

RF3_Team_WAL

Supported by IITP, StarLab.

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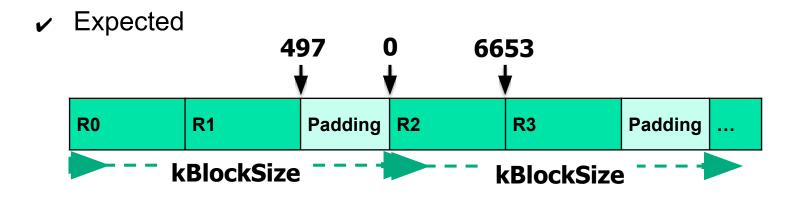
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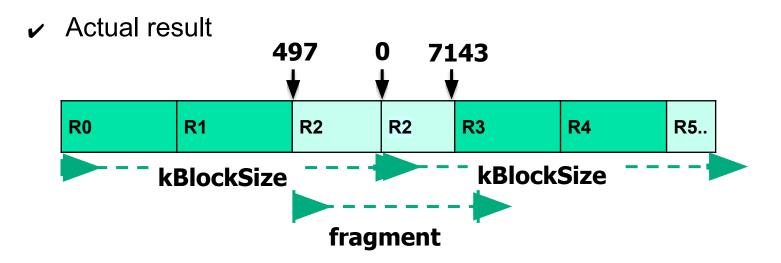




RocksDB Festival: kBlockSize

Previous experiment







- Experiment Info.
 - Insert padding in WAL
 - Purpose
 - Performance comparison according to existence of padding
 - Conditions
 - kBlockSize = 8KB, 16KB, 32KB
 - key_size = (16 128 512 1024)
 - value_size = (16 128 256 512 1024 2048 3072 4096 5120 6144 7168 8192)
 - benchmarks= fillrandom
 - threads=6
 - compression_type=NoCompression
 - entries = 100000 times





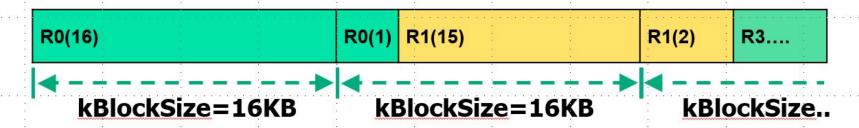
Hardware Environment : D

D		
CPU	1 * AMD Ryzen 5 3500X 6-Core	
os	Ubuntu 20.04.2 LTS	
SSD	mx500	

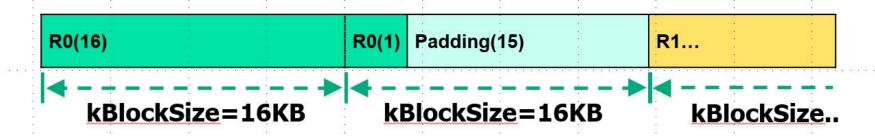




- Previous experiment (Hypothesis)
 - ✓ If kBlockSize = 16KB, record size = 17KB, num=100



- Padding is not exist, predict performance improve.
- ✓ But fragmentation may cause consistency issues.

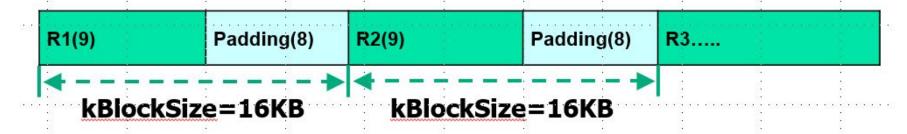


Padding is exist, overhead is too big.





- Previous experiment (Hypothesis)
 - ✓ If kBlockSize = 16KB, record size = 9KB, num=100



- ✓ Padding is extremely high size
 - Expected performance degradation
- ✓ Write in DB = 900KB, Write in Storage = 1600KB
 - Write Amplification is so high





Code revise

✓ log_writer.h append code

```
#define PADDINGINSERT // if you want disable padding insert, delete this line.
  std::string repeat_;
                         Max Padding size(kBlockSize)
  for(int i = 0; i < n; i++) {
    repeat += s;
  return repeat ;
std::string byte_padding = "\x00"; 패딩 \x00 삽입
#ifdef PADDINGINSERT
  std::string slice_data_ = repeat(byte_padding, kBlockSize);
delse
 std::string slice_data_ = repeat(byte_padding, 10);
#endif
const char *slice_data_ = slice_data_.c_str();
```





