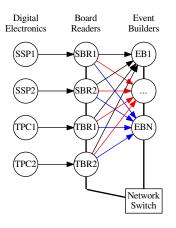
Online/Offline Buffer Options for protoDUNE/SP

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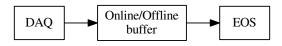
Model of Upstream DAQ



- DAQ as layered, directed acyclic graph.
- Figure glosses over distinctions in RCE + electronics (sorry).
- Board Readers route fragments from one trigger to Event Builders
- Event Builders concatenate fragments into contiguous readouts.
- Rates, bandwidths, processing time, switches, NICs all play major role in shape and size!

Data Scenario Implications on Buffer

Required to have 3 days buffer.



Recent upward revision of Data Scenarios Spreadsheet:

- DocDB 1086-v6 (let's put any and all updates here)
- \rightarrow 25-50Hz, 5ms, 6APA, 2-4× compression, 25-50M events.

Implications on buffer requirements

- 25-50TB buffer disk,
- 30-60 parallel HDD writes,
- 1.5-3.0 GByte/sec throughput.

Back of the envelope estimate

If assume 1Gbps NICs:

- 25 concurrent network streams
- NIC is bottleneck so minimum of 25 hosts writing 2-3 disks
- Beneficial side-effect: leaves plenty of idle CPUs to load up doing other, useful and prompt tasks.

If assume 10Gbps NICs:

- 3 concurrent network streams
- SATA is bottleneck so minimum of 6 hosts writing 10 disks
- Hosts: fewer, bigger, more expensive, less "off-the-shelf".
- CPUs loaded just doing I/O, little room for other processing.

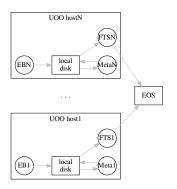
A **quantitative investigation** based on assumptions of rates, bandwidths, processing times, etc is being developed with the **Ersatz Simulation** package. Links: DAQ Sim presentation, GitHub

Buffer Design - Two Options

- **UOOB** Unified Online/Offline Buffer hosts
 - → Shared hosts, local disk for data hand-off.
- DOOB Dedicated Online/Offline Buffer hosts
 - → Separate layers, network for data hand-off.

Unified Online/Offline Buffer

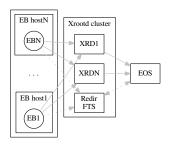
Online/offline interface: file-hierarchy on shared local disks.



Each UOOB computer must host:

- DAQ Event Builder node.
- File management glue scripts.
- File metadata producer.
- FTS instance.
- Logic to handle DAQ/FTS disk contention.
- Online/offline contract on detailed file locations.

Dedicated Online/Offline Buffer



Dedicated "layer" in distributed graph

- Each EB needs to know only about XRD redirector.
- Transfers still load-balanced.
- Single FTS instance, share XRD redir host.
- Direct XRD transfer to EOS, governed by FTS.
- Decoupled Online and Buffer specs.
- XRD_i nodes run metadata job after receive file.
- Interface specification = XRD URL namespace

Some pros/cons of each

- pro one less overall layer, file system hand-off maybe more familiar than XRD (?).
- con tight coupling in design and procurement, denser CPU requirement, O/O interface protocol more complex. Multiple FTS (not big problem), FTS→EOS mediated transfers.

DOOB:

- pro more distributed, leads to naturally more available CPU, decouples design and procurement, more fault-tolerant, O/O interface protocol simple. Direct XRD→EOS FTS-initiated transfers (EOS is native XRD).
- con requires extra layer, DAQ/EB needs XRD client lib (but not ROOT) or must use xrdcp unix command and thus its own local storage

Take Away

- We are looking to quantitatively understand the protoDUNE/SP online/offline needs.
- Likely major decision: 1Gbps vs. 10Gbps NICs
 - → want bottleneck at network or CPU/DISK?
- Two design options exist:
 - uoob concentrated complexity, tightly coupled, one less layer.
 - doob simpler, more distributed but one extra layer, requires XRD.

What are our external constraints? (how much \$\$\$?)