



Anagrams

str1 = "Baccio"

str2 = "Oface"

Is they are anagrams

* if all the characters of a string can be rearranged to form another string.

{'a', 'c', 'c', 'i', 'o'}



sort

{a, c, c, i, o}

{a, c, c, i, o}

Tc! O(N log₂ N)

Sc! O(N)

"accio" ← → "claoç"

Map

char	freq
✓ a	1 ✓
✓ c	✓ 2 ✓
✓ i	1 ✓
✓ o	1 ✓
↓	

Map 2

char	freq
wā	✓ 2
wē	1
wa	1
wō	1

TC : O(N) SC : O(26) ~ O(1)
=====

- ① if size are not same return false }
- ② verify all char has same freq }

~~0 0 0 1 0~~

a b c d e
0 1 2 3 4

1 2

1 1

z
25

$$\text{int pos} = (\text{ch} - 'a')$$

TC: O(N)

SC: O(26) ~ O(1)

Group Anagrams .

String [] arr = { "cat", "dog", "tac", "act", "god" }

{ { cat, tac, act }, { "dog", "god" } }



off

$\{$ "cat", "dog", "tac", "act", "god" $\}$
 ↑ ↑ ↑ ↑ ↑
 ↗

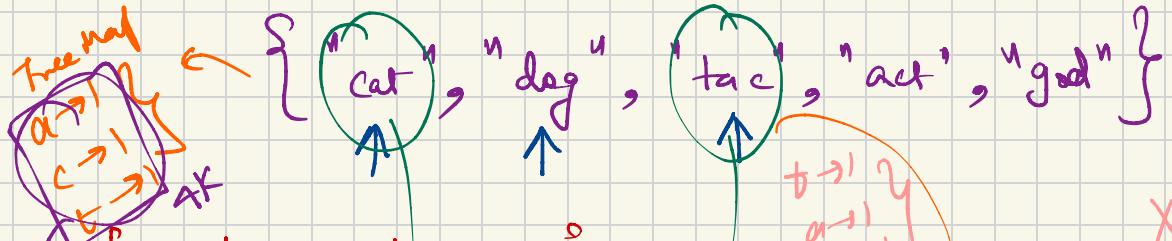
{ cat, tac, act }

{ dog, god }

Brute Force

$\{$ TC: $O(N^2) \times O(M)$
 $\approx O(N^2 M)$
 SC: $\underline{\underline{O(N)}}$

everyone will get a chance to start
 a group.
 ↗
 → avg. length of the { word }



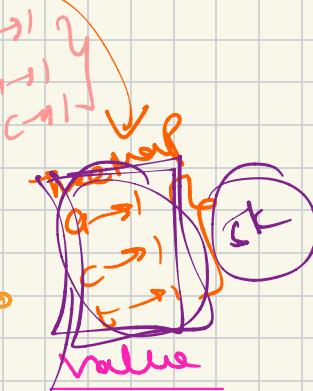
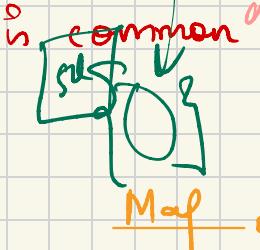
freq of each character

4K

key

freq

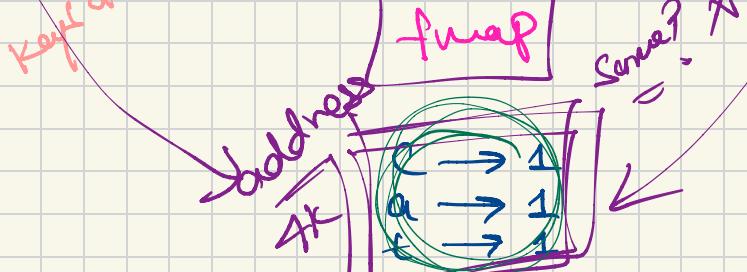
HashMap
Key are randomized.



X

`{ c → 1, a → 1, t → 1 }`

d → 1
o → 1
g → 1



d → 1
o → 1
g → 1

ArrayList<String>

{ cat, tac }

{ dog }

key

address?

HashMap <-, ->

ref variable
~~=~~

= new

ok?

