

# Conjugacy in Higman-Thompson groups

David Robertson

15<sup>th</sup> April, 2015

# Crash course on groups

- ▶ **Group**: set with a nice binary operation
  - ▶  $\mathbb{Z}$ ,  $\mathbb{R}$ ,  $\mathbb{C}$  with addition
  - ▶ non zero reals, invertible matrices with multiplication
  - ▶ invertible maps  $f: X \rightarrow X$  with function composition

# Crash course on groups

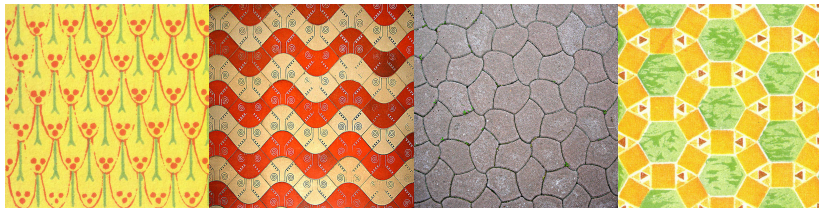
- ▶ **Group**: set with a nice binary operation
  - ▶  $\mathbb{Z}, \mathbb{R}, \mathbb{C}$  with addition
  - ▶ non zero reals, invertible matrices with multiplication
  - ▶ invertible maps  $f: X \rightarrow X$  with function composition
- ▶ **'Nice'**: operation has identity  $e$  and inverses  $x^{-1}$ 
  - ▶  $ex = xe = x$  e.g.  $0, 1, x \mapsto x$
  - ▶  $xx^{-1} = x^{-1}x = e$  e.g.  $-x, 1/x, f(x) \mapsto x$
  - ▶  $x(yz) = (xy)z$

# Crash course on groups

- ▶ **Group**: set with a nice binary operation
  - ▶  $\mathbb{Z}, \mathbb{R}, \mathbb{C}$  with addition
  - ▶ non zero reals, invertible matrices with multiplication
  - ▶ invertible maps  $f: X \rightarrow X$  with function composition
- ▶ **'Nice'**: operation has identity  $e$  and inverses  $x^{-1}$ 
  - ▶  $ex = xe = x$  e.g.  $0, 1, x \mapsto x$
  - ▶  $xx^{-1} = x^{-1}x = e$  e.g.  $-x, 1/x, f(x) \mapsto x$
  - ▶  $x(yz) = (xy)z$
- ▶ Used (e.g.) to describe/measure **symmetry**

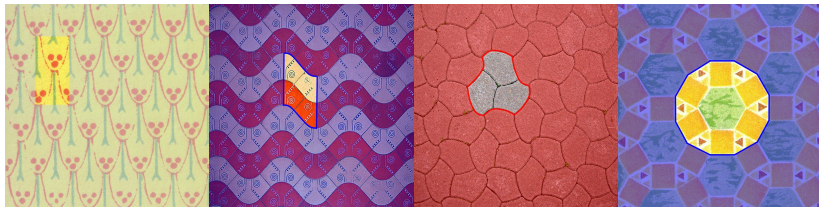
# Crash course on groups

- ▶ **Group**: set with a nice binary operation
  - ▶  $\mathbb{Z}$ ,  $\mathbb{R}$ ,  $\mathbb{C}$  with addition
  - ▶ non zero reals, invertible matrices with multiplication
  - ▶ invertible maps  $f: X \rightarrow X$  with function composition
- ▶ **'Nice'**: operation has identity  $e$  and inverses  $x^{-1}$ 
  - ▶  $ex = xe = x$  e.g.  $0, 1, x \mapsto x$
  - ▶  $xx^{-1} = x^{-1}x = e$  e.g.  $-x, 1/x, f(x) \mapsto x$
  - ▶  $x(yz) = (xy)z$
- ▶ Used (e.g.) to describe/measure **symmetry**



