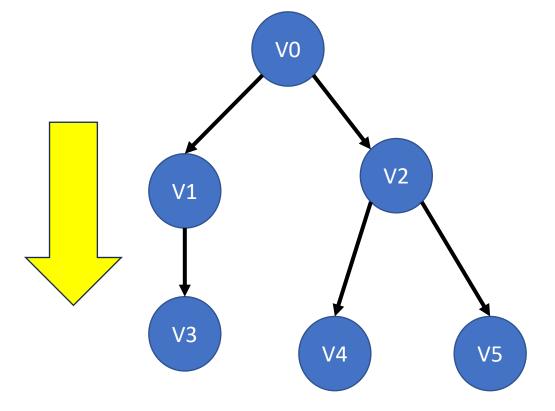
Path-finding

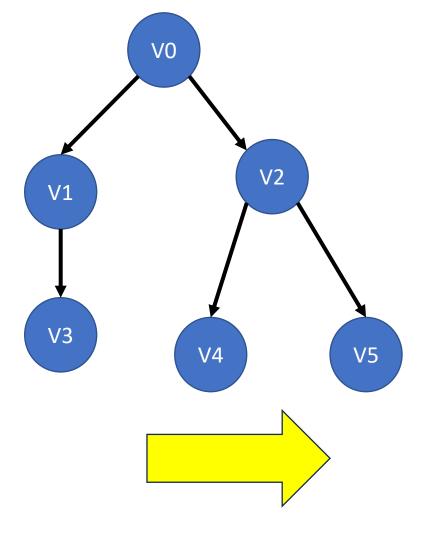
Daniel Nogueira

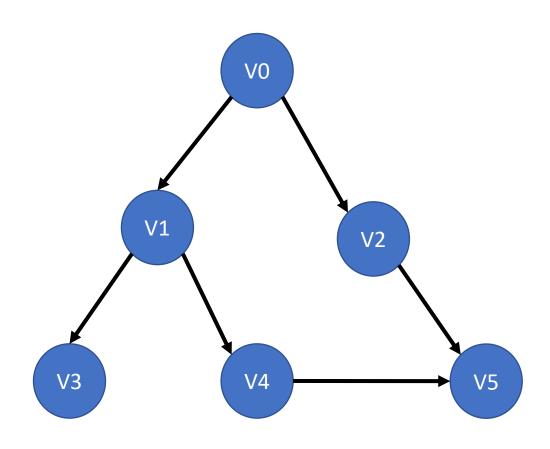
dnogueira@ipca.pt



Depth-First Search (DFS)



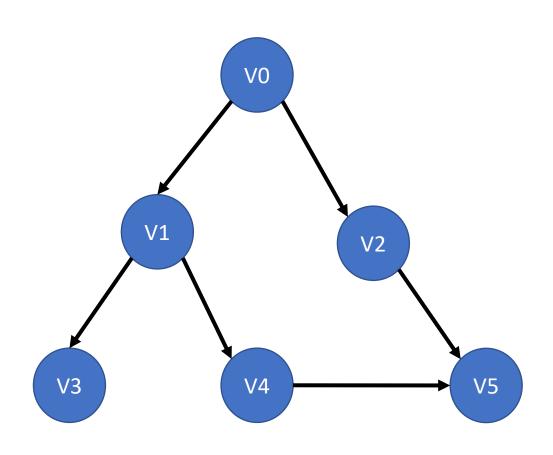




- 1. Set a start node
- 2. While this is not an objective or final node (node whose adjacency has already been visited):
 - Choose an adjacent node not yet visited
 - Visit it
- 3. If it is a non-objective end node:
 - Return to this father
- If there is a father, repeat. If there is no parent, choose another start node

