Chapter 8

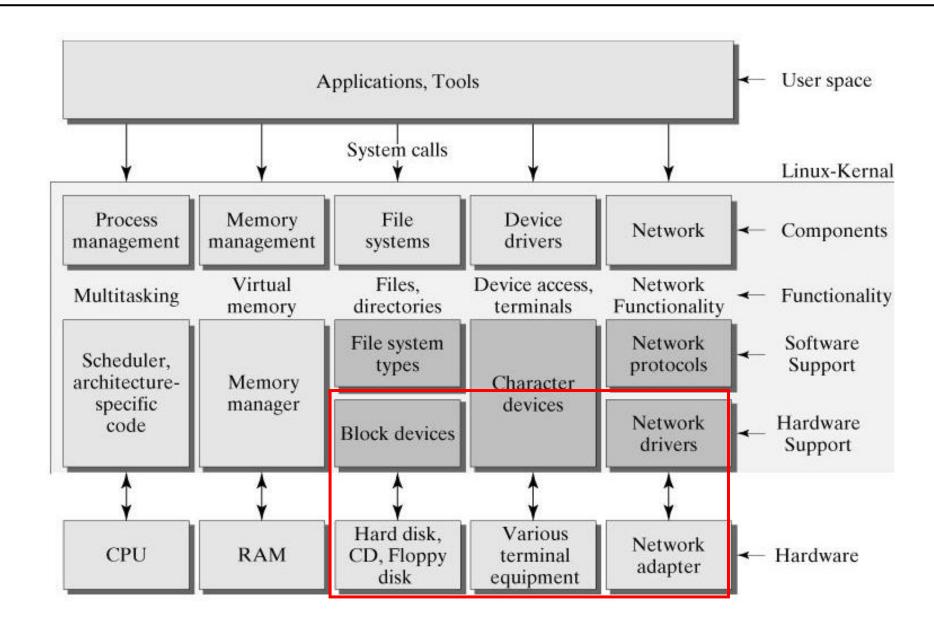
블록 디바이스 드라이버

15 February 2022 Minguk Choi koreachoi96@gmail.com

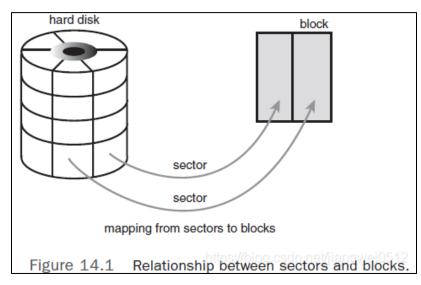
Contents

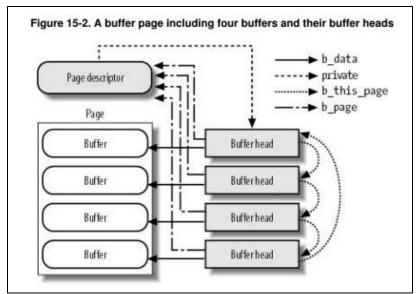
- Block Device Driver
 - Block Device
 - 구조체
 - I/O scheduler
- Block Device Driver 실습

Linux Kernel Architecture



Block Device





■ 섹터(sector)

- 블록 장치를 물리적으로 접근 가능한 최소 단위
- 일반적으로 512B

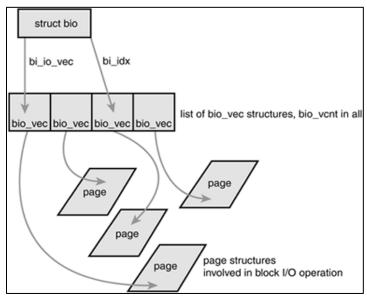
■ 블록(block)

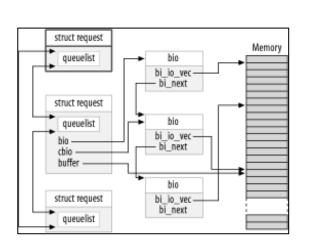
- 커널이 블록 장치를 논리적으로 접근 가능한 최소 단위
- 일반적으로 512B, 1KB, 4KB 많이 사용