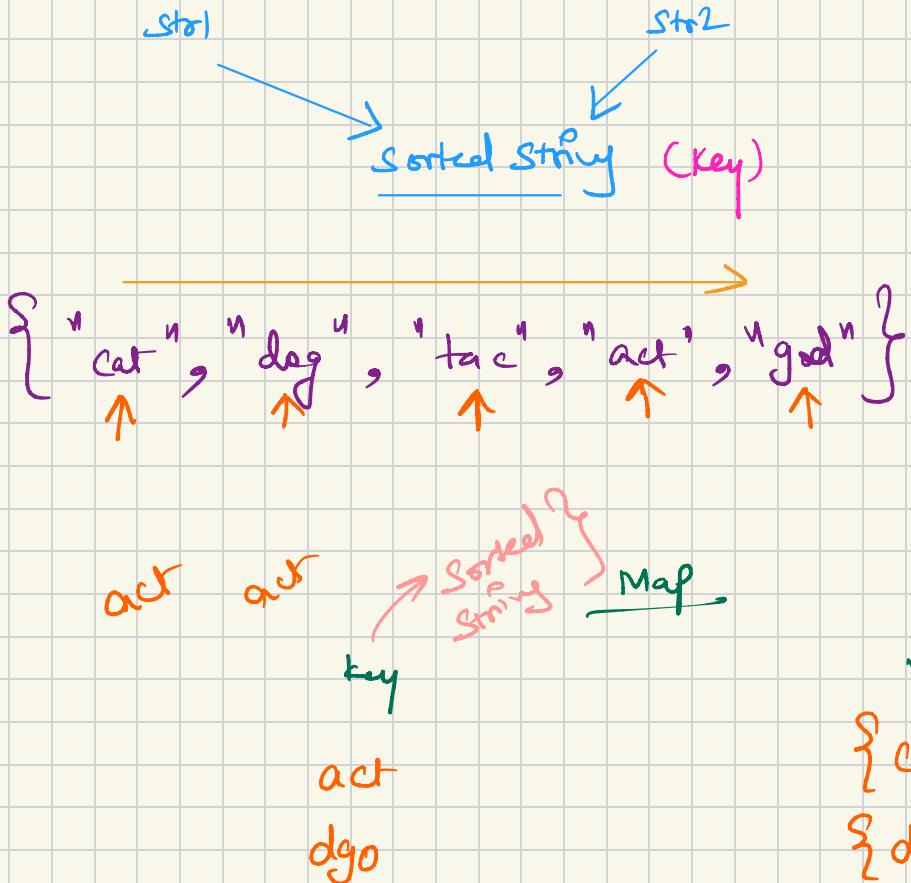
address?

allocate new memory

✓
OK

points

✓
OK



TC : $O(N) \times O(M \log_2 M)$

SC : $O(N)$

SC : $O(N \times M \log M)$

SC : $O(N)$

"Accio"

"Caior"

o encoding!

{
0 1 2
1 2

x
1

y
1

a1 c2 i1 o1

O(M)

$\{$ "cat", "dog", "tac", "act", "god" $\}$

$\alpha(|CT|)$

key

$\alpha(|CT|)$

$\alpha(g)$

Map

$T.C! \frac{\mathcal{O}(N \times M)}{\mathcal{O}(N)}$

value

$\{$ cat, tac, act $\}$

$\{$ dog, god $\}$

Minimum window Substring

str1: a d b a e b b a c a f d a

str2: c b a b d

Brute force

- ① generate all the substrings and try to find all char of str2 in window

Tc: $O(N^2 \times M)$ sc: $O(N)$

str1: ad b ac bb a ca f da
 ↓
 ↑
 inc

$$\begin{aligned}
 \text{len} &= \text{inc} - \text{exc} \\
 &= 6 - 5
 \end{aligned}$$

str2: cb ab d

freqMap

TC: $O(M) + O(N)$

a → 1 2 3

$\approx O(N+M)$

f → 1

SC: $O(1)$

b → 1

c → 1

d → 1

freqMap2

c	→	1
b	→	x 2
a	→	1
d	→	1

exc
↓

str1: a d b a c e b b a c a f d a



str2: c b a b d

dmcount = 5

freq Map²

c	→	1
b	→	2
a	→	1
d	→	1

mcnt = 1/2 / 3 / 4

freq Map¹

a	→	✓ / 3
f	→	1
b	→	1
c	→	1
d	→	1

A square root symbol contains the text "cmn = √ 5".

