

# Modeling Democracy

## Lecture 2 - **Voting rules and properties**



# Announcements

- **Assignment requirements:** PhD - 5 HW + project or presentation  
MS/Ugrad: 5 HW, where presentation can replace one HW and a project is fully optional  
After the first one, HW will be due Fridays on gradescope.
- **Slides and materials:** no Canvas, but we'll set up a github today. That will also include a link to a friendly intro-level social choice textbook.
- **Office Hours:** Alec will be available immediately after class on both Mon and Wed. Moon's OH TBD.
- **Fridays:** hold 1:30-2:30 (same as class time MW) for presentations. Attending these is "extra credit" for MS/Ugrad.
- **Skills:** overleaf, VoteKit. VoteKit trainings next week! Tue 10-12, Thu 1-3, Fri 1:30-3:30.

# Key concepts recap

- **Ballot** — for now, ranking: a (possibly truncated) permutation of  $m$  candidates
- **Preference profile** — table of ballots
- **Voting rule** — maps profiles to outcomes (e.g., rankings to winner sets)
- **Condorcet candidate** — one who beats all others H2H (head-to-head)
- **Pairwise comparison graph** (PWCG) / tournament graph — directed complete graph on  $\mathcal{C}$  with edges weighted by H2H margins
- **Condorcet cycle** — directed cycle in PWCG
- **Threshold** or **quota** — number/fraction of votes needed to be elected in IRV/STV

# Reminder: Exercises from last time

- A Condorcet 3-cycle exists on 3 candidates when some three cyclically permuted columns have been selected  $n_1, n_2, n_3$  times and those numbers satisfy the triangle inequality
- If a preference profile consists of complete rankings, then all margins in the PWCG have the same *parity* (all even or all odd)

# IRV/STV recap

- You fix a threshold or quota when electing for  $k$  seats — common choice is “Droop quota” of  $T = \frac{N}{k+1}$  (possibly rounded) for a profile with  $N$  valid ballots.
- An alternative is “Hare quota”  $N/k$  — variants include dynamical quota, changing by round
- Round 1: Suppose highest current candidate support is  $M$ .
  - If  $M > T$  change candidate status to “elected.” Each vote currently supporting them is transferred with weight  $\frac{M-T}{M}$  (share of support that was surplus).
  - Else, eliminate candidate with least support and transfer fully.
- Repeat in later rounds until  $k$  seats are filled.

Plurality	Positional (Borda)	Sequential
Pairwise (Lhull)	Condo-Borda	Top-Two
IRV	Coombs	STV
Beatpath (Schulze)	Smith	Smithified rules
Dodgson	Kemeny	“Secondality”

- **Plurality** — most first-place votes (FPV) wins
- **Positional (Borda)** — voter gives out points by position of ranking (e.g., 3 for 1st, 2 for 2nd, 1 for 3rd)
- **Sequential** — order fixed in advance, candidates face each other pairwise with one advancing based on head-to-head (H2H) comparison



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- **Pairwise** — each candidate gets 1 point per H2H win, 1/2 per tie
- **Condo-Borda** — Condorcet candidate if they exist, else Borda winner (a “Frankenstein” rule)
- **Top-Two** — two candidates with most FPV advance to instant runoff, decided by H2H



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- **IRV (instant runoff voting)** — successively eliminate those with fewest FPV, transferring their votes, until someone has a majority
- **Coombs**— same, but eliminate those with most last-place votes
- **STV (single transferable vote)** — set a threshold of election (typically  $\frac{1}{k+1}$  or so) and successively elect and eliminate until  $k$  are elected