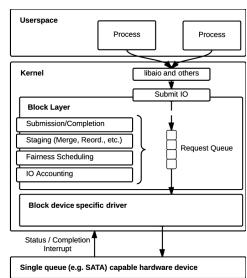
■ 버퍼

- 블록이 메모리 상에 존재할 때, 버퍼에 저장
- 커널은 버퍼 관련 정보를 버퍼 헤드에 저장
- 버퍼 헤드는 "버퍼-블록"의 연결관계를 나타냄

구조체

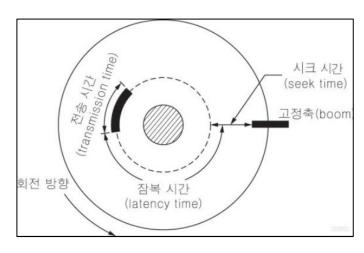


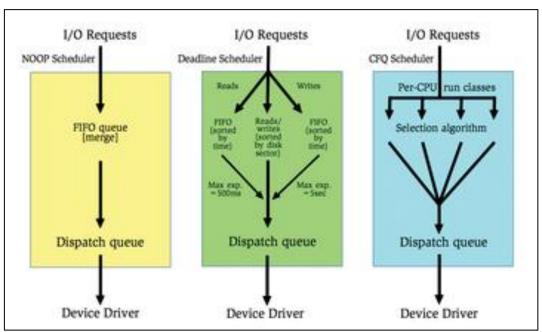




- struct bio
 - 커널 내부에서 블록 입출력을 전달하는 기본 장치
 - 현재 진행 중인 블록 입출력 동작을 세그먼트 리스트로 표현
- struct bio_vec: 세그먼트
- struct bio
- struct request
- struct request_queue

I/O scheduler





I/O scheduler

- 커널이 입출력 요청 순서대로 블록에 전달하면, 성능이 나쁨
- 현대 컴퓨터에서 가장 느린 동작 중 하나가 디스크 탐색
- 하드 디스크 헤드를 특정 블록이 있는 헤드로 옮겨야 함

■ 목적

• 디스크 탐색 시간을 최소화 하여 보다 나은 전체 성능 향상

■ 방법

- 병합
 - ✓ 인접 섹터를 읽는 요청을 병합
- 정렬
 - ✓ 섹터 순서에 따른 정렬 상태를 유치
 - ✓ 큐를 따라가며 탐색하는 동작이 디스크 섹터를 따라 순차적으로 일어나도록 함

■ 종류

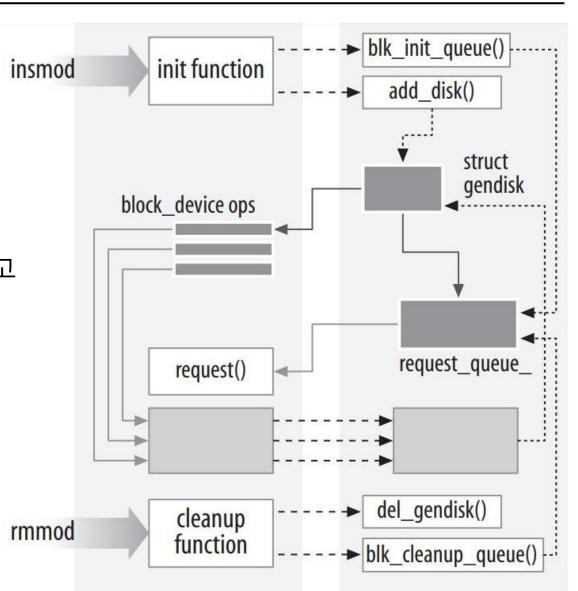
CFQ, Deadline, NOOP(SSD)

6

실습

- 리눅스 커널 내부구조
 - ch8 디바이스 드라이버
 - 블록 디바이스 드라이버 구현
- 구현 목표
 - 메모리의 일부를 가상 디스크 공간으로 할당하고
 - 해당 공간에 대한 I/O 제공

- 구현 내용
 - I/O request function
 - I/O operation function
 - ✓ open(), release(), ioctl()



실습

I/O request/operation function

.release =mydrv release;

.ioctl = mydrv_ioctl;

```
blk init queue()
                                                                             init function
                                                               insmod
                                                                                                     add disk()
static int MYDRV MAJOR = 0;
static char * mydrv data;
                                                                                                             struct
struct request_queue *mydrv_queue;
                                                                                                             gendisk
struct gendisk *mydrv disk;
                                                                          block_device ops
static int mydrv make request(struct request queue *q, struct bio *bio)
          bio for each segment(bvec, bio, i)
               case READ :
                      memcpy(buffer, data, bvec->bv len);
                      break:
                                                                                                       request_queue_
               case WRITE:
                                                                                request()
                      memcpy(dala, butter, bvec ->by Jen);
                      break;
static struct block device operations mydrv fops =
      .owner THIS MODULE;
                              int mydrv open(struct inode *inode, struct file *filp)
      .open = mydrv open;
```

