

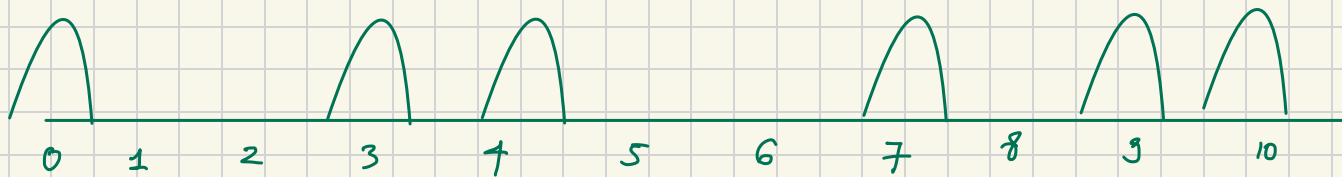


Aggressive Cows

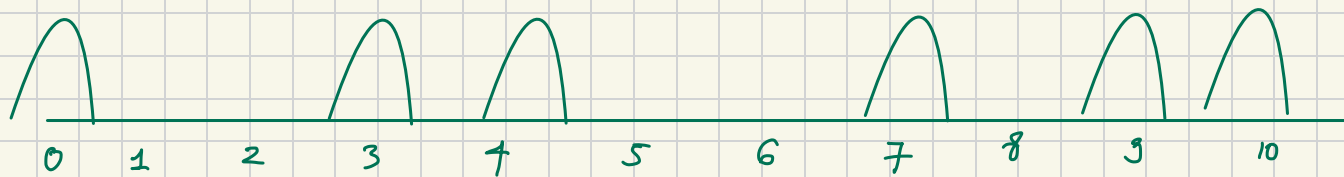
stalls $[] = \{ \overset{0}{0}, \overset{1}{3}, \overset{2}{9}, \overset{3}{7}, \overset{4}{10}, \overset{5}{4} \}$

cows = 4

inc. distance aggressive cows.



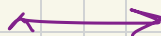
$\rightarrow \max^m \left\{ \min^m \text{distance b/w any two cows} \right\}$



