DM872 Math Optimization at Work

Preprocessing

Marco Chiarandini

Department of Mathematics & Computer Science University of Southern Denmark

Outline

1. Preprocessing

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Preprocessing rules

Consider
$$S = \{ \mathbf{x} : a_0 x_0 + \sum_{j=1}^n a_j x_j \le b, l_j \le x_j \le u_j, j = 0..n \}$$

Bounds on variables.
 If a₀ > 0 then:

$$x_0 \le \left(b - \sum_{j:a_j > 0} a_j I_j - \sum_{j:a_j < 0} a_j u_j\right) / a_0$$

and if $a_0 < 0$ then

$$x_0 \ge \left(b - \sum_{j: a_j > 0} a_j l_j - \sum_{j: a_j < 0} a_j u_j\right) / a_0$$

• Redundancy. The constraint $\sum_{i=0}^{n} a_i x_i \leq b$ is redundant if

$$\sum_{j:a_j>0} a_j u_j + \sum_{j:a_j<0} a_j l_j \le b$$

• Infeasibility: $S = \emptyset$ if (swapping lower and upper bounds from previous case)

$$\sum_{j:a_j>0}a_jI_j+\sum_{j:a_j<0}a_ju_j>b$$

• Variable fixing. For a max problem in the form

$$\max\{\boldsymbol{c}^{T}\boldsymbol{x}: A\boldsymbol{x} \leq \boldsymbol{b}, \boldsymbol{l} \leq \boldsymbol{x} \leq \boldsymbol{u}\}$$
 if $\forall i = 1..m: a_{ij} \geq 0, c_{j} < 0$ then fix $x_{j} = l_{j}$ if $\forall i = 1..m: a_{ii} < 0, c_{i} > 0$ then fix $x_{i} = u_{i}$

• Integer variables:

$$\lceil I_j \rceil \leq x_j \leq \lfloor u_j \rfloor$$

ullet Binary variables. Probing: add a constraint, eg, $x_2=0$ and check what happens

Example

$$\begin{array}{l} \max 2x_1 + \ x_2 - x_3 \\ \text{R1} : 5x_1 - 2x_2 + 8x_3 \leq 15 \\ \text{R2} : 8x_1 + 3x_2 - x_3 \geq 9 \\ \text{R3} : \ x_1 + x_2 + x_3 \leq 6 \\ 0 \leq x_1 \leq 3 \\ 0 \leq x_2 \leq 1 \\ x_3 \geq 1 \end{array}$$

R1:
$$5x_1 \le 15 + 2x_2 - 8x_3 \le 15 + 2 \cdot \underbrace{1 - 8 \cdot 1}^{u_2} = 9$$
 $\Rightarrow x_1 \le 9/5$
 $8x_3 \le 15 + 2x_2 - 5x_1 \le 15 + 2 \cdot 1 - 5 \cdot 0 = 17$ $\Rightarrow x_3 \le 17/8$
 $2x_2 \ge 5x_1 + 8x_3 - 15 \ge 5 \cdot 0 + 8 \cdot 1 = -7$ $\Rightarrow x_2 \ge -7/2, x_2 \ge 0$

$$\begin{array}{ll} \text{R2}: 8x_1 \geq 9 - 3x_2 + x_3 \geq 9 - 3 + 1 = 7 & \leadsto x_1 \geq 7/8 \\ \text{R1}: 8x_3 \geq 15 + 2x_2 - 5x_1 \leq 15 + 2 - 5 \cdot 7/8 = 101/8 & \leadsto x_3 \leq 101/64 \end{array}$$

 $\mathtt{R3}: x_1+x_2+x_3 \leq 9/5+1+101/64 < 6 \qquad \text{Hence R3 is redundant}$

Example

$$\begin{array}{l} \max 2x_1 + x_2 - x_3 \\ \text{R1} : 5x_1 - 2x_2 + 8x_3 \leq 15 \\ \text{R2} : 8x_1 + 3x_2 - x_3 \geq 9 \\ 7/8 \leq x_1 \leq 9/5 \\ 0 \leq x_2 \leq 1 \\ 1 \leq x_3 \leq 101/64 \end{array}$$

Increasing x_2 makes constraints satisfied $\rightsquigarrow x_2=1$ Decreasing x_3 makes constraints satisfied $\rightsquigarrow x_3=1$

We are left with:

$$\max\{2x_1: 7/8 \le x_1 \le 9/5\}$$

Preprocessing for Set Covering/Partitioning

1. if $e_i^T A = 0$ then the *i*th row can never be satisfied

2. if $e_i^T A = e_k$ then $x_k = 1$ in every feasible solution

$$\begin{bmatrix} 0 & 0 & \dots & 1 & \dots & 0 \end{bmatrix} \begin{bmatrix} -\frac{1}{0} & \frac{1}{0} & \frac{1}{0} & \frac{1}{0} \\ -\frac{1}{0} & \dots & \frac{1}{1} & \dots & 0 \\ -\frac{1}{0} & \dots & \frac{1}{1} & \dots & 0 \end{bmatrix} = \begin{bmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \\ 0 \end{bmatrix}$$

In SPP can remove all rows t with $a_{tk}=1$ and set $x_j=0$ (ie, remove cols) for all cols that cover t

3. if $e_t^T A \ge e_p^T A$ then we can remove row t, row p dominates row t (by covering p we cover t)



if ∑_{j∈S} Ae_j = Ae_k and ∑_{j∈S} c_j ≤ c_k then we can cover the rows by Ae_k more cheaply with S and set x_k = 0
 (Note, we cannot remove S if ∑_{j∈S} c_j ≥ c_k)

$$\left[\begin{array}{c|cccc} 1 & & & 1 \\ 1 & & & 1 \\ & 1 & & 1 \\ & 0 & 0 & & 0 \\ 1 & & & 1 \\ & 0 & 0 & & 0 \end{array}\right]$$

Preprocessing

Hot Topic

- MIP challenge 2024 on Presolving Reductions https://github.com/dominiqs81/MIPcc24
- Papilo, parallel presolving https://github.com/dominiqs81/MIPcc24

Preprocessing

Summary

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