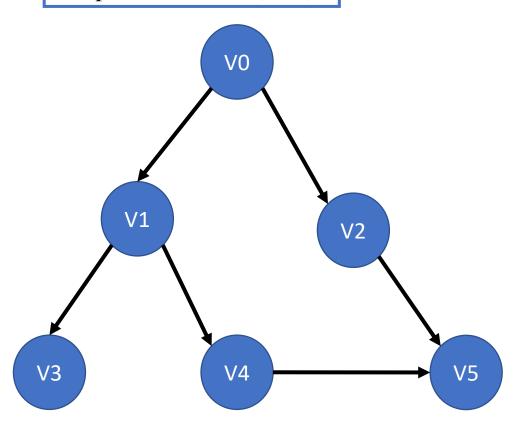
Path-finding

Daniel Nogueira

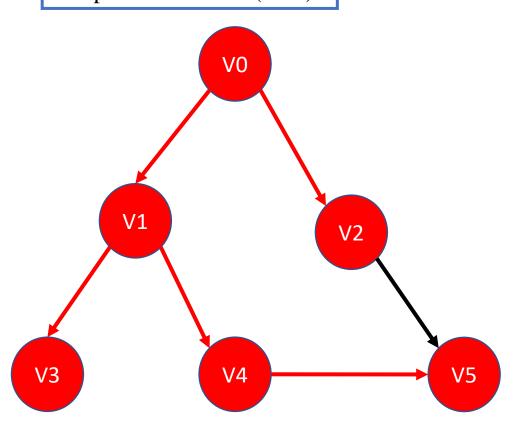
dnogueira@ipca.pt

Depth-First Search (DFS)



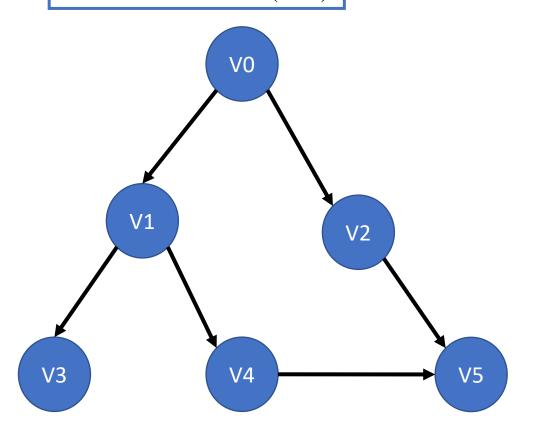
- 1.Set a start node
- 2. While this is not an objective or final node (node whose adjacency has already been visited):
 - Choose an adjacent node not yet visited
 - Visit it
- 3. If it is a non-objective end node:
 - Return to this father
- If there is a father, repeat. If there is no parent, choose another start node

Depth-First Search (DFS)



- 1.Set a start node
- 2. While this is not an objective or final node (node whose adjacency has already been visited):
 - Choose an adjacent node not yet visited
 - Visit it
- 3. If it is a non-objective end node:
 - Return to this father
- If there is a father, repeat. If there is no parent, choose another start node

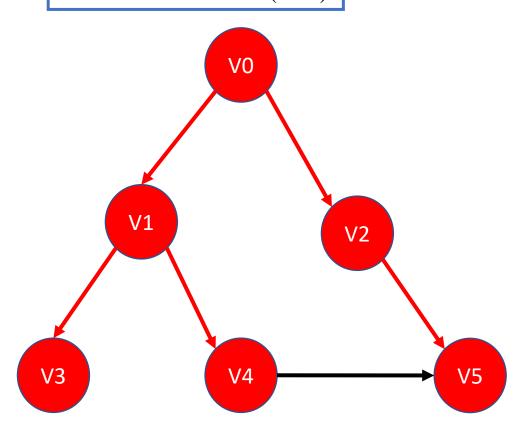
Breadth-First Search (BFS)



- 1.Define an initial node, marking it as explored
- 2. Put it on the list
- 3. As long as the queue is not empty:
 - Remove the 1st node from the queue, u
 - For each neighbour v of u:
 - * If v is not explored:
 - ** Mark v as explored
 - ** Put v at the end of the queue
- 4. Repeat from another starting node, if there is one



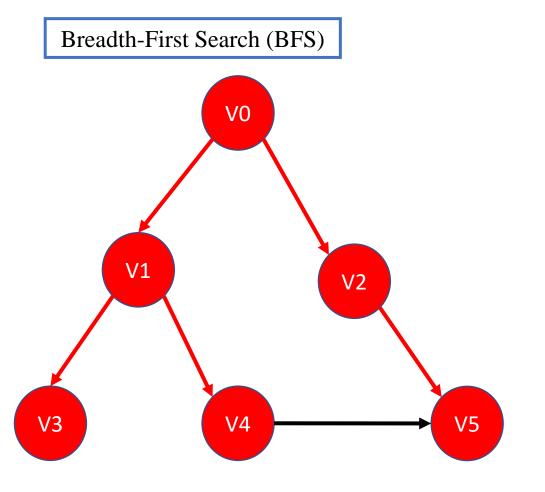
Breadth-First Search (BFS)

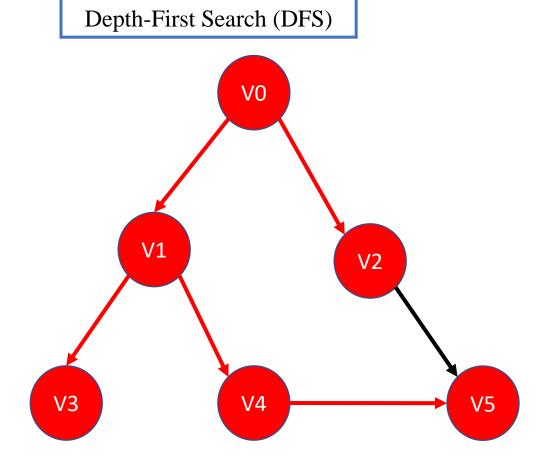


Following is the Breadth-First Search V0 V1 V2 V3 V4 V5

DFS V0 V1 V3 V4 V5 V2

BFS V0 V1 V2 V3 V4 V5





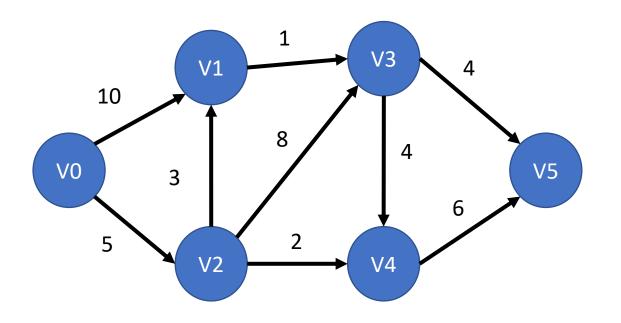
Dijkstra



Dijkstra's algorithm can be used to find the <u>shortest distance</u> from the source vertex to all other vertices in a weighted graph. For each vertex \underline{v} of the graph, we maintain an attribute $\underline{d(v)}$ which is an upper bound on the weight of the shortest path from the initial <u>node s to v</u>.

- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $\underline{v} \neq \underline{s}$, and p(v) = -1 for all \underline{v}
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
 - * Choose \underline{u} whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)

Dijkstra

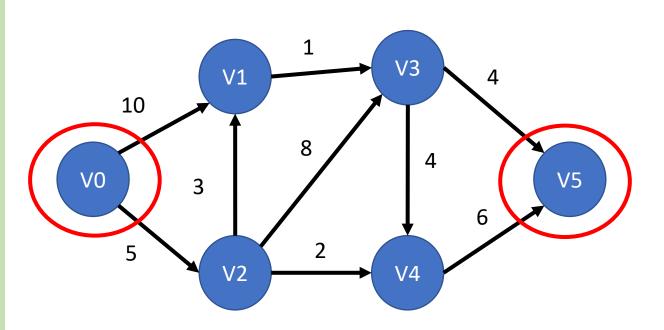


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose \underline{u} whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)

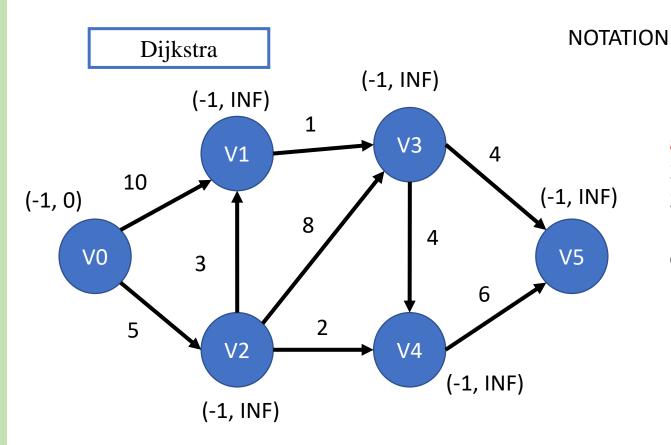
Dijkstra

NOTATION



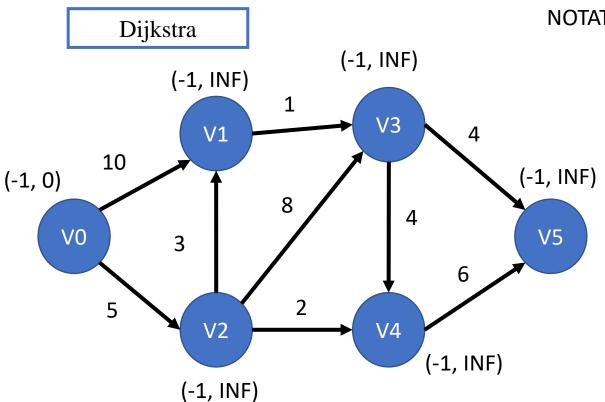


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose \underline{u} whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)





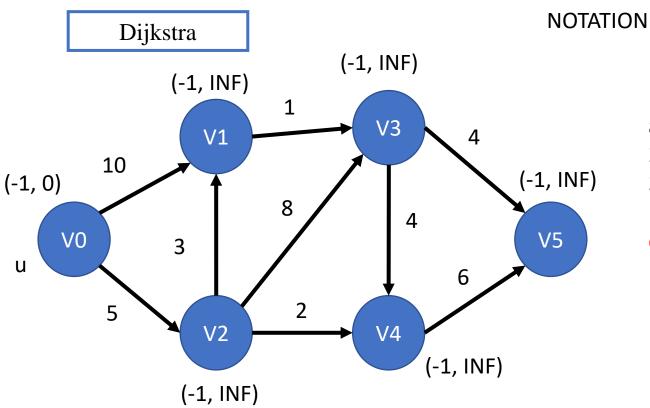
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose \underline{u} whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



NOTATION

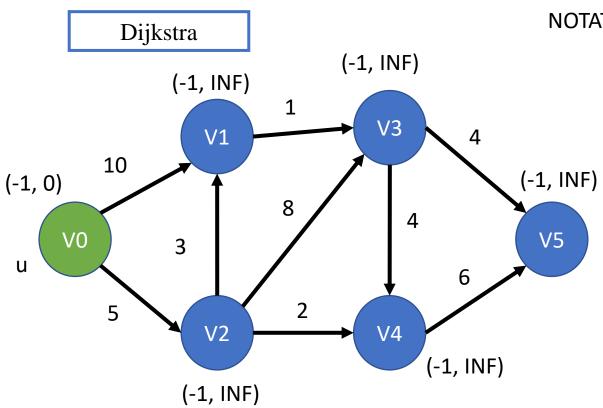


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose \underline{u} whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



TION

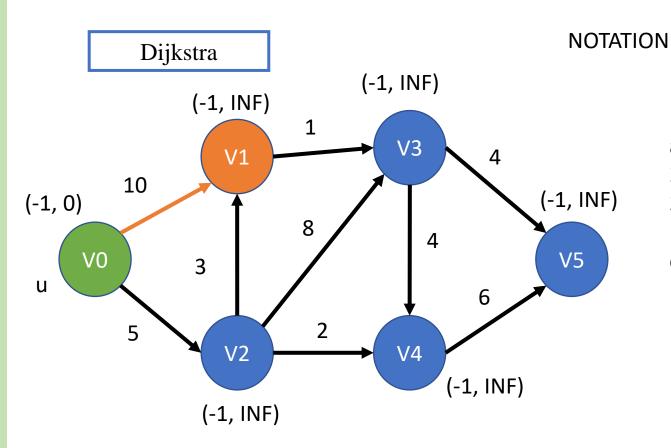
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