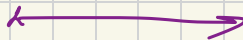
\min^n



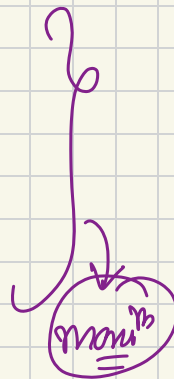
1



2

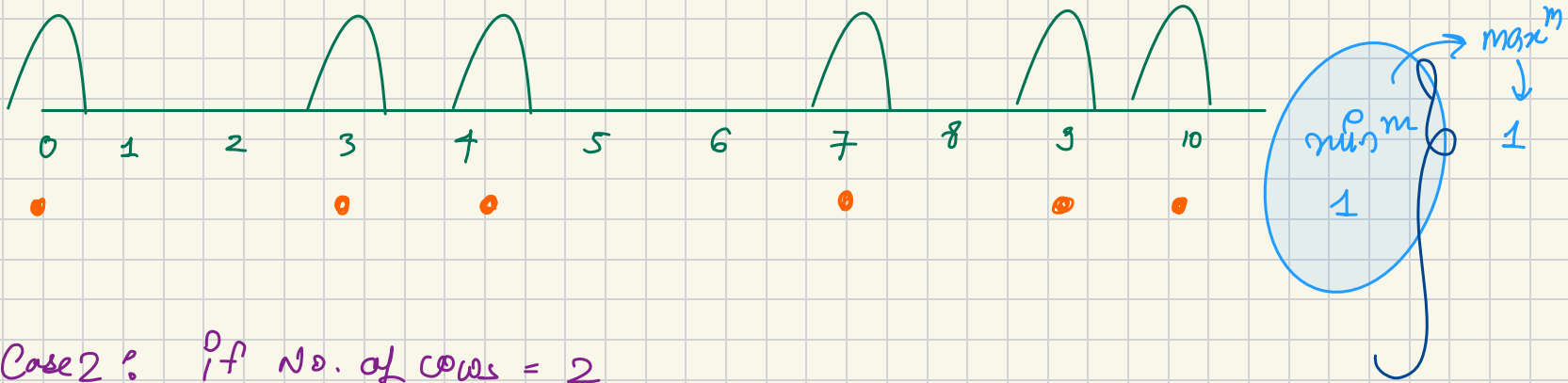


3

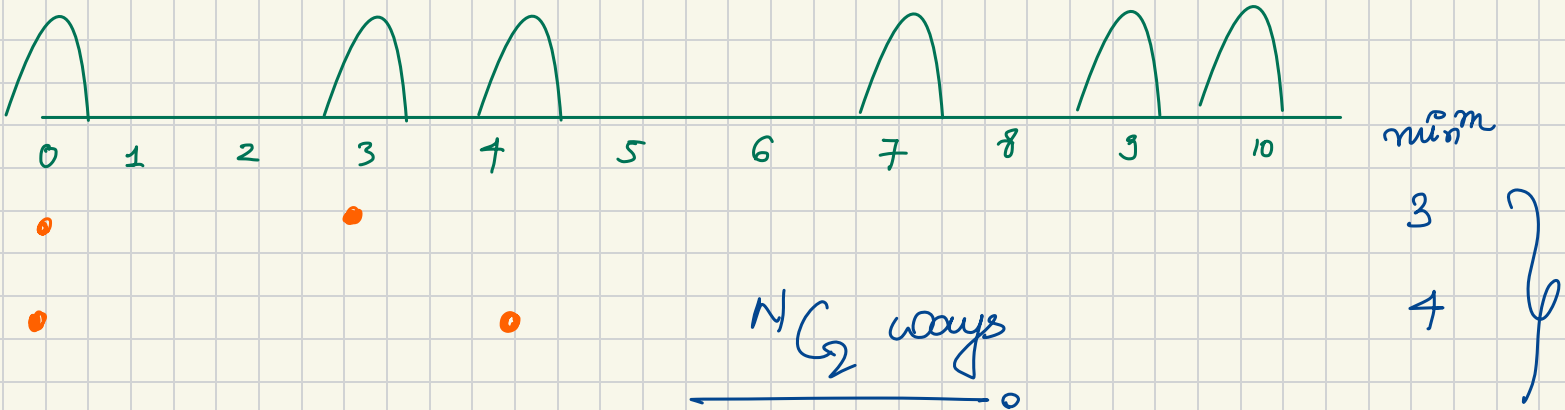


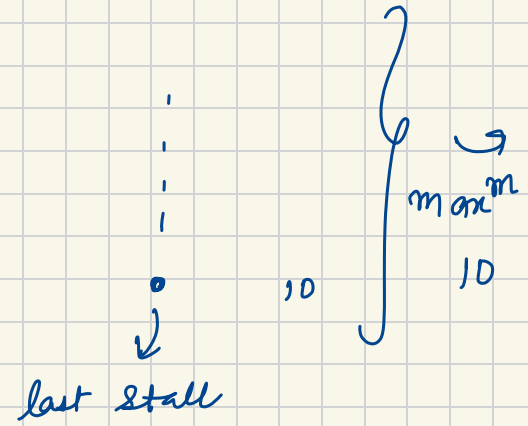
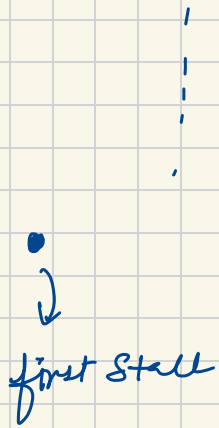
$N C_m$ ways to stone cows 0