# Crash course on groups

- ▶ **Group**: set with a nice binary operation
  - ▶  $\mathbb{Z}$ ,  $\mathbb{R}$ ,  $\mathbb{C}$  with addition
  - ▶ non zero reals, invertible matrices with multiplication
  - ▶ invertible maps  $f: X \rightarrow X$  with function composition
- ▶ **'Nice'**: operation has identity  $e$  and inverses  $x^{-1}$ 
  - ▶  $ex = xe = x$  e.g.  $0, 1, x \mapsto x$
  - ▶  $xx^{-1} = x^{-1}x = e$  e.g.  $-x, 1/x, f(x) \mapsto x$
  - ▶  $x(yz) = (xy)z$
- ▶ Used (e.g.) to describe/measure **symmetry**



# The conjugacy problem

Given a group  $G$  and elements  $x, y$  can you find a **conjugator**  $z$ ?

$$z^{-1}xz = y$$

# The conjugacy problem

Given a group  $G$  and elements  $x, y$  can you find a **conjugator**  $z$ ?

$$z^{-1}xz = y$$

[ $G, x, y$  given in terms of a presentation]

- ▶ This problem is **undecidable** (P. Novikov, 1954)
- ▶ Can't write a single program which correctly answers "Yes" or "No" in finite time for **every** group  $G$ .



# The conjugacy problem

Given a group  $G$  and elements  $x, y$  can you find a **conjugator**  $z$ ?

$$z^{-1}xz = y$$

[ $G, x, y$  given in terms of a presentation]

- ▶ This problem is **undecidable** (P. Novikov, 1954)
- ▶ Can't write a single program which correctly answers "Yes" or "No" in finite time for **every** group  $G$ .
- ▶ Cryptosystem?
- ▶ Can do this for **specific** groups, *i.e.* with  $G$  fixed.

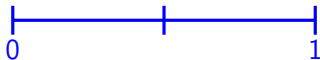
## A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .



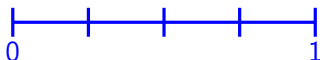
## A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .
- ▶ Chop it in half.



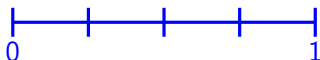
## A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .
- ▶ Chop it in half.
- ▶ If you like, chop some of the halves in half.



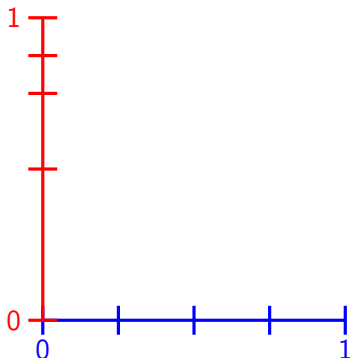
## A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .
- ▶ Chop it in half.
- ▶ If you like, chop some of the halves in half.
- ▶ Continue until you get bored.



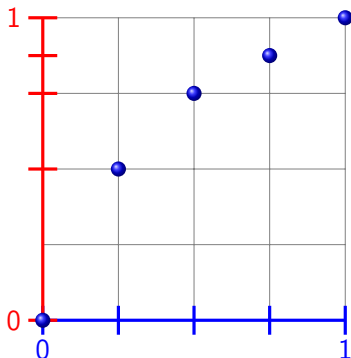
# A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .
  - ▶ Chop it in half.
  - ▶ If you like, chop some of the halves in half.
  - ▶ Continue until you get bored.
- 
- ▶ Do the same to a second copy of  $[0, 1]$ . Same number of chops!



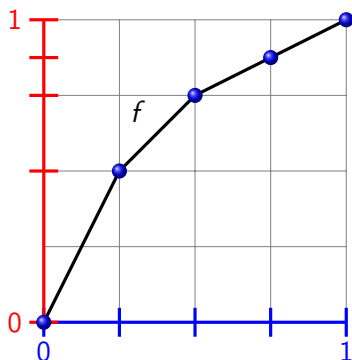
# A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .
  - ▶ Chop it in half.
  - ▶ If you like, chop some of the halves in half.
  - ▶ Continue until you get bored.
- 
- ▶ Do the same to a second copy of  $[0, 1]$ . Same number of chops!
  - ▶ Use the intervals as axes and join the dots.



