

# RF3\_Team\_WAL

Supported by IITP, StarLab.

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## RocksDB Festival

#### Content

- ✓ WAL Log File Format
- Experiment : is kBlockSize affected performance?
  - Info
  - Hypothesis
  - Result
  - Discussion1
  - Discussion2
- ✓ Next assignment : WAL performance according to value or key size





# RocksDB Festival: Log File Format

WAL Log File Format



#### Rn: variable size records

- Consists of a sequence of variable length records.
- ✓ Records are grouped by kBlockSize(32k).
- ✓ If a certain record cannot fit into the leftover space (leftover < Rn), then the leftover space is **padded** with empty (null) data.
- ✓ If record is bigger than kBlockSize, record occurs fragmentation.





# RocksDB Festival: Log File Format

The Legacy Record Format

- ✓ Record consists of CRC, Size, Type, Payload
  - CRC(Cyclic Redundancy Check): Verifies the integrity of the WAL
  - Size : Length of the record size
  - Type: kZeroType, kFullType, kFirstType, kLastType, kMiddleType



Payload : The actual value of the key-value is written





# RocksDB Festival: Log File Format

The Recyclable Record Format

CRC (4B) Size(2B) Type(1B) Log number (4B) Payload
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- Record consists of CRC, Size, Type, Log Number, Payload
  - CRC, Size, Payload : same as the components of the legacy record format.
  - Type: kRecyclableFullType, kRecyclableFirstType, kRecyclableMiddleType, kRecyclableLastType
  - Log Number: Distinguish between the previous log writer and the last one. (32bit)





Hardware Environment : D

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CPU	1 * AMD Ryzen 5 3500X 6-Core
OS	Ubuntu 20.04.2 LTS
SSD	mx500





- Experiment Info.
  - ✓ WAL overhead measurement according to the kBlockSize.
  - ✓ Because of the kBlockSize affect the size of padding, WAL overhead will change according to the kBlockSize.
  - ✓ Conditions
    - kBlockSize = 4KB, 8KB, 16KB, 32KB(default), 64KB

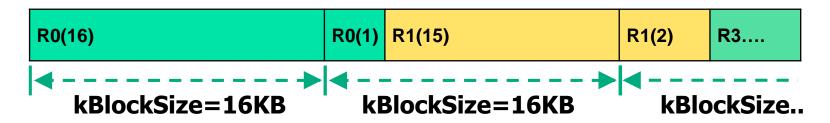
#### [db\_bench Option]

- benchmarks ="fillseq", "fillrandom"
- disable\_wal = false, true
- value\_size = 16byte, 32byte, 64byte, 128byte, 256byte, 512byte, 1024byte, 2048byte, 3072byte, 4096byte, 5120byte, 6144byte, 7168byte 8192byte 16384byte

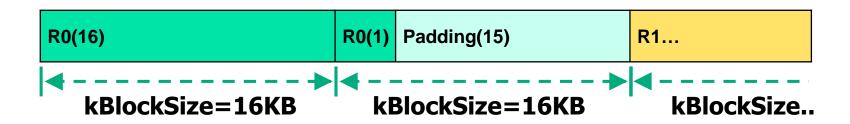




- Hypothesis If record size is bigger than kBlockSize
  - ✓ If kBlockSize = 16KB, record size = 17KB, num=100



- ✓ Padding is **not exist**, predict performance improve.
- But, Wouldn't fragmentation cause consistency issues?

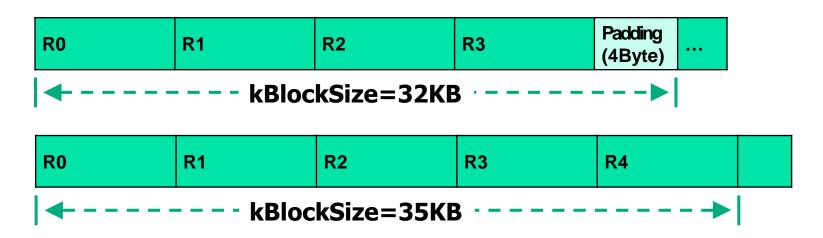


Padding is exist, overhead is too big.