Case 1: if No. of Cows = N ✓



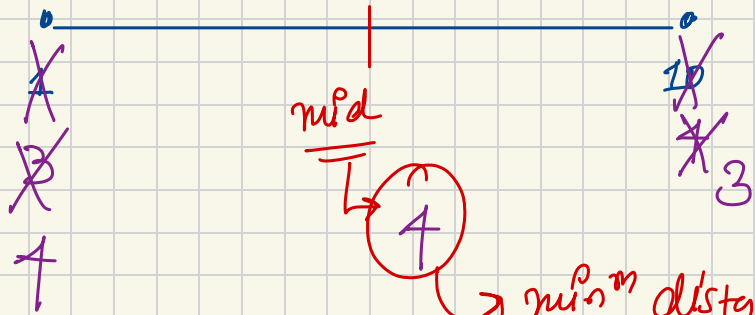
Case 2: if No. of cows = 2





got our range!

cols = 4

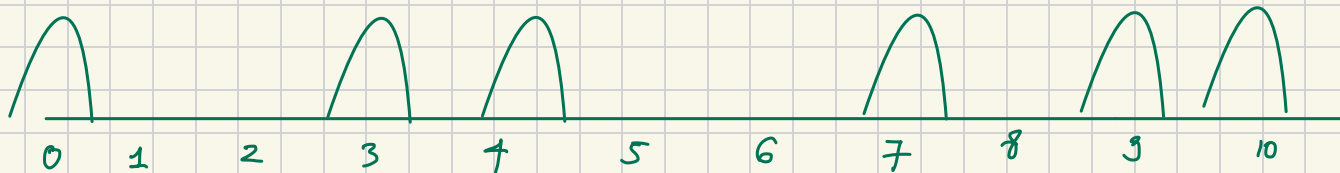


min distance b/w any two cows!

rows = 3

ans = 3

TC: $O(N \log_2 N)$
 SC: $O(1)$



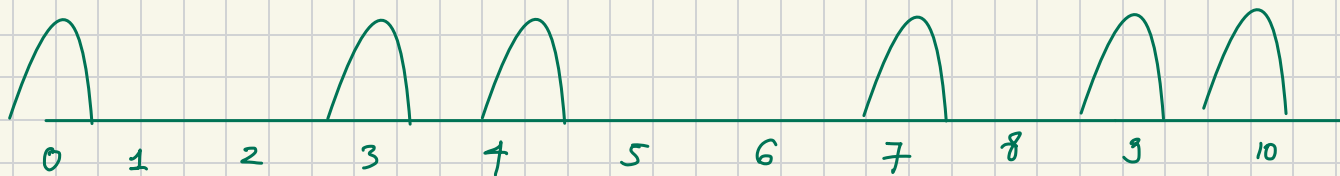
$$\text{numCap} = 5$$

•

C_1

•

C_2



$$\text{mid cap} = 2$$

•

C_1

•

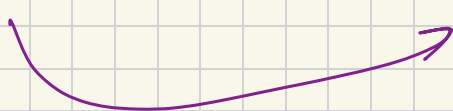
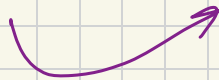
C_2