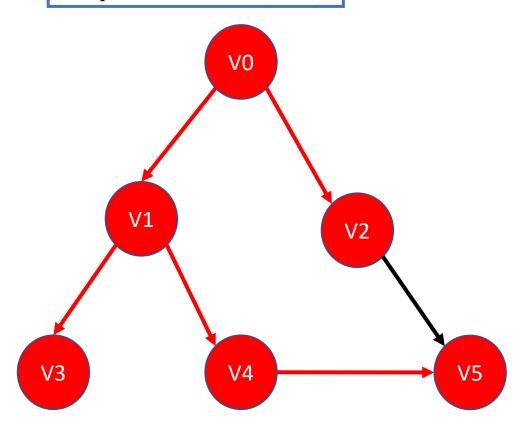
```
using System;
using System.Collections.Generic;
class Program
  static Dictionary<string, List<string>> graph = new Dictionary<string, List<string>>
    { "V0", new List<string> { "V1", "V2" } },
    { "V1", new List<string> { "V3", "V4" } },
    { "V2", new List<string> { "V5" } },
    { "V3", new List<string>() },
    { "V4", new List<string> { "V5" } },
    { "V5", new List<string>() }
  static HashSet<string> visited = new HashSet<string>();
  static void Main(string[] args)
    string startNode = "V0";
    Console.WriteLine("Following is the Depth-First Search:");
    DFS(startNode);
    Console.ReadLine();
  static void DFS(string node)
    if (!visited.Contains(node))
      Console.WriteLine(node);
      visited.Add(node);
      if (graph.ContainsKey(node))
        foreach (string neighbor in graph[node])
           DFS(neighbor);
```

```
private void DFSRecursive(string vertex, HashSet<string> visited)
    visited.Add(vertex);
    Console.WriteLine("Visiting vertex: " + vertex);
    foreach (string neighbor in adjacencyList[vertex])
      if (!visited.Contains(neighbor))
        DFSRecursive(neighbor, visited);
         } } }
class Program
  static void Main(string[] args)
    Graph graph = new Graph();
    // Adicionar vértices
    graph.AddVertex("V0");
    graph.AddVertex("V1");
    graph.AddVertex("V2");
    graph.AddVertex("V3");
    graph.AddVertex("V4");
    graph.AddVertex("V5");
    // Adicionar arestas
    graph.AddEdge("V0", "V1");
    graph.AddEdge("V0", "V2");
    graph.AddEdge("V1", "V3");
    graph.AddEdge("V1", "V4");
    graph.AddEdge("V2", "V5");
    graph.AddEdge("V4", "V5");
    Console.WriteLine("DFS starting from vertex V0:");
    graph.DFS("V0"); } }
```

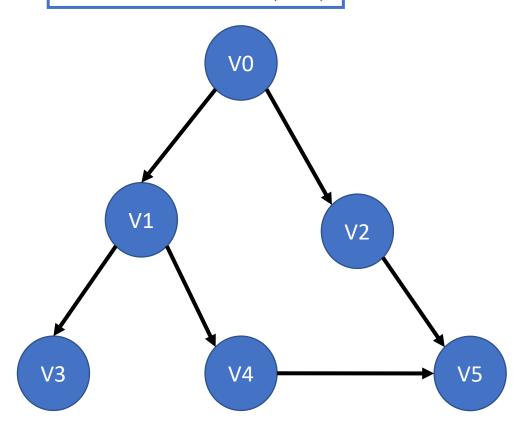


```
def dfs(graph, vertex, visited):
  # Marcar o vértice como visitado
  visited[vertex] = True
  print("Visitando vértice:", vertex)
  # Recursivamente visitar os vértices adjacentes não visitados
  for neighbor in graph[vertex]:
    if not visited[neighbor]:
       dfs(graph, neighbor, visited)
# Grafo representado como um dicionário de adjacências
graph = {
  'V0': ['V1', 'V2'],
  'V1': ['V3', 'V4'],
  'V2': ['V5'],
  'V3': [],
  'V4': ['V5'],
  'V5': []
# Inicializar um vetor de visitados
visited = {vertex: False for vertex in graph}
# Chamar o DFS a partir de todos os vértices não visitados
for vertex in graph:
  if not visited[vertex]:
    dfs(graph, vertex, visited)
```





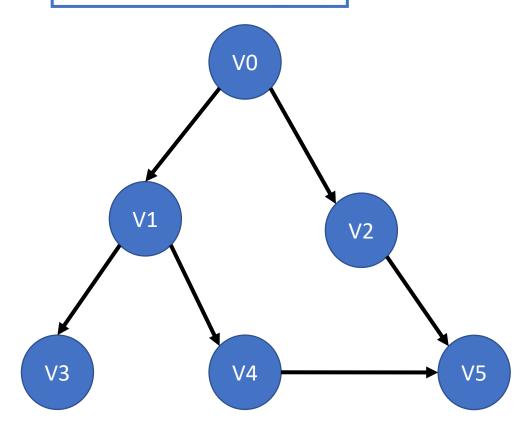
```
Following is the Depth-First Search
V0
V1
V3
V4
V5
```



- 1. Define an initial node, marking it as explored
- 2. Put it on the list
- 3. As long as the queue is not empty:
 - Remove the 1st node from the list, u
 - For each neighbour v of u:
 - * If v is not explored:
 - ** Mark v as explored
 - ** Put v at the end of the list
- 4. Repeat from another starting node, if there is one

```
using System;
using System.Collections.Generic;
class Program
  static Dictionary<string, List<string>> graph = new Dictionary<string, List<string>>
    { "V0", new List<string> { "V1", "V2" } },
    { "V1", new List<string> { "V3", "V4" } },
    { "V2", new List<string> { "V5" } },
    { "V3", new List<string>() },
    { "V4", new List<string> { "V5" } },
    { "V5", new List<string>() }
  static List<string> visited = new List<string>();
  static Queue<string> queue = new Queue<string>();
  static void Main(string[] args)
    string startNode = "V0";
    Console.WriteLine("Following is the Breadth-First Search:");
    BFS(startNode);
    Console.ReadLine();
```

