

Lagrangian Relaxation in Integer Programming

Original
Problem
(OP)

$$\begin{aligned} z = \min_{x \geq 0} \quad & c^T x \\ \text{subject to} \quad & Ax \leq b \\ & Dx \leq e \\ & x \geq 0 \\ & x \text{ integer} \end{aligned}$$

Lagrangian
Relaxation Problem
(LR)

$$\begin{aligned} z_{LR}(\lambda) = \min_{x \geq 0} \quad & c^T x + \lambda(Ax - b) \\ \text{subject to} \quad & Dx \leq e \\ & x \geq 0 \\ & x \text{ integer} \end{aligned}$$

z_{LP} : objective function value of linear
relaxation of OP

$$z_{LD} = \max_{\lambda \geq 0} z_{LR}(\lambda) \quad \text{Lagrangian dual
problem}$$

NB: In LR integrality constraint is not relaxed.

Facts:

- $Z_{LP} \leq Z$ because relaxation
- $Z_{LR} \leq Z$ because relaxation
- $Z_{LR} \leq Z_{LO}$ because of definition
- $Z_{LP} \leq Z_{LO}$ This is not trivial and important for motivating the use of Lagrangian Relaxation in Integer Programming:

motivation A: if $Z_{LP} < Z_{LO}$ then LR gives us a better bound to use in B&B

motivation B: if $Z_{LP} = Z_{LO}$ LR can still be worth because Z_{LO} can be found more easily than with LP

*under which conditions does it happen?
See Corollary below.*

motivation C: in any case LR gives us an alternative way to solve the problem. It is a heuristic way with the rare chance of getting also a dual bound and eventually

Proposition:

$$Z_{LP} \leq Z_{LD}$$

Proof:

There are two ways of proving this:

- = via the convexification argument as in the slides and previous classes, also presented in sec 16.4 of [AMO].
- = via the duality argument also presented in sec 8 of [Fi]

I report here the duality arg:

$$Z_{LD} = \max_{\lambda \geq 0} Z_{LR}(\lambda) =$$

$$= \max_{\lambda \geq 0} \left\{ \min_x \left\{ c^T x + \lambda^T (\Delta x - b) \mid \Delta x \leq e, x \geq 0 \text{ & integer} \right\} \right\}$$

$$\text{because}_{LP \text{ relaxation}} \geq \max_{\lambda \geq 0} \left\{ \min_x \left\{ c^T x + \lambda^T (\Delta x - b) \mid \Delta x \leq e, x \geq 0 \right\} \right\} =$$

$$\overbrace{\quad\quad\quad}^{\text{min } c^T x + \lambda^T (\Delta x - b)}$$

$$\mu: \Delta x \leq e$$

$$x \geq 0$$

Dual

$$\lambda^T \Delta + \mu^T \Delta \geq c$$

$$\mu \geq 0$$

$$\lambda \geq 0$$

Hence

$$= \max_{\lambda \geq 0} \left\{ \max_{\mu \geq 0} \left\{ \lambda^T b + \mu^T e \mid \lambda^T A + \mu^T D \geq c, \mu \geq 0 \right\} \right\} =$$

$\max \lambda^T b + \mu^T e$	$\min c^T x$
$x: \lambda^T A + \mu^T D \geq c$	$Ax \leq b$
$\mu \geq 0$	$Dx \leq e$
$\lambda \geq 0$	$x \geq 0$

Dual

$$= \min_x \{c^T x \mid Ax \leq b, Dx \leq e, x \geq 0\} =$$

$$= Z_{LP}$$



Corollary

$Z_{LD} = Z_{LP}$ when the LR problem has the integrality property.

Proof:

The only inequality introduced in the derivations of the proof above becomes equality as well.