{ Tc: O(N+m) Sc: O(2m) + O(2s) ~ O(1) }

Agenda :

- ① LRU Cache
- ② Snapshot Array
- ③ Longest Subarray with Equal freq of 0, 1, 2

Longest Subarray with equal freq of 0, 1 and 2

int arr = { 0, 1, 2, 0, 0, 1, 2, 0, 1, 2, 2, 2, 1, 0, 1 }

Brute Force

- o Calc Subarray, stone count 0, 1, 2
- o stone the maxlen

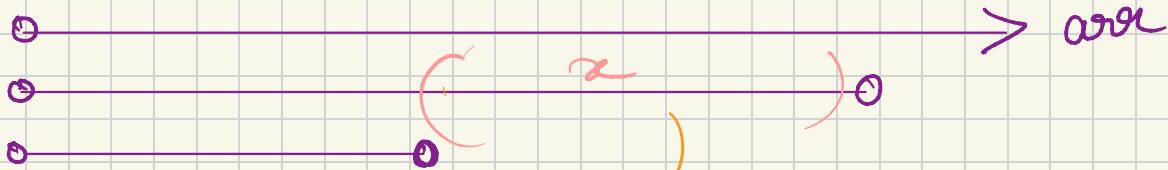
```
for (int i = 0 → n)
{   int cnt0, cnt1, cnt2
    for (int j = 0 → i)
    {
        if _____
```

TC: $O(N^2)$

SC: $O(1)$

$$\text{len} = \max(\text{cnt0})$$

$$^0_{\text{MTT}} \text{ aor} = \left\{ \begin{array}{c} 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14 \\ 0, 1, 2, 0, 0, 1, 2, 0, 1, 2, 2, 2, 2, 1, 0, 1 \end{array} \right\}$$



$$\begin{matrix} 0 & \rightarrow & \alpha \\ 1 & \rightarrow & \beta \\ 2 & \rightarrow & \gamma \end{matrix}$$

$$\begin{matrix} 0 & \rightarrow & \alpha' \\ 1 & \rightarrow & \beta' \\ 2 & \rightarrow & \gamma' \end{matrix} \quad \text{unknown}$$

$$\alpha' = \alpha + x \quad \text{①}$$

$$\beta' = \beta + x \quad \text{②}$$

$$\gamma' = \gamma + x \quad \text{③}$$

$$\textcircled{2} - \textcircled{1}$$

$$\beta' - \alpha' = \textcircled{\beta - \alpha} \longrightarrow \textcircled{4}$$

$$\textcircled{3} - \textcircled{2}$$

$$[\gamma' - \beta'] = \textcircled{\gamma - \beta} \longrightarrow \textcircled{3}$$

difference of current flow $\begin{cases} 1,0 \\ 210 \end{cases} \rightarrow \underline{\text{some}}$

$$[MT] \text{ arr} = \left\{ \begin{array}{cccccccccc|cccccccc} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 \\ 0, & 1, & 2, & 0, & 0, & 1, & 2, & 0, & 1, & 2, & 2, & 2, & 2, & 1, & 0, & 1 \end{array} \right\}$$

0	0	1	1	1	2	3	3	3	4	4	+
1	0	0	1	1	1	1	2	2	2	3	3
2	0	0	0	1	1	1	1	2	2	2	3
3	-1	0	-1	0	0	-1	-2	-1	-1	-2	-1
4	0	0	0	-1	0	0	0	-1	0	0	-1

$$\begin{pmatrix} a_1 - 0 \\ a_2 - 1 \end{pmatrix} \quad \boxed{0} \quad \boxed{-1} \quad \boxed{0} \quad \boxed{0} \quad -1 \quad -2 \quad -1 \quad -1 \quad -2 \quad -1 \quad \boxed{-1}$$