# Beatpaths and domsets

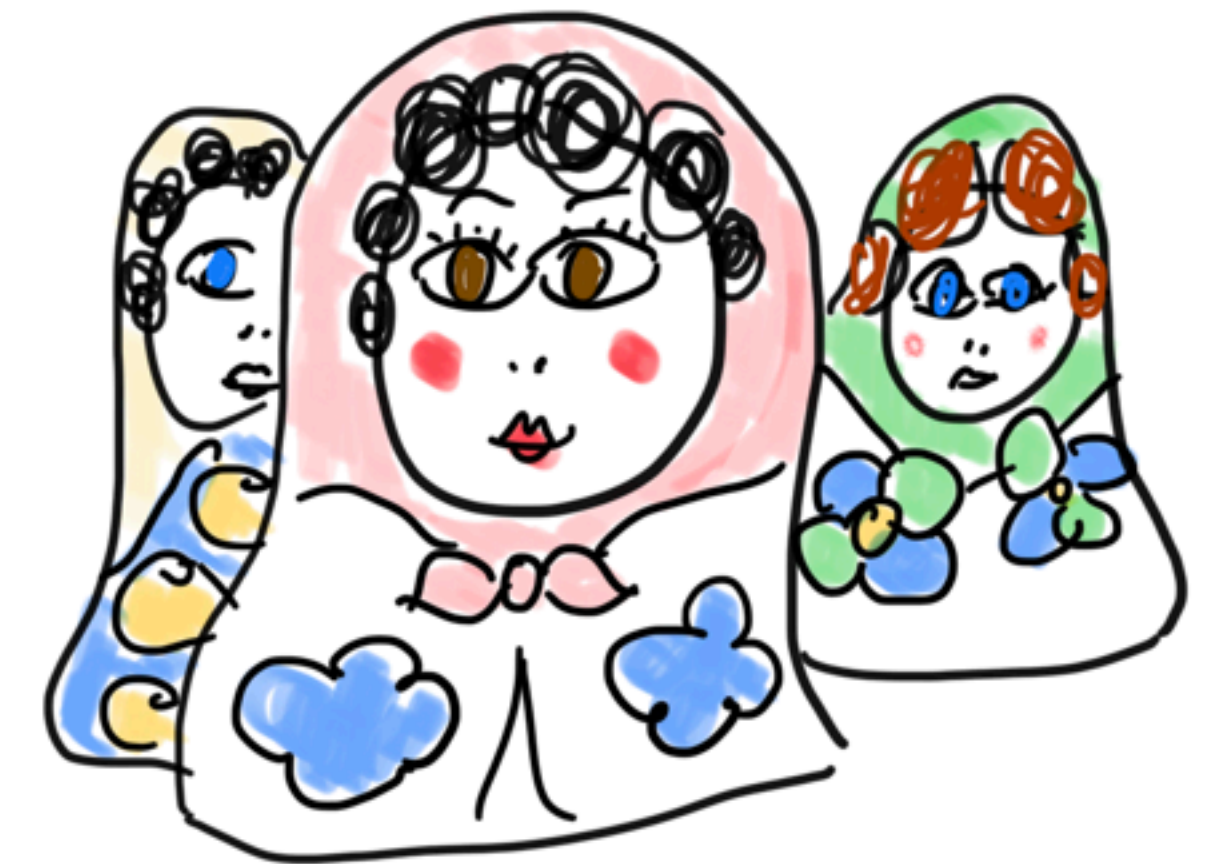
**Beatpath**: a directed path in the pairwise comparison graph

**Strength** of a beatpath: its lowest margin

Beatpath **elimination**: A eliminates B (denoted  $A \triangleright B$ ) if A has a beatpath to B (strictly) stronger than any beatpath from B to A

Dominating set ("**domset**"): a nonempty subset of candidates that beats everyone else head-to-head

Fact you will prove: Domsets are **nested**!  
So there is a smallest domset formed by intersecting them all.  
This is called the **Smith set**, and its members are called "strong" candidates (and others are called "weak").



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- **Beatpath** — candidate A eliminates B ( $A \triangleright B$ ) if there is a beatpath from A to B in pairwise comparison graph stronger than any path from B to A. Winner set is anyone not eliminated after all pairs are considered.
- **Smith**— candidates stratified into dominating tiers, where a dominating set has “all arrows pointing out”; Smith set / strong candidates/ winners are members of smallest domset
- **Smithified rules** — Smithified X: first consolidate to Smith set, then run method X

# Transitivity

Recall that a relation  $R$  is transitive if  $X_R Y, Y_R Z \implies X_R Z$ .

(Examples: *taller than*, *divisible by*. Nonexamples: *sibling of*.)

Head-to-head wins are not transitive — this is the **Condorcet paradox**. That means that if you tried to eliminate a candidate from contention because someone else beat them head-to-head, there would not be a well-defined outcome; it would depend on the order you considered them in.

However, as you will prove, beatpath elimination IS transitive, and this is enough to ensure that a beatpath election has at least one winner and that the winners are independent of the order of pairwise comparison.

# Swaps and such

Let's say that two ballots differ by a “**neighbor swap**” if one can be transformed to the other by exchanging candidates in neighboring positions, such as

$$\begin{pmatrix} A \\ B \\ C \\ D \end{pmatrix} \rightarrow \begin{pmatrix} A \\ C \\ B \\ D \end{pmatrix}$$

This can be considered an **elementary move** on ballots. (Another reasonable example of an elementary move is deleting a candidate name from the end of your ballot, or adding a name to the bottom of a partial ballot.)

