

# Midterm Exam 1

David Robinson

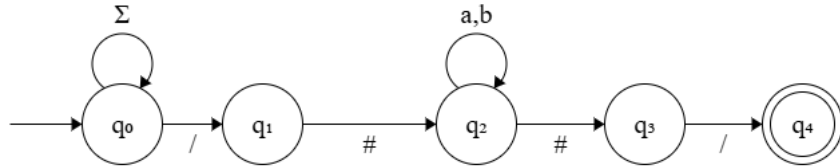
## Problem 1

$$(a \cup d)^* c (b (dd)^* \cup c \cup d)^* b$$

## Problem 2

Let  $M$  be a DFA that recognizes  $L$ . We can construct an NFA  $M'$  to recognize  $h(L)$  by replacing each transition  $\delta(q, a)$  in  $M$  with a sequence of transitions in  $M'$  that process  $h(a)$  instead of  $a$ . Since NFAs and DFAs both define the class of regular languages, the resulting language  $h(L)$  is regular. Therefore, regular languages are closed under homomorphism.

## Problem 3



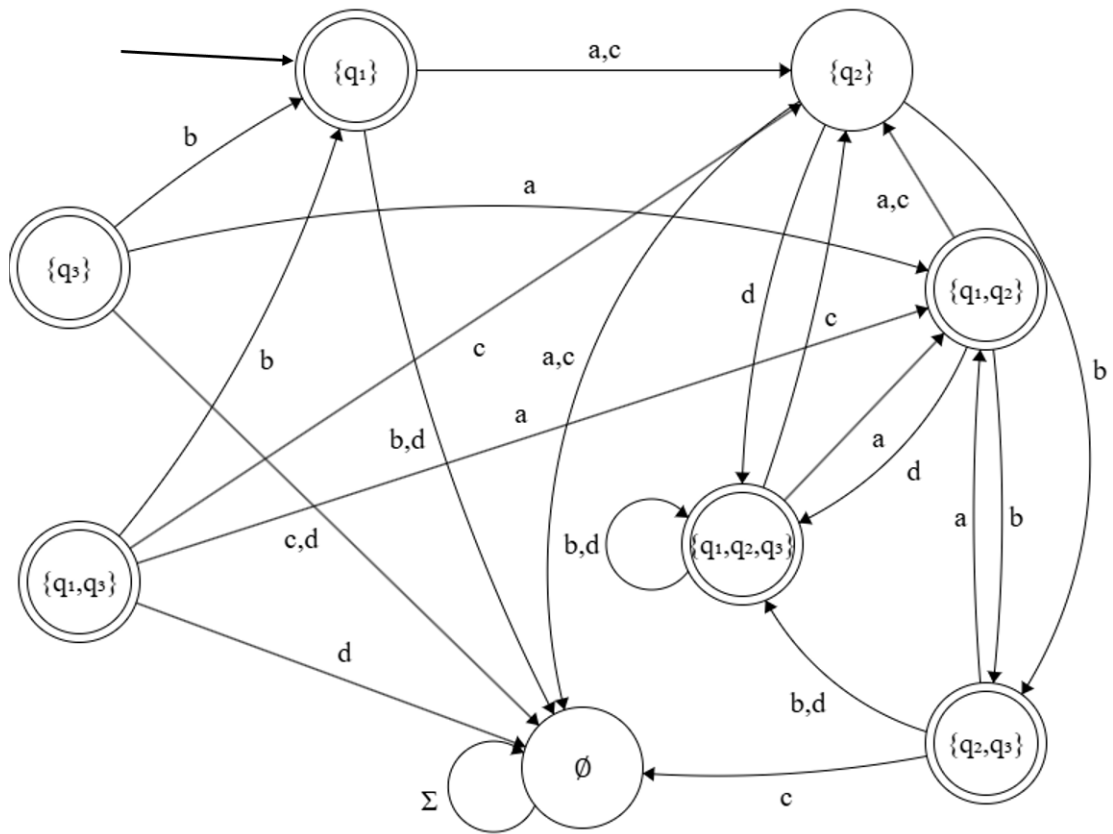
## Problem 4

If  $L = \{ccc\#cc\#c \mid c \in \{a,b\}^*\}$  is a regular language, then it must satisfy the Pumping Lemma, where there exists a pumping length  $p$  such that any string  $s$  in  $L$  with  $|s| \geq p$  can be split into  $s = xyz$  where:

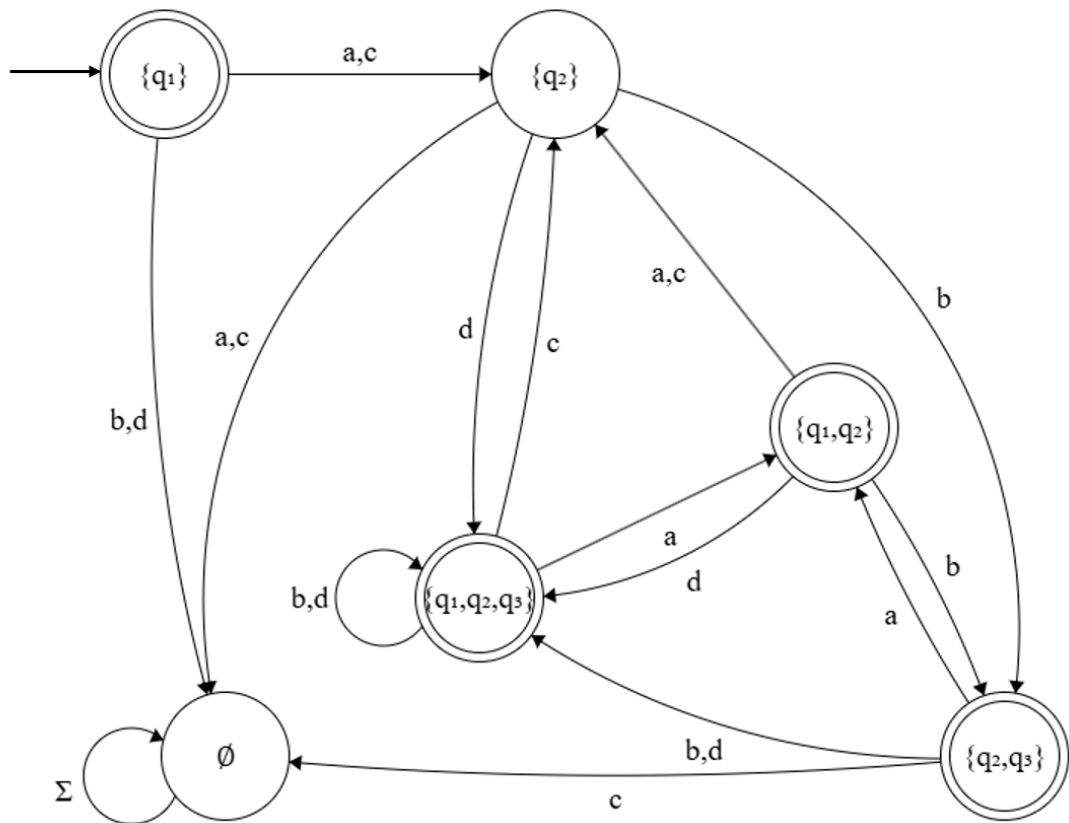
1.  $xy^kz \in L$  for all  $k \geq 0$
2.  $|y| > 0$
3.  $|xy| \leq p$

Consider a string  $s = a^p b^p a^p b^p a^p b^p \# a^p b^p a^p b^p \# a^p b^p$  in  $L$  where  $c = a^p b^p$ . Pumping  $y$  increases the number of  $a$ 's only in the first  $c$ , while the other instances of  $c$  are unchanged. Since the resulting string no longer follows the pattern  $ccc\#cc\#c$ , it contradicts the Pumping Lemma and  $L$  is not a regular language.

### Problem 5



### Problem 6



**Problem 7**