what is
a Hashmap

3
6
9

addresses same?
//

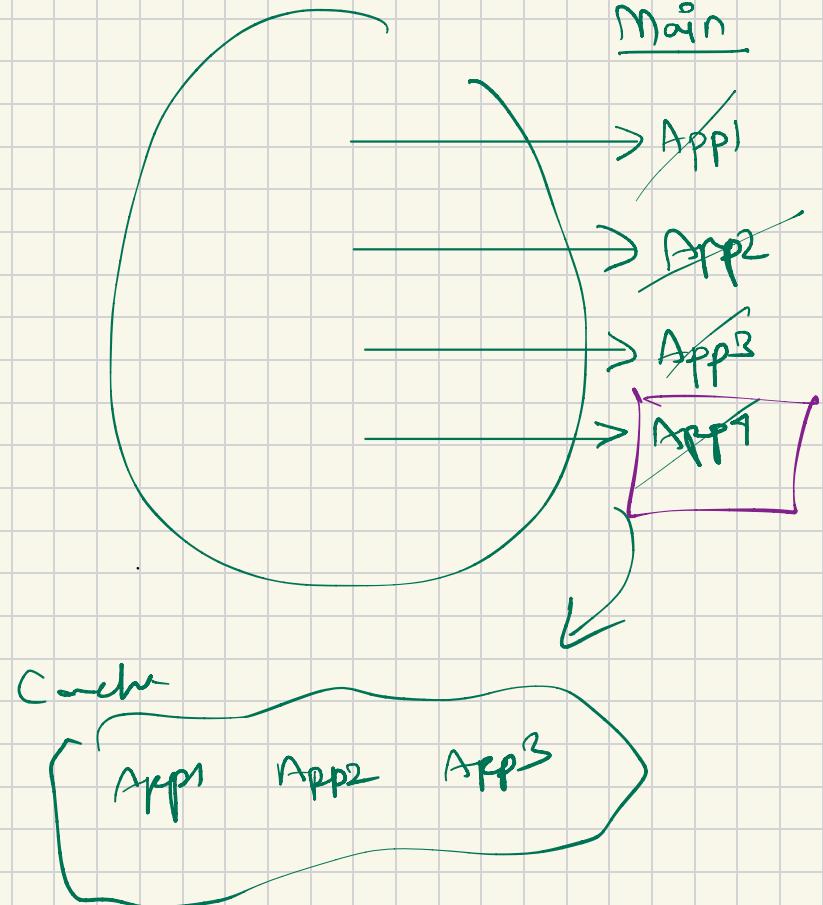
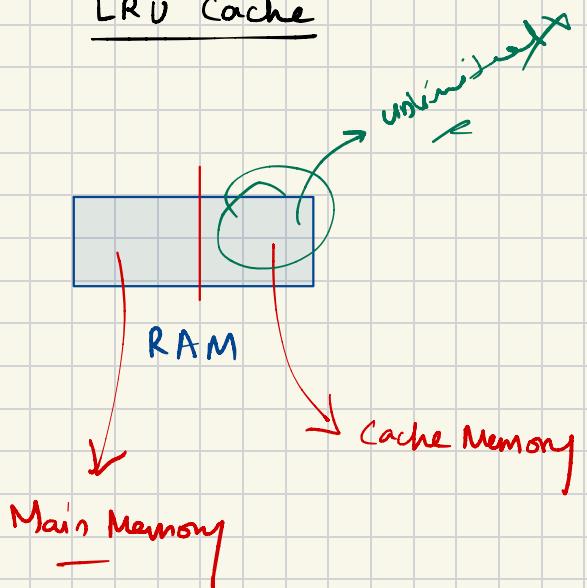
[class]

String = *key* \$ *keyL*

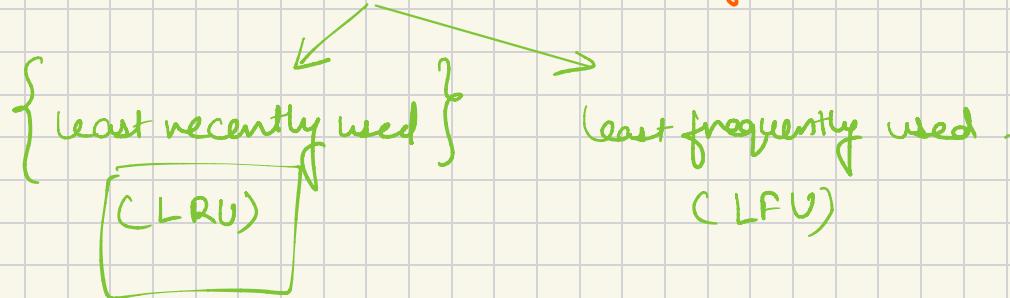
" 0 # -1 "