You can get from any ballot to any other ballot this way. For complete ballots, this can be phrased in mathematical language by noting that neighbor swaps are a **generating set** for the symmetric group  $S_n$ .



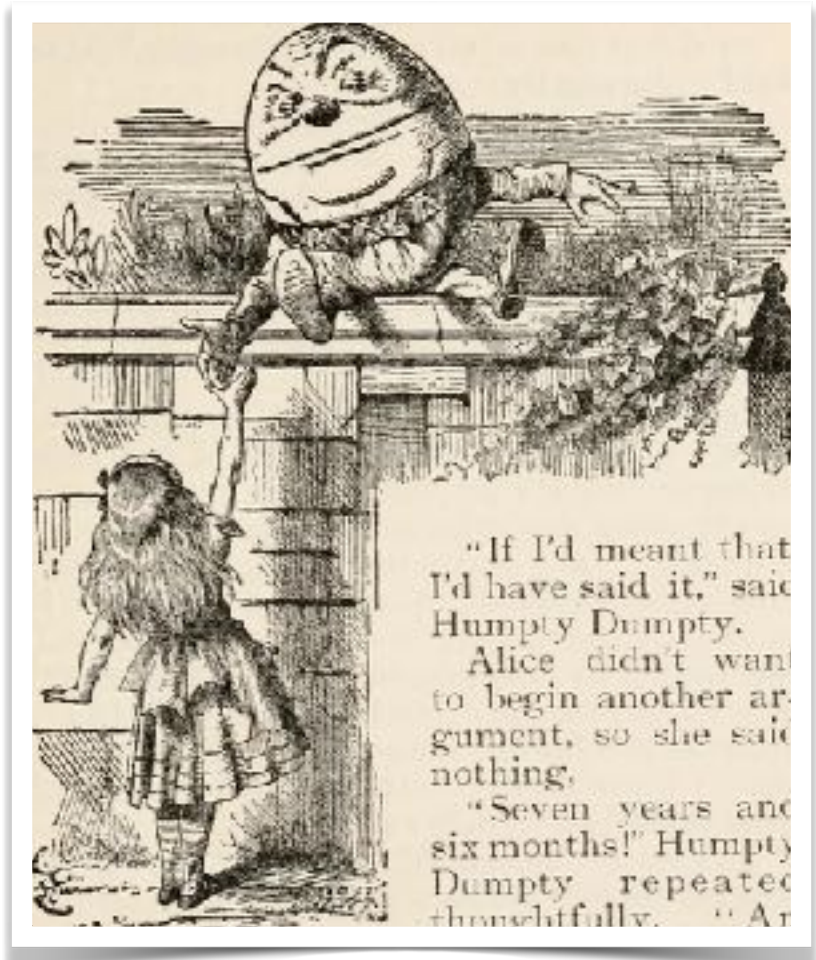
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- **Dodgson** — winner is candidate who requires the fewest voter ballot swaps to become Condorcet

- **Kemeny**— put a metric on ballots, say by swap distance/bubble sort. Winning ranking is minimizer of sum of distances to cast ballots.