- Code revise
 - ✓ log_writer.cc revise code

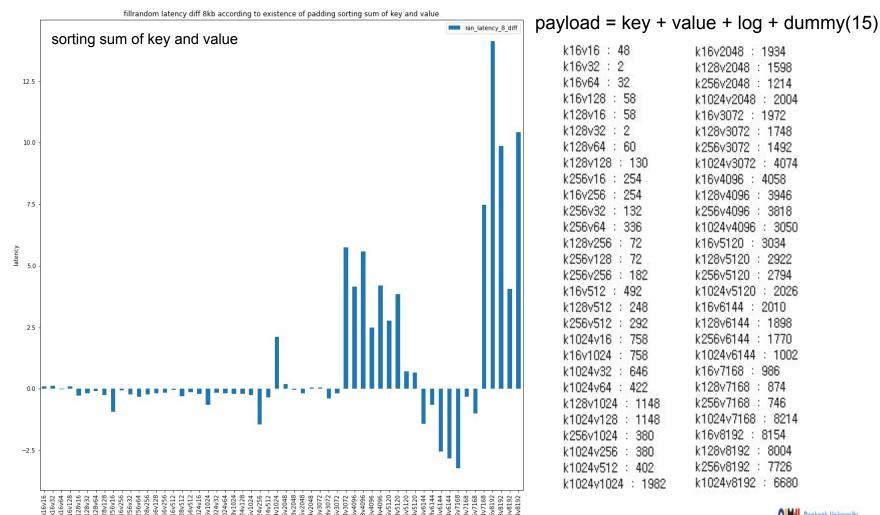
```
If padding exist
 bool begin = true;
                                        kBlock의 남은 크기가 record의
  const int64_t leftover = kBlockSize - block_offset_;
                                        남은 크기보다 작고 kBlock이
                                        비어있지 않을때
#ifdef PADDINGINSERT
   ispadding = (leftover < static_cast<int64_t>(left) && leftover != kBlockSize) || leftover < header_size ? true : false;
   ispadding = leftover < header_size ? true : false;
If padding not exist
  // fprintf(stdout, "leftover : %ld\n", leftover);
                                      kBlock의 남은 크기가
  assert(leftover >= 0);
                                      header의 크기보다 작을 때
if (ispadding) {
// Switch to a new block
                                             kBlock의 남은 크기만큼 max
   if (leftover > 0) {
                                             Padding size를 slicing해서
     // Fill the trailer (literal below relies on kHeaderSize and
     // kRecyclableHeaderSize being <= 11)
                                             padding(x00 -> 1byte)
     assert(header_size <= 11);
                                             slicing된 크기만큼 삽입한다.
// revised============
     // fprintf(stdout, "padding insert\n");
     s = dest_->Append(Slice(slice_data__, static_cast<size_t>(leftover)));
// -----
```





Result

✓ kBlockSize = 8KB, latency = padding insert – not padding insert (s)

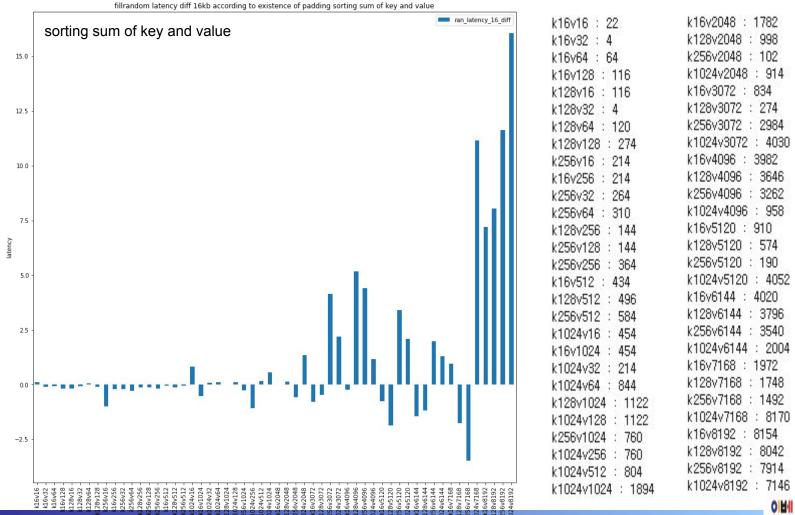






Result

✓ kBlockSize = 16KB, latency = padding insert – not padding insert (s)