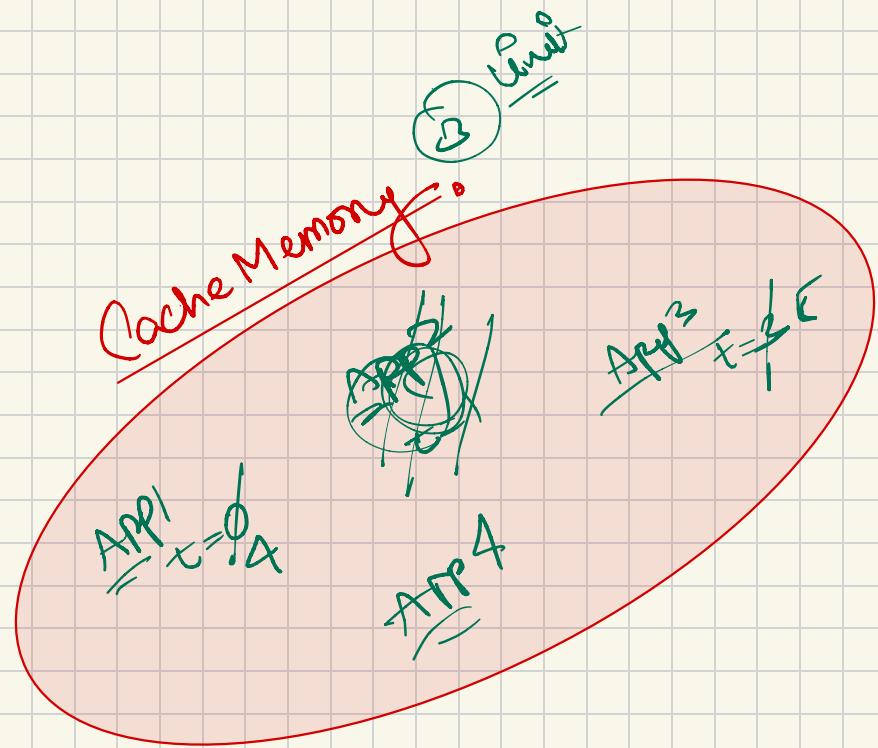
String → *key hashmap*

LRU Cache



Cache Memory Management Algo





App1 → opened
App2 → "
App3 →
App1 → pop up
App3 → updated
App4 → open

```
class LRUCache {  
    // your code here  
    public LRUCache(int capacity) {  
        // your code here  
    }  
  
    public int get(int key) {  
        // your code here  
    }  
  
    public void set(int key, int value) {  
        // your code here  
    }  
}
```