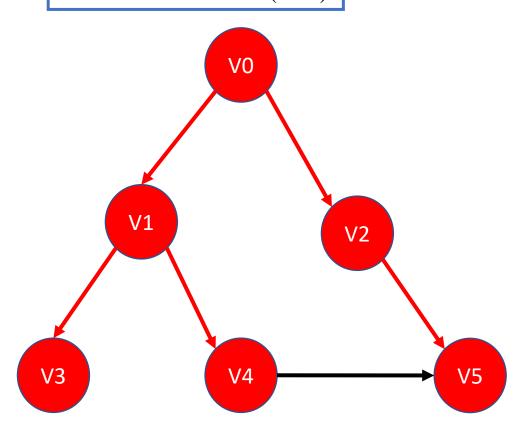
```
static void BFS(string node)
    visited.Add(node);
    queue.Enqueue(node);
    while (queue.Count > 0)
      string s = queue.Dequeue();
      Console.Write(s + " ");
      if (graph.ContainsKey(s))
        foreach (string neighbor in graph[s])
          if (!visited.Contains(neighbor))
            visited.Add(neighbor);
            queue.Enqueue(neighbor);
```



```
graph = {
    'V0' : ['V1','V2'],
    'V1' : ['V3', 'V4'],
    'V2' : ['V5'],
    'V3' : [],
    'V4' : ['V5'],
    'V5' : []
visited = [] # List to keep track of visited nodes.
queue = []
               #Initialize a queue
def bfs(visited, graph, node):
  visited.append(node)
  queue.append(node)
  while queue:
    s = queue.pop(0)
    print (s, end = " ")
    for neighbour in graph[s]:
      if neighbour not in visited:
        visited.append(neighbour)
        queue.append(neighbour)
# Driver Code
print("Following is the Breadth-First Search")
bfs(visited, graph, 'V0')
```



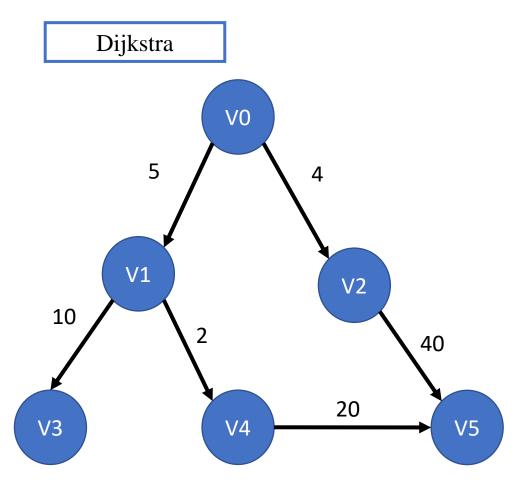
Breadth-First Search (BFS)



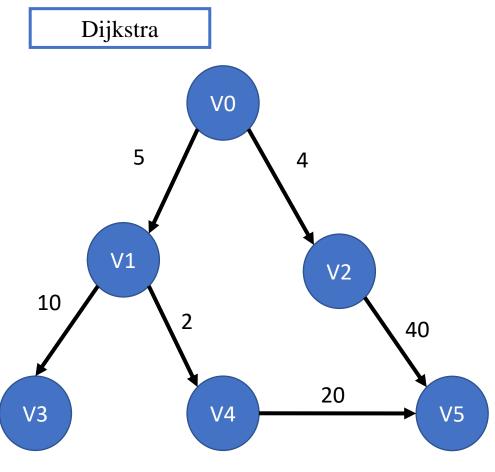
Following is the Breadth-First Search V0 V1 V2 V3 V4 V5

