

# Course Timetabling

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# University Course Timetabling Research

- Course Timetabling [ITC-2002, ITC-2007], [Di Gaspero et al. (2007)], [Bettinelli et al. (2015)]  
Curriculum-Based (CB-CTT)  $\rightsquigarrow$  Post-Enrolment-Based (PE-CTT)
- Solution Approaches: Modeling approach + solving algorithms:
  - Direct representation  
branch and bound, construction heuristics, metaheuristic methods
  - Mixed Integer Linear Programming (MILP)  
branch-and-cut
  - Satisfiability problems (SAT)  
backtracking-based algorithms
  - Constraint Satisfaction Problems (CSP)  
constraint propagation + backtracking-based algorithms

# University Course Timetabling Research

Common assumptions of CB-CTT and PE-CTT: weekly periodicity + classes of equal length

Inadequate in many institutions (SDU included)

- different requirements for the number of classes in each week of the semester
- different duration for each single class
- precedence constraints among classes, eg, introduction classes precede exercise classes

ITC2019: Course Timetabling Competition 2019 ([www.itc2019.org](http://www.itc2019.org)) including these features and student sectioning.

In the next slides we focus on the problem at SDU.

Open question: is there a possible reduction between the SDU problem and the problem at ITC2019?

# Outline

1. Problem Description

2. MILP Approach

# Timetabling at SDU

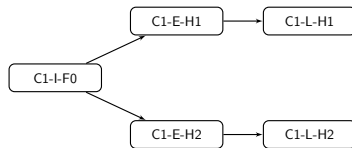
The timeline of timetabling activities for the Spring semester

week 37	39	40-52	49-50	52-2	6-23
Room negotiation	Teachers communicate organization of classes	Timetabling mandatory courses (CB-CTT)	Students register to courses	Timetabling elective courses (PE-CTT)	Semester

# Problem Description: Input

- A set of periods (timeslots)  $P = \{(h, d, w) \mid h \in \text{Hours}, d \in \text{Days}, w \in \text{Weeks}\}$
- A set of events  $E$  each event with a duration  $\ell(e)$ ,  $e \in E$  + precedence digraph  $D = (V, A)$

Course	Session	Section	Event
C1	Intro	F	C1-I-F
		H1	C1-E-H1
	Exercises	H2	C1-E-H2
		H1	C1-L-H1
		H2	C1-L-H2
		H2	C1-L-H2
	CompLab	H1	C1-L-H1

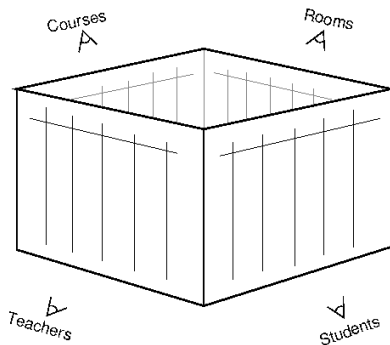


- A set of rooms  $R$  (seminar room + dummy)
- People: a set of students  $S$ , a set of teachers  $T$ 
  - A collection of enrollments  $\mathcal{Q} = \{E_s \subset E \mid s \in S\}$  that are events a student has subscribed to (post enrollment model)
  - A collection of teaching duties  $\mathcal{D} = \{D_t \subset E \mid t \in T\}$
  - Teacher unavailabilities  $\mathcal{U} = \{U_t \subset P \mid t \in T\}$
- Schedule of mandatory courses  $M = \{(e, r, p) \mid e \in E, r \in R, p \in P\} \equiv \text{preassignments}$

# Problem Description: Task

Schedule events in the semester such that the timetable is feasible and appealing from different **perspectives** (resources)

	Timeslots														
	8	9	10	11	12	13	14	15	16	17	8	9	...		
1	1	2					3					4			
2			6						5				7		
3						8									
4		9		10			11					12			
..		13						14			15				



# Problem Description

## Constraints (hard constraints)

Enforce All Events Scheduled

Prevent Room Conflicts

Prevent Staff Conflicts

Enforce Fixed Preassignments

Enforce Fixed Rooms

Enforce Max One Event  $\times$  Day  $\times$  Crs

Enforce Precedences

Enforce Banned Slots

Enforce Pairings

## Objective(s) (soft constraints):

Weekly Stability

Usage of seminar room

Student/Instructor Conflicts

Events  $\times$  Day  $\times$  Tch

Bad Slots

The classical approach:

Violations of these criteria are **penalized** with *appropriately* chosen **weights** in a (single) objective function to minimize

Collective welfare approach?