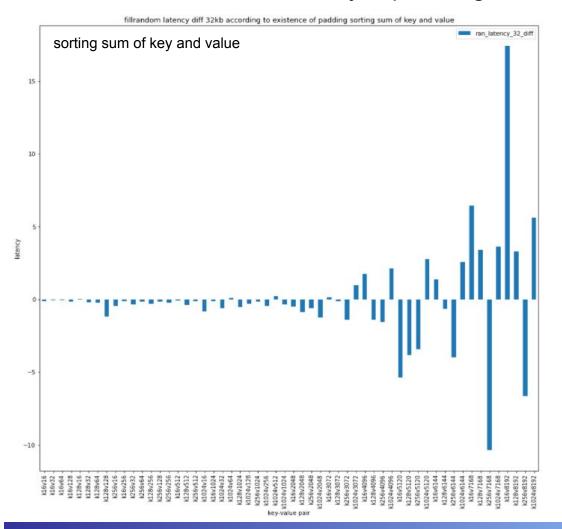






Result

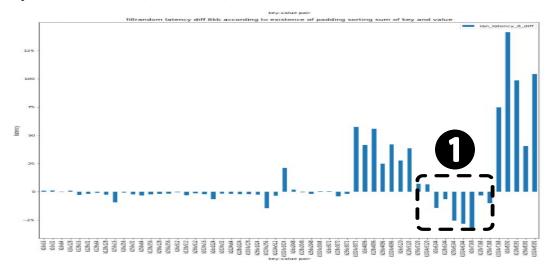
✓ kBlockSize = 32KB, latency = padding insert – not padding insert (s)



k16v16 : 42	k16v2048 : 1478
k16v32 : 8	k128v2048 : 1996
k16v64 : 26	k256v2048 : 204
k16v128 : 66	k1024v2048 : 1828
k128v16 : 66	k16v3072 : 1668
k128v32 : 8	k128v3072 : 548
k128v64 : 26	k256v3072 : 2618
k128v128 : 242	k1024v3072 : 3942
k256v16 : 134	k16v4096 : 3830
k16v256 : 134	k128v4096 : 3046
k256v32 : 218	k256v4096 : 2150
k256v64 : 278	k1024v4096 : 1916
k128v256 : 288	k16v5120 : 1820
k256v128 : 288	k128v5120 : 1148
k256v256 : 194	k256v5120 : 380
k16v512 : 318	k1024v5120 : 1938
k128v512 : 330	k16v6144 : 1858
k256v512 : 378	k128v6144 : 1298
k1024v16 : 908	k256v6144 : 658
k16v1024 : 908	k1024v6144 : 4008
k1024v32 : 428	k16v7168 : 3944
k1024v64 : 578	k128v7168 : 3496
k128v1024 : 1070	k256v7168 : 2984
k1024v128 : 1070	k1024v7168 : 8126
k256v1024 : 218	k16v8192 : 8078
k1024v256 : 218	k128v8192 : 7742
k1024v512 : 50	k256v8192 : 7358
k1024v1024 : 1718	k1024v8192 : 5054
K1024V1024 - 1710	FILE Deals



- Discussion
 - ✓ kBlockSize = 8KB



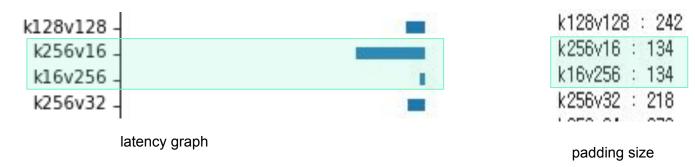
- Despite insert padding, latency is lower than no padding result (performance better).
 Why?
- 2. The larger the padding, the worse the performance(latency) but exception(compression, memtable, etc....) is exist







- Discussion
 - ✓ kBlockSize = 8KB



Padding sizes are same, different performance



Discussion