DFS V0 V1 V3 V4 V5 V2

BFS V0 V1 V2 V3 V4 V5

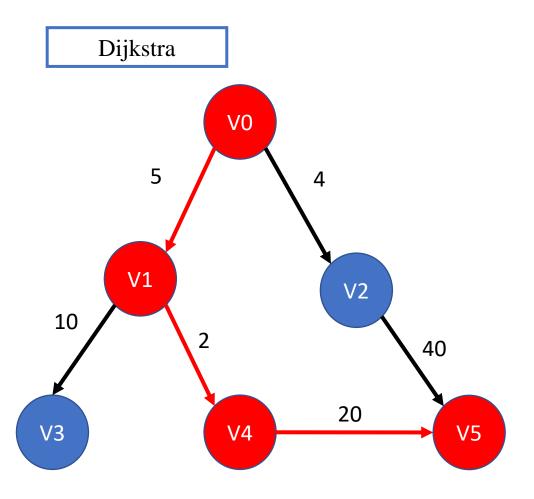


- 1. Initialize the graph with d(s) = 0, d(v) = INF, for all $\underline{v} \neq \underline{s}$, and p(v) = -1 for all \underline{v}
- 2. Make open(v) = True for every v in the graph
- 3. As long as there is an open vertex:
- * Choose \underline{u} whose estimate is the smallest among the open
 - * Close <u>u</u>
 - * For every open node \underline{v} adjacent to \underline{u} : relax edge ($\underline{u},\underline{v}$)



```
def dijkstra_algorithm(graph, start_node):
    unvisited nodes = list(graph.get nodes())
    shortest path = {}
   previous nodes = {}
    # We'll use max value to initialize the "infinity" value of the unvisited nodes
    max value = sys.maxsize
    for node in unvisited_nodes:
        shortest path[node] = max value
    # However, we initialize the starting node's value with 0
    shortest path[start node] = 0
    while unvisited nodes:
        current min node = None
        for node in unvisited_nodes: # Iterate over the nodes
            if current_min_node == None:
                current min node = node
            elif shortest path[node] < shortest path[current min node]:</pre>
                current min node = node
        # The code block below retrieves the current node's neighbors and updates their distances
       neighbors = graph.get_outgoing_edges(current_min_node)
        for neighbor in neighbors:
            tentative_value = shortest_path[current_min_node] + graph.value(current_min_node, neighbor)
           if tentative_value < shortest_path[neighbor]:</pre>
                shortest path[neighbor] = tentative value
                # We also update the best path to the current node
                previous_nodes[neighbor] = current_min_node
       unvisited_nodes.remove(current_min_node)
    return previous_nodes, shortest_path
```





We found the following best path with a value of 27. V0 -> V1 -> V4 -> V5

Dijkstra

```
using System;
using System.Collections.Generic;
class Program
  static void Main(string[] args)
    Dictionary<string, Dictionary<string, int>> initGraph = new Dictionary<string, Dictionary<string, int>>()
      { "V0", new Dictionary<string, int> { { "V1", 5 }, { "V2", 4 } } },
      { "V1", new Dictionary<string, int> { { "V3", 10 }, { "V4", 2 } } },
      { "V2", new Dictionary<string, int> { { "V5", 40 } } },
      { "V3", new Dictionary<string, int>() },
      { "V4", new Dictionary<string, int> { { "V5", 20 } } },
      { "V5", new Dictionary<string, int>() }
    Graph graph = new Graph(initGraph);
    string startNode = "V0";
    Dictionary<string, int> shortestPath = DijkstraAlgorithm(graph, startNode);
    PrintResult(shortestPath, startNode, "V5");
    Console.ReadLine();
```

```
class Graph
{
    private Dictionary<string, Dictionary<string, int>> graph;

    public Graph(Dictionary<string, Dictionary<string, int>> initGraph)
    {
        graph = new Dictionary<string, Dictionary<string, int>>(initGraph);
    }

    public bool ContainsNode(string node)
    {
        return graph.ContainsKey(node);
    }

    public Dictionary<string, int> GetEdges(string node)
    {
        return graph[node];
    }
}
```

Dijkstra

```
static Dictionary<string, int> DijkstraAlgorithm(Graph graph, string startNode)
    Dictionary<string, int> shortestPath = new Dictionary<string, int>();
   HashSet<string> visited = new HashSet<string>();
   foreach (var node in graph.GetEdges(startNode).Keys)
     shortestPath[node] = int.MaxValue;
   shortestPath[startNode] = 0;
   while (true)
     string currentNode = null;
     int minDistance = int.MaxValue;
     foreach (var node in graph.GetEdges(startNode).Keys)
        if (!visited.Contains(node) && shortestPath[node] < minDistance)
          currentNode = node;
          minDistance = shortestPath[node];
     if (currentNode == null)
        break;
     visited.Add(currentNode);
```

```
foreach (var kvp in graph.GetEdges(currentNode))
      string neighbor = kvp.Key;
      int weight = kvp.Value;
      if (shortestPath[currentNode] + weight < shortestPath[neighbor])</pre>
        shortestPath[neighbor] = shortestPath[currentNode] + weight;
  return shortestPath;
static void PrintResult(Dictionary<string, int> shortestPath, string startNode, string endNode)
 List<string> path = new List<string>();
 string currentNode = endNode;
  while (currentNode != startNode)
    path.Add(currentNode);
    foreach (var kvp in shortestPath)
      if (kvp.Key == currentNode)
        currentNode = kvp.Value.ToString();
        break;
  path.Add(startNode);
  path.Reverse();
  Console.WriteLine($"Shortest Path from {startNode} to {endNode} with a value of {shortestPath[endNode]}:");
  Console.WriteLine(string.Join(" -> ", path));
```

A*



A* is an <u>informed search</u> algorithm. <u>Informed Search</u> signifies that the algorithm has extra information, to begin with. It is a complete as well as an optimal solution for solving path and grid problems.