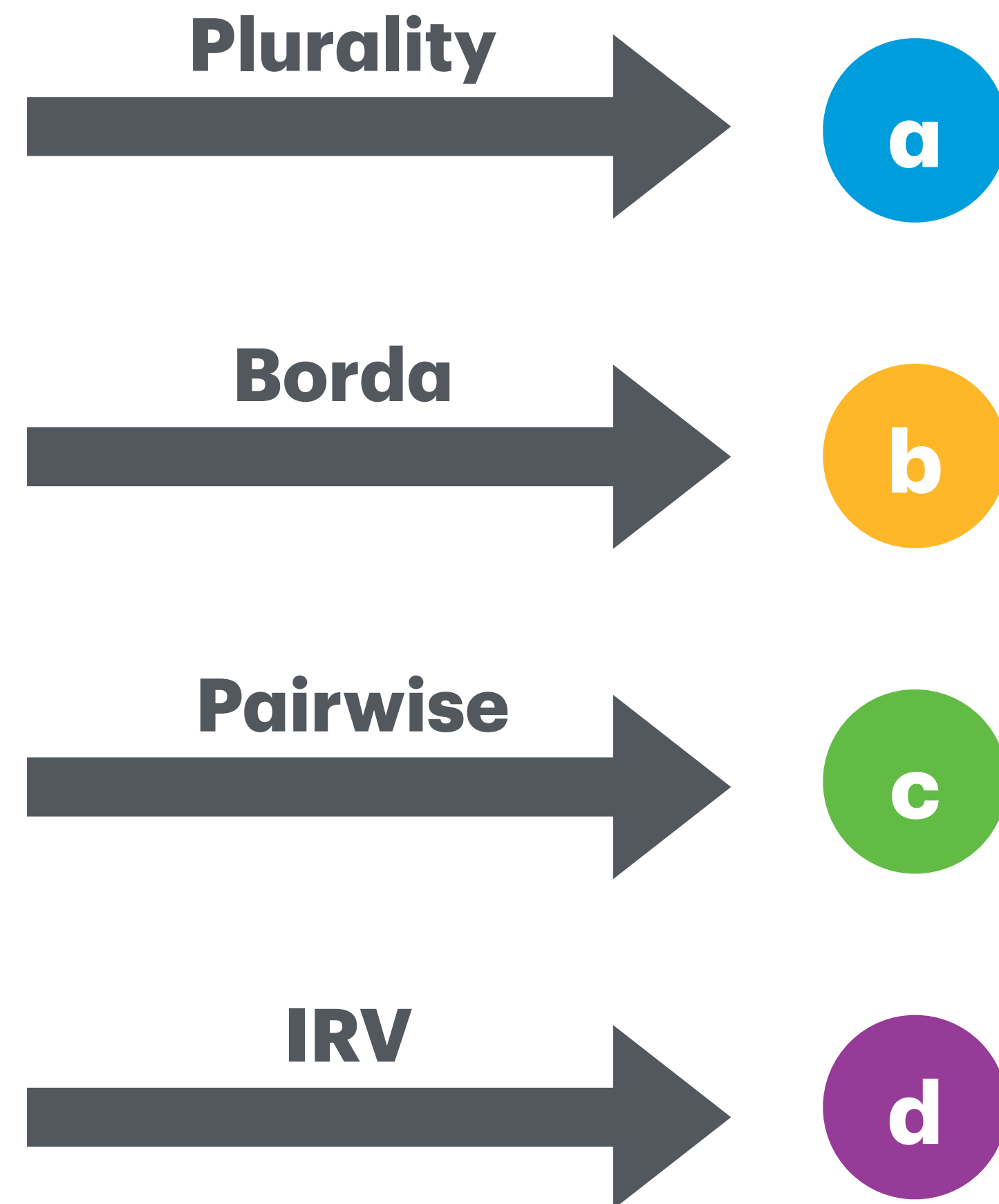
- **Secondality** —most second-place votes wins!



ok so who *should*  
win this one:



33 voters	16 voters	3 voters	8 voters	18 voters	22 voters
<i>a</i>	<i>b</i>	<i>c</i>	<i>c</i>	<i>d</i>	<i>e</i>
<i>b</i>	<i>d</i>	<i>d</i>	<i>e</i>	<i>e</i>	<i>c</i>
<i>c</i>	<i>c</i>	<i>b</i>	<i>b</i>	<i>c</i>	<i>b</i>
<i>d</i>	<i>e</i>	<i>a</i>	<i>d</i>	<i>b</i>	<i>d</i>
<i>e</i>	<i>a</i>	<i>e</i>	<i>a</i>	<i>a</i>	<i>a</i>



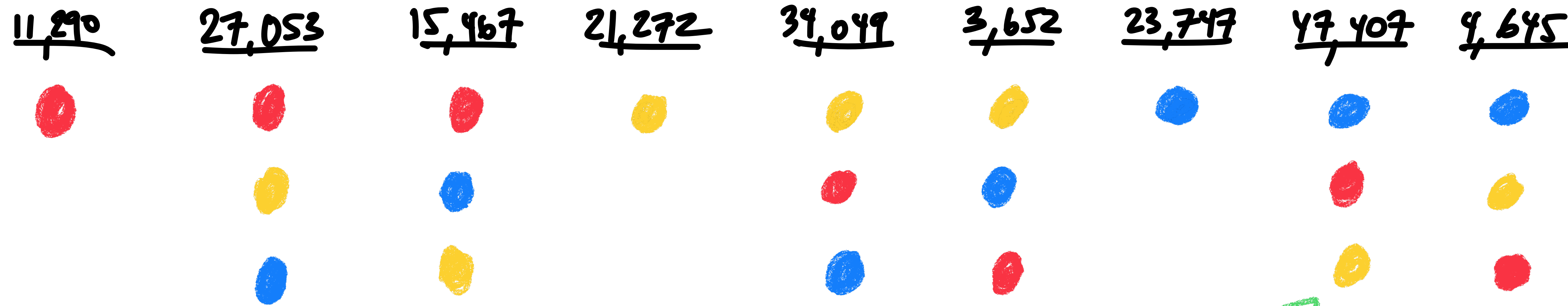
# What are we after? Axioms of good voting rules

