

# Domain-Specific Languages of Mathematics

## Course codes: DAT326 / DIT983

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<b>Results</b>	Announced within 15 workdays
<b>Exam check</b>	Tue 2025-04-08, 12.15-12.45 in EDIT 6128
<b>Aids</b>	One textbook of your choice (Domain-Specific Languages of Mathematics, or Beta - Mathematics Handbook, or Rudin, or Adams and Essex, or ...). No printouts, no lecture notes, no notebooks, etc.
<b>Grades</b>	To pass you need <b>a minimum of 5p on each question (1 to 4)</b> and also reach these grade limits: 3: $\geq 48$ p, 4: $\geq 65$ p, 5: $\geq 83$ p, max: 100p

Remember to write legibly. Good luck!

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For reference: the learning outcomes. Some are tested by the hand-ins, some by the written exam.

- Knowledge and understanding
  - design and implement a DSL (Domain-Specific Language) for a new domain
  - organize areas of mathematics in DSL terms
  - explain main concepts of elementary real and complex analysis, algebra, and linear algebra
- Skills and abilities
  - develop adequate notation for mathematical concepts
  - perform calculational proofs
  - use power series for solving differential equations
  - use Laplace transforms for solving differential equations
- Judgement and approach
  - discuss and compare different software implementations of mathematical concepts

1. [25p] **Algebraic structure: a DSL for vector spaces**

A vector space over  $\mathbb{R}$  is a set  $V$  together with a constant (or nullary) operation  $0 : V$  (called “zero”), an operation  $(+) : V \rightarrow V \rightarrow V$  (called “add”), an operation  $(-) : V \rightarrow V$  (called “negate”) and an operation  $(\cdot) : \mathbb{R} \rightarrow V \rightarrow V$  (called “scale”), such that

$\forall v \in V.$	$v + 0 = 0 + v = v$	-- additive zero
$\forall v_1, v_2, v_3 \in V.$	$(v_1 + v_2) + v_3 = v_1 + (v_2 + v_3)$	-- associative (+)
$\forall v \in V.$	$v + (-v) = (-v) + v = 0$	-- additive inverse
$\forall v_1, v_2 \in V.$	$v_1 + v_2 = v_2 + v_1$	-- commutative (+)
$\forall x_1, x_2 \in \mathbb{R}, v \in V.$	$x_1 \cdot (x_2 \cdot v) = (x_1 * x_2) \cdot v$	-- repeated scaling
$\forall v \in V.$	$1 \cdot v = v$	-- unit scaling
$\forall x \in \mathbb{R}, v_1, v_2 \in V.$	$x \cdot (v_1 + v_2) = x \cdot v_1 + x \cdot v_2$	-- scaling left-distributive
$\forall x_1, x_2 \in \mathbb{R}, v \in V.$	$(x_1 + x_2) \cdot v = x_1 \cdot v + x_2 \cdot v$	-- scaling right-distributive

*Remark:*  $(*)$  denotes the standard multiplication in  $\mathbb{R}$ .

- Define a type class *Vector* that corresponds to the structure “vector space over  $\mathbb{R}$ ”.
- Define a datatype *VecSyn a* for the language of vector space expressions (with variables of type *a*) and define a *Vector* instance for it. (These are expressions formed from applying the vector operations to the appropriate number of arguments, e.g., all the left hand sides and right hand sides of the above equations.)
- Find and implement two other instances of the *Vector* class. Make sure the laws are satisfied.
- Give a type signature for, and define, a general evaluator (on the basis of an assignment function) from *VecSyn a* expressions to any semantic type with a *Vector* instance.
- Specialise the evaluator to the two *Vector* instances defined in (1c). Take three vector expressions (of type *VecSyn String*), give the appropriate assignments and compute the results of evaluating, in each case, the three expressions.

Each question carries 5pts.