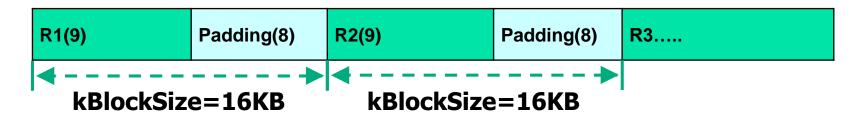


- Hypothesis If kBlockSize is bigger than record size
  - ✓ Size of Padding = kBlockSize % Size of Recode(fixed)
  - ✓ Ex. Size of record = 7KB, num=20
    - kBlockSize = 32KB, Size of Padding = 4KB
      - $\rightarrow$  Total 160KB = 140KB + **20KB** (Higher overhead)
    - kBlockSize = 35KB, Size of Padding = 0KB
      - $\rightarrow$  Total 140KB = 140KB + **0KB** (Lower overhead)





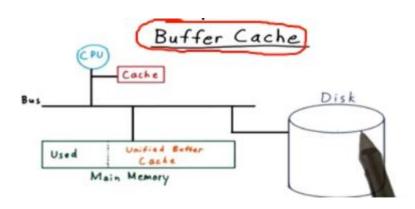
- Hypothesis Extreme situations
  - ✓ If kBlockSize = 16KB, record size = 9KB, num=100



- ✓ Padding is extremely high size
  - Expected performance degradation
- ✓ Write in DB = 900KB, Write in Storage = 1600KB
  - Write Amplification is so high



- Method for decreasing padding size
  - ✓ Predict payload size
  - Determine kBlockSize considering payload and padding size
- Despite of disadvantage, Why kBlockSize is used in Log File?
  - For delayed write in OS, managing static size is efficiently for processing (buffer cache)
  - ✓ Like Paging!!