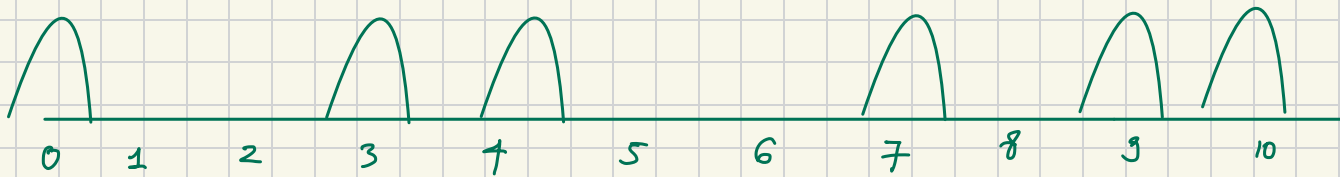
•

C_3

•

C_4





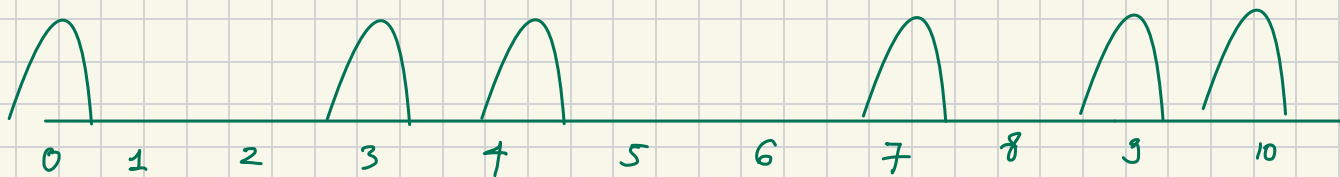
maxlap
= 3

C_1

C_2

C_3

C_4



maxlap = 4

C_1

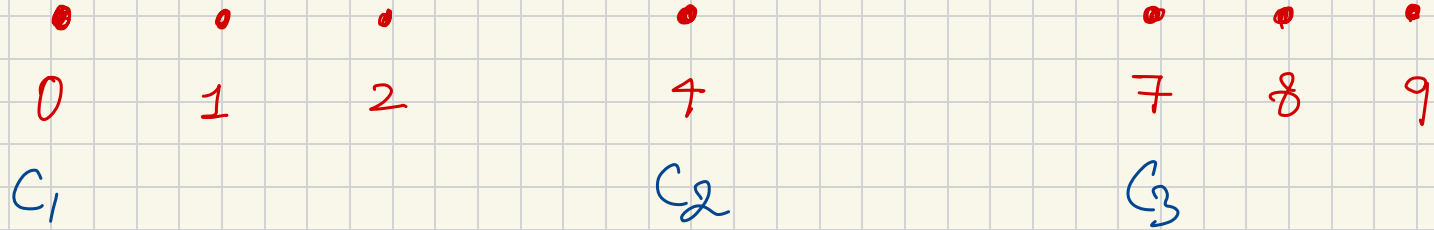
C_2

C_3

Binary Search ✓

impossible ✓





$$\text{min}^m = 3$$

$$\rightarrow \text{lastPositionPlaced} = \cancel{0} \cancel{1} 7$$

$\{0, 1, 2, 5, 7, 9, 10, 14\}$

min^m dist = 4

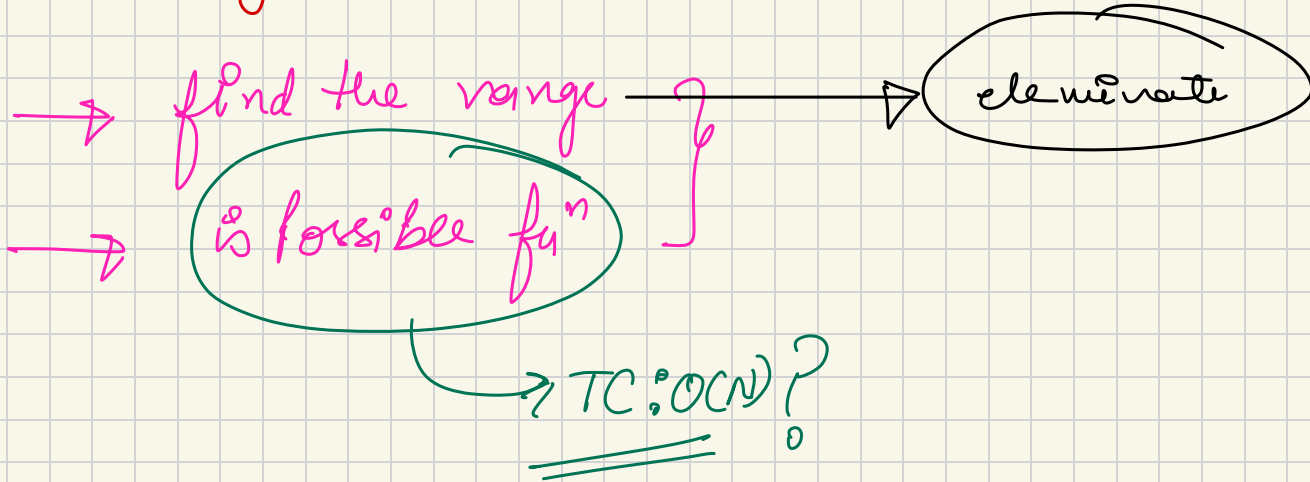
$\overset{\cdot}{i} \rightarrow \rightarrow \overset{\cdot}{i} \rightarrow \rightarrow \overset{\cdot}{i} \rightarrow \rightarrow \overset{\cdot}{i}$