NOTATION

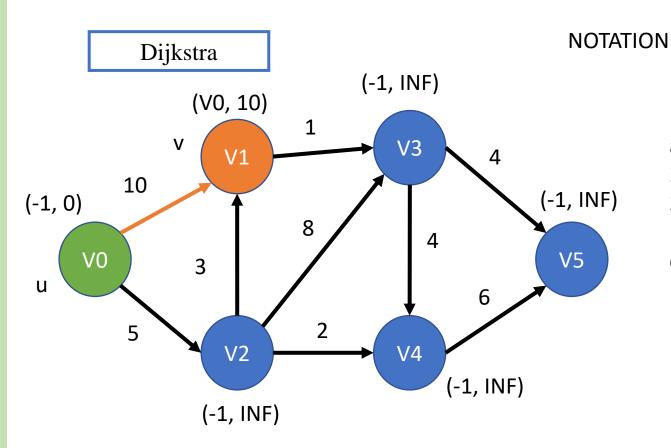


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



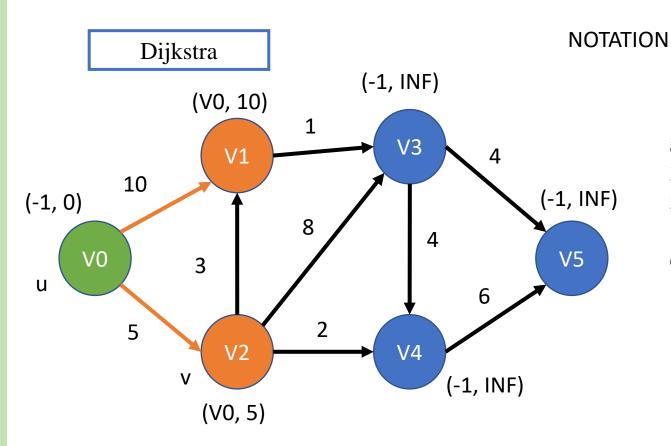


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



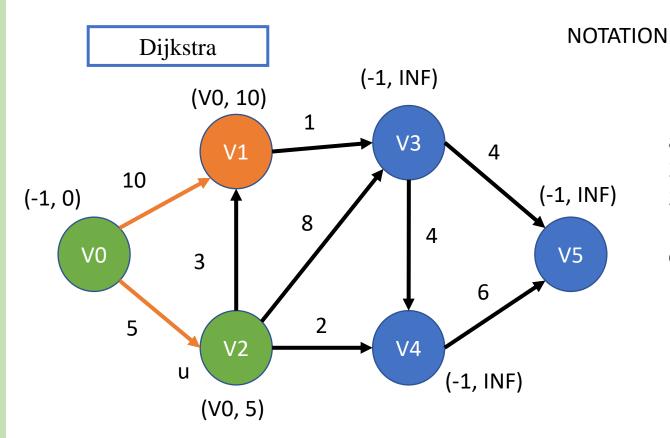


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



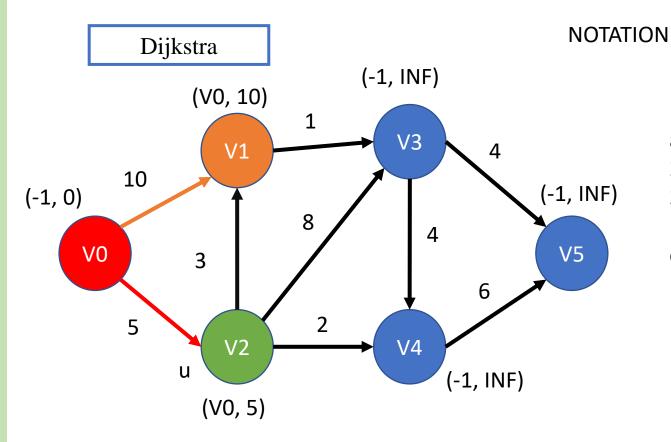


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



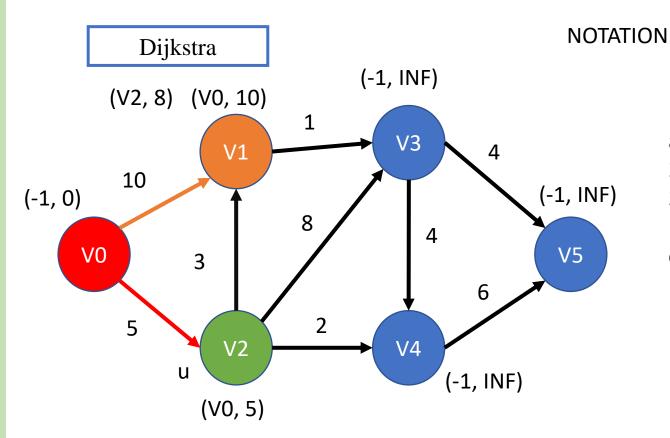


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



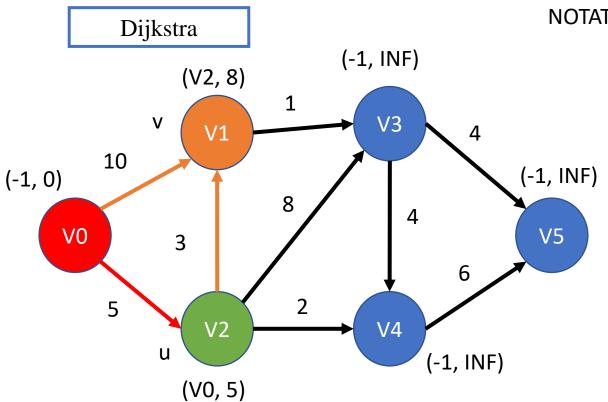


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)





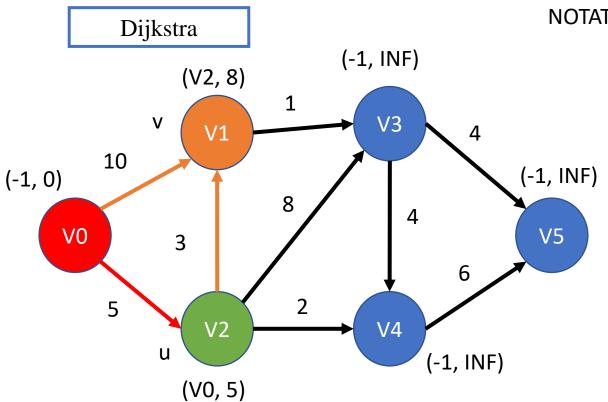
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



NOTATION



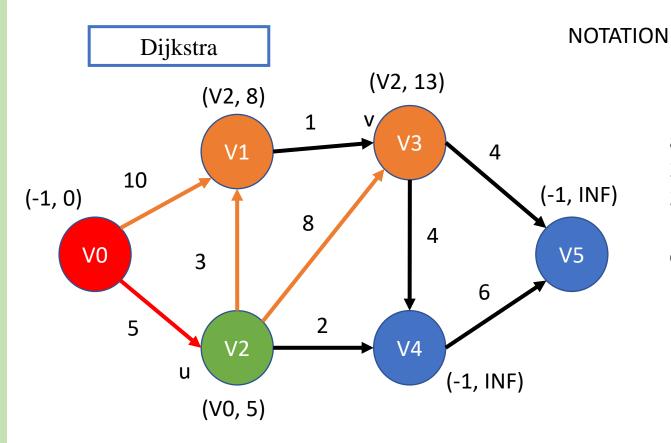
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



