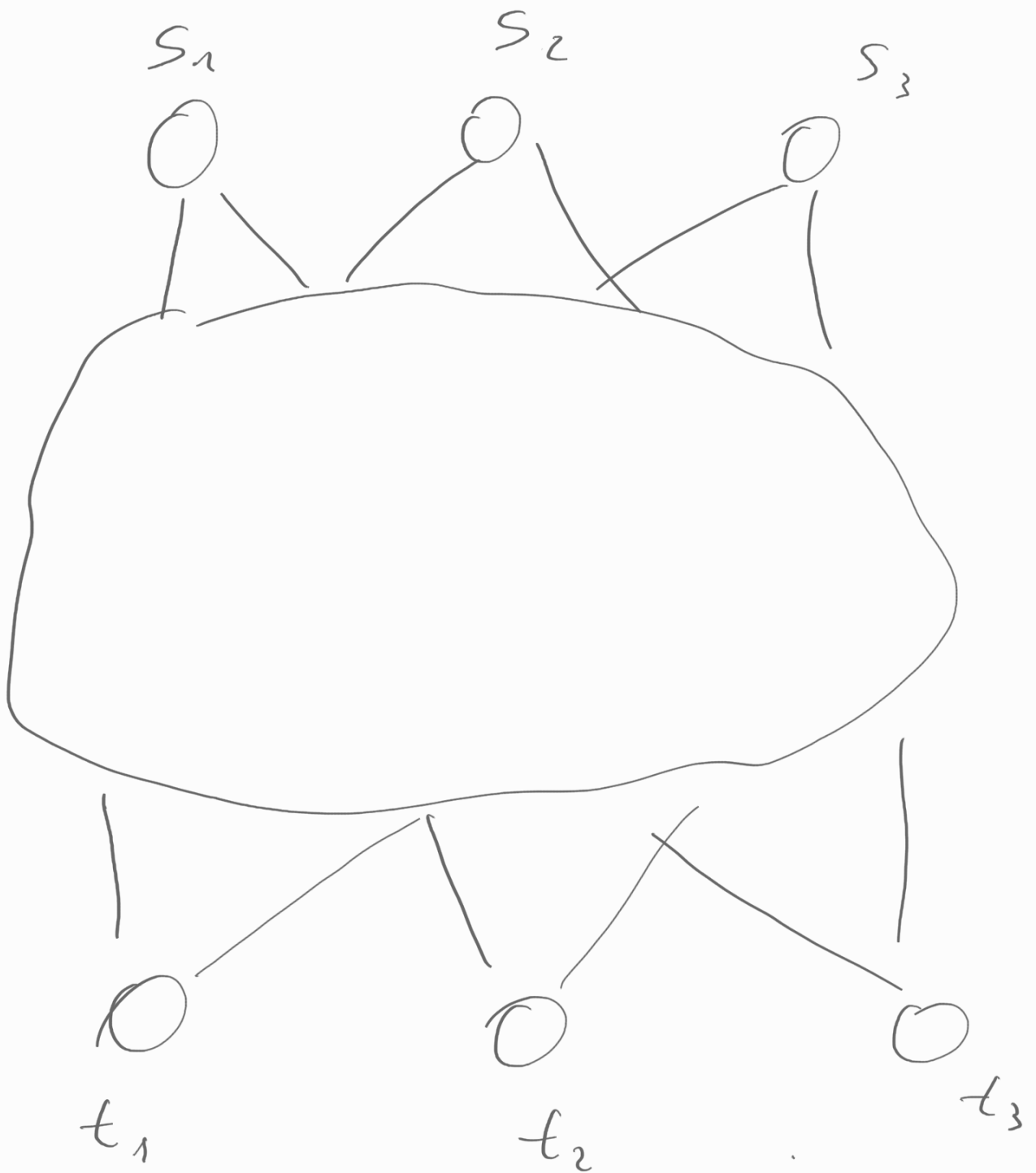


MULTICOMMODITY NETWORK FLOW



$$\min c^T x$$

$$Ax = b$$

$$0 \leq x \leq u$$

$$|z| = \sum_{k=1}^n c_k^T x^k$$

$$z = \min_k \sum_k c_k x_k$$

$$N x^k = b$$

$$\forall k = 1..3$$

$$\sum_k x_{ij}^k \leq u_{ij}$$

$$\forall k = 1..3 \quad ij \in A$$

$$0 \leq x_{ij}^k \leq u_{ij}^k$$

$$\forall ij \in A \quad k = 1..3$$

(x_{ij}^k is integral)

