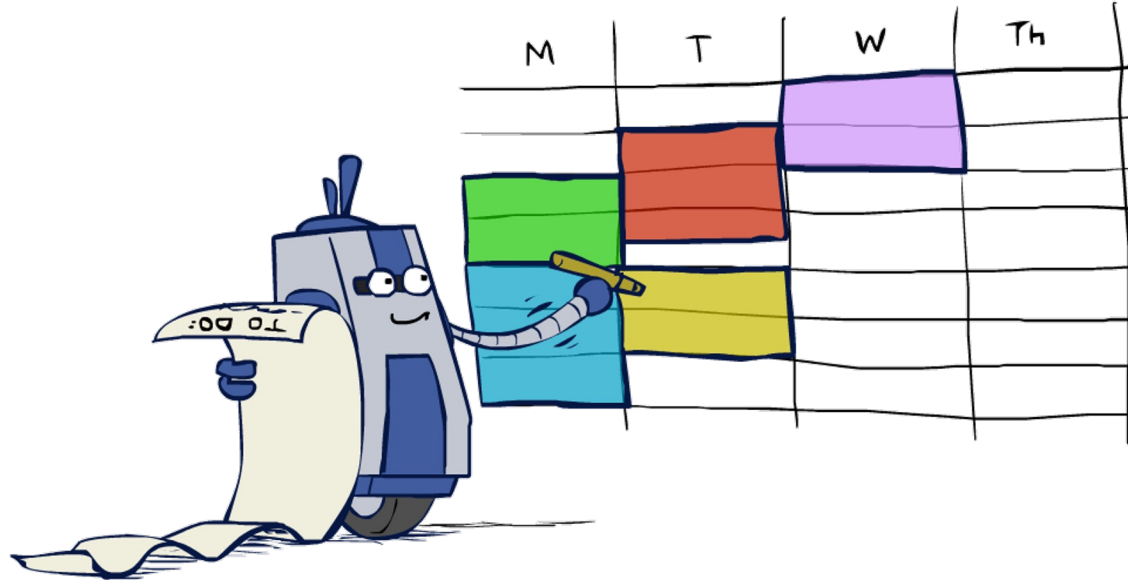


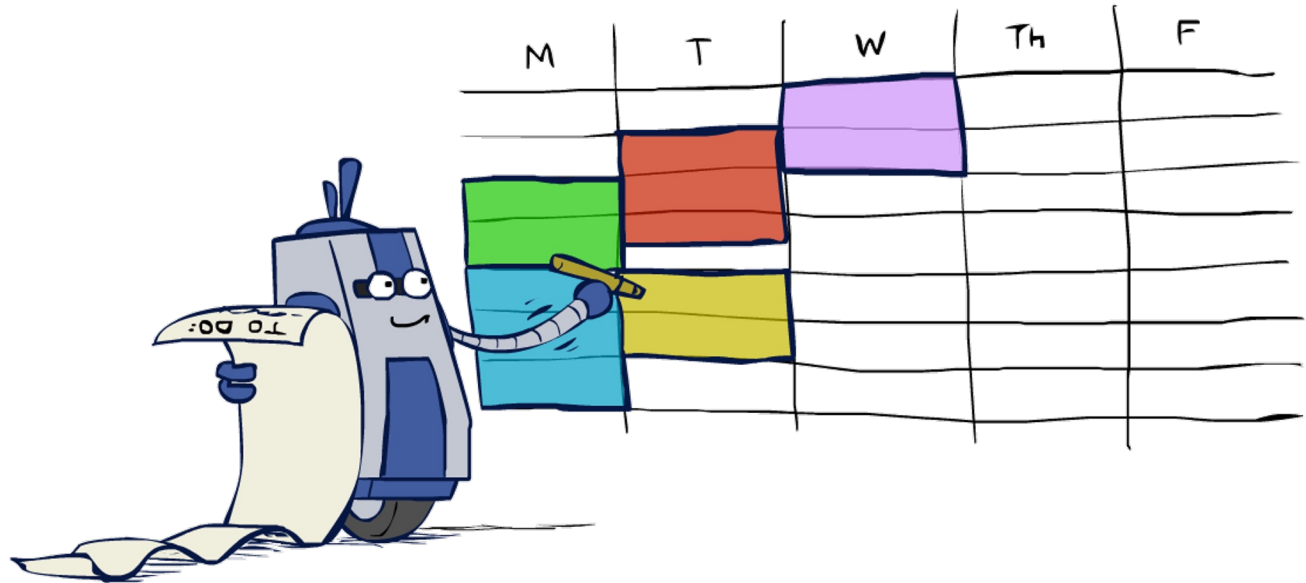
Varieties of Constraints

- Varieties of Constraints
 - Unary constraints involve a single variable (equivalent to reducing domains), e.g.: $SA \neq \text{green}$
 - Binary constraints involve pairs of variables, e.g.: $SA \neq WA$
 - Higher-order constraints involve 3 or more variables
- Preferences (soft constraints):
 - E.g., red is better than green
 - Often representable by a cost for each variable assignment
 - Gives constrained optimization problems