NOTATION

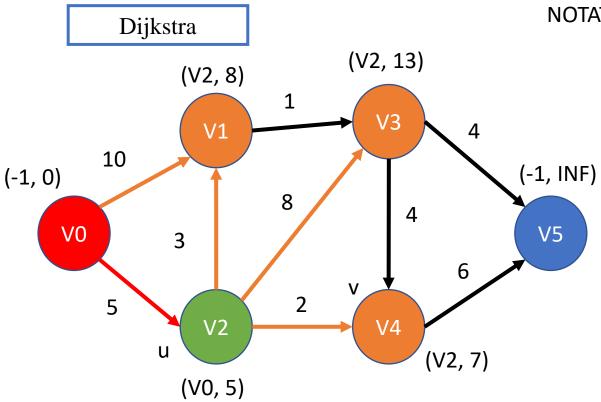


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



N _

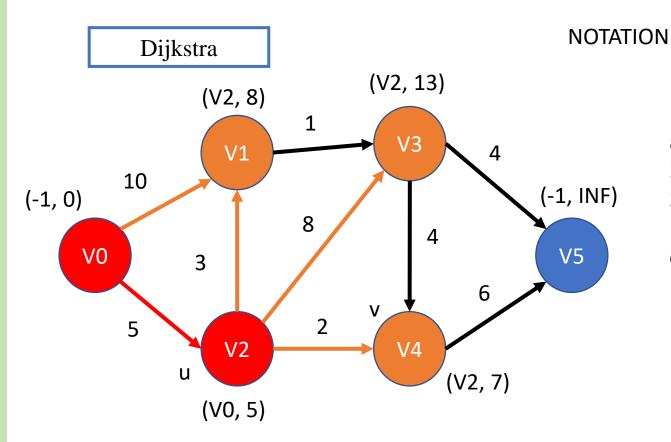
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



NOTATION

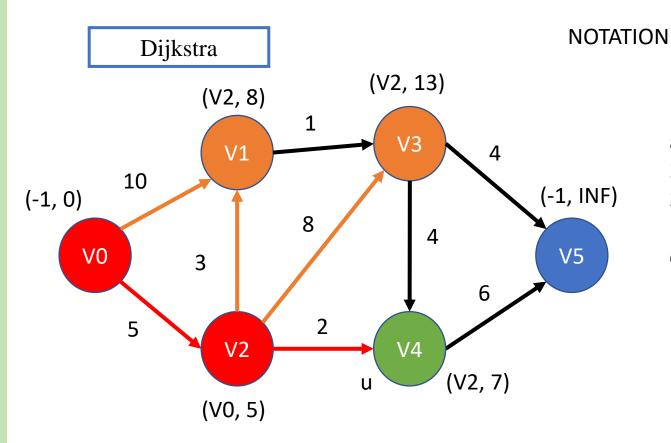


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



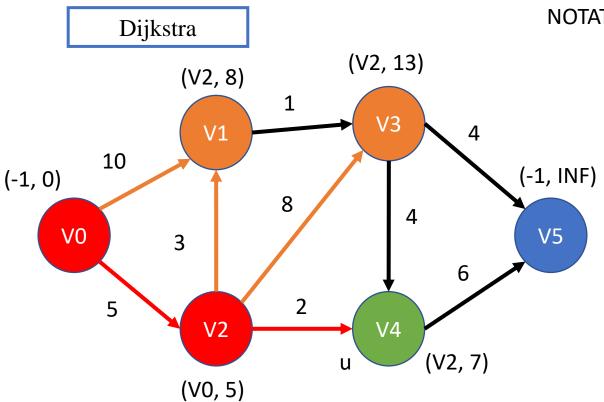


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)





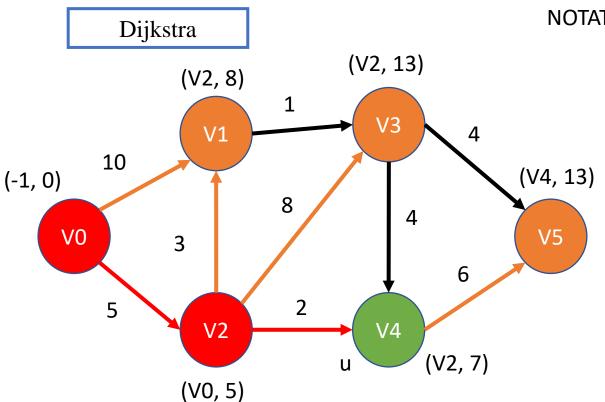
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



NOTATION



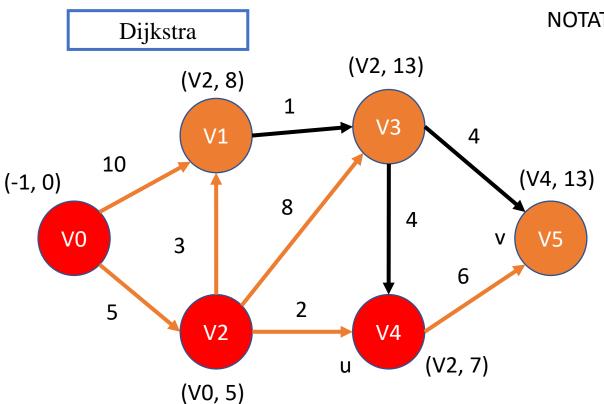
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



NOTATION



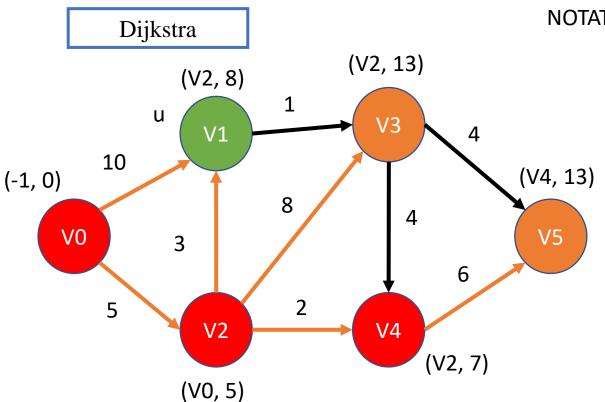
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



NOTATION



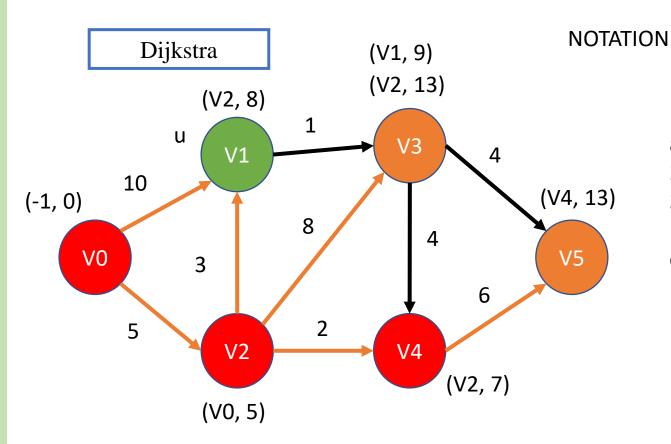
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