no. of cows placed = $\cancel{0} \cancel{1} \cancel{2} \cancel{5} 7 \checkmark$

last position placed = $\cancel{0} \cancel{1} \cancel{2} \cancel{5} 14$

TC: $O(N)$ is possible function!

Binary Search Ours Solⁿ



$$\log_2(10^8) = 8 \times \log_2 10^3 = 24$$

$$\log_2(10^4) = 4 \times \log_2 10^3 = 12$$

$$\log_2(100) = 2 \times \log_2 10^3 = 6$$



i

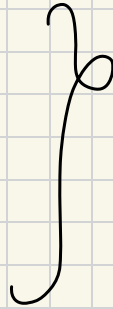
10

dist \rightarrow

$\left\{ \begin{array}{l} \min^m \text{dist} = 1 \quad \{ \text{adjacent to each other} \} \\ \max^m \text{dist} = 10^6 \quad \{ \text{far apart} \} \end{array} \right.$



- Sorted Array
- Rotated Sorted Array
- Mountain Array



• Infinite Array

✓ Binary Search!

Binary Search

$$\text{arr} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

→ define range!

→ remove half.

$$\begin{aligned} \text{TC: } &O(\log_2 N) \\ \text{SC: } &O(1) \end{aligned}$$

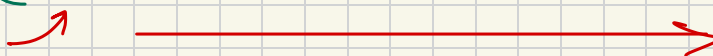
Search region

$$N \rightarrow \frac{N}{2} \rightarrow \frac{N}{4} \rightarrow \dots \rightarrow 1$$

Suppose,

you got a infinite sorted array!

{ 1, 3, 10, 11, 15, 17, 20, 25, - - - - - } ^{infinite ele.}



target = 10/4

Brute force
Linear Search

$\{ \overset{0}{1}, \overset{1}{3}, \overset{2}{10}, \overset{3}{11}, \overset{4}{15}, \overset{5}{17}, \overset{6}{20}, \overset{7}{25}, \overset{8}{30}, \overset{9}{35}, \dots, \overset{16}{103}, \infty \}$

target = 90

lo

hi

Range $[0, 1] \xrightarrow{\times} 2 \quad \{ arr[lo] \leq target \leq arr[hi] \}$

Range $[2, 2] \rightarrow 1 \sim 2^0$

Range $[3, 4] \rightarrow 2 \sim 2^1$

Range $[5, 8] \rightarrow 4 \sim 2^2$

Range $[9, 16] \rightarrow 8 \sim 2^3$

$lo = hi + 1$

$hi = 2 * hi$

$\{ \overset{0}{1}, \overset{1}{3}, \overset{2}{10}, \overset{3}{11}, \overset{4}{15}, \overset{5}{17}, \overset{6}{20}, \overset{7}{25}, \overset{k^{th}}{\text{...}} \dots \infty \}$