Real-World CSPs

- Scheduling problems: e.g., when can we all meet?
 - Timetabling problems: e.g., which class is offered when and where?
 - Assignment problems: e.g., who teaches what class
 - Hardware configuration
 - Transportation scheduling
 - Factory scheduling
 - Circuit layout
 - Fault diagnosis
 - ... lots more!
 - Many real-world problems involve real-valued variables...
- 
- The image shows a cartoon robot character with a blue body and a single wheel, holding a large yellow scroll. It is pointing with its right arm towards a grid that represents a schedule. The grid has four columns labeled 'M', 'T', 'W', and 'Th' at the top. The rows are represented by horizontal lines. Colored rectangles are placed in the grid: a green rectangle in the first row under 'M', a red rectangle in the second row under 'T', a purple rectangle in the third row under 'W', and a yellow rectangle in the fourth row under 'T'. The robot is standing on a small patch of ground.

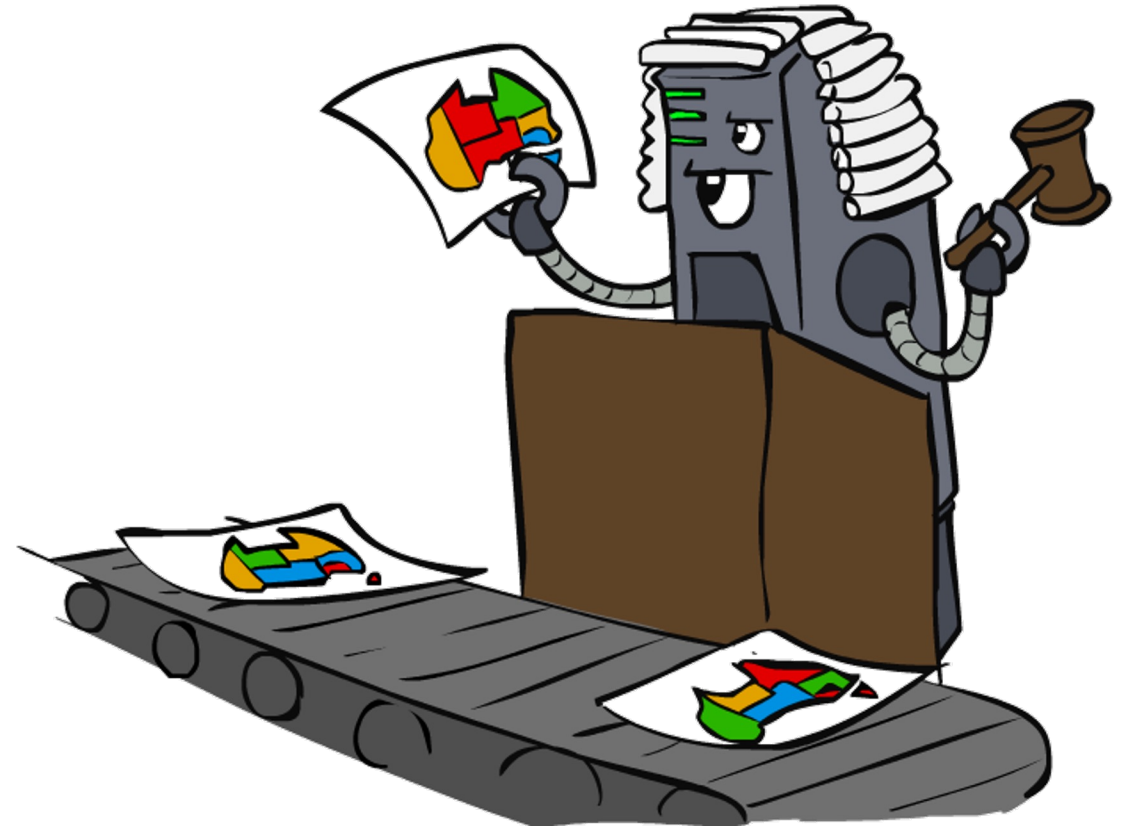


Solving CSPs



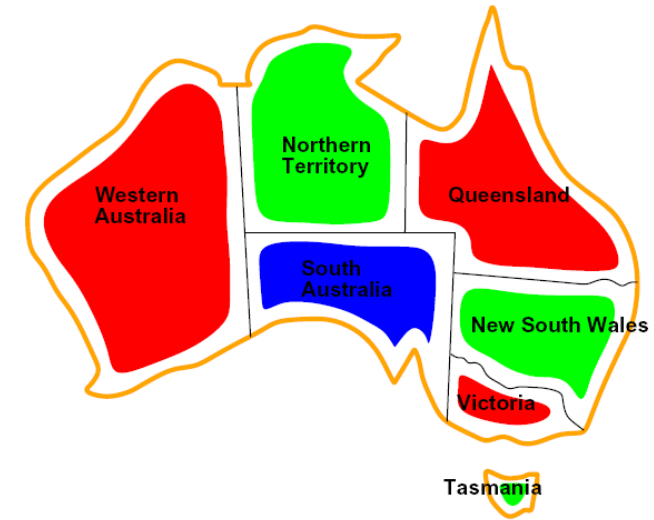
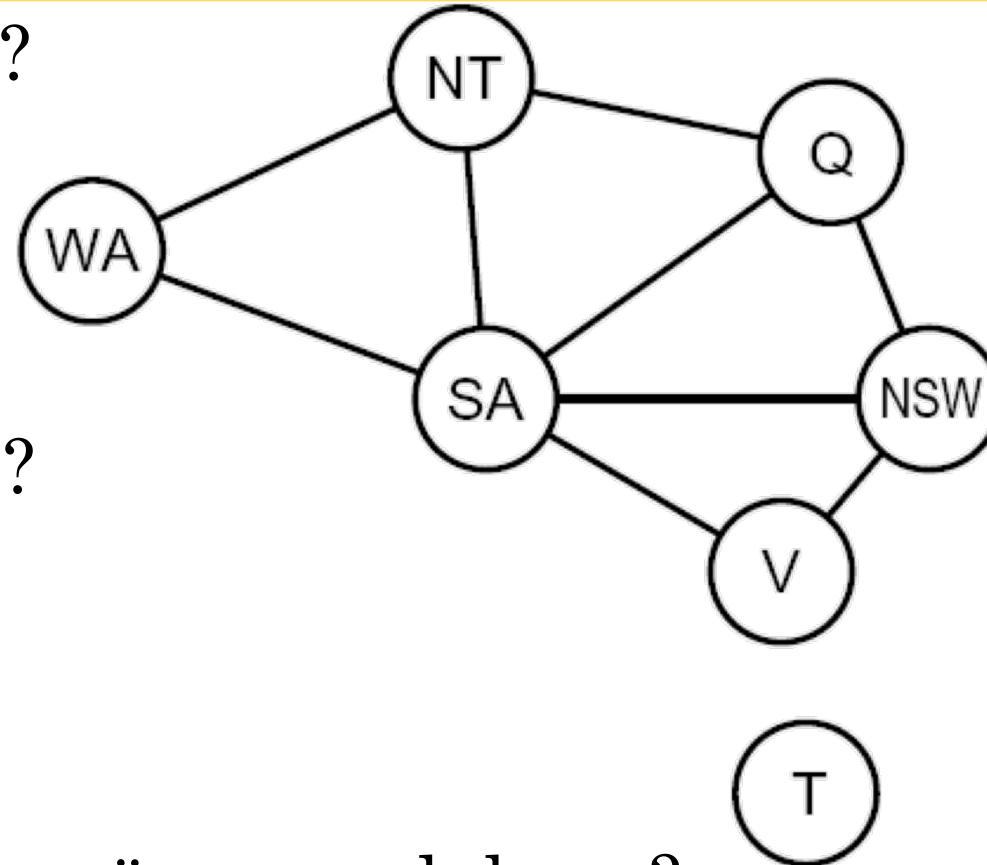
Standard Search Formulation

- Standard search formulation of CSPs
- States defined by the values assigned so far (partial assignments)
 - Initial state: the empty assignment, $\{\}$
 - Successor function: assign a value to an unassigned variable
 - Goal test: the current assignment is complete and satisfies all constraints
- We'll start with the straightforward, naïve approach, then improve it