# Outline

1. Problem Description

2. MILP Approach

# Mixed Integer Linear Programming Formulations (1/2)

- Compact MILP formulation (aka Monolithic) [Burke et al. (2008, 2010a)]  
variables:  $x_{erp} \in \{0, 1\}$  whether event  $e \in E$  placed in room  $r \in R$  in timeslot  $p \in P$
- Two-stage formulation [Lach and Lübbecke (2012)]
  - ① Variables:  $x_{ep} \in \{0, 1\}$  whether event  $e \in E$  placed in timeslot  $p \in P$
  - ② One-sided perfect matching in a bipartite graph with additional constraints
- Divide-and-conquer approach [Hao and Benlic (2011)]
  - ① Generate a partition  $E_i$  of the set of events.
  - ② Solve each subproblem  $P(E_i), i = 1, \dots, k$ , with a MILP solver to compute lower bound  $LB_i$ .
  - ③ Sum up the  $k$  values  $LB_i$  to achieve the final lower bound to CB-CTT.
- Column generation approach [Cacchiani et al. (2013)]: Two Weekly Schedule Types
  - ① A vector  $\mathbf{x} \in B^{|E| \times |R| \times |P|}$  made of components  $x_{erp} \in \mathbb{B}$  representing the assignment of event  $e \in E$  to room  $r \in R$  at timeslot  $p \in P$ .  
Used to consider penalties for room capacity and room stability.
  - ② A vector  $\boldsymbol{\theta} \in B^{|E| \times |P|}$  made up of components  $\theta_{ep} \in \mathbb{B}$  that indicates if an event  $e \in E$  is scheduled at timeslot  $p \in P$ .  
Used to consider penalties for curriculum compactness and minimum working days.

# Mixed Integer Linear Programming Formulations (2/2))

- A Resource Constrained Project Scheduling Model

- ① Variables: Time indexed variables of starting time of an event

$$x_{erp} \in \{0, 1\} \quad \forall e \in E, r \in R, p = (h, d, w) \in P$$

denote whether an event  $e \in E$  is located in room  $r \in R$  and starts at timeslot  $p = (h, d, w) \in P$

- Patterns

- ① Patterns are generated: In a pattern all meetings of a class start at the same time, run for the same number of slots and are placed in the same room.
- ② Variables  $x_{ci} \in \{0, 1\}$  selects a pattern for a class,  $y_{cr}$  selects a room for a class.

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```
<class id="40" limit="34" parent="39">
  <room id="16" penalty="4"/>
  <room id="21" penalty="0"/>
  <room id="22" penalty="4"/>
  <room id="3" penalty="0"/>
  <room id="13" penalty="0"/>
  <room id="25" penalty="4"/>
  <room id="27" penalty="0"/>
  <room id="7" penalty="0"/>
  <room id="17" penalty="0"/>
  <time days="0001000" start="96" length="22" weeks="01111111111110" penalty="0"/>
  <time days="0000100" start="96" length="22" weeks="01111011111110" penalty="0"/>
  <time days="0001000" start="120" length="22" weeks="01111111111110" penalty="6"/>
  <time days="0000100" start="120" length="22" weeks="01111011111110" penalty="0"/>
  <time days="0001000" start="144" length="22" weeks="01111111111110" penalty="6"/>
  <time days="0000100" start="144" length="22" weeks="01111011111110" penalty="0"/>
  <time days="0001000" start="168" length="22" weeks="01111111111110" penalty="6"/>
  <time days="0000100" start="168" length="22" weeks="01111011111110" penalty="2"/>
  <time days="0001000" start="192" length="22" weeks="01111111111110" penalty="0"/>
  <time days="0000100" start="192" length="22" weeks="01111011111110" penalty="8"/>
```

# The Resource Constrained Project Scheduling Model

## Rooms

Hard constraints:

- Prevent Room Conflicts + Room Availability + Suitable Events

$$\sum_{e \in E} \sum_{h=\max\{8, s-\ell(e)\}}^s x_{erhdw} \leq a_{rhdw} \quad \forall r \in R, s = 8, \dots, 17$$

	Monday																Tuesday															
	16	17	8	9	10	11	12	13	14	15	16	17	8	9	10	11	12	13	14	15	16	17										
Event 1																																
Event 2																																
Event 3																																
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$$x_{1,\cdot,1,10,w} + x_{1,\cdot,1,11,w} + x_{2,\cdot,1,10,w} + x_{2,\cdot,1,11,w} + x_{3,\cdot,1,9,w} + x_{3,\cdot,1,10,w} + x_{3,\cdot,1,11,w} \leq 1$$