# A recipe for $F$

- ▶ Take the interval  $[0, 1]$ .
- ▶ Chop it in half.
- ▶ If you like, chop some of the halves in half.
- ▶ Continue until you get bored.
- ▶ Do the same to a second copy of  $[0, 1]$ . Same number of chops!
- ▶ Use the intervals as axes and join the dots.

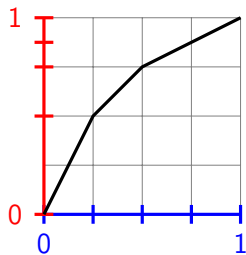


Functions  $f$  like this  
are the elements of  
Thompson's group  $F$ .



## Thompson's other groups $T$ and $V$

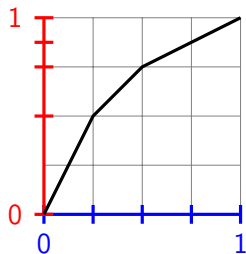
$F$



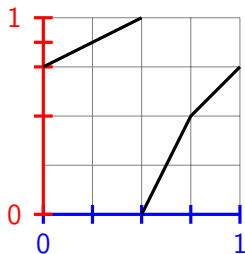
$F$ : increasing  
functions

## Thompson's other groups $T$ and $V$

$$F < T$$



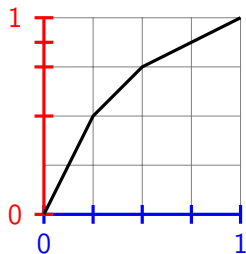
$F$ : increasing  
functions



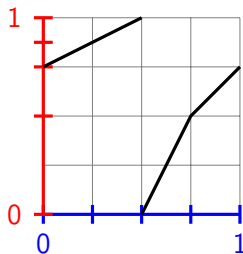
$T$ : don't have to start  
at  $(0,0)$

# Thompson's other groups $T$ and $V$

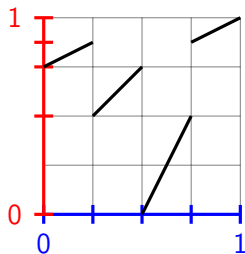
$$F < T < V$$



$F$ : increasing functions



$T$ : don't have to start at  $(0, 0)$



$V$ : don't have to be continuous

## People find these groups interesting

$T$  and  $V$  are finitely presented, infinite simple groups (**rare!**)

- ▶ **Finitely Presented:** described by finitely many equations.
- ▶ **Simple:** can't 'compress' the group to get a smaller one.

## People find these groups interesting

$T$  and  $V$  are finitely presented, infinite simple groups (**rare!**)

- ▶ **Finitely Presented:** described by finitely many equations.
- ▶ **Simple:** can't 'compress' the group to get a smaller one.

$V$  contains a copy of every finite group  $G$ :

- ▶  $G \hookrightarrow \mathcal{S}_n \hookrightarrow \mathcal{S}_{2^m}$
- ▶ Take two full binary trees of depth  $m$ , each with  $2^m$  leaves.
- ▶  $\sigma \in \mathcal{S}_{2^m}$  corresponds to a permutation of the leaves.

# People find these groups interesting

$T$  and  $V$  are finitely presented, infinite simple groups (**rare!**)

- ▶ **Finitely Presented:** described by finitely many equations.
- ▶ **Simple:** can't 'compress' the group to get a smaller one.

$V$  contains a copy of every finite group  $G$ :

- ▶  $G \hookrightarrow \mathcal{S}_n \hookrightarrow \mathcal{S}_{2^m}$
- ▶ Take two full binary trees of depth  $m$ , each with  $2^m$  leaves.
- ▶  $\sigma \in \mathcal{S}_{2^m}$  corresponds to a permutation of the leaves.