Search Methods

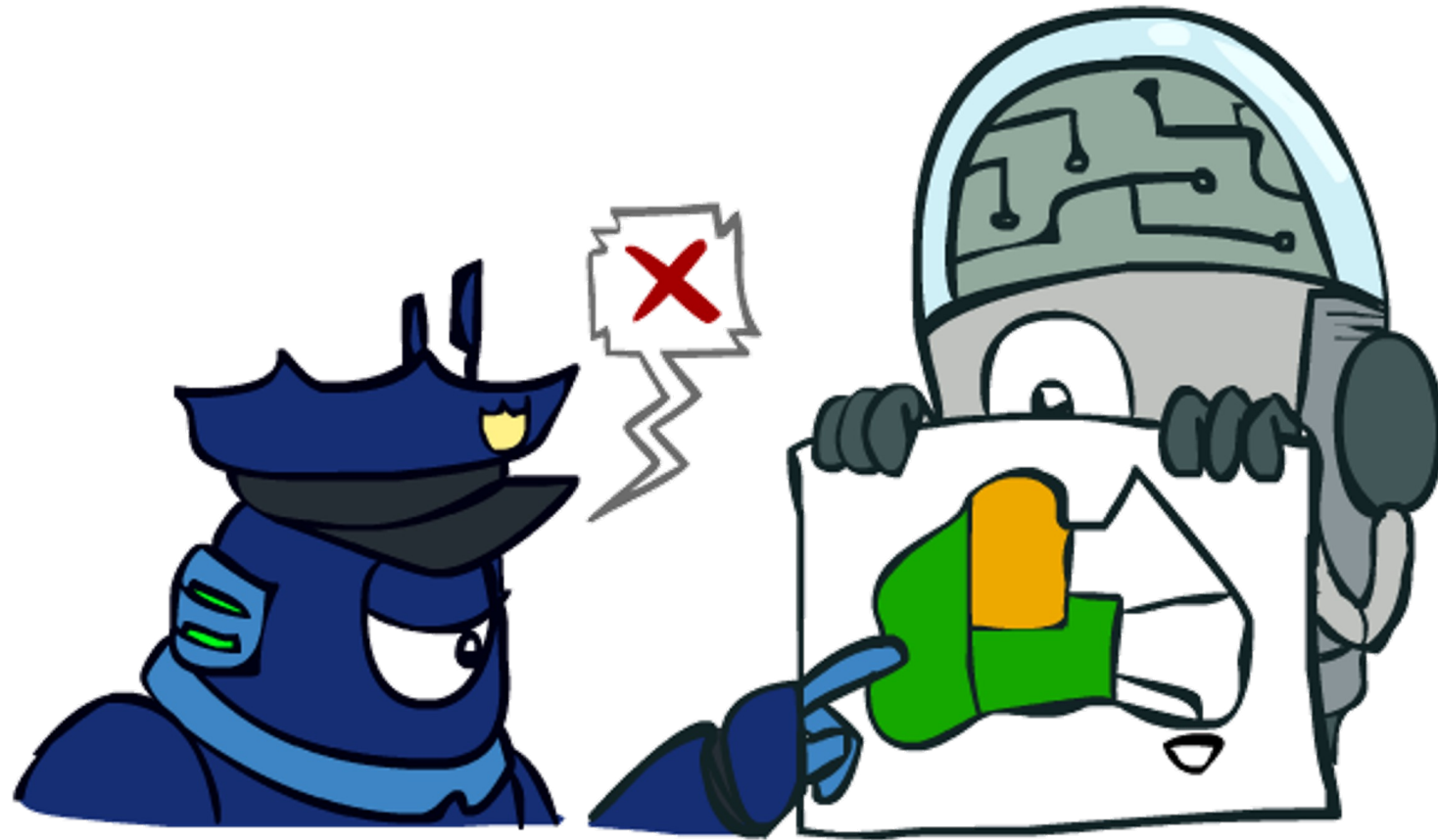
- What would BFS do?



- What would DFS do?

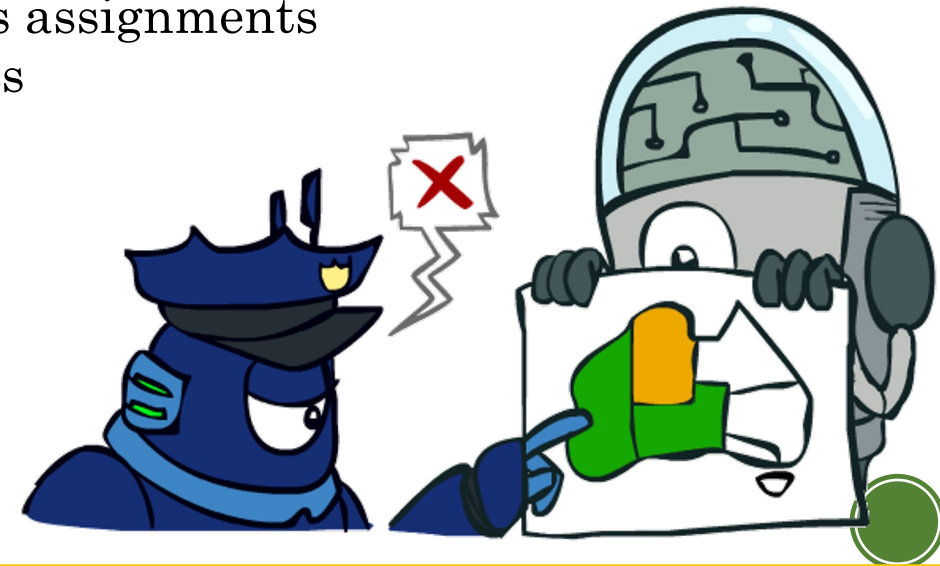
- What problems does naïve search have?

