

Understanding the problem

18 February 2025

02:08

— Monotonic Stack: —

Q. What is monotonic stack.

Before this we stack to store the elements

↳ But not store the elements in a specific order

◦ it may be rise

◦ \rightarrow — base.

◦ \rightarrow — (any other order)

One of Stack. 496. Next Greater Element I

Attempted

Easy

Topics

Companies

The next greater element of some element x in an array is the first greater element that is to the right of x in the same array.

You are given two **distinct 0-indexed** integer arrays `nums1` and `nums2`, where `nums1` is a subset of `nums2`.

For each $0 \leq i < \text{nums1.length}$, find the index j such that `nums1[i] == nums2[j]` and determine the **next greater element** of `nums2[j]` in `nums2`. If there is no next greater element, then the answer for this query is `-1`.

Return an array `ans` of length `nums1.length` such that `ans[i]` is the **next greater element** as described above.

Example

Input: nums1 = [4,1,2], nums2 =
[1,3,4,2]

Output: [-1,3,-1]

→ this actual question

But sir tougher Q:-

Example - [6, 0, 8, 2, 3] ↗ O/P
O/P [8, 8, -1, 3, -1]

Method 1 solve this problem

New
ans: →

```

nge = [n]
for i = 0 → n-1
{
    for j = i+1 → n-1
    {
        if (arr[i] > arr[j])
            nge[i] = arr[j], break
    }
}
return nge

```

TC - $O(N^2)$

SC - $O(N)$ for store answer.

↳ Not happy to reduce time complexity

$O(N^2) \rightarrow O(N)$

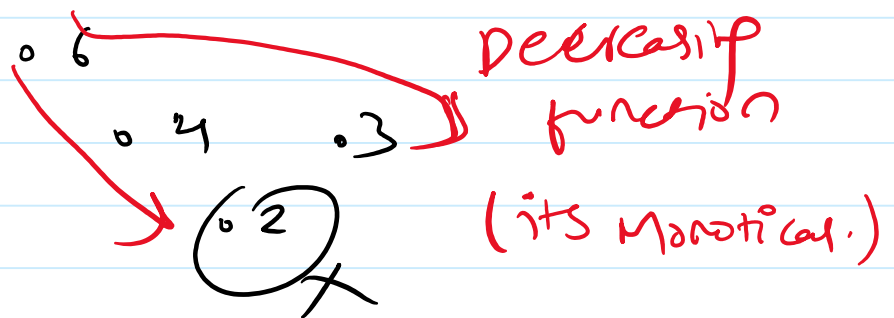
Example to use the monotonic stack

(4, 12, 5, 3, 1, 2, 5, 3, 1, 2, 4, 6)

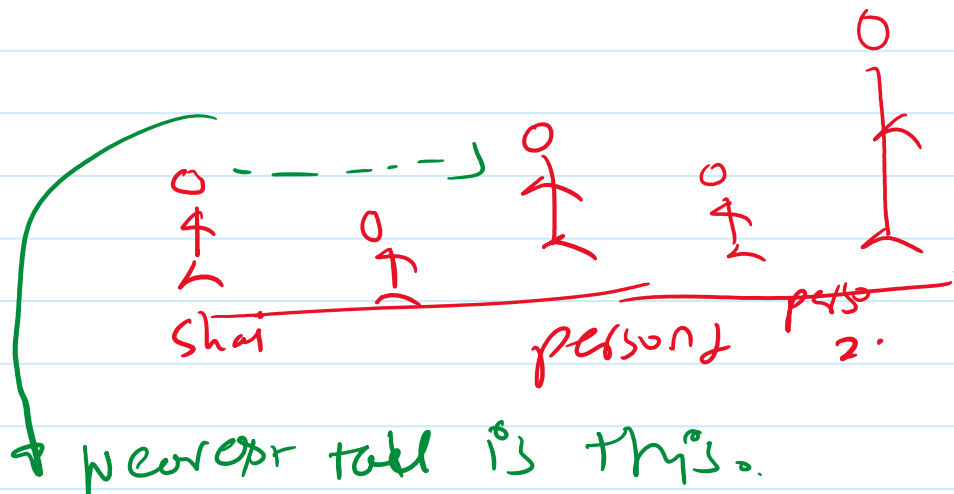
$$O(N^2) \rightarrow O(N)$$

4	12	5	3	1	2	5	3	1	2	4	6
↑	↑	↑	↑	↑	↑	↑	↑	↑	↑	↑	↑
ngc[] → 12	-1	6	5	2	5	6	4	2	4	6	-1

instead of storing the complete value on stack
use monotonic function.



Consider similar to —



Code:-

How to intuition come



```
list<int> findNGE(arr[])
{
    nge[n], stack st
    for (i = n-1 → 0) → O(N)
        while (!st.empty() && st.top() <= arr[i])
            st.pop(); → O(N)
        if (st.empty()) nge[i] = -1
        else nge[i] = st.top()
        st.push(arr[i])
    }
    return nge;
}
```

TC → $O(2N)$ → worst case
 SC → $O(N)$ + $O(N)$ → Ans
 → Stack

Solution :

Method -1.

```

class Solution {
    public int[] nextGreaterElement(int[] nums1, int[] nums2) {
        int[] ans = new int[nums1.length];
        Arrays.fill(ans, -1);
        for (int i = 0; i < nums1.length; i++) {
            for (int j = 0; j < nums2.length; j++) {
                if (nums1[i] == nums2[j]) {
                    for (int k = j + 1; k < nums2.length; k++) {
                        if (nums2[j] < nums2[k]) {
                            ans[i] = nums2[k];
                            break;
                        }
                    }
                }
            }
        }
        return ans;
    }
}

```

TC $\rightarrow O(N^3)$

Method 2.

```

19 public static int[] nGE(int[] arr){
20     int[] ans = new int[arr.length];
21     Stack<Integer> stack = new Stack<>();
22     for (int i = arr.length - 1; i >= 0; i--) {
23         while (stack.isEmpty() == false && arr[i] >= stack.peek())
24             stack.pop();
25
26         if (stack.isEmpty())
27             ans[i] = -1;
28         else {
29             ans[i] = stack.peek();
30         }
31         stack.push(arr[i]);
32     }
33
34     return ans;
35 }

```

```

2 public int[] nextGreaterElement(int[] nums1, int[] nums2) {
3     int[] ans = new int[nums1.length];
4
5     int[] arr = nGE(nums2);
6     for (int i = 0; i < nums1.length; i++) {
7         for (int j = 0; j < nums2.length; j++) {
8             if (nums1[i] == nums2[j]) {
9                 {
10                     ans[i] = arr[j];
11                     break;
12                 }
13             }
14         }
15     }
16     return ans;
17 }

```