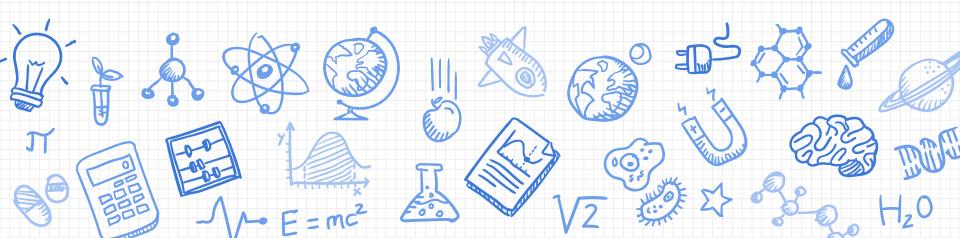
Advanced Programming Crash Course

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What will we learn?

- Arrays
- 2D Arrays (Matrices)
- Array Sorting
- Collection Classes
- Recursion



Arrays

 Location in memory that stores multiple elements under the same TYPE and NAME

```
// Ways to declare an array
String stringArray[] = new String[2];
String stringArray2[] = new String[] { "Element 1", "Element 2" };
String stringArray3[];
stringArray3 = new String[2];

System.out.println("Method 1 of outputting using a loop");
for (int i = 0; i < stringArray2.length; i++) {
    System.out.println(stringArray2[i]);
}

System.out.println("Method 2 of outputting using a loop");
for (String string : stringArray2) {
    System.out.println(string);
}</pre>
```



2D Arrays (Matrices)

- Some data is stored in multiple dimensions
 - Example: Tables, Maps, Images

Name	Age
Shalee	17
Campos	25



- A lot harder to manipulate the data if its in a
 1D array
- We need to call each row by 2 indices

In this example:

Row 1 is index 0 and index 1



- Easier if its 2D since we can call each Row by 1 index

	Shalee, 17],
	(amp 05, 25]

Sha	lee	

Campos

In this example:

Row 1 is index 0



17

Index order for an individual element is [Row][Col]

Name	Age
Shalee	17
Campos	25

Name	Age
Row 1, Col 1 (0, 0)	Row 1, Col 2 (0, 1)
Row 2, Col 1 (1, 0)	Row 2, Col 2 (1, 1)



```
int intArray[][] = new int[2][2];
int intArray2[][] = new int[][] { { 1, 2, 3 }, { 4, 5, 6 } };
int intArray3[][];
intArray3 = new int[2][2];
System.out.println("Method 1 of outputting using a loop");
    for (int j = 0; j < intArray2[i].length; j++) {</pre>
        System.out.print(intArray2[i][j] + " ");
    System.out.println();
System.out.println("Method 2 of outputting using a loop");
for (int[] is : intArray2) {
        System.out.print(is2 + " ");
    System.out.println();
```



Array Sorting

- Array Sorting is crucial to organizing data
 - Ascending order, Descending order, Alphabetical, etc
- Algorithms we will discuss
 - Bubble Sort
 - Selection Sort
 - Insertion Sort



- Bubble Sort
 - Compare the first two elements
 - Swap if out of order
 - Continue this until you reach the end of the array

$$\frac{N(n!)}{2} \rightarrow O(n^2)$$



- Example: Sorting in Ascending order
 - Array: [5, 3, 8, 4, 6]
 - [5, 3, 8, 4, 6] -> Compare first and second (Swap)
 - [3, 5, 8, 4, 6] -> Compare second and third (Don't swap)
 - [3, 5, 8, 4, 6] -> Compare third and fourth (Swap)
 - [3, 5, 4, 8, 6] -> Compare fourth and fifth (Swap)
 - [3, 5, 4, 6, 8] -> Repeat this process until no more swaps are required

- Selection Sort
 - Search for the smallest element and swap it into the first index
 - Search for the second smallest element and swap it into the second index
 - Continue this process until the array is fully sorted

$$\frac{\mathcal{N}(n\cdot 1)}{2} \rightarrow \mathcal{O}(n^2)$$



- Insertion Sort
 - Takes advantage of pre-sorting (only 1 pass)
 - Requires fewer comparisons than the other 2 Algorithms
 - Merge based (merging unsorted data to the sorted side)
- View pseudocode here: Click me