Backtracking Search

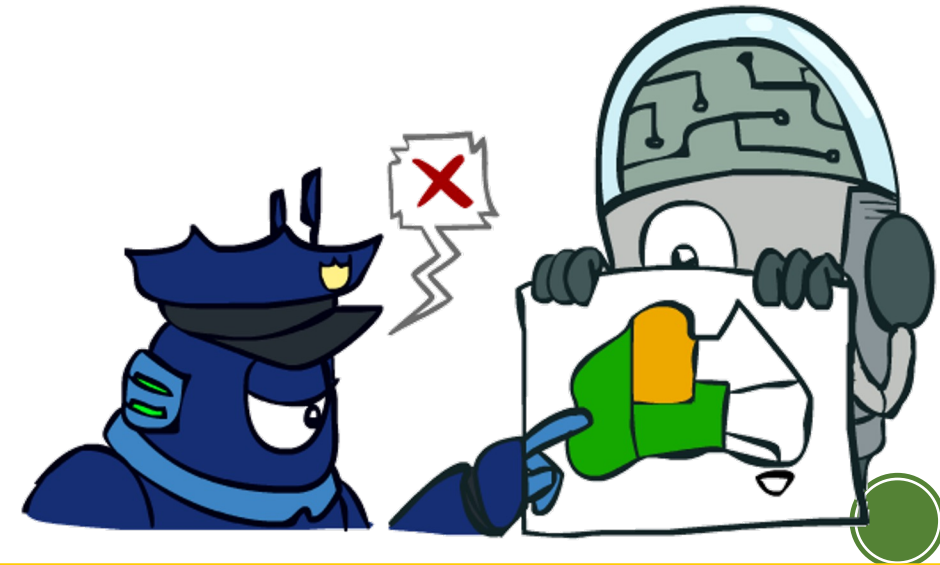
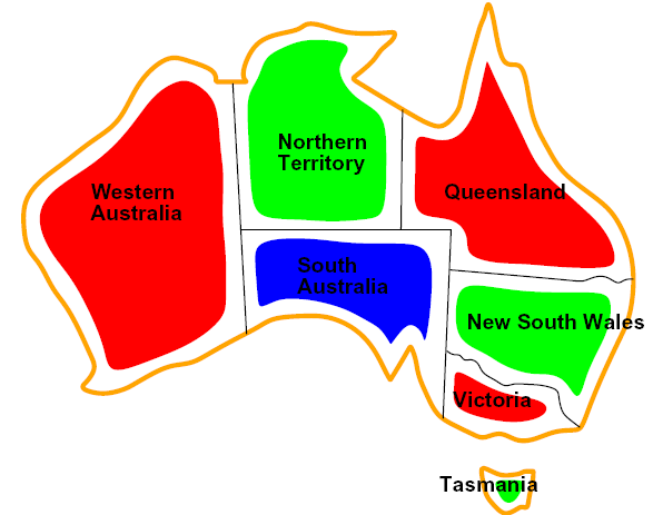
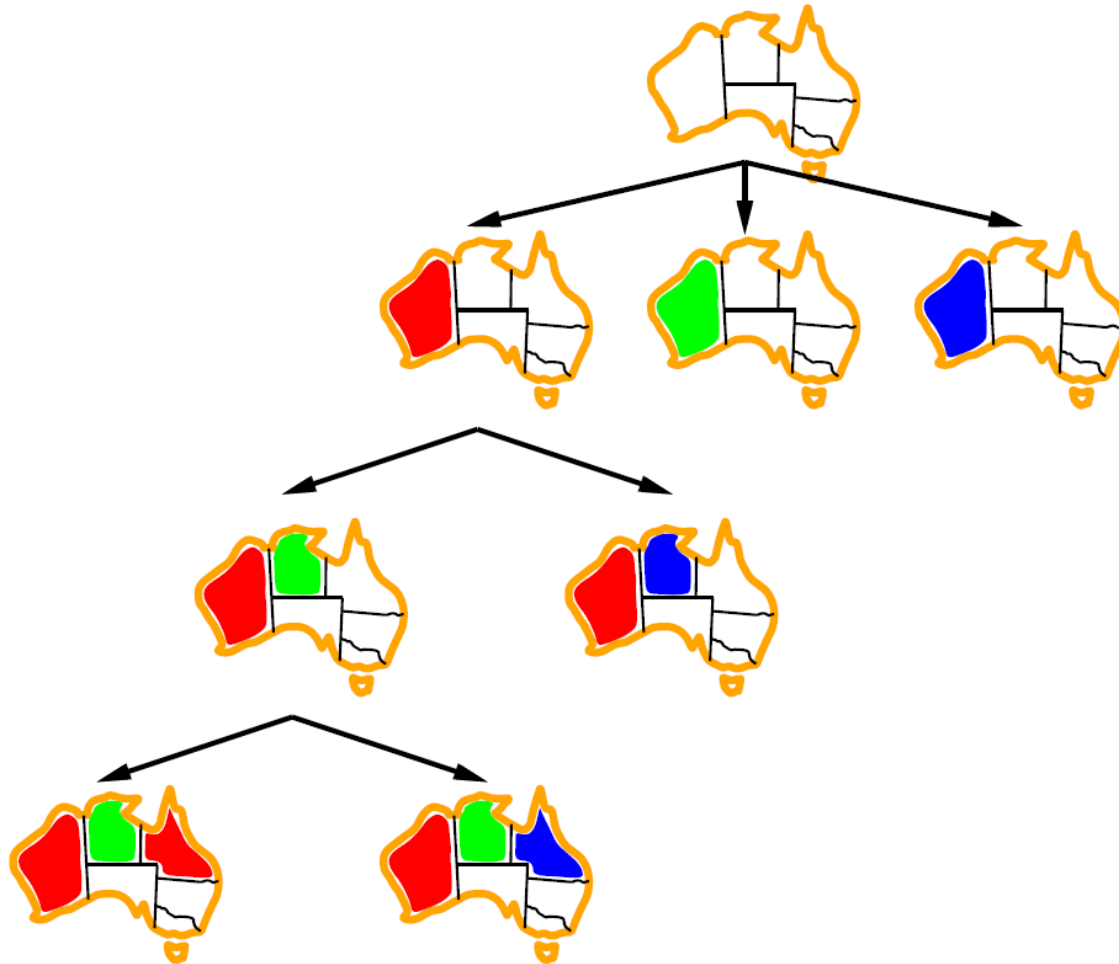