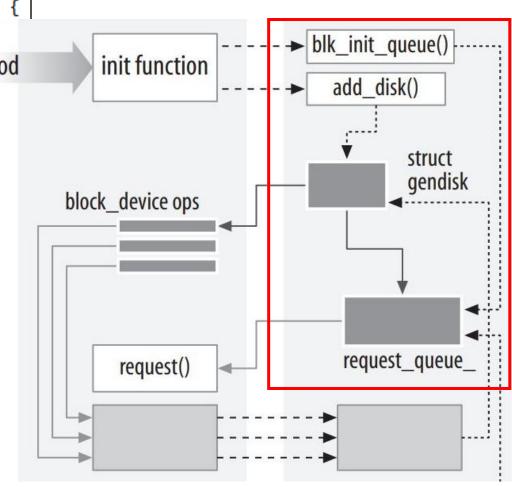
int mydrv release (struct inode *inode, struct file *filp)

int mydrv ioctl(struct inode *inode, struct file *filp, uns

실습

Init block device driver

```
int mydrv_init(void)
    if((MYDRV_MAJOR = register_blkdev(MYDRV_MAJOR, DEVICE_NAME)) < 0) {</pre>
    if ((mydrv_data = vmalloc(MYDRV_MAX_LENGTH)) == NULL) {
                                                                  insmod
   if ((mydrv_disk = alloc_disk(1)) == NULL) {
   if((mydrv queue = blk alloc queue(GFP KERNEL)) == NULL)
    blk_queue_make_request(mydrv_queue, &mydrv_make_request);
    blk_queue_hardsect_size(mydrv_queue, MYDRV_BLK_SIZE);
    mydrv disk->major = MYDRV MAJOR;
    mydrv_disk->first_minor = 0;
    mydrv_disk->fops = &mydrv_fops;
    mydrv disk->queue = mydrv queue;
    set_capacity(mydrv_disk, MYDRV_TOTAL_BLK);
    add_disk(mydrv_disk);
    return 0;
```



컴파일 에러

컴파일 에러 원인

- 리눅스 커널 내부구조
 - Kernel: Linux 2.6 (2003년)
- VirtualBox
 - Kernel: Linux 5.11.0
- \$ hostnamectl

Index of /pub/linux/kernel/v2.6/

<u> </u>	18-Dec-2003 03:27	32M
<u> </u>	18-Dec-2003 03:27	40M
<u>linux-2.6.0.tar.sign</u>	08-Aug-2013 19:25	665
<u>linux-2.6.0.tar.xz</u>	18-Dec-2003 03:27	25M
<u> inux-2.6.1.tar.bz2</u>	09-Jan-2004 07:31	32M
<u> </u>	09-Jan-2004 07:31	40M
<u>linux-2.6.1.tar.sign</u>	08-Aug-2013 19:25	665
<u>linux-2.6.1.tar.xz</u>	09-Jan-2004 07:31	25M
<u> </u>	<u> </u>	ocu

```
mingu@mingu-VirtualBox:~$ hostnamectl
   Static hostname: mingu-VirtualBox
        Icon name: computer-vm
        Chassis: vm
        Machine ID: a6346fa91d5b404da2c8238150f4d85b
        Boot ID: 81b0f244e1b04725a914912e95a3d0cb
   Virtualization: oracle
   Operating System: Ubuntu 20.04.3 LTS
        Kernel: Linux 5.11.0-46-generic
   Architecture: x86-64
```