Capacity
of
Cache

return value
of the
app. i.e LRU

update/
and

$\text{set}(1, 10)$

$\text{set}(2, 30)$

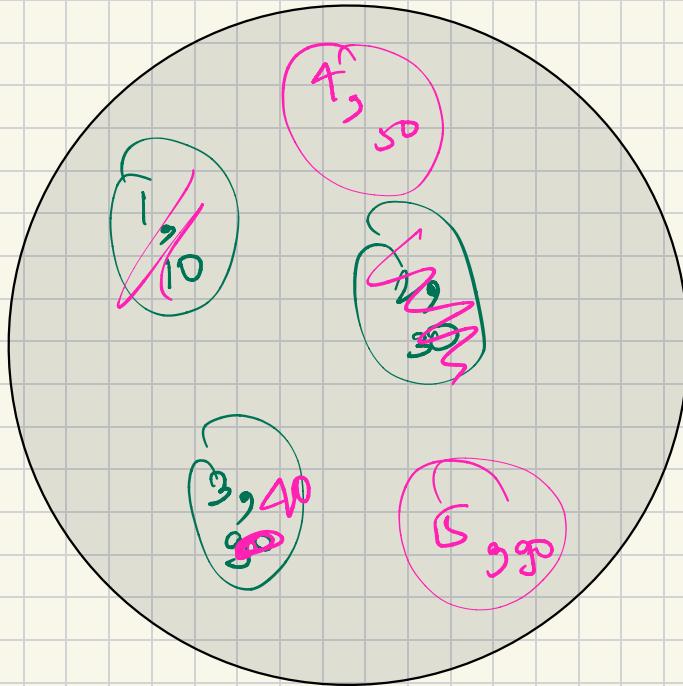
$\text{set}(3, 90)$

$\text{get}(1) \rightarrow 10$

$\text{set}(4, 50)$

$\text{set}(3, 40)$

$\text{set}(5, 90)$

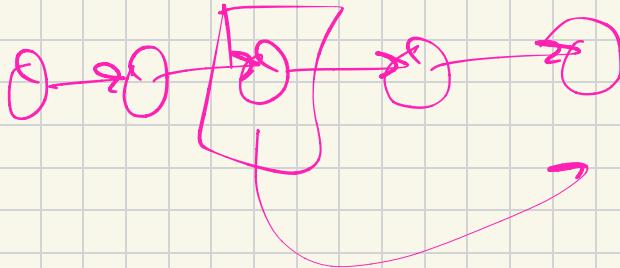


MRV

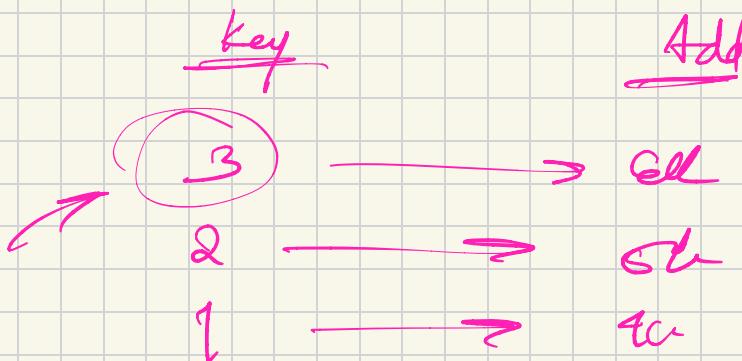
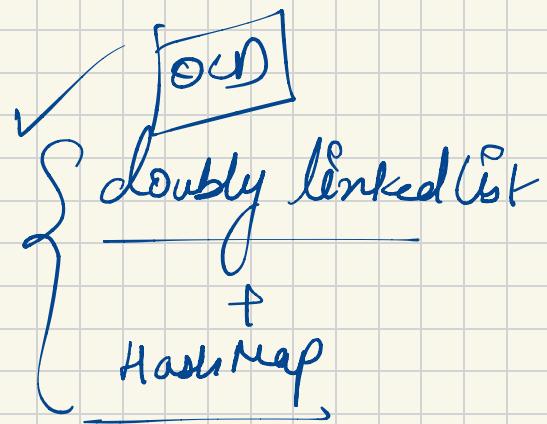
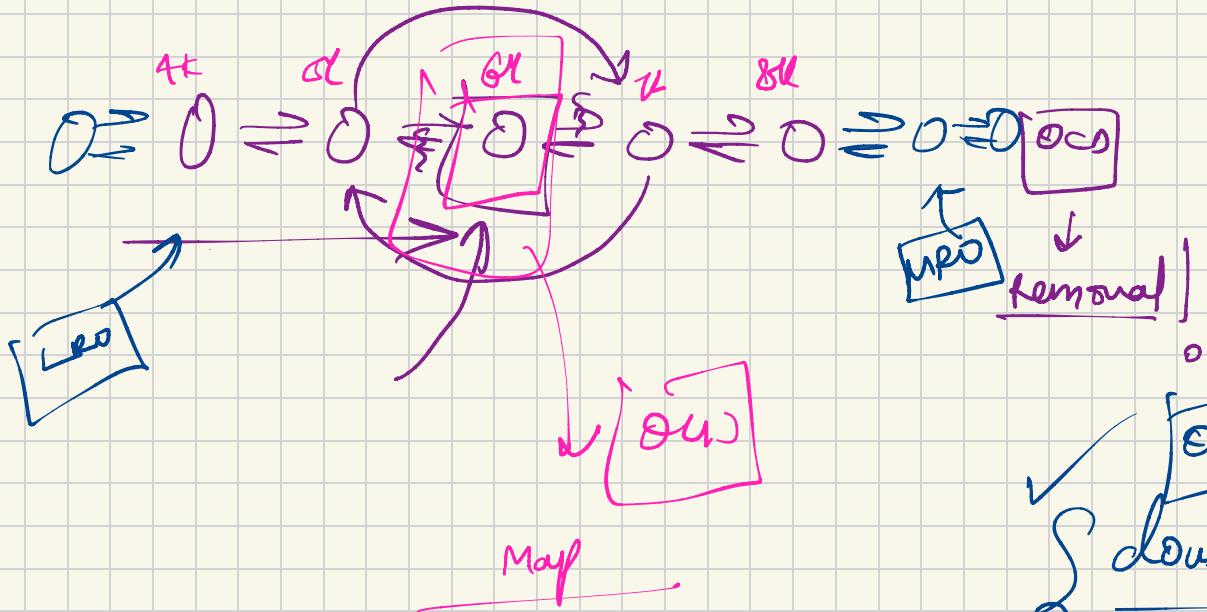
LRU