People are also interested in whether  $F$  is amenable or not,

**whatever that means...**

# Conjugacy algorithms in $V$

1974 Higman described a conjugacy algorithm in a series of lectures.

- ▶ Works for  $V = G_{2,1}$  and generalisations  $G_{n,r}$ .

# Conjugacy algorithms in $V$

1974 Higman described a conjugacy algorithm in a series of lectures.

- ▶ Works for  $V = G_{2,1}$  and generalisations  $G_{n,r}$ .
- ▶ Different tools used to address conjugacy since:

1997 Guba, Sapir: [diagram groups](#)

2007 Belk, Matucci: [strand diagrams](#)

2010 Salazar-Díaz: [revealing tree pairs](#)

2011 Bleak *et al.*: [train tracks and flow graphs](#)



# Conjugacy algorithms in $V$

1974 Higman described a conjugacy algorithm in a series of lectures.

- ▶ Works for  $V = G_{2,1}$  and generalisations  $G_{n,r}$ .
- ▶ Different tools used to address conjugacy since:

1997 Guba, Sapir: [diagram groups](#)

2007 Belk, Matucci: [strand diagrams](#)

2010 Salazar-Díaz: [revealing tree pairs](#)

2011 Bleak *et al.*: [train tracks and flow graphs](#)

2014 Barker used Higman's ideas to solve **power** conjugacy in  $V$ .

# Conjugacy algorithms in $V$

1974 Higman described a conjugacy algorithm in a series of lectures.

- ▶ Works for  $V = G_{2,1}$  and generalisations  $G_{n,r}$ .
- ▶ Different tools used to address conjugacy since:

1997 Guba, Sapir: [diagram groups](#)

2007 Belk, Matucci: [strand diagrams](#)

2010 Salazar-Díaz: [revealing tree pairs](#)

2011 Bleak *et al.*: [train tracks and flow graphs](#)

2014 Barker used Higman's ideas to solve **power** conjugacy in  $V$ .

Given  $x, y$  can you solve  $z^{-1}x^a z = y^b$ ?

# Conjugacy algorithms in $V$

1974 Higman described a conjugacy algorithm in a series of lectures.

- ▶ Works for  $V = G_{2,1}$  and generalisations  $G_{n,r}$ .
- ▶ Different tools used to address conjugacy since:

1997 Guba, Sapir: [diagram groups](#)

2007 Belk, Matucci: [strand diagrams](#)

2010 Salazar-Díaz: [revealing tree pairs](#)

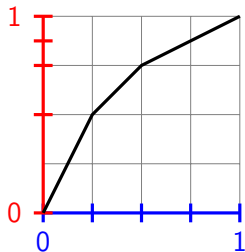
2011 Bleak *et al.*: [train tracks and flow graphs](#)

2014 Barker used Higman's ideas to solve **power** conjugacy in  $V$ .

Given  $x, y$  can you solve  $z^{-1}x^a z = y^b$ ?

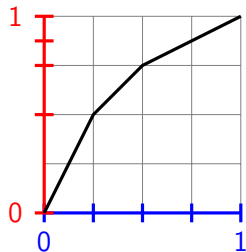
2015 **Barker, Duncan and R.** Generalisation to  $G_{n,r}$  and corrections.  
[Proof of concept implementation.](#)

## Interval partitions $\rightarrow$ trees



How can we store this in memory?

## Interval partitions $\rightarrow$ trees

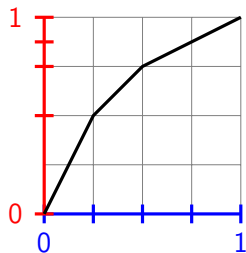


$[0, 1]$

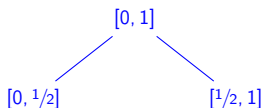
How can we store this in memory?