NOTATION

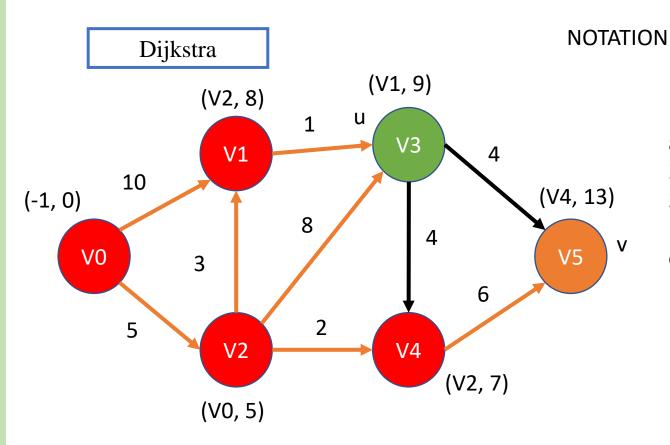


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



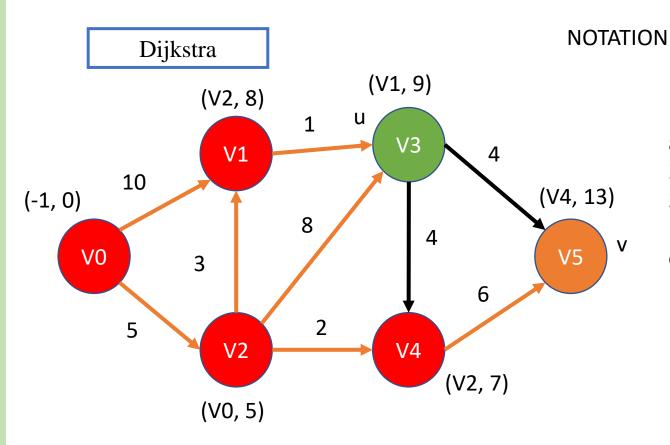
N

- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



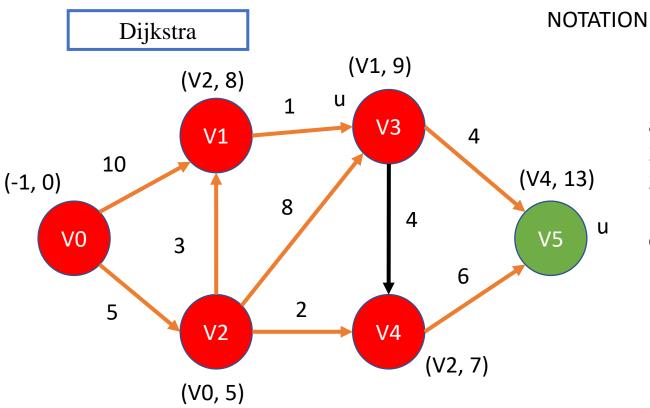


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$



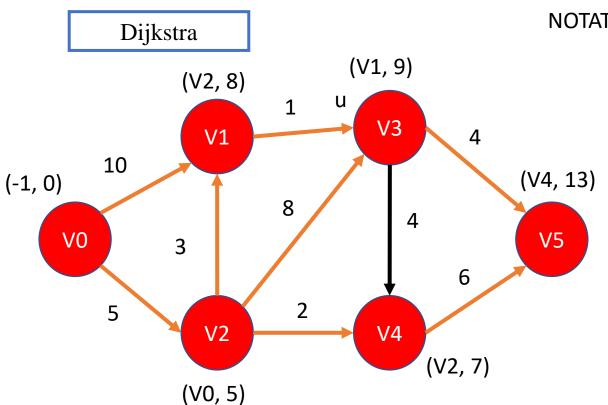


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge $(\underline{u},\underline{v})$





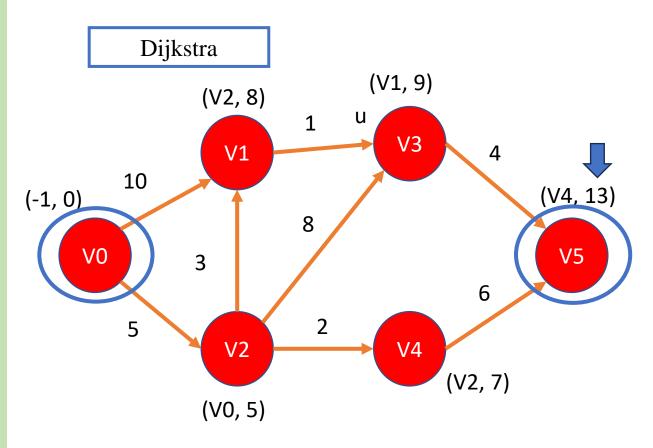
- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close *u*
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



NOTATION



- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $v \neq s$, and p(v) = -1 for all $v \neq s$.
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose u whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)

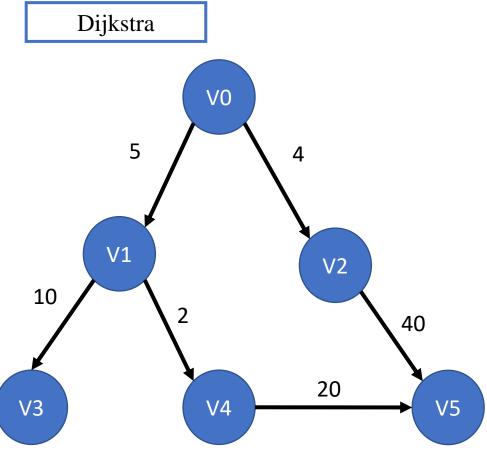


the shortest distance V0 => V5



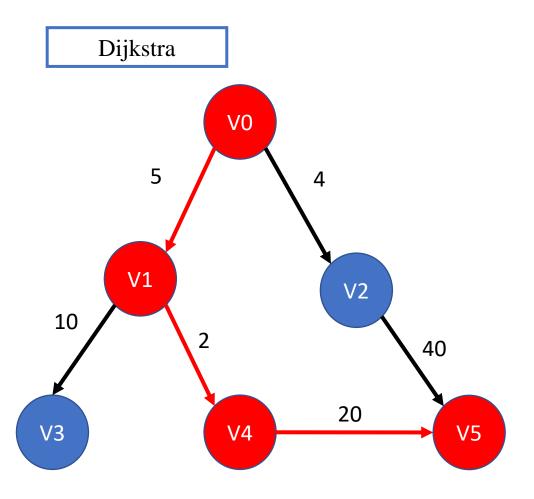
V5, V4, V2, V0

TOTAL = 13



```
def dijkstra_algorithm(graph, start_node):
 unvisited nodes = list(graph.get nodes())
 shortest path = {}
previous nodes = {}
 # We'll use max value to initialize the "infinity" value of the unvisited nodes
 max value = sys.maxsize
 for node in unvisited_nodes:
     shortest path[node] = max value
 # However, we initialize the starting node's value with 0
 shortest path[start node] = 0
 while unvisited nodes:
     current min node = None
     for node in unvisited_nodes: # Iterate over the nodes
         if current_min_node == None:
             current min node = node
         elif shortest path[node] < shortest path[current min node]:</pre>
             current min node = node
     # The code block below retrieves the current node's neighbors and updates their distances
    neighbors = graph.get_outgoing_edges(current_min_node)
     for neighbor in neighbors:
         tentative_value = shortest_path[current_min_node] + graph.value(current_min_node, neighbor)
        if tentative_value < shortest_path[neighbor]:</pre>
             shortest path[neighbor] = tentative value
             # We also update the best path to the current node
             previous_nodes[neighbor] = current_min_node
    unvisited_nodes.remove(current_min_node)
 return previous_nodes, shortest_path
```





We found the following best path with a value of 27. V0 -> V1 -> V4 -> V5

 A^*



A* is an informed search algorithm. Informed Search signifies that the algorithm has extra information, to begin with. It is a complete as well as an optimal solution for solving path and grid problems.

