

# Vehicle Scheduling: Models and Algorithms

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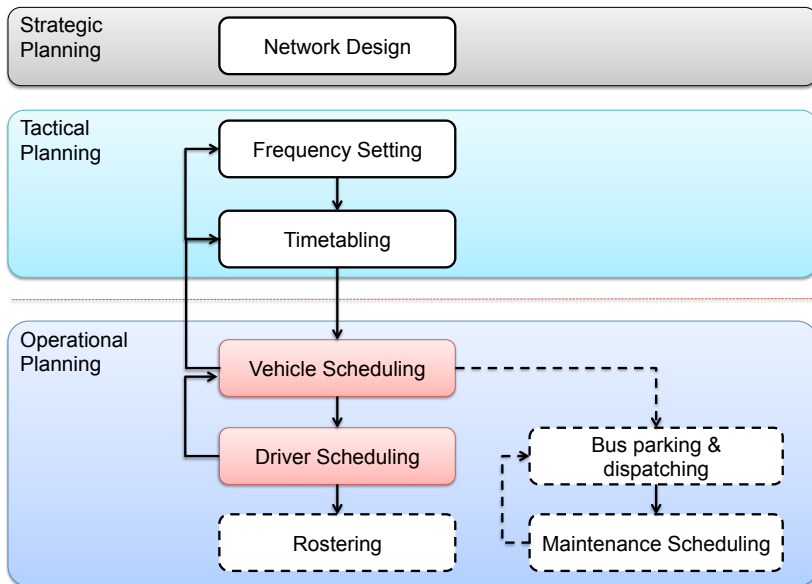
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## 1 Introduction

## 2 Vehicle Scheduling (VS)

# Overview of Planning Activities

(Desaulniers&Hickman2007)



# Strategic Planning: Network Design (Urban)



# Strategic Planning: Network Design (Regional)



# Tactical Planning: Frequency Setting and Timetabling

41

42

Odense Banegård - Syddansk Universitet (SDU)

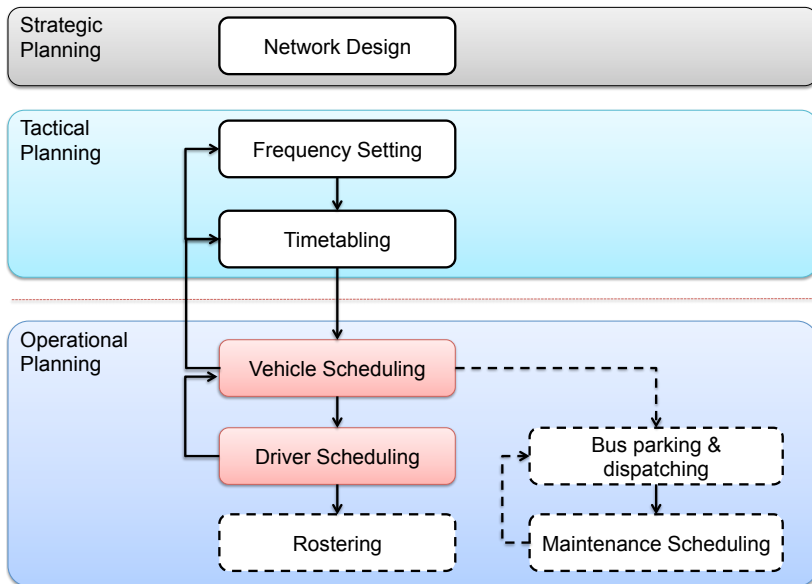
Hverdag

Rullevr.	OBC Plads B	Hans Mules Gade	Nyborgvej / Frederikssgade	Palmekæsevej	Nansenegade	Epbygade	L.A. Ringe Vej	Rosengårdc. / Gul Indg.	Blåka	Rosengårdc. / Østbævej	Niels Bohr's Allé	Campusvej	SDU
42	05.25	05.26	05.28	05.29	05.30	05.31	05.33			05.34	05.36	05.37	05.40
41	06.23	06.24	06.26	06.27	06.28	06.29	06.31	06.33	06.35			06.37	06.40
41	07.20	07.21	07.23	07.24	07.25	07.26	07.29	07.31	07.34			07.36	07.40
42	07.28	07.30	07.32	07.33	07.34	07.35	07.39			07.41	07.44	07.46	07.50
42	07.36	07.38	07.40	07.41	07.42	07.44	07.46			07.48	07.51	07.53	07.56
42	07.44	07.46	07.48	07.49	07.50	07.51	07.55			07.57	08.00	08.02	08.06
42	07.52	07.54	07.56	07.56	07.57	07.58	08.01			08.03	08.06	08.07	08.12
151P	08.00	08.02	08.04	08.04	08.05	08.06	08.09			08.11	08.14	08.15	08.20
41	08.10	08.12	08.14	08.15	08.16	08.18	08.21	08.23	08.27			08.29	08.33
41	08.25	08.27	08.29	08.30	08.31	08.32	08.35	08.37	08.41			08.43	08.47
41	08.35	08.37	08.39	08.40	08.41	08.42	08.45	08.47	08.51			08.53	08.57
41	08.45	08.47	08.49	08.50	08.51	08.52	08.55	08.57	09.01			09.03	09.07
41	08.55	08.57	08.59	09.00	09.01	09.02	09.05	09.07	09.11			09.13	09.17
41	09.05	09.07	09.09	09.10	09.11	09.12	09.15	09.17	09.21			09.23	09.27
41	09.15	09.17	09.19	09.20	09.21	09.22	09.25	09.27	09.31			09.33	09.37