$x = (p-1)^{th}$
 lo hi

$$lo = x + 1 = p$$

$$hi' = 2 * x = 2(p-1)$$

p^{th} Index

$O(\log_{f_2} P)$

$$\text{No. of Ele} = 2(p-1) - p + 1 = (p-1) \text{ Elements}$$

$$1^{\text{st}} \text{ Range} = 2$$

$$2^{\text{nd}} \text{ Range} = 1 = 2^0$$

$$3^{\text{rd}} \text{ Range} = 2^1$$

$$4^{\text{th}} \text{ Range} = 2^2$$

,

,

,

,

$$k^{\text{th}} \text{ Range} = 2^{(k-2)}$$

$$2^{(k-2)} = p - 1$$

$$\overset{k}{\cancel{(k-2)}} = \log_2 \overset{p}{\cancel{(p-1)}}$$

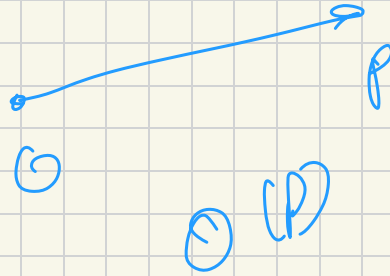
$$\checkmark k = \log_2(p)$$

Time Complexity

$$O(K) + O(\log_2 P)$$

$$O(\log_2 P) + O(\log_2 P)$$

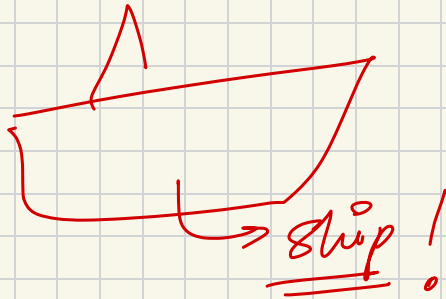
TC! $O(\log_2 P)$ ✓



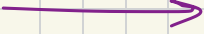
orders of ans.

Capacity of ship packages within B days.

$$\text{arr}[] A = \left\{ \overset{0}{2}, \overset{1}{2}, \overset{2}{2}, \overset{3}{4}, \overset{4}{1}, \overset{5}{4} \right\} \quad \text{days} = 3$$



$\min^m \left\{ \max^m \text{ capacity ship to deliver} \right\}$
these packages within
B days

$$\text{arr}[] A = \left\{ \overset{0}{2}, \overset{1}{2}, \overset{2}{2}, \overset{3}{4}, \overset{4}{1}, \overset{5}{4} \right\}$$




Min^m limit of balls in A Bag.

$$\text{arr}[] = \left\{ \overset{0}{2}, \overset{1}{4}, \overset{2}{8}, \overset{3}{2} \right\} \quad \text{opt} = 4$$

min { max^m balls in a bag }

$$\text{arr}[] = \{ 2, \overset{0}{\underset{\downarrow}{4}}, \overset{1}{8}, \overset{2}{2} \} \quad \text{opt} = 4$$

way 1

$(3, 1) (2, 2)$

opt 1

$$\{ 2, 3, 1, \overset{\downarrow}{8}, 2 \}$$

$(1, 7) (2, 6) (3, 5) (4, 4)$

opt 2

$$\{ 2, 3, 1, 3, \overset{\downarrow}{5}, 2 \}$$

$(1, 4) (2, 3)$

opt 3

$$\{ 2, 3, 1, 3, 2, \overset{\downarrow}{3}, 2 \}$$

$(1, 2)$

opt 4

$$\{ 2, 3, 1, 3, 2, 1, 2, 2 \}$$

$\max^m = 3$

$$\text{arr}[] = \left\{ \overset{0}{2}, \overset{1}{4}, \overset{2}{8}, \overset{3}{2} \right\}$$

Case 1 $\text{opt} = \infty$

$$\{1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1\}$$

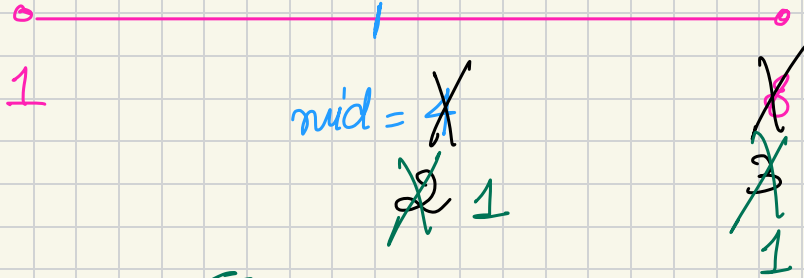
$$\text{max}^m = 1 \rightarrow \text{min}$$