실습 목표

■ 커널 함수를 업데이트 하자!

```
static int mydrv make request(struct request queue *q, struct bio *bio)
       struct gendisk *bdisk = bio->bi disk;
       struct block device *bdev = bio->bi bdev;
       sector t sector = bio->bi iter.bi sector;
       sector t sector = bio->bi sector;
       unsigned long len = bio->bi iter.bi size >> 9;
       unsigned long len = bio->bi size >> 9;
       blk_mq_init_queue(&mydrv_make_request);
       blk queue make request(mydrv queue, &mydrv make request);
       blk queue logical block size(mydrv queue, MYDRV BLK SIZE);
       blk queue hardsect size(mydrv queue, MYDRV BLK SIZE);
```

blk_queue_make_request()

```
int mydrv_init(void)
    if((MYDRV_MAJOR = register_blkdev(MYDRV_MAJOR, DEVICE_NAME)) < 0) {</pre>
   if ((mydrv data = vmalloc(MYDRV MAX LENGTH)) == NULL) {
   if ((mydrv disk = alloc disk(1)) == NULL) {
   if((mydrv queue = blk alloc queue(GFP KERNEL)) == NULL)
   blk queue make request(mydrv queue, &mydrv make request);
   blk queue hardsect size(mydrv queue, MYDRV BLK SIZE);
   mydrv disk->major = MYDRV MAJOR;
   mydrv disk->first minor = 0;
   mydrv disk->fops = &mydrv fops;
   mydrv disk->queue = mydrv queue;
   set capacity(mydrv disk, MYDRV TOTAL BLK);
   add disk(mydrv disk);
   return 0;
```

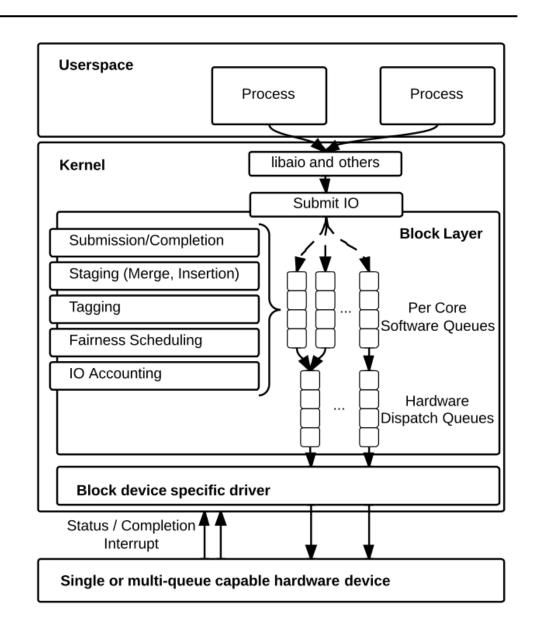
- blk_queue_make_request()
 - define an alternate make_request function for a device

Multi-Queue Block I/O (blk-mq)

- Multi-Queue Block I/O
 - why? single-queue is bottleneck
 - Multi-processor, SSD

blk-mq (Multi-Queue Block IO Queueing Mechanism) is a new framework for the Linux block layer that was introduced with Linux Kernel 3.13, and which has become feature-complete with Kernel 3.16.^[1] Blk-mq allows for over 15 million IOPS with high-performance

Diagram).^[3] How much longer this request_fn based approach will exist in the Linux kernel is currently unclear (July 2014).^{[4][5]}



Linux Kernel 5.0 Updated 10 March 2019



torvalds

I'm not getting any more than that 4.x numbers started and toes.

(I repeat once again – our releases are not tied to any specific features, so the number of the new version 5.0 means only that for numbering versions 4.x I don't have enough fingers and toes)

It's all about the multi-queue block layer (blk-mq). There are plenty of introductory articles about him on the Internet, so let's get straight to the point. The transition to blk-mq was started a long time ago and slowly advanced. Multi-queue scsi appeared (kernel parameter scsi_mod.use_blk_mq), new mq-deadline schedulers, bfq, etc. appeared ...

https://prog.world/linux-kernel-5-0-we-write-simple-block-device-under-blk-mg/

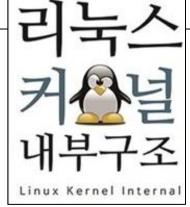
Linux Kernel 5.0

Updated 10 March 2019

The number of block device drivers, which work in the old manner, has been reduced. And in 5.0, the function blk_init_queue () was removed as unnecessary. And now the old glorious code lwn.net/Articles/58720 from 2003 is not only not going, but has lost its relevance. Moreover, new distributions that are being prepared for release this year, in the default configuration, use multiqueue block layer. For example, on the 18th Manjaro, the kernel, though version 4.19, but blk-mq by default.