*(Classic)  
fairness criteria*

Majority-fair	Condorcet-fair	Pareto efficient
Unanimity-fair	Monotonic	Strongly monotonic
Strategy-proof	No (weak) spoilers	Smith-fair
Voter anonymity	Candidate anonymity	Independence of irrelevant alternatives

- **Majority-fair** — in a single-winner contest, a strict majority of FPV guarantees a win
- **Condorcet-fair** — if there is a Condorcet candidate, they win
- **Pareto efficient** — if everyone has same first-place choice, they win (or are ranked first if the output is a ranking)

2022 Special Election by IRV for U.S. Congress from Alaska



Begich

Palin

Peltola

- Begich is Condorcet
- Begich is first to be eliminated
- Peltola easily defeats Palin H2H and wins





# IOP FELLOWS



**Dan Caldwell**  
Former Advisor to  
the U.S. Secretary  
of Defense



**Yangyang Cheng**  
Research Scholar  
in Law and Fellow  
at Yale Law  
School's Paul  
Tsai China Center



**Mary Peltola**  
Former U.S.  
Representative  
from Alaska



**David Pressman**  
Former U.S.  
Ambassador to  
Hungary



**Alex Wagner**  
Former Advisor to  
the U.S. Secretary  
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**Vince Warren**  
Executive Director  
of the Center for  
Constitutional  
Rights

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## SEMINAR SCHEDULE

MONDAYS	TUESDAYS	WEDNESDAYS	THURSDAYS
		12:30 – 1:45 PM DAN CALDWELL	12:30 – 1:45 PM ALEX WAGNER
3:30 – 4:45 PM DAVID PRESSMAN	3:30 – 4:45 PM VINCE WARREN		

WEEK OF OCT 20

MARY PELTOLA

WEEKS OF NOV 3 & NOV 10

YANGYANG CHENG



REGISTER FOR  
SEMINARS



SIGN UP FOR  
ONE-ON-ONE  
OFFICE HOURS





# More axioms

(Classic)  
fairness criteria

Majority-fair	Condorcet-fair	Pareto efficient
<b>Unanimity-fair</b>	<b>Monotonic</b>	<b>Strongly monotonic</b>
Strategy-proof	No (weak) spoilers	Smith-fair
Voter anonymity	Candidate anonymity	Independence of irrelevant alternatives

- **Unanimity property**— for single-winner, if everyone ranks  $X > Y$  then  $Y \notin \mathcal{W}$ ; for ranking output, if everyone ranks  $X > Y$  then society ranks  $X > Y$
- **Monotonic** — if  $X \in \mathcal{W}$  and profile changes only by moves favorable to  $X$ , then  $X \in \mathcal{W}'$
- **Strongly monotonic** — if  $X \in \mathcal{W}$  and profile changes only by moves favorable to or neutral to  $X$ , then  $X \in \mathcal{W}'$

# More axioms

(Classic)  
fairness criteria

Majority-fair	Condorcet-fair	Pareto efficient
Unanimity-fair	Monotonic	Strongly monotonic
<b>Strategy-proof</b>	<b>No (weak) spoilers</b>	<b>Smith-fair</b>
Voter anonymity	Candidate anonymity	Independence of irrelevant alternatives

- **Strategy-proof** — given a profile  $P$  with ballot  $B$ , there should be no  $B', P'$  so that  $f(P')$  is preferable according to  $B$
- **Spoilers** —  $X$  is **not** a spoiler if, when  $P, P'$  differ by removal of candidate  $X$  from  $P$ , one of these holds. If none of these, then  $X$  is a spoiler.  
(a)  $\mathcal{W} = \mathcal{W}'$ ,   (b)  $\mathcal{W} = \{X\}$ ,   (c)  $\mathcal{W}' = \mathcal{W} \setminus \{X\}$ .

- **Smith-fair** — only the strong can win. ( $\mathcal{W} \subseteq \mathcal{S}$ )

“weak spoiler” — a spoiler who is outside the Smith set



# More axioms

*(Classic)  
fairness criteria*

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Strategy-proof	No (weak) spoilers	Smith-fair
<b>Voter anonymity</b>	<b>Candidate anonymity</b>	<b>Independence of irrelevant alternatives</b>

- **Voter anonymity** — independence of voter names or order
- **Candidate anonymity** — independence of candidate names or order
- **IIA** (for rankings) — whether  $X > Y$  in final ranking depends only on whether  $X > Y$  on ballots