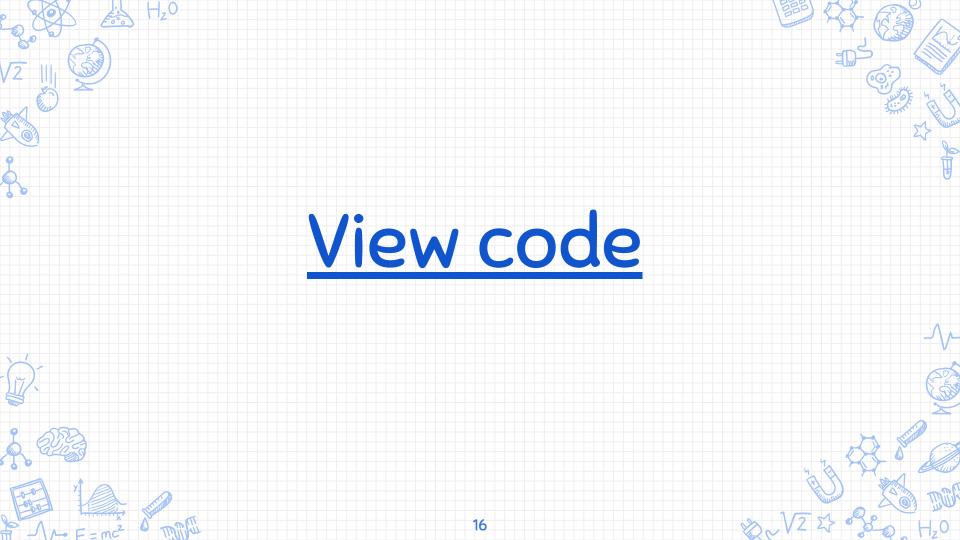
		Time Complexity					
		Best	Worst	Avg.	Space	Stable	Comments
Comparison Sort	Bubble Sort	O(n^2)	O(n^2)	O(n^2)	O(1)	Yes	For each pair of indices, swap the elements if they are out of order
	Modified Bubble Sort	O(n)	O(n^2)	O(n^2)	O(1)	Yes	At each Pass check if the Array is already sorted. Best Case-Array Already sorted
	Selection Sort	O(n^2)	O(n^2)	O(n^2)	O(1)	Yes	Swap happens only when once in a Single pass
	Insertion Sort	O(n)	O(n^2)	O(n^2)	O(1)	Yes	Very small constant factor even if the complexity is O(n^2). Best Case: Array already sorted Worst Case: sorted in reverse order
	Quick Sort	O(n.lg(n))	O(n^2)	O(n.lg(n))	O(1)	Yes	Best Case: when pivot divide in 2 equal halves Worst Case: Array already sorted - 1/n-1 partition
0	Randomized Quick Sort	O(n.lg(n))	O(n.lg(n))	O(n.lg(n))	0(1)	Yes	Pivot chosen randomly
	Merge Sort	O(n.lg(n))	O(n.lg(n))	O(n.lg(n))	O(n)	Yes	Best to sort linked-list (constant extra space). Best for very large number of elements which cannot fit in memory (External sorting)
	Heap Sort	O(n.lg(n))	O(n.lg(n))	O(n.lg(n))	0(1)	No	
Non- Comparison Sort	Counting Sort	O(n+k)	O(n+k)	O(n+k)	O(n+2^k)	Yes	k = Range of Numbers in the list
	Radix Sort	O(n.k/s)	O(2^s.n.k/s)	O(n.k/s)	O(n)	No	
	Bucket Sort	O(n.k)	O(n^2.k)	O(n.k)	O(n.k)	Yes	



Collection Classes

- Lists
 - An ordered collection of elements
 - Elements can be accessed through indices
 - May Contain duplicate values
- Sets
 - An unordered collection of elements
 - No indices
 - All values must be unique





Recursion

- An Algorithm where a method calls itself
 - Form of looping where a method calls itself repeatedly to solve simpler versions of a problem
- Properties
 - Must contain a base case
 - Has to be reached eventually to prevent a StackOverFlow error
 - Contains decision structures rather than conventional loops



Recursion - Continued

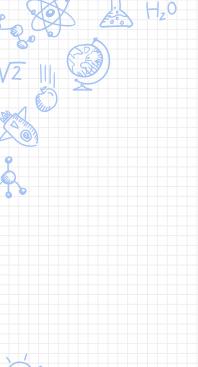
- Process
 - Involves the usage of a Stack
 - Stack = Data Abstraction where:
 - New data is "pushed" to the top of the stack
 - Old data is "popped" or removed from the top of the stack
 - The current computation is suspended and placed onto a Stack
 - Once the base case is reached the method unwinds and the Stack is popped



Recursion - Continued Part 2

```
private static int power(int base, int n) {
   if (n <= 1) {
        return base;
   return base * power(base, n - 1);
Run | Debug
public static void main(String[] args) {
   int value = power(2, 3); // Calling the power method
   System.out.println(value); // Outputting the answer
```





View PPT Code