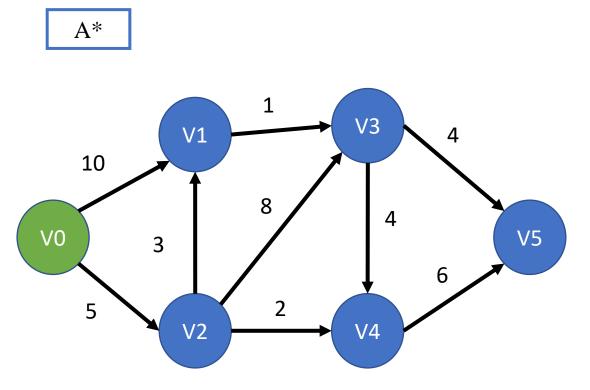
Optimal – find the least cost from the starting point to the ending point. Complete – It means that it will find all the available paths from start to end.

The Algorithm

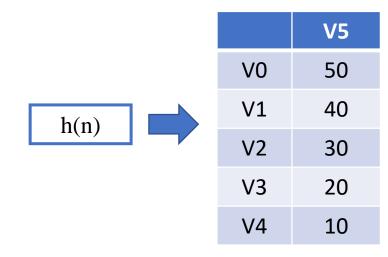
1. Place the starting node into OPEN and find its f (n) value.

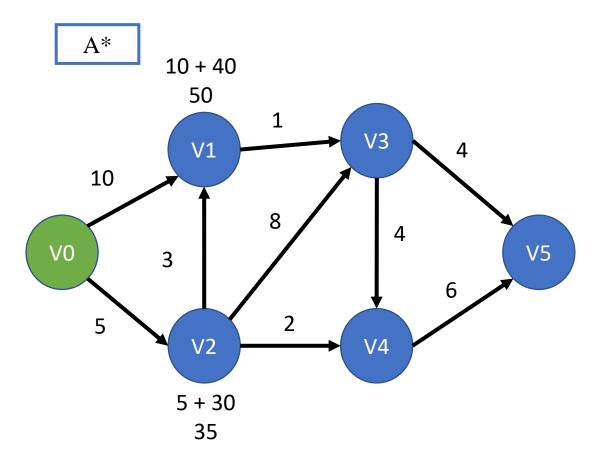
f(n) = g(n) + h(n)

- 2. Remove the node from OPEN, having the smallest f (n) value.
 - * If it is a goal node, then stop and return to success.
 - * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.

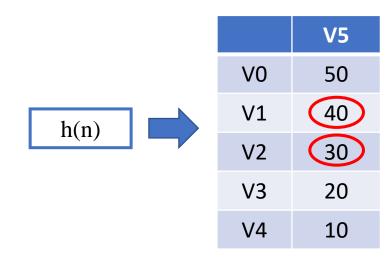


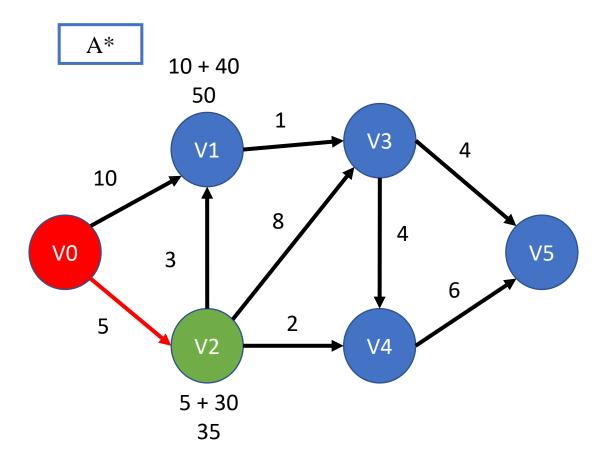
- 1. Place the starting node into OPEN and find its f (n) value.
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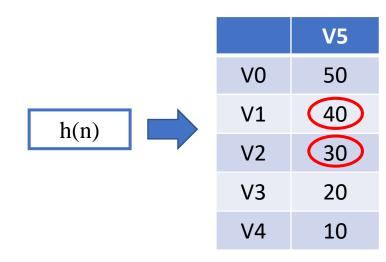