 A^*



A* is an informed search algorithm. Informed Search signifies that the algorithm has extra information, to begin with. It is a **complete** as well as an **optimal** solution for solving path and grid problems.

Complete – It means that it will find all the available paths from start to end.

Optimal – find the least cost from the starting point to the ending point.

A*

A* is an <u>informed search</u> algorithm. <u>Informed Search</u> signifies that the algorithm has extra information, to begin with. It is a <u>complete</u> as well as an <u>optimal</u> solution for solving path and grid problems.

$$f(n) = g(n) + h(n)$$

The Algorithm

- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
 - * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.

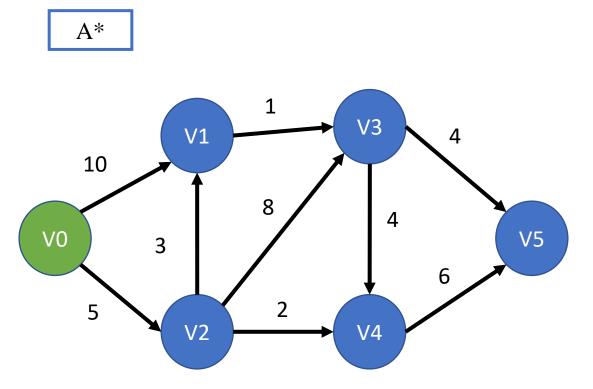
A*

Algorithms

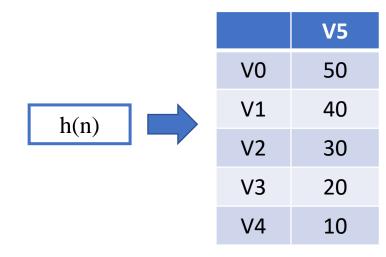
manhattan((x1, y1), (x2, y2)) = |x1 - x2| + |y1 - y2|euclidean((x1, y1), (x2, y2)) = $sqrt(x^2 + y^2)$

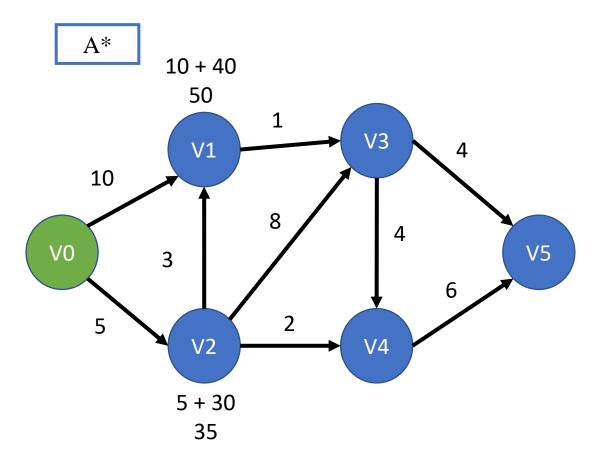
h(n)

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
5	4	3	2	1
4	3	2	1	0

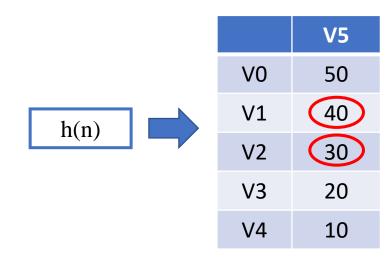


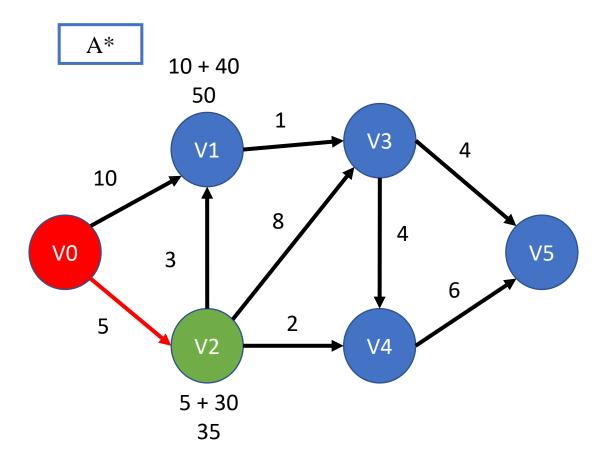
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



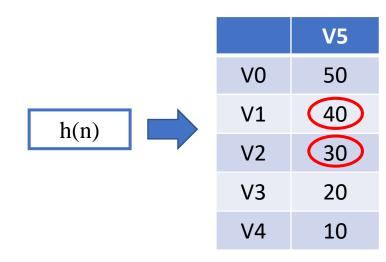


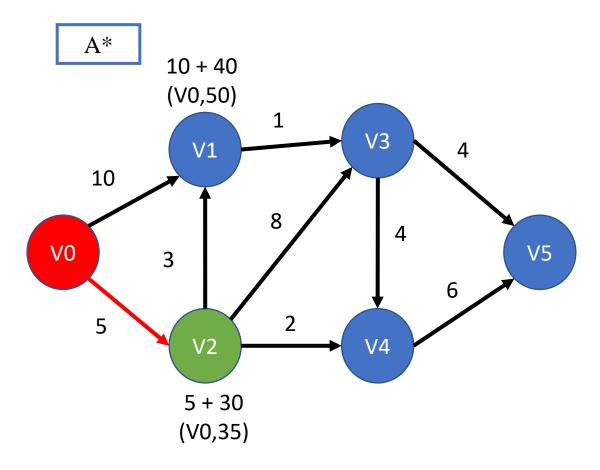
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



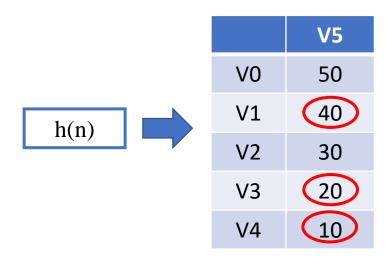


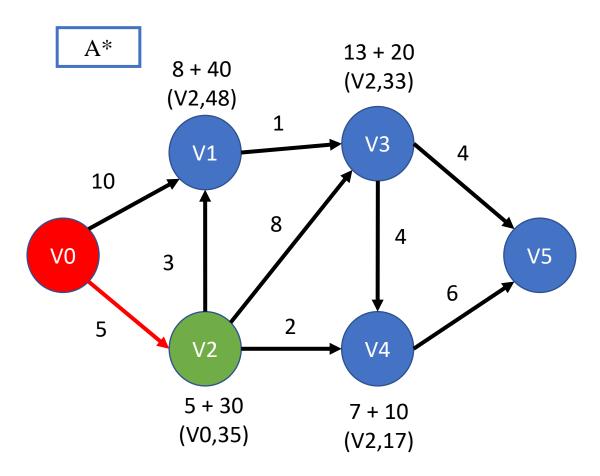
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