$[0, 1]$

## Interval partitions $\rightarrow$ trees

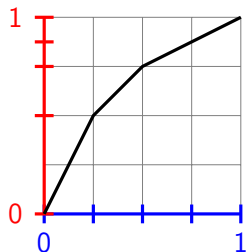


How can we store this in memory?

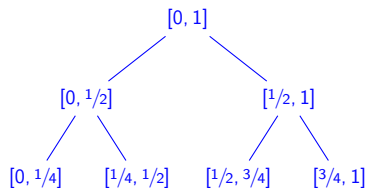


$[0, 1]$

# Interval partitions $\rightarrow$ trees

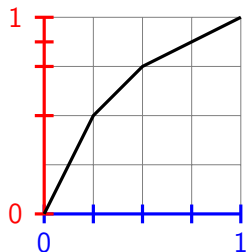


How can we store this in memory?

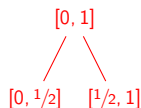
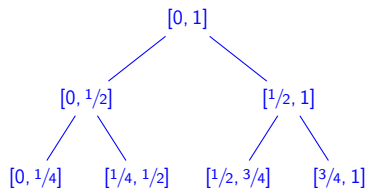


$[0, 1]$

# Interval partitions $\rightarrow$ trees

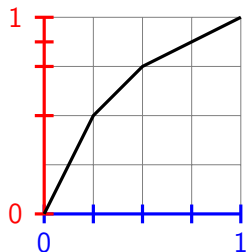


How can we store this in memory?

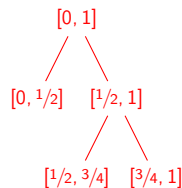
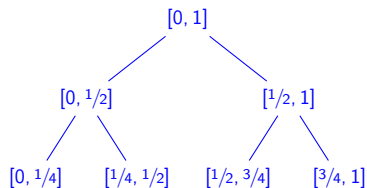




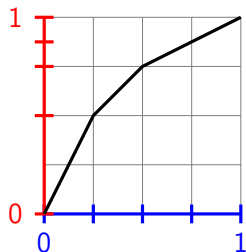
# Interval partitions $\rightarrow$ trees



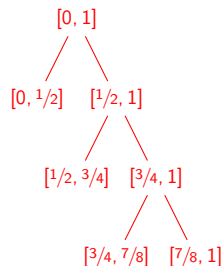
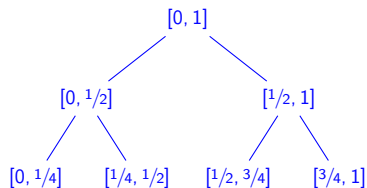
How can we store this in memory?



# Interval partitions $\rightarrow$ trees



How can we store this in memory?



## Trees → paths and words

- ▶ Trees aren't always easy to work with
- ▶ Can only tell if you're at the top (**root**) or bottom (**leaf**)
- ▶ Recursively delegate to children

# Trees → paths and words

- ▶ Trees aren't always easy to work with
- ▶ Can only tell if you're at the top (**root**) or bottom (**leaf**)
- ▶ Recursively delegate to children
- ▶ Quickly becomes confusing! for me, at least...
- ▶ Where does  $f$  send  $[^{34}/_{128}, ^{35}/_{128}]$ ?

# Trees $\rightarrow$ paths and words

- ▶ Trees aren't always easy to work with
- ▶ Can only tell if you're at the top (**root**) or bottom (**leaf**)
- ▶ Recursively delegate to children
- ▶ Quickly becomes confusing! for me, at least...
- ▶ Where does  $f$  send  $[^{34/128}, ^{35/128}]$ ?