Backtracking Search

- Backtracking search is the basic uninformed algorithm for solving CSPs
- **Idea 1: One variable at a time**
 - Variable assignments are commutative, so fix ordering
 - I.e., [WA = red then NT = green] same as [NT = green then WA = red]
 - Only need to consider assignments to a single variable at each step
- **Idea 2: Check constraints as you go**
 - I.e. consider only values which do not conflict with previous assignments
 - Might have to do some computation to check the constraints
 - “Incremental goal test”
- Depth-first search with these two improvements is called *backtracking search* (not the best name)
- Can solve n-queens for $n \approx 25$



Backtracking Example



Backtracking Search

```
function BACKTRACKING-SEARCH(csp) returns solution/failure
  return RECURSIVE-BACKTRACKING({ }, csp)

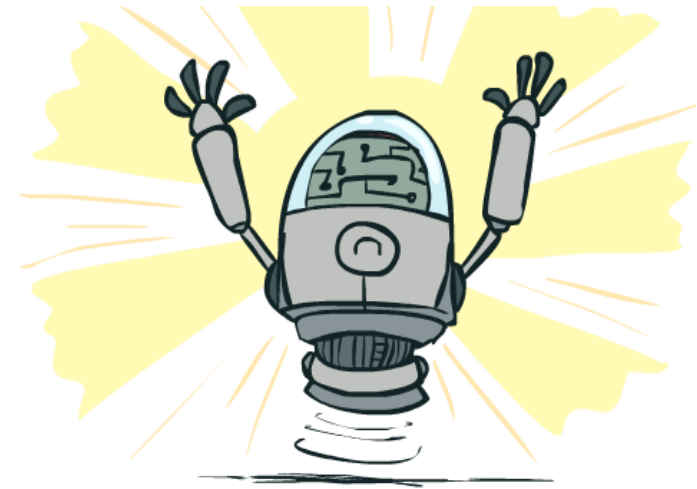
function RECURSIVE-BACKTRACKING(assignment, csp) returns soln/failure
  if assignment is complete then return assignment
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
    if value is consistent with assignment given CONSTRAINTS[csp] then
      add {var = value} to assignment
      result ← RECURSIVE-BACKTRACKING(assignment, csp)
      if result ≠ failure then return result
      remove {var = value} from assignment
  return failure
```

- Backtracking = DFS + variable-ordering + fail-on-violation
- What are the choice points?