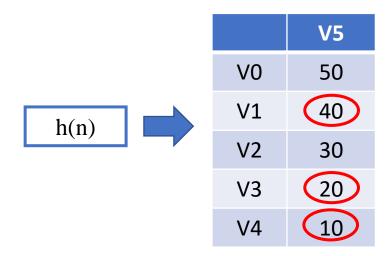


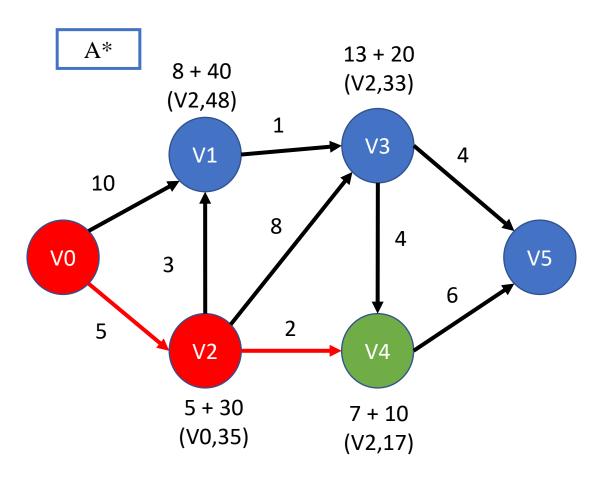
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



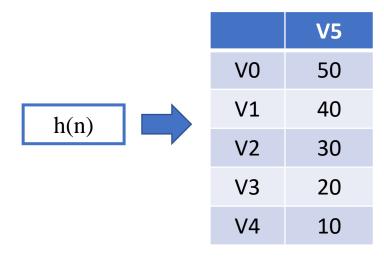


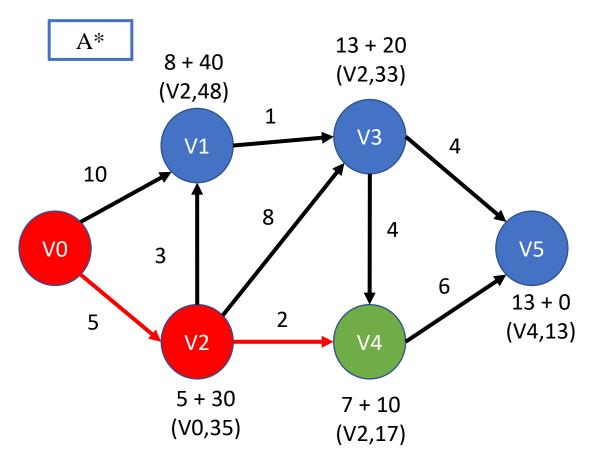
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



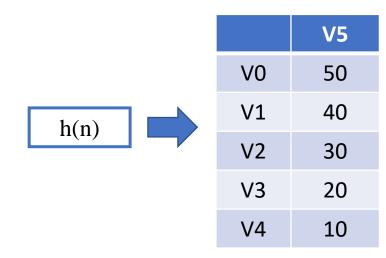


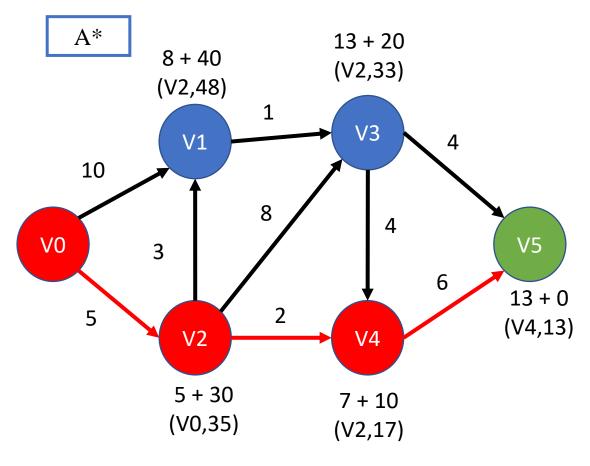
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



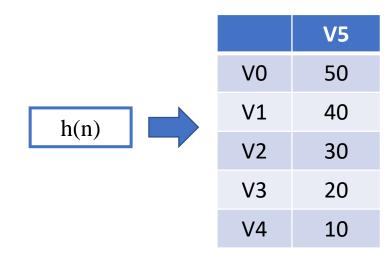


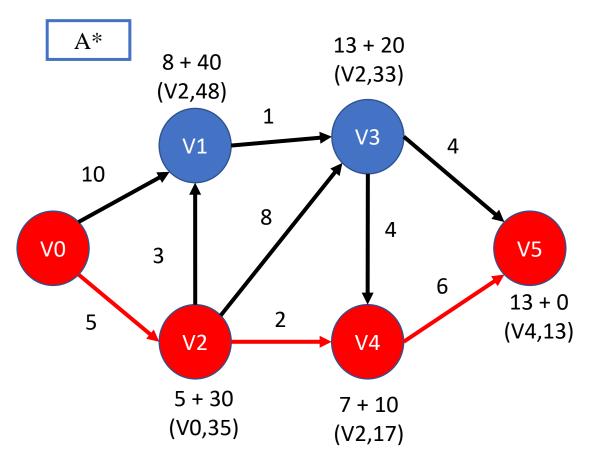
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



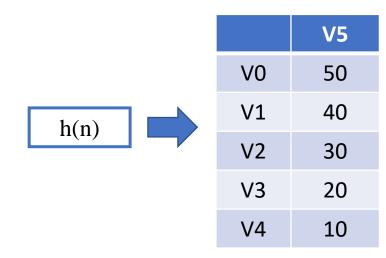


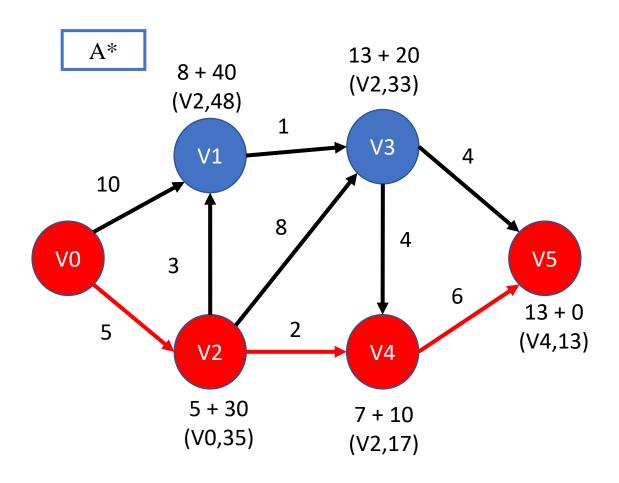
- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.