2. [25p] **Laplace**

Consider the following differential equation:

$$f'' + 20f' + 99f = -99, \quad f(0) = 0, \quad f'(0) = -1$$

- [10p] Solve the equation assuming that  $f$  can be expressed by a power series  $fs$ , that is, use *integ* and the differential equation to express the relation between  $fs$ ,  $fs'$ ,  $fs''$ . What are the first three coefficients of  $fs$ ? Explain how you compute them.
- [15p] Solve the equation using the Laplace transform. You should need this formula (note that  $\alpha$  can be a complex number) and the rules for linearity + derivative:

$$\mathcal{L}(\lambda t. e^{\alpha * t}) s = 1/(s - \alpha)$$

Show that your solution does indeed satisfy the three requirements.

3. [25p] **Typing:** Lattice subgroups

Consider the following (rephrased) quote from Wikipedia:

In geometry and group theory, a lattice in the real coordinate space  $\mathbb{R}^n$  is a set of points  $L$  in this space with the properties that coordinate-wise addition or subtraction of two points in the lattice produces another lattice point, that there is a minimum distance  $d_1$  such that any two lattice points  $a, b$  have at least distance  $d_1$ , and that every point  $x$  in  $\mathbb{R}^n$  is within some maximum distance  $d_2$  of a lattice point.

- (a) [5p] Give types for  $L, d_1, a, b, d_2$ . Explain your reasoning and be careful to interpret “distance” and “two lattice points” in a way that makes the definition non-trivial.
- (b) [5p] Define (in first-order logic) the predicate  $\text{isLattice}(U, n)$  which is true if and only if the set  $U$  is a lattice in  $\mathbb{R}^n$ . You may assume a function  $d : \mathbb{R}^n \rightarrow \mathbb{R}^n \rightarrow \mathbb{R}$  computing the distance of two points in  $\mathbb{R}^n$ .
- (c) [5p] Prove  $\neg \text{isLattice}(\{(x, 0) \in \mathbb{R}^2 \mid x \in \mathbb{Z}\}, 2)$ .
- (d) [5p] Let  $a > 0$  be a positive real number. Prove  $\text{isLattice}(\{(x, y) \in \mathbb{R}^2 \mid ax, ay \in \mathbb{Z}\}, 2)$ .
- (e) [5p] Find an example of a lattice in  $\mathbb{R}^2$  which is not of the form given above. Explain your reasoning.

4. [25p] **Computational proof:** Syntactic derivatives

Consider the following Haskell code for a DSL of 1-argument functions. The evaluator  $eval$  and the function  $der_2$  are structural homomorphisms from the syntax type  $F$ . Your task is to implement some of the missing parts of the “deep copy and derivative” function  $der_2$ , and prove some properties about it.

```
data F = Zero | One | X | Sub F F | Mul F F | Div F F deriving (Eq, Show)
eval :: Field a => F -> (a -> a)
eval Zero  _ = zero
eval One   _ = one
eval X     _ = x
eval (Sub fe ge) x = eval fe x - eval ge x
eval (Mul fe ge) x = eval fe x * eval ge x
eval (Div fe ge) x = eval fe x / eval ge x
der2 :: F -> (F, F)
der2 Zero    = der2Zero
der2 One     = der2One
der2 X       = der2X
der2 (Sub fe ge) = der2Sub (der2 fe) (der2 ge)
der2 (Mul fe ge) = der2Mul (der2 fe) (der2 ge)
der2 (Div fe ge) = der2Div (der2 fe) (der2 ge)
```

Let the property  $P(fe)$  be “Let  $fe_2$  be the first, and  $fe'$  the second, component of the pair returned by  $der_2 fe$ . Then  $fe_2 == fe$  and  $eval fe' == D (eval fe)$ , where  $D$  is the differentiation operator.” The specification of  $der_2$  is then  $\forall fe. P(fe)$ .

- (a) [5p] Give types for, and implement,  $der2One$ ,  $der2X$ , and  $der2Sub$ .
- (b) [5p] In an induction proof of correctness of  $der_2$ , one of the base cases is  $P(X)$  (where  $X$  is a constructor in the datatype  $F$ ). Prove this case using equational reasoning, carefully motivating each step as in “by def.  $eval One$ ”, “by def.  $der_2 X$ ”, etc.
- (c) [5p] Give the type for, and implement,  $der2Div$ .
- (d) [10p] One of the inductive step cases is  $\forall ge. \forall he. P(ge) \wedge P(he) \Rightarrow P(Div ge he)$ . Prove this step using equational reasoning.