Case 2 $\text{opt} = 0$

$$\{2, 4, 8, 2\}$$

$$\text{max}^m = 8 \rightarrow \text{min}$$

opt = 4



2 ans = ~~4~~ 2

$$\text{arr}[] = \{ \overset{0}{2}, \overset{1}{4}, \overset{2}{8}, \overset{3}{2} \}$$

\nearrow \nearrow \nearrow \nearrow

$(4, 4)$

$\text{maxCap} = 4$

opt = 1

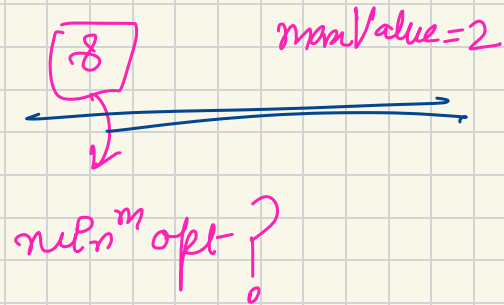
$$\text{arr}[] = \{ \overset{0}{2}, \overset{1}{4}, \overset{2}{8}, \overset{3}{2} \}$$

\nearrow \nearrow \nearrow \nearrow

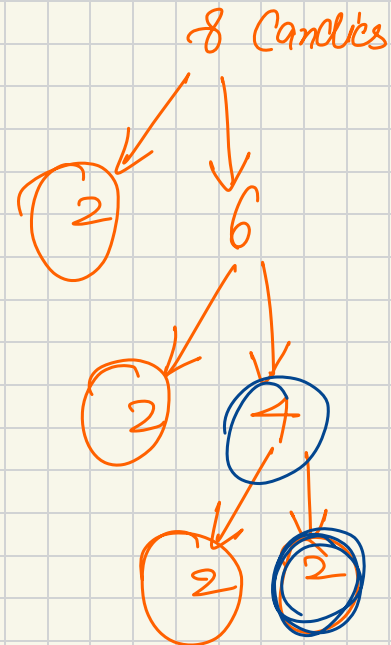
$(2, 2)$ $(2, 6)$ $(2, 4)$ $(2, 2)$

$\text{opt} = \cancel{4}$

$\text{maxCap} = 2$



$$\left\{ \frac{8}{2} - 1 \right\}$$

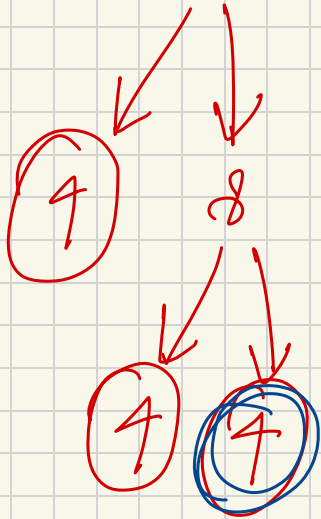


max^m 2 Candies

3 opt

12 Candies

maxValue = 4



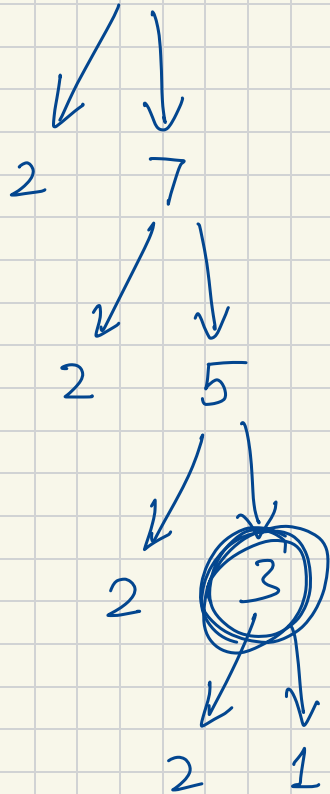
opt = 2

$$\left\{ \frac{12}{4} - 1 \right\}$$

$$\text{No. of opt} = \frac{\text{Balls}}{\text{maxValue}} - 1 \quad \checkmark$$

9 Candies

maxValue = 2



$$9/2 = 4 \checkmark$$

opt = Balls
None Value