- 1. Place the starting node into OPEN and find its f (n) value.
- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.





the shortest distance V0 => V5



V5, V4, V2, V0

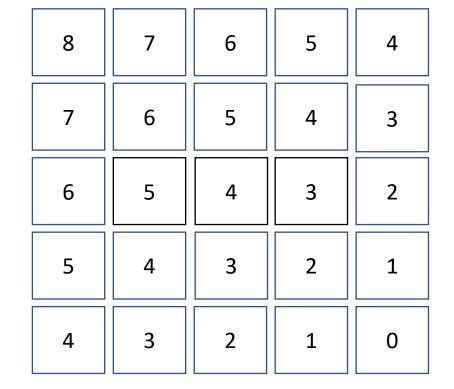
 A^*

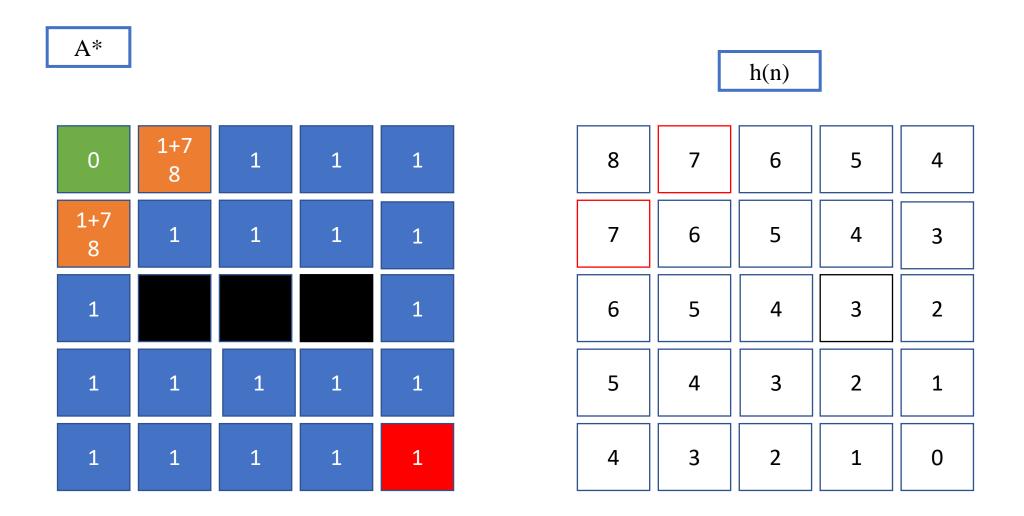
Algorithms

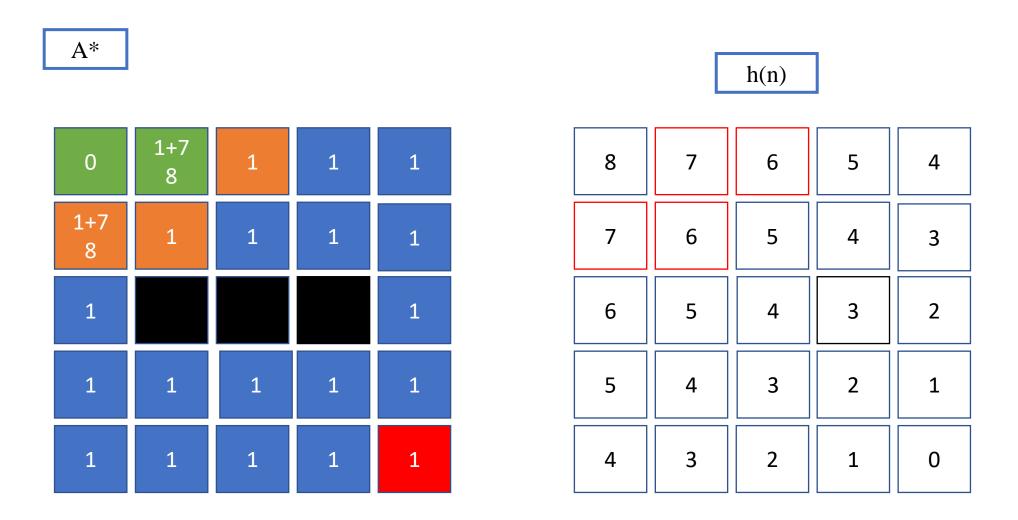
manhattan((x1, y1), (x2, y2)) = |x1 - x2| + |y1 - y2|euclidean((x1, y1), (x2, y2)) = $sqrt(x^2 + y^2)$

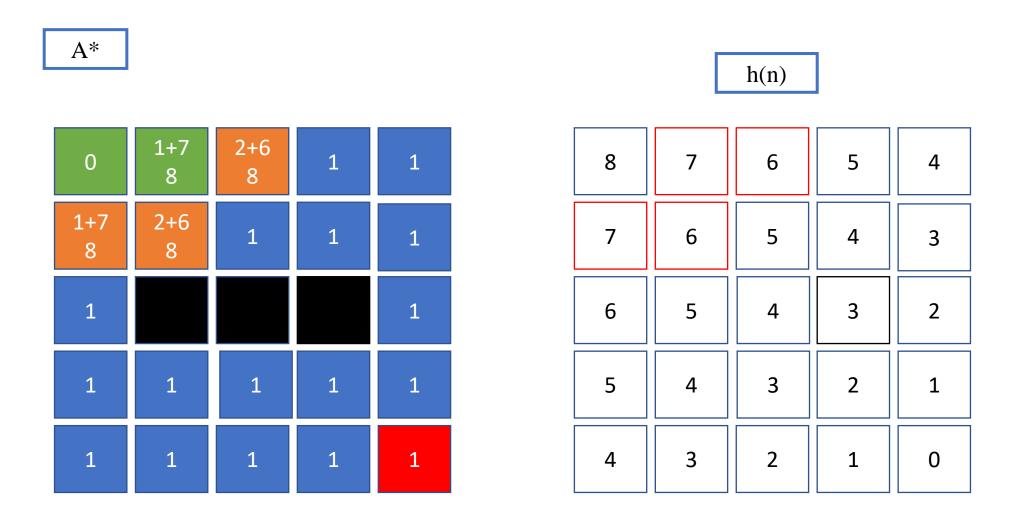
h(n)

0	1	1	1	1
1	1	1	1	1
1				1
1	1	1	1	1
1	1	1	1	1











 A^* h(n) 1+7 2+6 3+5 8 5 4 1+7 2+6 3+5 7 6 4 3 6 5 3 4 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 8 5 4 1+7 2+6 3+5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 1+7 2+6 3+5 4+4 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 1+7 2+6 3+5 4+4 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6+2 6 5 3 4 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6+2 6 5 3 7+1 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6+2 6 5 3 4 8+2 7+1 5 2 4 1 10 8+0 4 0 8



Daniel Nogueira

dnogueira@ipca.pt