Motivations for
Lagrangian Relaxation

$$\sum_k x_{ij}^k + s_{ij} = u_{ij}$$

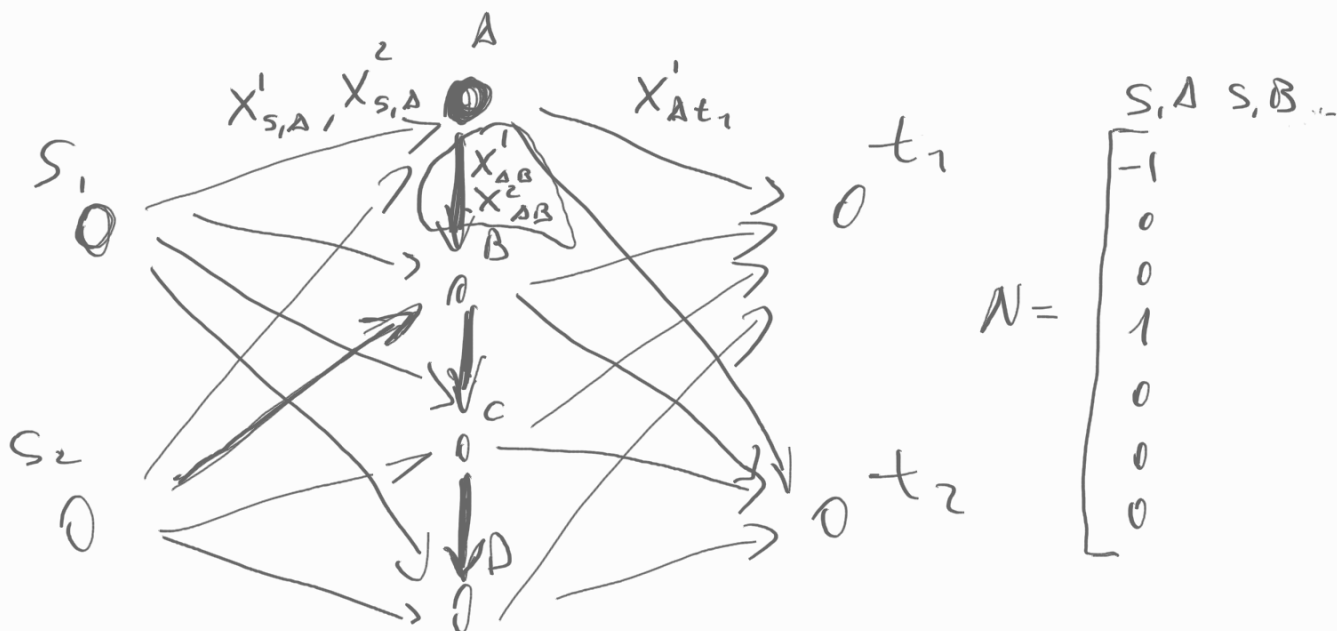
Divisible goods \Rightarrow Linear Prog. Problem

- Solve the problem without simplex and possibly faster

Indivisible goods \Rightarrow Integer Lin Prog. Problem

- As a banding procedure in B2B
- To heuristic solutions

Application example: Warehouse problem



$$\Delta: X'_{s_2A} + X'_{s_1A} - X_{AR} - X_{At_1} - X_{At_2} = 0$$

$$\Delta B: X'_{AB} + X'_{AB} \leq u_{AB}$$

$$X'_{s_2A} + X'_{s_2A} \leq u_{s_1A}$$

• LR(λ)

Lagrangian
Relaxation
problem
(on x vars)

$$z_{LR1} = \min_x d^T x^n$$

$$\sum_k x_{ij}^k \leq u_{ij}$$

$$0 \leq x_{ij}^k \leq u_{ij}^k$$

$$z_{LR2} = \min_x$$

$$e^T x^k$$

$$N x^k = b$$

$$0 \leq x_{ij}^k \leq u_{ij}^k$$

$$\lambda \geq 0$$

$$c^T x^k + \lambda (\sum x^k - u)$$

$$\underbrace{x^k (c^T + \lambda)}_{e^T} - \underbrace{\lambda u}_{const}$$

Two alternative
relaxations.
The second one
leads to k
shortest paths
problems and
hence seems
easier to solve

• LD

$$z_{LD} = \max_{\lambda \geq 0} LR(\lambda)$$

Lagrangian
Dual Problem
(on λ vars)

Relations between z , z_{LR} , z_{LD}

$$z \perp$$

$$z_{LR} \leq z_{LD} \quad \left(\text{since it is the max of } z_{LR}(\lambda) \right)$$

$$\begin{array}{c} z_{LD} \\ z_{LR}(\lambda) \end{array}$$

$$z_{LD} \leq z$$

Proof:

Assume x^* optimal sol for the original problem Then

$$\sum_n (x_{ij}^*)^n \leq u_{ij}$$

and

$$z_{LR}(\lambda) = c^T(x^*)^n + \lambda \underbrace{\left(\sum_n (x_{ij}^*)^n - u_{ij} \right)}_{\substack{\geq 0 \\ \leq 0}} \leq c^T(x^*)^n = z$$

So $z_{LR}(\lambda) \leq z$ for all $\lambda \geq 0$
and hence also

$$z_{LD} \leq z$$

The relationship above holds always.

In some cases $z_{LD} = z$. Two of these cases are:

- 1) when the original problem is an LP problem
- 2) when the original problem is an ILP problem and the Lagrangian relaxation problem has the integrality property.

Proof of 1) (SMO Th 16.6)

Assume x^* is opt sol for original problem
Opt conditions: complementary slackness th.

$$\pi^* \left(\sum_n x^{*n} - u \right) = 0$$

↑
dual vars associated with linking constr.

Consider $LR(\lambda)$ with $\lambda = \pi^*$. x^* is feasible in LR .

$$z_{LR}(\pi^*) = C^T x^* + \pi^* \left(\sum_n x^{*n} - u \right) = C^T x^* = z^*$$

Proof of 2)

A consequence of 1).