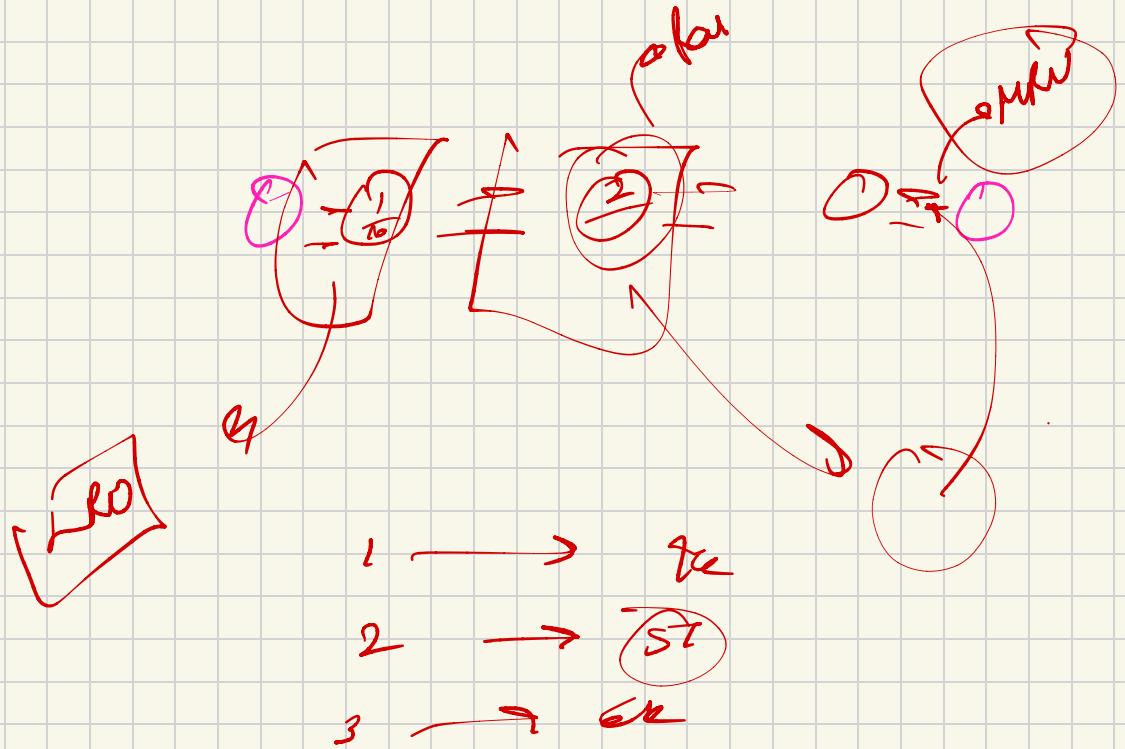
(3, 90) (2, 70) (1, 10)

[~~0 CND~~] P



δ(w) X





Snapshot Array

int[] arr = { 0, 1, 2, 3, 4, 5, 6 }
0, 0, 8, 0, 3, 0, 0 }

set(2, 10)

Snapshot = 0

0, 1, 2, 3, 4, 5, 6
0, 0, 10, 0, 3, 0, 0 }

set(4, 3)

snapshot()

set(2, 8)