Improving Backtracking

- General-purpose ideas give huge gains in speed
- Ordering:
 - Which variable should be assigned next?
 - In what order should its values be tried?
- Filtering: Can we detect inevitable failure early?
- Structure: Can we exploit the problem structure?

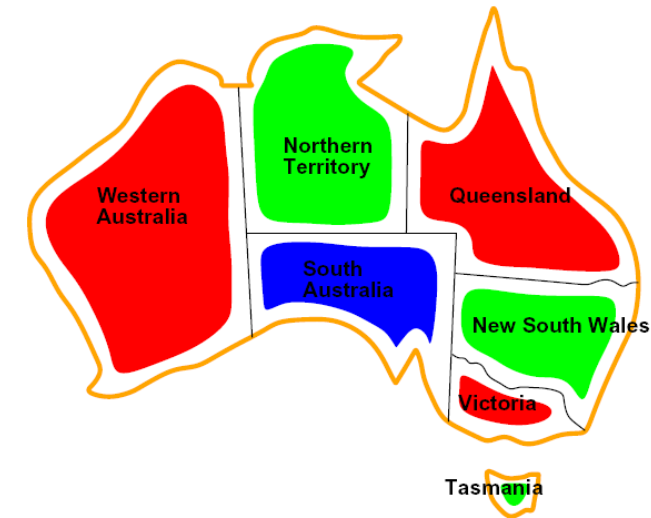
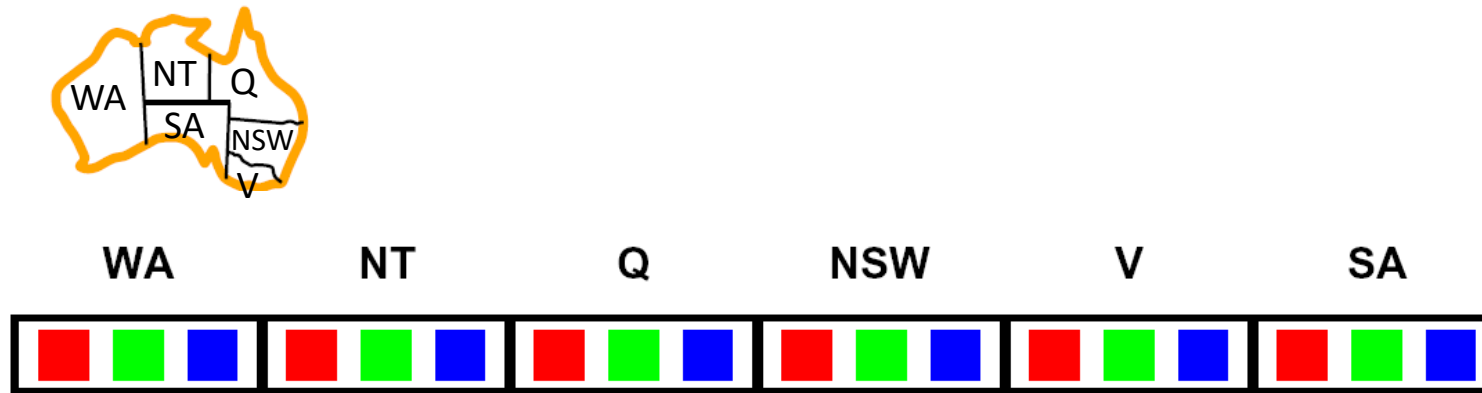


Filtering



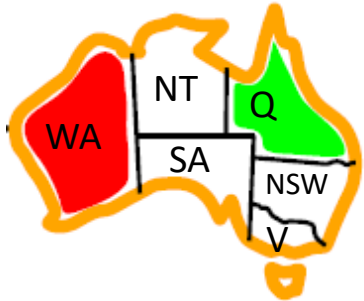
Filtering: Forward Checking

- Filtering: Keep track of domains for unassigned variables and cross off bad options
- Forward checking: Cross off values that violate a constraint when added to the existing assignment



Filtering: Constraint Propagation

- Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:



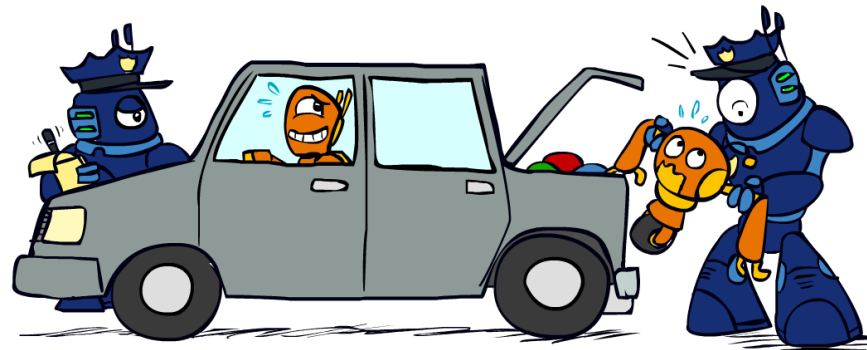
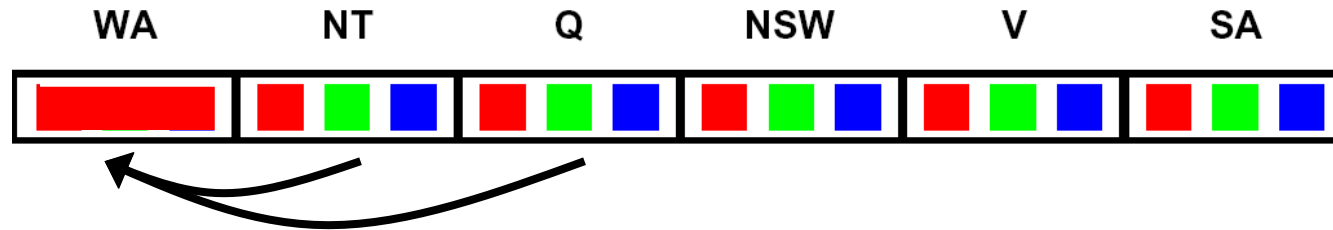
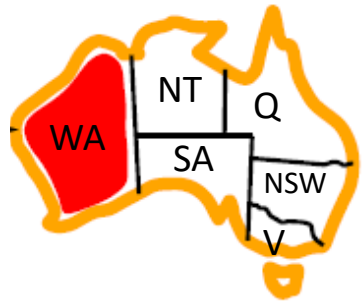
WA	NT	Q	NSW	V	SA
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- NT and SA cannot both be blue!
- Why didn't we detect this yet?
- Constraint propagation*: reason from constraint to constraint



Consistency of A Single Arc

- An arc $X \rightarrow Y$ is **consistent** iff for *every* x in the tail there is *some* y in the head which could be assigned without violating a constraint



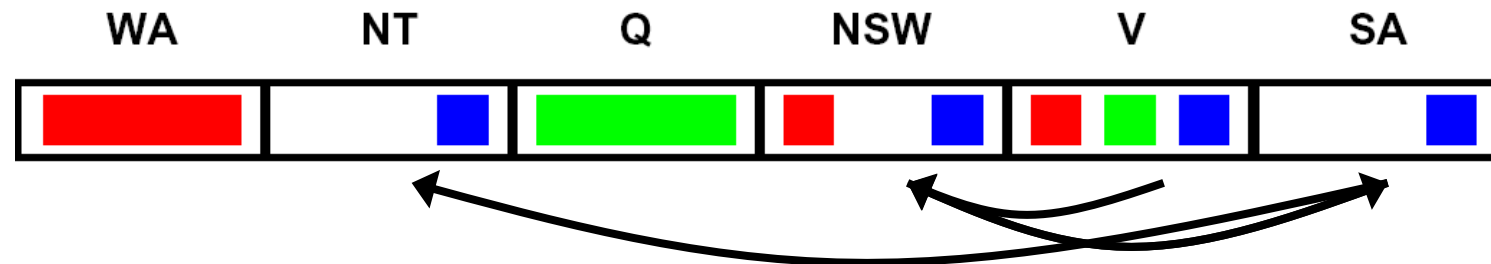
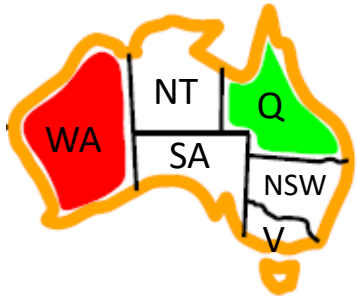
Delete from the tail!

- Forward checking: Enforcing consistency of arcs pointing to each new assignment



Arc Consistency of an Entire CSP

- A simple form of propagation makes sure **all** arcs are consistent:



- Important: If X loses a value, neighbors of X need to be rechecked!
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment
- What's the downside of enforcing arc consistency?

*Remember:
Delete from
the tail!*