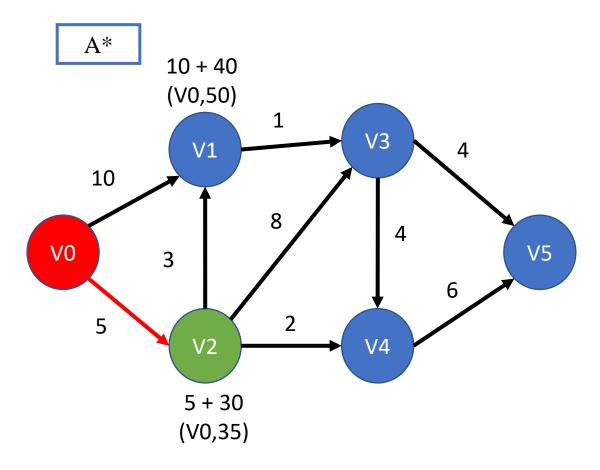
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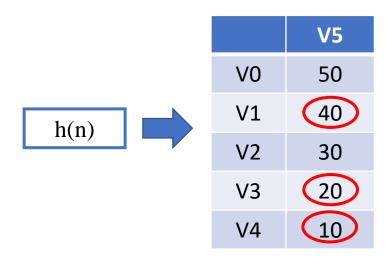


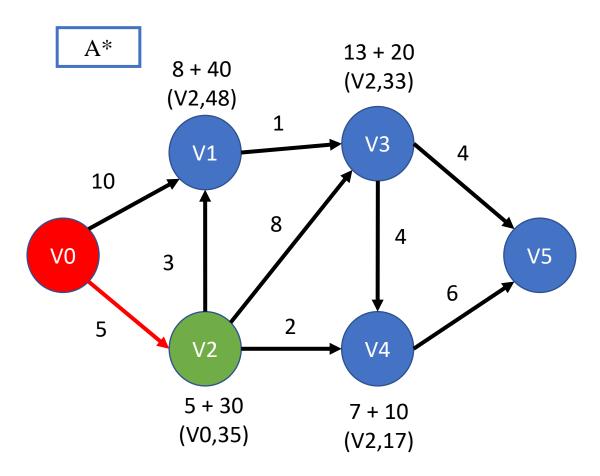
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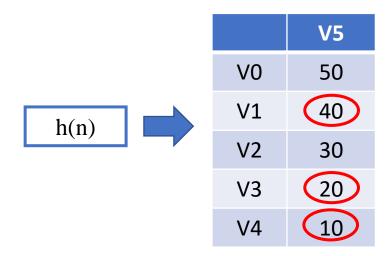


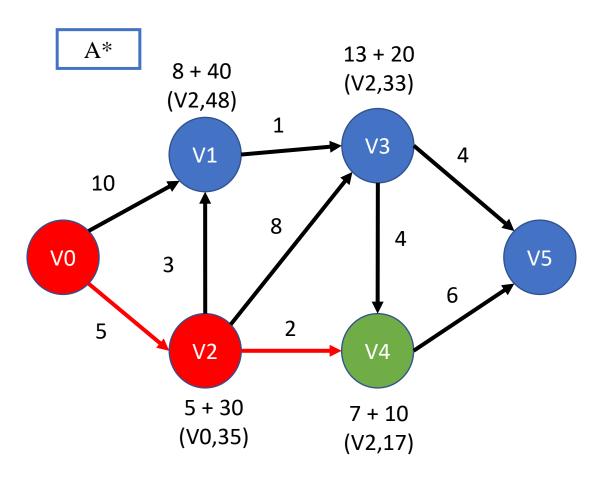
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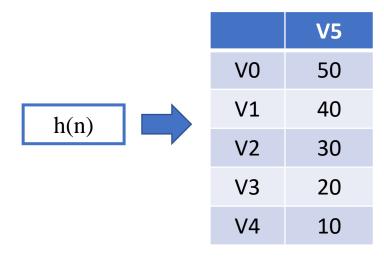


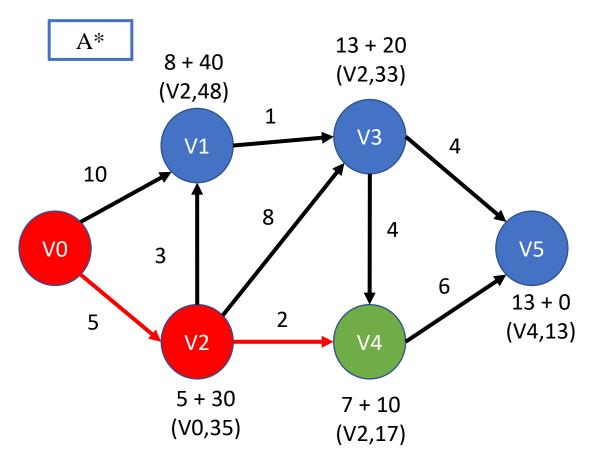
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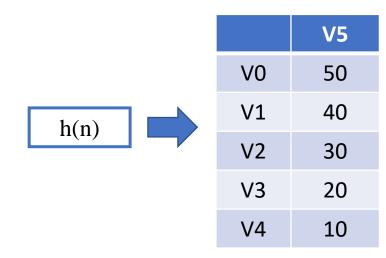