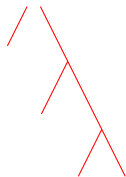
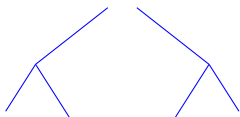
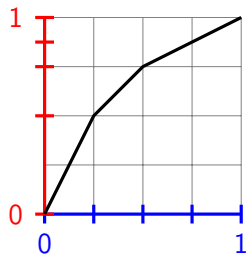
Higman described paths in the tree using an algebra. Introduce labels:

Root  $\mapsto x$

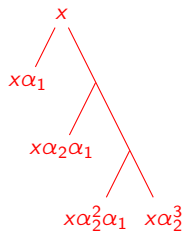
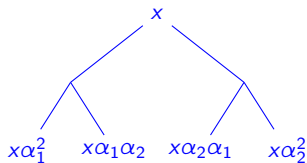
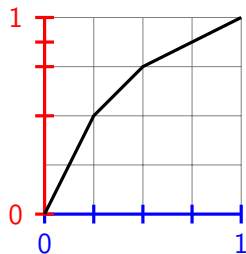
left  $\mapsto \alpha_1$

right  $\mapsto \alpha_2$

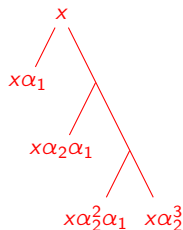
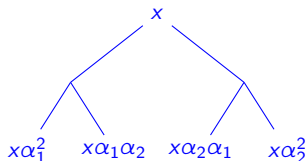
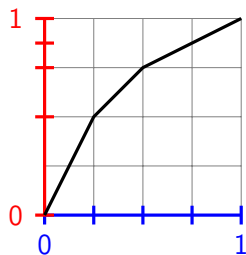
# Trees $\rightarrow$ paths and words



# Trees $\rightarrow$ paths and words



# Trees $\rightarrow$ paths and words



Maps specified by lists of **domain** and **range** words.

$$x\alpha_1^2 \mapsto x\alpha_1$$

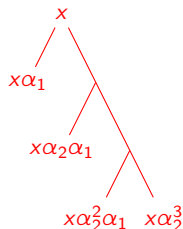
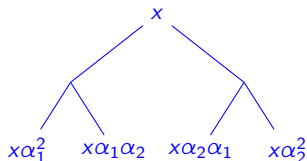
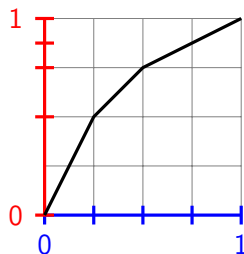
$$x\alpha_1\alpha_2 \mapsto x\alpha_2\alpha_1$$

$$x\alpha_2\alpha_1 \mapsto x\alpha_2^2\alpha_1$$

$$x\alpha_2^2 \mapsto x\alpha_2^3$$



# Trees $\rightarrow$ paths and words



Maps specified by lists of **domain** and **range** words.

$$\begin{array}{ll}
 x\alpha_1^2 & \mapsto x\alpha_1 \\
 x\alpha_1\alpha_2 & \mapsto x\alpha_2\alpha_1 \\
 x\alpha_2\alpha_1 & \mapsto x\alpha_2^2\alpha_1 \\
 x\alpha_2^2 & \mapsto x\alpha_2^3
 \end{array}$$

Easier to compute **components**, e.g.

$$\begin{array}{llll}
 \dots & \mapsto & x\alpha_1^3 & \mapsto & x\alpha_1^2 & \mapsto & x\alpha_1 \\
 \dots & \mapsto & [0, 1/8] & \mapsto & [0, 1/4] & \mapsto & [0, 1/2]
 \end{array}$$

## Components ( $\approx$ orbits)

- ▶ Pick your favourite word e.g.  $w = x\alpha_1^2 \leftrightarrow [0, 1/4]$ .
- ▶ Compute component of  $w$  until you can't any more:

$$\dots, f^{-2}, f^{-1}(w), w, f(w), f^2(w), \dots$$

# Components ( $\approx$ orbits)

- ▶ Pick your favourite word e.g.  $w = x\alpha_1^2 \leftrightarrow [0, 1/4]$ .
- ▶ Compute component of  $w$  until you can't any more:

$$\dots, f^{-2}, f^{-1}(w), w, f(w), f^2(w), \dots$$

Components come in five different shapes:



# Components and Conjugacy

1. Break down a map  $f \in V$  into components.

# Components and Conjugacy

1. Break down a map  $f \in V$  into components.
2. If  $g \in V$  is conjugate to  $f$ , the components of  $g$  must match the components of  $f$ .
  - e.g. periodic  $\mapsto$  periodic, with the same period

# Components and Conjugacy

1. Break down a map  $f \in V$  into components.
2. If  $g \in V$  is conjugate to  $f$ , the components of  $g$  must match the components of  $f$ .  
e.g. periodic  $\mapsto$  periodic, with the same period
3. Only finitely many matchings—check them all!

# Components and Conjugacy

1. Break down a map  $f \in V$  into components.
2. If  $g \in V$  is conjugate to  $f$ , the components of  $g$  must match the components of  $f$ .  
e.g. periodic  $\mapsto$  periodic, with the same period
3. Only finitely many matchings—check them all!
4. If one works: we get a conjugator  $h$ .

# Components and Conjugacy

1. Break down a map  $f \in V$  into components.
2. If  $g \in V$  is conjugate to  $f$ , the components of  $g$  must match the components of  $f$ .  
e.g. periodic  $\mapsto$  periodic, with the same period
3. Only finitely many matchings—check them all!
4. If one works: we get a conjugator  $h$ .
5. If none of them work: no conjugator exists.



# Implementation

- ▶ My job: get a computer to do this.
- ▶ Like programming a fancy calculator!
- ▶ Wrote in Python: something I knew and quick to work with.

# Implementation

- ▶ My job: **get a computer to do this.**
- ▶ Like programming a fancy calculator!
- ▶ Wrote in Python: something I knew and quick to work with.
- ▶ Summer project that was more like a semester project. . .

# Implementation

- ▶ My job: **get a computer to do this.**
- ▶ Like programming a fancy calculator!
- ▶ Wrote in Python: something I knew and quick to work with.
- ▶ Summer project that was more like a semester project. . .
- ▶ Other tools **exist** to do calculations in  $V$ , but not to solve the conjugacy problem.

Code is on **GitHub**. Come and find me if you want a demo!

## Sphinx: comments in source code

```
def format(word):  
    """Turns a sequence of integers representing a *word* into [...]  
    >>> format([2, -1, 2, -2, 0])  
    'x2 a1 x2 a2 L'  
    >>> format([])  
    The Spanish Inquisition  
    """  
    if len(word) == 0:  
        return "<the empty word>"  
    return " ".join(_char(i) for i in word)
```

Sphinx generates nice HTML documentation and runs tests based on  
"""comments like this""".

## Sphinx: doctest

```
H:\thompsons_v\docs>make doctest
```

```
[...]
```

```
*****
```

```
File "thompson.word.rst", line 10, in default
```

```
Failed example:
```

```
    format([])
```

```
Expected:
```

```
    The Spanish Inquisition
```

```
Got:
```

```
    '<the empty word>'
```

```
*****
```

```
1 items had failures:
```

```
    1 of 100 in default
```

```
100 tests in 1 items.
```

```
99 passed and 1 failed.
```

```
***Test Failed*** 1 failures.
```

## Other lessons learned

- ▶ Small test suites—catch bugs before they happen
- ▶ Generate random examples
- ▶ Immutable words
- ▶ Document the code

# Future Work

## Code

- ▶ More testing
- ▶ Complexity analysis

## Theory

- ▶ **Simultaneous** conjugacy

Given  $x_1, \dots, x_n; y_1, \dots, y_n$  find a **single** conjugator  $z$  such that

$$z^{-1}x_iz = y_i, \quad \forall i$$

- ▶ Try to solve different kinds of equations?
- ▶ Transfer to more general Thompson-like groups  $V(\Sigma)$ ?