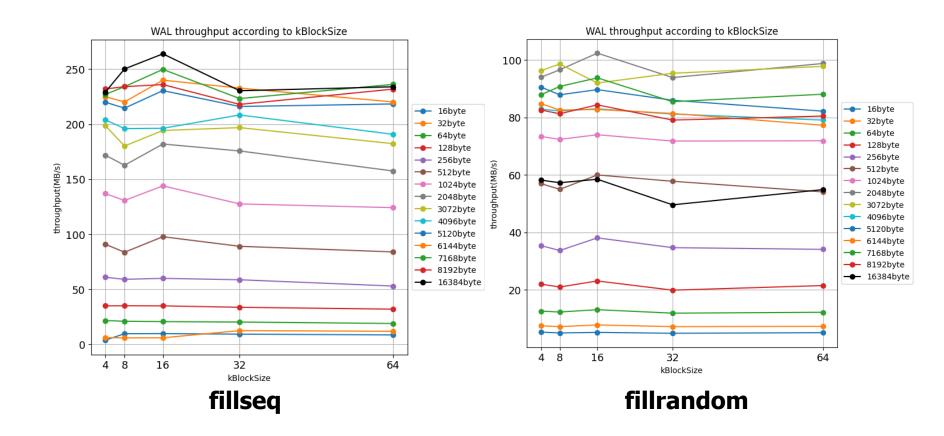






#### Result 1

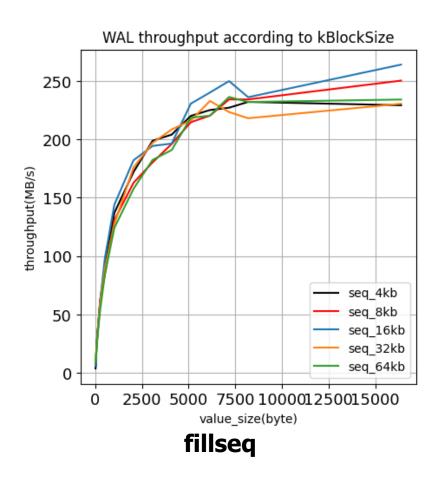
kBlockSizes do not affect WAL overhead

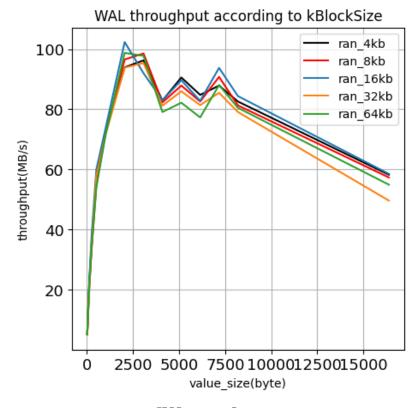




#### Result 2

√ kBlockSizes do not affect WAL overhead





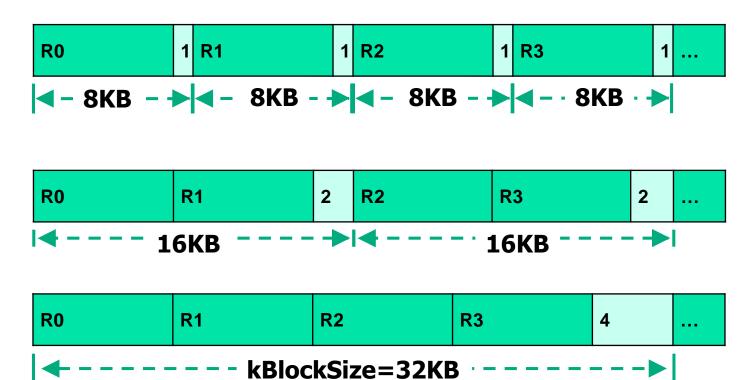






#### Discussion1

- ✓ The sum of padding sizes is constant.
- ✓ Size of record = 7KB





#### Discussion2

- ✓ Padding is not exist!
- Additional experiments to observe padding

```
// is empty, we still want to iterate once to emit a single
// zero-length record
IOStatus s;
bool begin = true;
do {
 const int64 t leftover = kBlockSize - block offset ;
 fprintf(stdout,"leftover : %ld\n", leftover);
  assert(leftover >= 0);
 if (leftover < header size) {</pre>
   // Switch to a new block
   if (leftover > 0) {
     // Fill the trailer (literal below relies on kHeaderSize and
     // kRecyclableHeaderSize being <= 11)
     assert(header size <= 11);
     static_cast<size_t>(leftover)));
```

log\_writer.cc





#### Discussion2

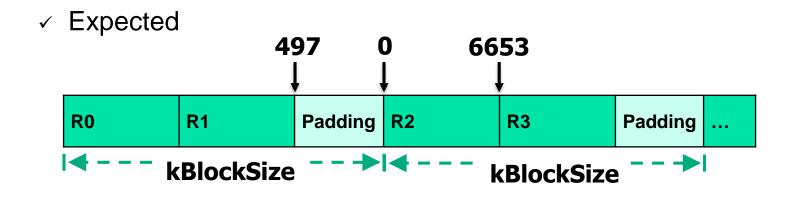
- ✓ Kblocksize: 8KB, Key Size: 16byte, Value Size: 1500byte
- ✓ leftover = kBlockSize block\_offset\_;

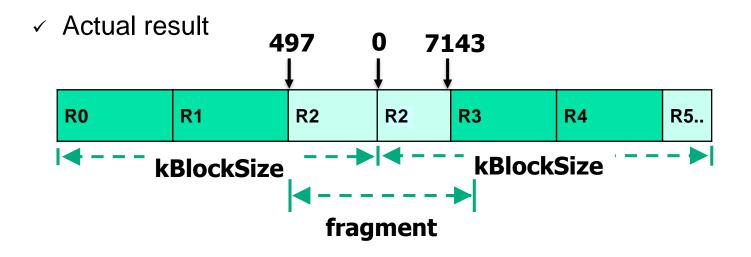
```
DB path: [./mnt]
    leftover: 8192
                        8192 - 6653 = 1539
    leftover : 6653
kBlock-1
                        6653 - 5114 = 1539
                                               Size of record: 1539byte
    leftover : 5114
                        5114 - 3575 = 1539
    leftover : 3575
    leftover
              : 2036
    leftover : 497
                              Fragmentation!
     leftover :
    leftover
             : 7143
    leftover : 5604
                              8192 - 7143 = 1049
kBlock-2
    leftover: 4065
                              497 + 1049 = 1539 + 7 = 1546 byte
    leftover : 2526
    leftover : 987
    leftover :
                              Where is padding??
     leftover : 7633
```





Discussion2 - Is padding existed in kBlock?







# RocksDB Festival: key / value size

- Next assignment
  - ✓ WAL performance according to value or key size.



Size(2B)	Type(1B) Payload (Variable Length)
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- ✓ db\_bench options
  - --disable\_wal=[boolean]
  - --key\_size=[int value]
  - --value\_size=[int value]



# Discussion