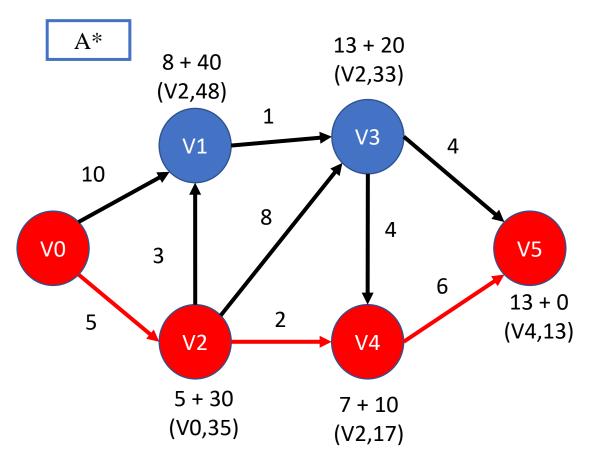
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- 2. Remove the node from OPEN, having the smallest f (n) value.
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- * Else remove the node from OPEN, and find all its successors.
- * Find the f (n) value of all the successors, place them into OPEN, and place the removed node into CLOSE.



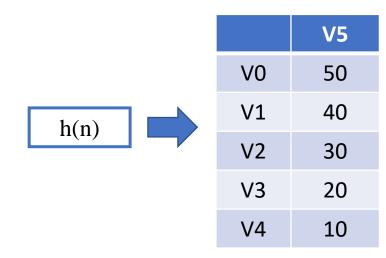


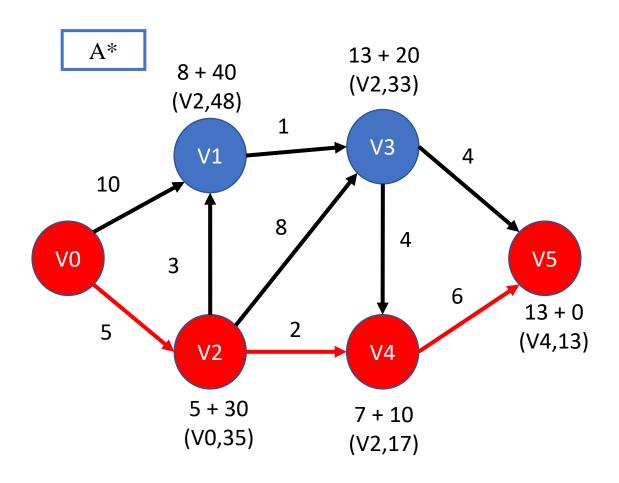
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 - * If it is a goal node, then stop and return to success.
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the shortest distance V0 => V5