OBS:  
Ikke gyldig mellem  
13. maj - 23. august.  
Se sommerkøreplanen  
på side 70-71.

# Overview of Planning Activities

(Desaulniers&Hickman2007)



# Leuthardt Survey

(Leuthardt 1998, Kostenstrukturen von Stadt-, Überland- und Reisebussen, DER NAHVERKEHR 6/98, pp. 19-23.)

<i>bus costs (DM)</i>	<i>urban</i>	<i>%</i>	<i>regional</i>	<i>%</i>
crew	349,600	73.5	195,000	67.5
depreciation	35,400	7.4	30,000	10.4
calc. interest	15,300	3.2	12,900	4.5
materials	14,000	2.9	10,000	3.5
fuel	22,200	4.7	18,000	6.2
repairs	5,000	1.0	5,000	1.7
other	34,000	7.1	18,000	7.2
total	475,500	100.0	288,900	100.0



## 1 Introduction

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## Smart models start small

Posted on **SEPTEMBER 9, 2013** Written by **MARC-ANDRE**  [LEAVE A COMMENT](#)

There is only one good way to build large-size or complex optimization models: to start by a small model and adding elements gradually until you get the model you wanted in the first place. I have seen so many people (including myself) try to build large-size, complex models from scratch, only to spend countless frustrating hours trying to debug all kinds of problems. It just doesn't work.

A better approach is to start with the simplest version of the model. On or two

# Vehicle Scheduling

Given a timetable as a set  $V = \{v_1, \dots, v_n\}$  of **trips**, where for each trip  $v_i$  we have:

$t_i$  : departure time

$a_i$  : arrival time

$o_i$  : origin (departure terminal)

$d_i$  : destination (arrival terminal)

$v_i$	$t_i$	$a_i$	$o_i$	$d_i$
$v_1$	7:10	7:30	$T_a$	$T_b$
$v_2$	7:20	7:40	$T_c$	$T_d$
$v_3$	7:40	8:05	$T_b$	$T_a$
$v_4$	8:00	8:30	$T_d$	$T_c$
$v_5$	8:35	9:05	$T_c$	$T_d$

Given the **deadheading trips** (i.e. trips without passengers) of duration  $h_{ij}$  between every pair of terminals

$h_{ij}$	$T_a$	$T_b$	$T_c$	$T_d$
$T_a$	0	15	20	20
$T_b$	15	0	25	10
$T_c$	20	25	0	15
$T_d$	20	10	15	0

## Definition (Compatible Trips)

A pair of trips  $(v_i, v_j)$  is **compatible** if and only if  $a_i + h_{ij} \leq t_j$

# Vehicle Scheduling

## Definition (Vehicle Duty)

A subset  $C = \{v_{i_1}, \dots, v_{i_k}\}$  of  $V$  is a **vehicle duty (or block)** if  $(v_{i_j}, v_{i_{j+1}})$  is a **compatible pair of trips**, for  $j = 1, \dots, k - 1$

## Definition (Vehicle Schedule)

A collection  $C_1, \dots, C_r$  of *vehicle duties* such that each trip  $v$  in  $V$  belongs to exactly one  $C_j$  with  $j \in \{1, \dots, r\}$  is said to be a **Vehicle Schedule**

# Vehicle Scheduling: Example

$v_i$	$t_i$	$a_i$	$o_i$	$d_i$
$v_1$	7:10	7:30	$T_a$	$T_b$
$v_2$	7:20	7:40	$T_c$	$T_d$
$v_3$	7:40	8:05	$T_b$	$T_a$
$v_4$	8:00	8:30	$T_d$	$T_c$
$v_5$	8:35	9:05	$T_c$	$T_d$

$h_{ij}$	$T_a$	$T_b$	$T_c$	$T_d$
$T_a$	0	15	20	20
$T_b$	15	0	25	10
$T_c$	20	25	0	15
$T_d$	20	10	15	0

**Example:** These 5 trips can be scheduled with 2 vehicle duties:

- $C_1 = \{v_1, v_3\}$
- $C_2 = \{v_2, v_4, v_5\}$

## Further Features of the Problem

- Limited number of vehicles
- Minimize fleet size (number of vehicles)
- Minimize operational costs (given by pull-out and pull-in from depots and deadheading trips)
- Multiple depots
- Different types of vehicles with different operational costs located at a single depot