#### **Braille Auto-Correct**

This document details a streamlined approach to auto-correct Braille input using only the Levenshtein distance for matching, along with optimizations and trade-offs.

#### 1. Objective

Provide a simple yet effective real-time suggestion system for QWERTY-based Braille input that:

- Corrects substitutions, insertions, and deletions.
- Requires minimal setup and dependencies.
- Delivers low latency for small-to-medium dictionaries.

#### 2. Approach

- 1. Input Normalization
  - o Map keys s, d, f, j, k,  $l \rightarrow dots 1-6$ .
  - o Sort and join into a string pattern ("sdf"  $\rightarrow$  "123").
- 2. Dictionary Preprocessing
  - o Precompute pattern for each word in the dictionary once at startup.
- 3. Brute-Force Matching
  - o For each query, iterate all dictionary patterns:
    - Compute Levenshtein distance with early exit.
    - Collect words whose distance ≤ max\_dist.
- 4. Ranking & Suggestion
  - o Sort by (distance, word) to resolve ties.
  - o Return the top-K candidates (e.g., 3).

## 3. Core Algorithm Components

- 3.1. QWERTY-to-Pattern Conversion
  - Mapping Dict: {'s':1,'d':2,'f':3,'j':4,'k':5,'l':6}
  - Function: Filter valid keys, map, sort, join.
- 3.2. Levenshtein Distance
  - DP Table: Compute in  $O(n \cdot m)$  time, where n, m are string lengths.
  - Space: O(min(n,m)) by storing only the previous row.
  - Early Exit: Abort when the current row's minimal edit cost exceeds max\_dist.

## 4. Optimizations

- Early-Abandon: Greatly reduces work on dissimilar patterns by cutting DP loops early.
- One-Time Preprocessing: Patterns computed just once, reducing per-query overhead.
- Compact Patterns: Using short digit strings speeds up string comparisons and distance calls.
- Threshold Tuning: Adjusting max\_dist to balance recall vs. computation (e.g., 1–2 for typical Braille patterns).

# 5. Trade-Offs Analysis

Choice	Benefit	Drawback
Brute-Force Search	Simplicity; no additional data structures	$\mathrm{O}(N^{\centerdot}n^{\centerdot}m)$ per query may be slow for large $N$
Levenshtein	Handles all typo types uniformly	Quadratic with pattern length
Pure Python	Portable; no external deps	Single-threaded performance ceiling
Early-Termination	Cuts off costly comparisons early	Requires careful threshold setting

### 6. Suitability

- When to Use: Prototyping, educational demos, dictionaries up to ~50K entries.
- Scalability: For larger datasets or microsecond latency, layer in an index (e.g., BK-tree) or compile hotspots in C/C++.