Snapshot = 1

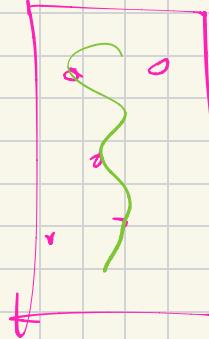
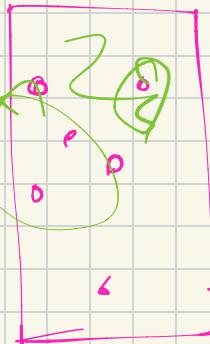
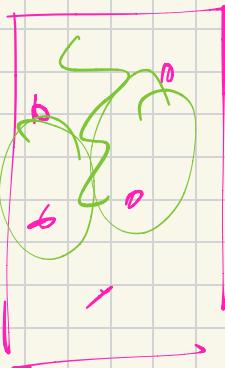
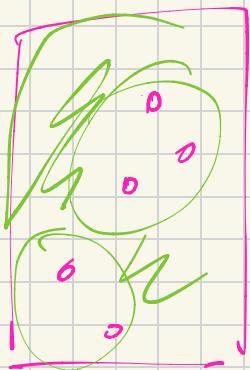
0, 1, 2, 3, 4, 5, 6
0, 0, 8, [0], 3, 0, 0 }

snapshot()

get(2, 0) → [10]

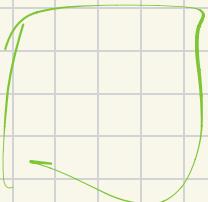


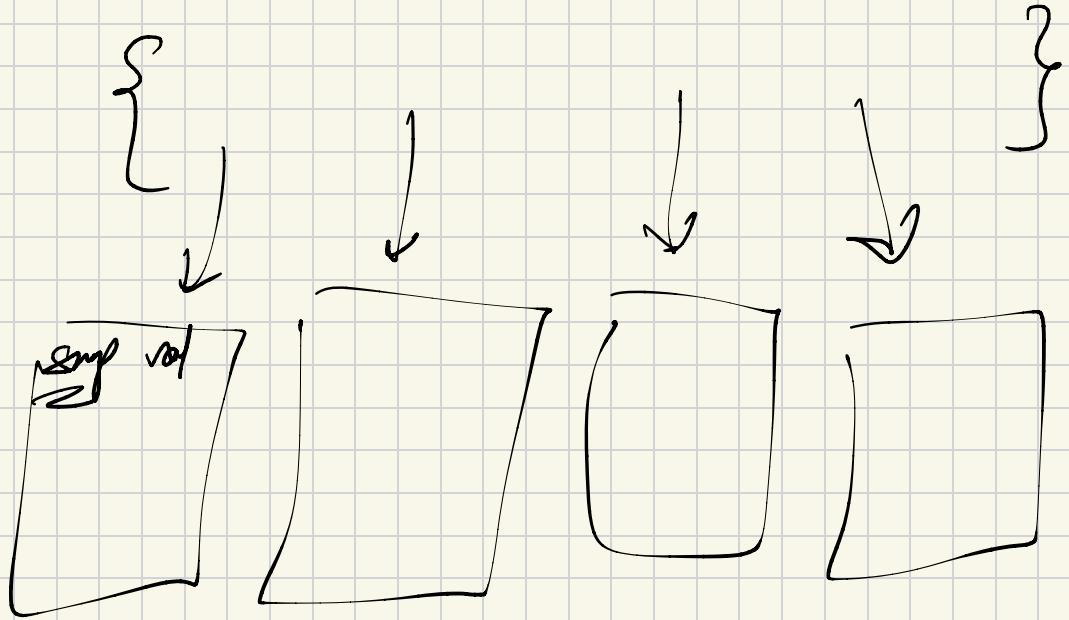
get(3, 0) → [0]



4K

⇒ Separable
~~+~~





$[0, \cancel{0}, \cancel{0}, 0, 0]$

$\text{set}(1, 3)$

$\text{set}(2, 5)$

$\text{snip}(1 \rightarrow \cancel{0})$

$\text{set}(1, 9)$

$\text{snip}(0 \rightarrow \cancel{1})$

$\text{get}(1, 0)$

$\text{get}(2, 10)$

$\text{get}(6, 5)$

$\text{snip}(1 \rightarrow 0/2)$