Therefore, we can assume that in 5.0 the transition to blk-mq is completed. And for me this is an important event that will require rewriting the code and additional testing. That in itself promises the appearance of bugs large and small, as well as several fallen servers (It is necessary, Fedya, it is necessary! (C)).

https://prog.world/linux-kernel-5-0-we-write-simple-block-device-under-blk-mg/





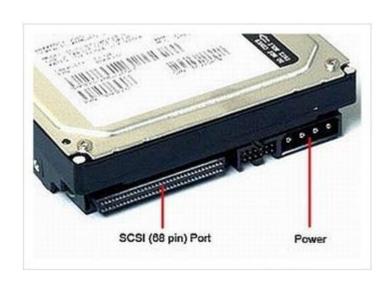
MQ processing structure flow diagram

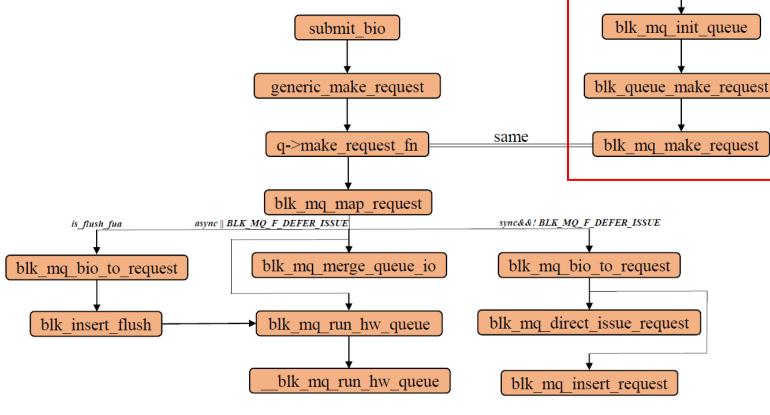
Kernel 3.xx

blk_queue_make_request() -> blk_mq_make_request()

SCSI(Small Computer System Interface)

Fast but Expensive I/O Interface for Sever/Workstation





scsi alloc sdev

scsi mq alloc queue

What's different?

```
static int init blkdev init(void) {
           if ((ret = blkdev add device()) != SUCCESS)
                                                                                 static in do request(struct request *rq, unsigned int *nr bytes) {
                                                                                        int ret = SUCCESS:
static int blkdev add device(void) {
                                                                                        struct bio vec bvec;
        int ret = SUCCESS:
                                                                                        struct req iterator iter;
                                                                                        block_dev_t *dev = rq->q->queuedata;
        struct gendisk *disk;
                                                                                        loff t pos = blk rq pos(rq) << SECTOR SHIFT;</pre>
        struct request queue *queue;
                                                                                        loff t dev size = (loff t)(dev->capacity << SECTOR SHIFT);</pre>
        block dev t *dev = kzalloc(sizeof(block dev t), GFP KERNEL);
                                                                                        printk("request start from sector %lld\n", blk_rq_pos(rq));
        dev->tag set.ops → &mg ops:
                                                                                        rq for each segment(bvec, rq, iter) {
        queue = blk mg init queue(&dev->tag set);
                                                                                               unsigned long b_len = bvec.bv_len;
                                                                                               void *b buf = page address(bvec.bv page) + bvec.bv offset;
        ret = blk mg alloc tag set(&dev->tag set);
                                                                                               if ((pos + b len) > dev size)
                                                                                                      b len = (unsigned long)(dev size - pos);
                                                                                               if (b len < 0)
static struct blk mq ops mq ops = {
                                                                                                      b len = 0;
        . queue rq = queue rq
                                                                                               if (rq data_dir(rq) == WRITE)
                                                                                                      memcpy(dev->data + pos, b buf, b len);
                                                                                               else
                                                                                                      memcpy(b buf, dev->data + pos, b len);
static blk status t queue rq(struct blk mq hw ctx *hctx,
                                                                                               pos += b len;
                        const struct blk_mq_queue data *bd) {
                                                                                               *nr bytes += b len;
        unsigned int nr bytes = 0;
        blk status t status = BLK STS OK;
                                                                                        return ret;
        struct request *rq = bd->rq;
        blk mg start request(rg);
        if (do_request(rq, &nr_bytes) != 0) status = BLK_STS_IOERR;
```

Q & A

Block Device Driver

- Block Device Driver 구현 실습
 - 컴파일 에러

Multi-Queue Block I/O

Linux kernel 5

- Block Device Driver 구현 실습
 - Linux kernel 5