V5, V4, V2, V0

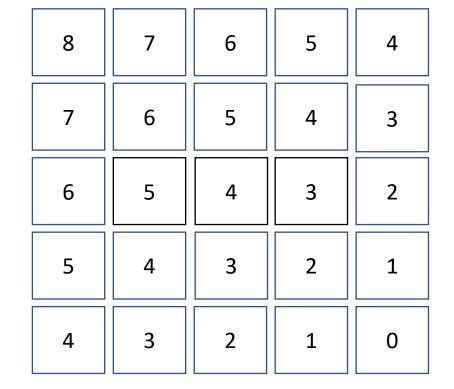
A*

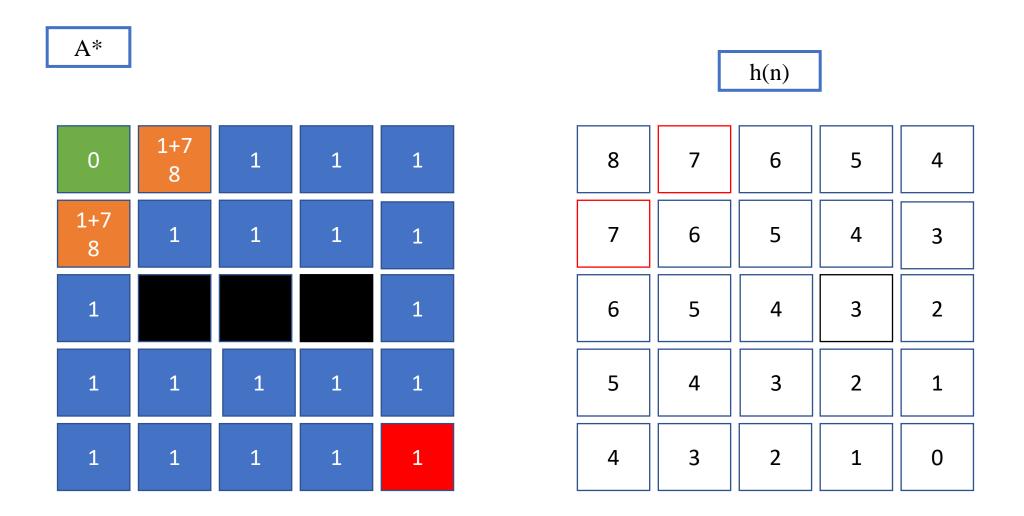
Algorithms

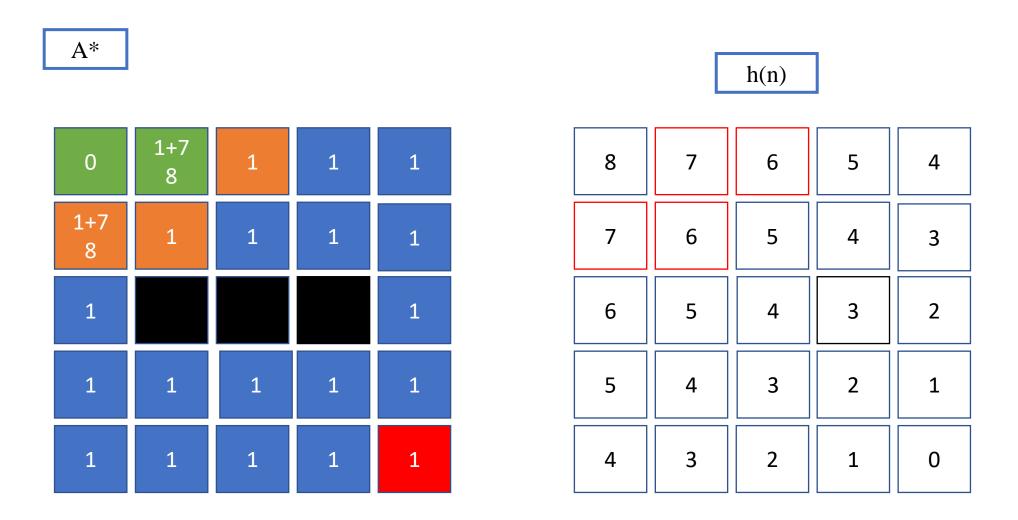
manhattan((x1, y1), (x2, y2)) = |x1 - x2| + |y1 - y2|euclidean((x1, y1), (x2, y2)) = $sqrt(x^2 + y^2)$

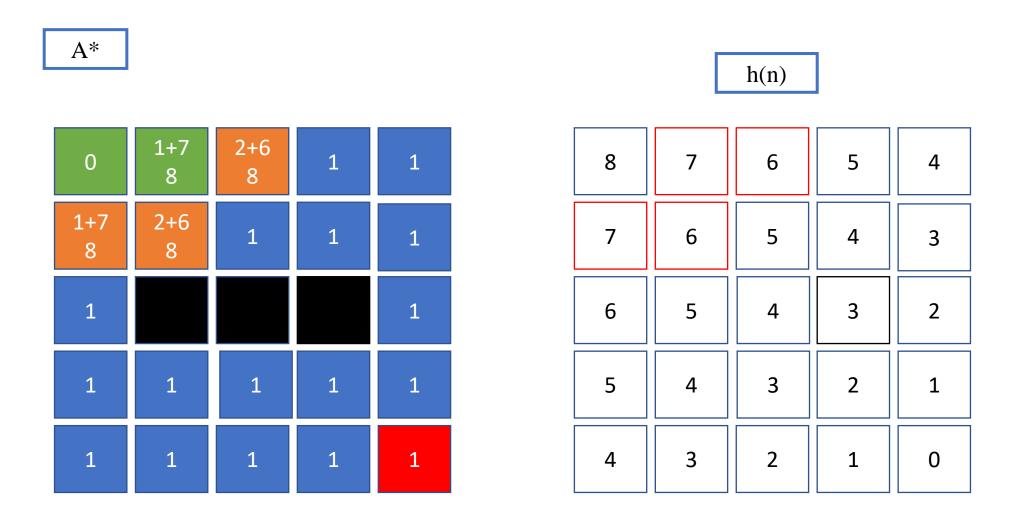
h(n)

| 0 | 1 | 1 | 1 | 1 |
|---|---|---|---|---|
| 1 | 1 | 1 | 1 | 1 |
| 1 | | | | 1 |
| 1 | 1 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 |











 A^* h(n) 1+7 2+6 3+5 8 5 4 1+7 2+6 3+5 7 6 4 3 6 5 3 4 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 8 5 4 1+7 2+6 3+5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 1+7 2+6 3+5 4+4 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 1+7 2+6 3+5 4+4 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6 5 3 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6+2 6 5 3 4 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6+2 6 5 3 7+1 5 2 4 4 0

 A^* h(n) 1+7 2+6 3+5 4+4 8 5 4 4+4 1+7 2+6 3+5 5+3 5 7 6 4 3 6+2 6 5 3 4 8+2 7+1 5 2 4 1 10 8+0 4 0 8



Daniel Nogueira

dnogueira@ipca.pt